

Stochastic optimization

Chance constrained programming - SAA

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We consider the problem

$$\begin{aligned} \min_{x \in \mathcal{X}} f(x) \\ \text{s.t. } P[G(x, \xi) \leq 0] \geq 1 - \alpha, \end{aligned}$$

where $G : \mathcal{R}^n \times \Xi \rightarrow \mathcal{R}^m$.

This allows to consider separate and joint chanced constrained.

We can reformulate this problem as

$$\begin{aligned} \min_{x \in \mathcal{X}} f(x) \\ \text{s.t. } P[G(x, \xi) > 0] \leq \alpha, \end{aligned}$$

If $P[G(x, \xi) \leq 0]$ has no closed form, we can always approximate it by sample average approximation.

With the indicator function of $(0, \infty)$

$$\mathcal{I}_{(0,\infty)} = \begin{cases} 1 & \text{if } t > 0, \\ 0 & \text{if } t \leq 0, \end{cases}$$

we can rewrite the chance constraint as

$$p(x) = E [\mathcal{I}_{(0,\infty)} (G(x, \xi))] \leq \alpha.$$

Let ξ_1, \dots, ξ_N be an independent identically distributed (iid) sample of N realizations of ξ . We approximate the chance constraint by

$$\hat{p}_N = \frac{1}{N} \sum_{j=1}^N \mathcal{I}_{(0,\infty)} (G(x, \xi_j)) \leq \gamma$$

$\hat{p}_N(x)$ is equal to the proportion of times that $G(x, \xi_j) > 0$.
 γ is not necessarily equal to α , especially when we want to enlarge the feasible of the approximate problem.

The approximate problem is

$$\begin{aligned} \min_{x \in \mathcal{X}} & f(x) \\ \text{s.t. } & \hat{p}_N(x) \leq \gamma. \end{aligned}$$

For theoretical justifications, see

B.K. Pagnoncelli, S. Ahmed, and A. Shapiro, "Sample Average Approximation Method for Chance Constrained Programming: Theory and Applications", Journal of Optimization Theory and Applications 142 (2), pp. 399–416, 2009.

Especially, under some regularity conditions, we have

$$\sup_{x \in C} |\hat{p}_N(x) - p(x)| \rightarrow 0,$$

w.p.1, as $N \rightarrow \infty$.

Assume that there is an optimal solution \bar{x} of the true problem such that for any $\epsilon > 0$ there is $x \in \mathcal{X}$ with $\|x - \bar{x}\| \leq \epsilon$ and $p(x) < \alpha$.

Theorem

Suppose $\gamma = \alpha$, \mathcal{X} compact, and $f(x)$ continuous. Then $\hat{v}_N \rightarrow v^$, where \hat{v}_N and v^* are the optimal values of the approximate and true problems, respectively, and $D(\hat{S}_N, S) \rightarrow 0$ almost surely as $N \rightarrow \infty$.*

Application: call centers

Work in progress with Thuy Anh Ta and Pierre L'Ecuyer. The next slides are due Thuy Anh Ta.

A call center: an office used for receiving or transmitting customers requests by **telephone**.

- A **call** represents a customer and it is classified by **type** (K types).
- Customers are served by **agents**.
- Agents sharing a common set of tasks form a **group** of agents (I groups).
- A day is divided into **periods** (P periods).
- The work schedule of an agent is determined by her **shift** (Q shift configurations).

Performance measure

- **Performance measures (PM)** allow to assess the quality of service (QoS) and efficiency of a call center, e.g., Service level (SL) and average waiting time (AWT), abandonment ratio, occupancy ratio, etc.
- **Service level (SL)**: the percentage of calls that are answered in a defined waiting threshold.
- **Average waiting time (AWT)**: the average (or mean) time a customer waited to have a service.
- QoS can be defined in the long run (expectations), or over a given time period (random variable).

Staffing, Scheduling and Routing

- The objective is to minimize the operating cost of a call center under a set of constraints on certain PMs (SL or AWT).
- Staffing: deciding the number of agents $y_{i,p}$ of each group for each period in a day.
- Scheduling: determining how many agents $x_{i,q}$ to assign to a set of predefined shifts.
- Routing: determining which agents are allowed to handle the call, and how agents are chosen when several agents are free.

Modeling a call center

- The calls arrive according to arbitrary stochastic processes that could be non-stationary (e.g. homogeneous Poisson process), and perhaps doubly stochastic (e.g. Poisson gamma process, the arrival rate is a random variable of gamma distribution).
- Service times: exponential, log-normal distributions.
- Delay time, patience time, etc.
- When number of skills is low: Arrivals are served in a first come first serve (FCFS) order and/or a longest idle server first rules.
- Otherwise, **skills based routing** may be used to route calls to appropriate agents.

Staffing and scheduling with predictable arrival rates

$$g(y) = \frac{\mathbb{E}[\text{\# of served call that waited at most } \tau]}{\mathbb{E}[\text{\# Total \# of calls}]}$$

Scheduling problem

$$\min c^T x = \sum_{i=1}^I \sum_{q=1}^Q c_{i,q} x_{i,q}$$

subject to:

$$\begin{aligned} Ax &\geq y \\ g_{k,p}(y) &\geq l_{k,p} \quad \forall k, p, \\ g_p(y) &\geq l_p \quad \forall p, \\ g_k(y) &\geq l_k \quad \forall k, \\ g(y) &\geq l, \\ x, y &\geq 0 \text{ and integer.} \end{aligned} \tag{P1}$$

Staffing problem

$$\min_y c'^T y = \sum_{i=1}^I \sum_{p=1}^P c'_{i,p} y_{i,p}$$

subject to:

$$\begin{aligned} g_{k,p}(y) &\geq l_{k,p} \quad \forall k, p, \\ g_p(y) &\geq l_p \quad \forall p, \\ g_k(y) &\geq l_k \quad \forall k, \\ g(y) &\geq l, \\ y &\geq 0 \text{ and integer.} \end{aligned} \tag{P2}$$

SL $S(\tau, y)$ and AWT $W(y)$ are considered as random variables in a given period.

The chance-constrained scheduling problem for multiskill call centers is formulated as:

$$\min c^T x = \sum_{i=1}^I \sum_{q=1}^Q c_{i,q} x_{i,q}$$

subject to:

$$\begin{aligned} Ax &\geq y \\ \mathbb{P}[S_{k,p}(\tau_{k,p}, y) \geq s_{k,p}] &\geq r_{k,p} \quad \forall k, p \\ \mathbb{P}[W_{k,p}(y) \leq w_{k,p}] &\geq v_{k,p} \quad \forall k, p \\ x, y &\geq 0 \text{ and integer.} \end{aligned} \tag{P3}$$

Sample average approximate problem

$$\min c^T x = \sum_{i=1}^I \sum_{q=1}^Q c_{i,q} x_{i,q}$$

subject to:

$$\begin{aligned} Ax &\geq y \\ \frac{1}{N} \sum_{d=1}^N \mathbb{I}[\hat{S}_{N,k,p}^d(\tau_{k,p}, y) \geq s_{k,p}] &\geq r_{k,p} \quad \forall k, p \\ \frac{1}{N} \sum_{d=1}^N \mathbb{I}[\hat{W}_{N,k,p}^d(y) \leq w_{k,p}] &\geq v_{k,p} \quad \forall k, p \\ x, y &\geq 0 \text{ and integer.} \end{aligned} \tag{SAA}$$

Under some specific conditions, the optimal value and the set of optimal solutions of the SAA problem converge to those of the true problem with probability one as N approaches infinity.