

Nash Equilibria with Irradical Probabilities

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Abstract

We present for every $n \geq 4$ an n -player game in normal form with payoffs in $\{0, 1, 2\}$ that has a unique, fully mixed, Nash equilibrium in which all the probability weights are irradical (i.e., algebraic but not closed form expressible even with m -th roots for any integer m).

Keywords: Nash equilibrium, algebraic, irrational

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1. Introduction

Every Nash equilibrium (NE) in a 2-player normal form game with rational payoffs is a solution to a linear program with rational weights (?). Therefore, if σ is a NE then there is a NE τ with rational probability weights (henceforth we refer to them as *probabilities*) and support $\text{supp } \sigma$. However, this conclusion fails to hold in n -player games when $n \geq 3$. This is because more generally, NEs are solutions of polynomial equations and inequalities in n variables and degree at most $n - 1$, and these might not have solutions in the rationals. A concrete example is found in ?: a 3-player game with integer payoffs in $\{0, 1, 2, 3\}$, possessing a unique NE, and its probabilities are $((x_1, 1 - x_1), (x_2, 1 - x_2), (x_3, 1 - x_3))$ where $x_1 = \frac{25 - \sqrt{409}}{12}$, $x_2 = \frac{13 + \sqrt{409}}{60}$, $x_3 = \frac{21 - \sqrt{409}}{2}$. These 6 probabilities are all irrational. A different known example is the 3-player poker game in ?.

Such a result is possible for every $n \geq 4$ as well but there seems to be no previously known game with this property for any such n . In this paper we

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present for every $n \geq 4$ an n -player game G_n that has 2 actions per player and the following properties: (1) The payoffs are in $\{0, 1, 2\}$; (2) G_n has a unique NE; (3) that NE is fully mixed, and (4) the probability weights of this NE are all inexpressible with radicals (henceforth we call such numbers *irradicals*). Inexpressibility with radicals is an even stronger requirement than irrationality – it means that these probabilities do not have closed form expressions using rational arithmetic and m -th roots for any integer m .

To play a mixed NE, a player must randomly select an action from the underlying (discrete) probability distribution. Exact sampling from discrete distributions with rational probabilities is a well understood problem solved with efficient algorithms (??). There, the computations are done using integer arithmetic which be done precisely on a modern digital computer. But sampling from distributions with irrational (or irradical) probabilities using rational arithmetic is a less studied problem. Irrational numbers can only be approximated by finite integers. With current known techniques, exact sampling from a distribution with irrational weights can only be done with arbitrary precision arithmetic using a power series expansion of the probabilities (??), which is computationally and pragmatically more involved than sampling from distributions with rational probabilities. In our context, all the probability weights are algebraic numbers, so the sampling problem appears to require computing roots of polynomials with potentially very large degrees and coefficients (the polynomials associated with G_5 have degree 26 and 10-digit coefficients). This fact highlights a challenge that players face in this context.

For example, our 4-player game has a unique NE, $((x_1, 1 - x_1), (x_2, 1 - x_2), (x_3, 1 - x_3), (x_4, 1 - x_4))$ such that x_1, x_2, x_3, x_4 are these roots of the polynomials

$$\begin{aligned} P_1(z) &= 7z^6 - 42z^5 + 89z^4 - 83z^3 + 40z^2 - 10z + 1, \\ P_2(z) &= 4z^6 - 27z^5 + 70z^4 - 79z^3 + 45z^2 - 13z + 1, \\ P_3(z) &= 140z^6 - 511z^5 + 701z^4 - 454z^3 + 141z^2 - 19z + 1, \\ P_4(z) &= 5z^6 - 44z^5 + 143z^4 - 163z^3 + 85z^2 - 21z + 2, \end{aligned}$$

that are approximately $x_1 \approx 0.529$, $x_2 \approx 0.846$, $x_3 \approx 0.523$, $x_4 \approx 0.320$. x_1, x_2, x_3, x_4 cannot be expressed with radicals because P_1, P_2, P_3, P_4 are all irreducible polynomials over the rationals whose Galois groups are all S_6 which is not solvable (see ? for reference).

To obtain the main result we first construct the mentioned 4-player game

(Proposition ??) and a 5-player game (Proposition ??) with a unique NE whose probabilities are irrational. We then use these two games to construct such a game for any $n \geq 6$ (Proposition ??). More specifically, for $n \geq 6$ we juxtapose independent copies of the 4-player and 5-player games, and to one of the game copies possibly attach 2 additional players who force each other to mimic the mixed strategy of one player in the game copy.

We complement this result by showing that every $2 \times 2 \times 2$ game with integer payoffs has a NE whose probabilities can be expressed using only rationals and square roots of rationals (Proposition ??). We also present a simple 3-player game with integer payoffs in only $\{0, 1, 2\}$ whose unique NE has only irrational probabilities (Proposition ??). This game is simpler than that of ?. Its unique NE $((x_1, 1 - x_1), (x_2, 1 - x_2), (x_3, 1 - x_3))$ where $x_1 = \frac{7-\sqrt{13}}{6}$, $x_2 = \frac{7+\sqrt{13}}{18}$, $x_3 = \frac{-3+\sqrt{13}}{2}$. These are approximately $x_1 \approx 0.566$, $x_2 \approx 0.589$, $x_3 \approx 0.303$.

Note that given such a 3-player game it is easy to extend it to an n -player game where *some* of the players need to play mixed actions with irrational probabilities in every NE – simply add players with one dominating action each. But then only the first 3 players will need to play NE strategies with irrational probabilities. We are interested in forcing *all* players to play NE strategies with irrational probabilities. For this reason we are interested only in fully mixed equilibria.

It is also possible to use such a 3-player game to generate n -player games for every $n \geq 5$ where all players need to play NE strategies with irrational probabilities. This is achieved by applying our results in Section ?? to such a 3-player game. However, this approach alone leaves unresolved the case of 4 players, and all the probabilities in the other cases will not be irrational.

There are some related results in the literature. First is the characterization of ? of (mixed) strategy profiles which can be the unique NEs of a suitable normal form game. In particular, any fully mixed strategy profile for n players with 2 actions each is the unique NE of some n -player game (with 2 actions per player). However, the payoffs in the constructed game will involve the probabilities in that strategy profile. In particular, if the probabilities are irrational then the payoffs found will be irrational. Here we are looking for games with integer payoffs, since these are simpler to store on a computer than irrational payoffs.

Another result is by ? who constructs for every algebraic number α a 3-player game with integer payoffs such that α is the payoff to some player in

the NE of the game. However, in the constructed game player 3 has the fully mixed uniform distribution (which is comprised of rational probabilities) as an equilibrium strategy.

Another related result is the universality of NEs by ?. This result asserts that every semialgebraic set is isomorphic to the set of NEs of some game. However, an isomorphism in this context only preserves the topological structure of the semialgebraic set rather than its exact members. In particular, in ?, a semialgebraic set consisting of a single point with irrational entries will be mapped to a NE profile with rational probabilities.

Paper organization. In Section ?? we define the setup and provide algebra background. In Section ?? we present our results for 3-player, $2 \times 2 \times 2$ games. In Section ?? and Section ?? we present the 4-player game and 5-player game. In Section ?? we use the latter two games to construct the remaining n -player games for every $n \geq 6$.

2. Setup and lemmas

The games we define are in the usual normal form. They will be presented as a 2-column table of rows in the format

$$[\text{action tuple} \mid \text{payoffs}].$$

The action tuple and the payoffs in each row are ordered by the players $1, \dots, n$.

In the games we present each player has 2 actions called 0 and 1. Each mixed action profile is defined by $x = (x_1, \dots, x_n)$ where $x_i \in [0, 1]$ is the probability player i gives to playing action 0. The payoff function of player i is u_i . We then define for every i the polynomial

$$f_i(x_{-i}) = u_i(x_1, \dots, x_{i-1}, 1, x_{i+1}, \dots, x_n) - u_i(x_1, \dots, x_{i-1}, 0, x_{i+1}, \dots, x_n). \quad (1)$$

This is the expected payoff advantage action 0 has over action 1 for player i given the mixed actions x_{-i} of the other players. We sometimes write $f_i(x)$ or f_i for simplicity.

Given a game G let $\text{NE}(G)$ be the set of NEs of G and $\text{PNE}(G)$ the set of pure NEs of G . For reference we state the following lemma which can be proved by definition.

Lemma 2.1. *If $x \in \text{NE}(G)$ then for every i ,*

$$\begin{aligned} f_i(x_{-i}) > 0 &\Rightarrow x_i = 1, \\ f_i(x_{-i}) < 0 &\Rightarrow x_i = 0, \\ 0 < x_i < 1 &\Rightarrow f_i(x_{-i}) = 0. \end{aligned} \tag{2}$$

2.1. Algebra background

Here we cover the necessary algebra background for this paper (see ? for a detailed reference). Let $\mathbb{Q}[x]$ be the ring (set) of polynomials in the variable x with rational coefficients. $P \in \mathbb{Q}[x]$ is irreducible if it cannot be written as $P_1 P_2$ for $P_1, P_2 \in \mathbb{Q}[x]$ where $0 < \deg P_1, \deg P_2 < \deg P$. A number a is called *algebraic* if there is a $0 \neq P \in \mathbb{Q}[x]$ such that $P(a) = 0$. Every algebraic number has a *minimal polynomial*: a unique, monic polynomial of minimum degree $P_a \in \mathbb{Q}[x]$ such that $P_a(a) = 0$. For example, $P_{\sqrt{2}} = x^2 - 2$. Here we will not insist on monicity in our search of these polynomials.

One way to prove a number a is irrational is by showing that $\deg P_a \geq 2$. To verify that a polynomial at hand P is P_a , it is sufficient to show that $P(a) = 0$ and P is irreducible in $\mathbb{Q}[x]$ (and that P is monic).

An irrational algebraic number can often be represented in a compact way, using *radicals* – m -roots ($\sqrt[m]{}$ for an integer m) of rationals. For example $\sqrt{2}$ represents the positive root of $x^2 - 2$. Another example is $\sqrt{2} + 3\sqrt[3]{5}$. Given a polynomial with a root a , we want to know when it is possible to construct such a representation for a . This will be relevant in Section ?? and Section ??.

Let P be an irreducible polynomial with degree n and roots a_1, \dots, a_n . The field $F = \mathbb{Q}[a_1, \dots, a_n]$ is defined to be the smallest field that contains \mathbb{Q} and a_1, \dots, a_n .¹ For example, $\mathbb{Q}[\sqrt{2}]$ is the set of all numbers $a + b\sqrt{2}$ where $a, b \in \mathbb{Q}$. An *automorphism of F that extends \mathbb{Q}* is a function $f : F \rightarrow F$ that is a bijection, preserves the field operations (e.g., $f(a+b) = f(a) + f(b)$) and maps every rational to itself. It can be shown that f is uniquely defined by its action on the roots a_1, \dots, a_n . Furthermore, $f(\{a_1, \dots, a_n\}) = \{a_1, \dots, a_n\}$, since for every i it holds that $P(f(a_i)) = f(P(a_i)) = f(0) = 0$, meaning $f(a_i)$ is a root of P . The set $\text{Aut}(F)$ is called the *Galois group* of P .

¹ $\mathbb{Q}[a_1, \dots, a_n]$ is a field while $\mathbb{Q}[x]$ and $\mathbb{Q}[x_1, \dots, x_n]$ are rings because x_1, \dots, x_n are symbolic variables with no defined inverses, while a_1, \dots, a_n are real numbers that do have their inverses $\frac{1}{a_1}, \dots, \frac{1}{a_n}$ in $\mathbb{Q}[a_1, \dots, a_n]$.

A fundamental result in algebra states that a number a can be represented with radicals if and only if the Galois group of its minimal polynomial P_a is *solvable*. The formal definition of solvability does not matter here. The point is that to determine if a can be written with radicals we only need to determine if the Galois group of P_a is solvable (and wlog we can ignore monicity). A particular type of group that is *unsolvable* is S_n for $n \geq 5$. S_n is the set of all $n!$ permutations of $\{1, \dots, n\}$ with the group operation of permutation composition.

Take two examples. The first is $\sqrt{2}$. Its minimal polynomial is $x^2 - 2$. Its Galois group is called \mathbb{Z}_2 , which is simply the set $\{0, 1\}$ equipped with the XOR operation. This group is solvable, and indeed $\sqrt{2}$ is a radical representation of the positive root of $x^2 - 2$. The second example is any root of $P = x^5 - x - 1$. It can be shown that the Galois group of P is S_5 . It is unsolvable, and indeed the roots of P cannot be represented with radicals. We call such numbers *irradicals*.

Note that the pair $(x_i, 1 - x_i)$ satisfies $x_i \in \mathbb{Q} \iff 1 - x_i \in \mathbb{Q}$, and similarly for expressibility with radicals. Therefore in our proofs of irrationality and irrationality we only analyze x_1, \dots, x_n .

3. Three players

Define the $2 \times 2 \times 2$ game G_3 :

$$\left[\begin{array}{ccc|ccc} 0 & 0 & 0 & 0 & 0 & 2 \\ 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 & 1 \end{array} \right]$$

Figure 1: The game G_3

Proposition 3.1. G_3 has a unique NE $((x_1, 1 - x_1), (x_2, 1 - x_2), (x_3, 1 - x_3))$

given by

$$x_1 = \frac{7 - \sqrt{13}}{6}, \quad x_2 = \frac{7 + \sqrt{13}}{18}, \quad x_3 = \frac{-3 + \sqrt{13}}{2}.$$

The probabilities in the NE are all irrational.

Proof. First observe that $\text{PNE}(G_3) = \emptyset$ using this table of profitable deviations:

action profile	unsatisfied players
(0, 0, 0)	2
(0, 0, 1)	3
(0, 1, 0)	1, 3
(0, 1, 1)	1
(1, 0, 0)	2, 3
(1, 0, 1)	1
(1, 1, 0)	3
(1, 1, 1)	2

Now the polynomials f_1, f_2, f_3 from Equation ?? are

$$\begin{aligned} f_1(x_2, x_3) &= -x_2x_3 + 2x_2 - 1, \quad f_2(x_1, x_3) = x_1x_3 - x_1 - 2x_3 + 1, \\ f_3(x_1, x_2) &= 3x_1x_2 - 1. \end{aligned}$$

We show using Lemma ?? that if $x_i \in \{0, 1\}$ for some i then $x \notin \text{NE}(G_3)$. There are 6 cases.

Case 1: $x_1 = 0$. then $f_2 = -2x_3 + 1$, $f_3 = -1 < 0$, so $x_3 = 0$ so $f_2 = 1 > 0$ so $x_2 = 1$. Therefore $x = (0, 1, 0) \notin \text{NE}(G_3)$.

Case 2: $x_1 = 1$. then $f_2 = -x_3$, $f_3 = 3x_2 - 1$. If $x_3 > 0$ then $f_2 < 0$ so $x_2 = 0$ so $f_3 = -1 < 0$ so $x_3 = 0$ – a contradiction. Therefore $x_3 = 0$, so $f_3 \leq 0$ so $x_2 \leq \frac{1}{3}$ so $f_1 \leq -\frac{1}{3} < 0$ so $x_1 = 0$ – a contradiction.

Case 3: $x_2 = 0$. then $f_1 = -1$ so $x_1 = 0$ so by case 1 $x \notin \text{NE}(G_3)$.

Case 4: $x_2 = 1$. then $f_1 = -x_3 + 1$, $f_3 = 3x_1 - 1$. By cases 1 and 2, $x_1 \in (0, 1)$ so $f_1 = 0$ so $x_3 = 1$ so $f_2 = -1 < 0$ so $x_2 = 0$ – a contradiction. Therefore $x \notin \text{NE}(G_3)$.

Case 5: $x_3 = 0$. then $f_2 = -x_1 + 1$. By cases 3 and 4, $x_2 \in (0, 1)$ so $f_2 = 0$ so $-x_1 + 1 = 0$ so $x_1 = 1$ so by case 2 we conclude that $x \notin \text{NE}(G_3)$.

Case 6: $x_3 = 1$. then $f_1 = x_2 - 1$. Again $0 = f_1 = x_2 - 1$ so $x_2 = 1$, so by case 4 we conclude that $x \notin \text{NE}(G_3)$.

Therefore the only NEs are fully mixed. This means that $x \in \text{NE}(G_3)$ satisfies $f_i(x) = 0$ for every i . By f_3 we have $x_1 = \frac{1}{3x_2}$. Then by f_2 we have $\frac{x_3}{3x_2} - \frac{1}{3x_2} - 2x_3 + 1 = 0$ so $(6x_3 - 3)x_2 = x_3 - 1$. $x_3 \neq \frac{1}{2}$ because $x_3 \neq 1$ so $x_2 = \frac{x_3 - 1}{6x_3 - 3}$. Then by f_1 we have $\frac{(x_3 - 1)(2 - x_3)}{6x_3 - 3} = 1$ so $-x_3^2 + 3x_3 - 2 = 6x_3 - 3$ so $x_3^2 + 3x_3 - 1 = 0$ so because $x_3 > 0$ we get $x_3 = \frac{-3 + \sqrt{13}}{2}$, $x_2 = \frac{-5 + \sqrt{13}}{-24 + 6\sqrt{13}} = \frac{7 + \sqrt{13}}{18}$ and $x_1 = \frac{6}{7 + \sqrt{13}} = \frac{7 - \sqrt{13}}{6}$.

$x_1, x_2, x_3 \notin \mathbb{Q}$ because they are expressions involving only rationals and $\sqrt{13}$ which is irrational. Their minimal polynomials in $\mathbb{Q}[y]$ are

$$P_1(y) = 3y^2 - 7y + 3, \quad P_2(y) = 9y^2 - 7y + 1, \quad P_3(y) = y^2 + 3y - 1.$$

These are all indeed irreducible by the rational root theorem. \square

We conclude with a proof that square roots of rationals (along with rational arithmetic) are always enough to represent a NE of a $2 \times 2 \times 2$ game.

Proposition 3.2. *Let G be a $2 \times 2 \times 2$ game with integer payoffs. There is a NE $((x_1, 1 - x_1), (x_2, 1 - x_2), (x_3, 1 - x_3))$ such that for every $i \in \{1, 2, 3\}$, $x_i = a_i + b_i\sqrt{c_i}$ for some $a_i, b_i, c_i \in \mathbb{Q}$.*

Proof. In full generality, the polynomials f_1, f_2, f_3 have the form

$$\begin{aligned} f_1(x_2, x_3) &= A_1x_2x_3 + B_1x_2 + C_1x_3 + D_1, \\ f_2(x_1, x_3) &= A_2x_1x_3 + B_2x_1 + C_2x_3 + D_2, \\ f_3(x_1, x_2) &= A_3x_1x_2 + B_3x_1 + C_3x_2 + D_3, \end{aligned}$$

where the coefficients A_i, B_i, C_i, D_i are all integers, as the payoffs in the game are integers.

Let $x \in \text{NE}$. There are 3 cases up to permuting actions and payoffs:

Case 1: $x_1 = 0$ and $x_2, x_3 \in (0, 1)$. Then $f_1 \leq 0$, $f_2 = C_2x_3 + D_2 = 0$, $f_3 = C_3x_2 + D_3 = 0$.

From $f_1 \leq 0$ we get $(A_1x_3 + B_1)x_2 \leq -C_1x_3 - D_1$. If $A_1x_3 + B_1 = 0$ then either $x_3 = -\frac{B_1}{A_1} \in \mathbb{Q}$ or $A_1 = B_1 = 0$, in which case x_2, x_3 are defined by linear equations and inequalities with rational coefficients, so x_2, x_3 can be rational in a NE.

Otherwise, $x_2 \leq -\frac{C_1x_3 + D_1}{A_1x_3 + B_1}$.

If $C_i \neq 0$ for $i \in \{2, 3\}$ then $x_{5-i} = -\frac{D_i}{C_i} \in \mathbb{Q}$. If $i = 2$ then $f_1 = \left(-\frac{A_1D_2}{C_2} + B_1\right)x_2 - \frac{C_1D_2}{C_2} + D_1 \leq 0$. Since this inequality holds for some value of x_2 , the inequality also holds for some rational value of x_2 . In total, x_1, x_2, x_3 can all be rational (in a NE).

Otherwise $C_2 = C_3 = 0$. Then the only constraints on x_2, x_3 are being in $(0, 1)$ and the inequality $x_2 \leq -\frac{C_1x_3 + D_1}{A_1x_3 + B_1}$. This inequality is in particular satisfied for some rational values of x_2, x_3 , meaning that x_1, x_2, x_3 can all be rational.

Case 2: $x_1 = x_2 = 0$ and $x_3 \in (0, 1)$. Then $f_1 = C_1x_3 + D_1 \leq 0$, $f_2 = C_2x_3 + D_2 \leq 0$, $f_3 = D_3 = 0$. So x_3 can be rational in a NE with $x_1 = x_2 = 0$.

Case 3: $x_1, x_2, x_3 \in (0, 1)$. Then $f_1 = f_2 = f_3 = 0$. Then $(A_1x_3 + B_1)x_2 = -C_1x_3 - D_1$. Suppose that (*) $A_1x_3 + B_1 \neq 0$. Then $x_2 = -\frac{C_1x_3 + D_1}{A_1x_3 + B_1}$. From $f_3 = 0$ we then get $\left(-\frac{A_3(C_1x_3 + D_1)}{A_1x_3 + B_1} + B_3\right)x_1 - \frac{C_3(C_1x_3 + D_1)}{A_1x_3 + B_1} + D_3 = 0$. So

$$((B_3A_1 - A_3C_1)x_3 + B_3B_1 - A_3D_1)x_1 = (C_3C_1 - D_3A_1)x_3 + C_3D_1 - D_3B_1.$$

Suppose that (**) $(B_3A_1 - A_3C_1)x_3 + B_3B_1 - A_3D_1 \neq 0$. Then

$$x_1 = \frac{(C_3C_1 - D_3A_1)x_3 + C_3D_1 - D_3B_1}{(B_3A_1 - A_3C_1)x_3 + B_3B_1 - A_3D_1}.$$

Now $f_2 = 0$ simplifies into a quadratic equation in x_3 involving rational coefficients. So x_3 can be written with rational arithmetic and (non-nested) square roots. The same then holds for x_1, x_2 by their expressions as ratios of linear functions of x_3 .

Now suppose that (*) and (**) are not both true. There are 2 cases:

- (*) is false, so $A_1x_3 + B_1 = 0$. Then $C_1x_3 + D_1 = 0$. If $A_1 = C_1 = 0$ then $B_1 = D_1 = 0$ and f_1 is always 0. If f_2 is always 0 then x_3 can

take any value in $(0, 1)$ (particularly rational ones), and one can verify that x_1, x_2 can take rational values. So f_2, f_3 are not always 0, and because of symmetry we can swap f_1 and f_2 to have one of $A_1 \neq 0$ and $C_1 \neq 0$ being true which implies that $x_3 \in \mathbb{Q}$. It can then be verified that x_1, x_2 can both be rational.

- f_1, f_3 are not always 0 and $A_1x_3 + B_1 \neq 0$, $(*)$ is true and $(**)$ is false, so $(B_3A_1 - A_3C_1)x_3 + B_3B_1 - A_3D_1 = 0$. Then $(C_3C_1 - D_3A_1)x_3 + C_3D_1 - D_3B_1 = 0$. If $B_3A_1 - A_3C_1 = C_3C_1 - D_3A_1 = 0$ then $C_3D_1 - D_3B_1 = B_3B_1 - A_3D_1 = 0$ and conditioned on $f_1(x_2, x_3) = 0$, we get that f_3 is always 0 in contradiction to our assumption.

So one of $B_3A_1 - A_3C_1$ and $C_3C_1 - D_3A_1$ is not 0. Then $x_3 \in \mathbb{Q}$. It can then be verified that x_1, x_2 can both be rational. \square

4. Four players

Define the $2 \times 2 \times 2 \times 2$ game G_4 :

$$\left[\begin{array}{cccc|cccc} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 2 \\ 0 & 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 \end{array} \right]$$

Figure 2: The game G_4

Proposition 4.1. G_4 has a unique NE $((x_1, 1 - x_1), (x_2, 1 - x_2), (x_3, 1 - x_3), (x_4, 1 - x_4))$ given by

$$\begin{aligned} x_1 &\approx 0.529270752820, \quad x_2 \approx 0.846414728986, \\ x_3 &\approx 0.523440476515, \quad x_4 \approx 0.320065197645. \end{aligned}$$

The probabilities in the NE are all irrational and irradical.

Proof. First observe that $\text{PNE}(G_4) = \emptyset$ using this table of profitable deviations:

action profile	unsatisfied players
(0, 0, 0, 0)	3
(0, 0, 0, 1)	2
(0, 0, 1, 0)	1
(0, 0, 1, 1)	1
(0, 1, 0, 0)	1, 3
(0, 1, 0, 1)	4
(0, 1, 1, 0)	4
(0, 1, 1, 1)	2
(1, 0, 0, 0)	2, 3, 4
(1, 0, 0, 1)	1
(1, 0, 1, 0)	2
(1, 0, 1, 1)	3
(1, 1, 0, 0)	3, 4
(1, 1, 0, 1)	1, 2
(1, 1, 1, 0)	1, 4
(1, 1, 1, 1)	1, 3

Now the polynomials f_1, f_2, f_3, f_4 from Equation ?? are

$$\begin{aligned} f_1(x_2, x_3, x_4) &= x_2x_3x_4 + 2x_2x_3 - 2x_2 - 2x_3x_4 + 1, \\ f_2(x_1, x_3, x_4) &= 3x_1x_3x_4 - 3x_1x_3 + x_1 - x_3x_4 + x_3 - x_4, \\ f_3(x_1, x_2, x_4) &= x_1x_4 - x_1 - 2x_4 + 1, \\ f_4(x_1, x_2, x_3) &= -x_1x_2x_3 + 3x_1x_3 - x_2x_3 + x_2 - 1. \end{aligned}$$

We show using Lemma ?? that if $x_i \in \{0, 1\}$ for some i then $x \notin \text{NE}(G_4)$. There are 8 cases.

Case 1: $x_1 = 0$. then $f_2 = -x_3x_4 + x_3 - x_4$, $f_3 = -2x_4 + 1$, $f_4 = -x_2x_3 + x_2 - 1$. If $x_4 = \frac{1}{2}$ then $f_2 = \frac{x_3-1}{2}$. If $x_3 = 1$ then $f_4 = -1 < 0$ so $x_4 = 0$ – a contradiction. Therefore $x_3 < 1$ so $f_2 < 0$ so $x_2 = 0$. Since $x_4 = \frac{1}{2}$ we get $0 = f_4 = -1$ – a contradiction.

Therefore $x_4 \neq \frac{1}{2}$. If $x_4 < \frac{1}{2}$ then $f_3 > 0$ so $x_3 = 1$ so $f_4 = -1$ so $x_4 = 0$ so $f_2 = 1$ so $x_2 = 1$. Therefore $x = (0, 1, 1, 0) \notin \text{PNE}(G_4)$. If $x_4 > \frac{1}{2}$ then $f_3 < 0$ so $x_3 = 0$ so $f_2 = -x_4 < 0$ so $x_2 = 0$ so $f_4 = -1$ so $x_4 = 0$ – a contradiction.

Overall x_4 cannot take any value in $[0, 1]$, so $x \notin \text{NE}(G_4)$.

Case 2: $x_1 = 1$. then $f_2 = 2x_3x_4 - 2x_3 - x_4 + 1$, $f_3 = -x_4$, $f_4 = -2x_2x_3 + x_2 + 3x_3 - 1$. If $x_4 > 0$ then $f_3 < 0$ so $x_3 = 0$ so $f_2 = -x_4 + 1$ and $f_4 = x_2 - 1$. If $x_4 < 1$ then $f_2 > 0$ so $x_2 = 1$ so $f_1 = -1$ so $x_1 = 0$ – a contradiction. Therefore, still assuming that $x_4 > 0$ we get $x_4 = 1$ so $f_4 = x_2 - 1 \geq 0$ so $x_2 = 1$ so $f_1 = -1$ so $x_1 = 0$ – a contradiction.

Therefore $x_4 = 0$. So $f_2 = -2x_3 + 1$. If $x_3 = \frac{1}{2}$ then $f_4 = \frac{1}{2}$ so $x_4 = 1$ – a contradiction. Therefore $x_3 \neq \frac{1}{2}$. If $x_3 > \frac{1}{2}$ then $f_2 < 0$ so $x_2 = 0$ so $f_4 = 3x_3 - 1 > 0$ so $x_4 = 1$ – a contradiction. Therefore $x_3 < \frac{1}{2}$. Then $f_2 > 0$ so $x_2 = 1$ so $f_1 = 2x_3 - 1 < 0$ so $x_1 = 0$ – a contradiction.

Overall $x_1 \neq 1$.

Case 3: $x_2 = 0$. then $f_1 = -2x_3x_4 + 1$, $f_4 = 3x_1x_3 - 1$. By cases 1 and 2, $x_1 \in (0, 1)$ so $f_1 = 0$ so $x_3x_4 = \frac{1}{2}$ so $x_3, x_4 \neq 0$ so $f_3, f_4 \geq 0$. Therefore $x_3 = \frac{1}{2x_4}$ so $f_4 = \frac{3x_1}{2x_4} - 1 \geq 0$ so $x_1 \geq \frac{2x_4}{3}$. Also, $f_3 = x_1x_4 - x_1 - 2x_4 + 1 \geq 0$ so $(1-x_4)x_1 \leq 1-2x_4$. If $x_4 = 1$ then $f_3 = -1 < 0$ – a contradiction. So $x_4 < 1$, so $\frac{2}{3}x_4 \leq x_1 \leq \frac{1-2x_4}{1-x_4}$. The quadratic inequality is then $2x_4^2 - 8x_4 + 3 \geq 0$, and in particular $x_4 < \frac{1}{2}$ so $x_3 > 1$ – a contradiction to $x_3 \leq 1$.

Overall x_4 cannot take on any value, concluding this case.

Case 4: $x_2 = 1$. then $f_1 = -x_3x_4 + 2x_3 - 1$, $f_4 = 2x_1x_3 - x_3$. Again $f_1 = 0$ so $x_3 = \frac{1}{2-x_4}$. If $x_4 = 1$ then $x_3 = 1$ and $f_3 = -1$ so $x_3 = 0$ – a contradiction. Therefore $x_4 < 1$ so $x_3 < 1$ so $f_3, f_4 \leq 0$. Also $x_3 > 0$ and by $f_4 \leq 0$ we get $x_1 \leq \frac{1}{2}$. By $f_3 \leq 0$ we get $x_1 \geq \frac{1-2x_4}{1-x_4}$, so in total $1-x_4 \geq 2-4x_4$ so $x_4 \geq \frac{1}{3}$. Therefore $x_3 \geq \frac{3}{5}$. This means that $f_3 = f_4 = 0$, so $x_1 = \frac{1}{2}$ so $f_3 = -\frac{3}{2}x_4 + \frac{1}{2} = 0$ so $x_4 = \frac{1}{3}$ so $x_3 = \frac{3}{5}$. Therefore $f_2 = -\frac{1}{30}$ so $x_2 = 0$ – a contradiction to $x_2 = 1$. This establishes case 4.

Case 5: $x_3 = 0$. then $f_1 = -2x_2 + 1$, $f_2 = x_1 - x_4$, $f_4 = x_2 - 1$. By the previous cases $x_1, x_2 \in (0, 1)$ so $f_1 = f_2 = 0$ so $x_2 = \frac{1}{2}$ and $x_1 = x_4$. We also get $f_4 < 0$ so $x_4 = 0$ so $x_1 = 0$ – a contradiction.

Case 6: $x_3 = 1$. then $f_1 = x_2x_4 - 2x_4 + 1$, $f_2 = 3x_1x_4 - 2x_1 - 2x_4 + 1$, $f_4 = -x_1x_2 + 3x_1 - 1$. Again $f_1 = f_2 = 0$ so $x_4 = \frac{1}{2-x_2}$ so $f_2 = \frac{3x_1}{2-x_2} - 2x_1 - \frac{2}{2-x_2} + 1 = 0$ so $\frac{2x_2-1}{2-x_2}x_1 = \frac{x_2}{2-x_2}$ so $(2x_1 - 1)x_2 = 1$ so $x_2 > 1$ because $x_1 \in (0, 1)$ – a contradiction to $x_2 < 1$.

Case 7: $x_4 = 0$. then $f_3 = -x_1 + 1$. By the previous cases, $x_1, x_2, x_3 \in (0, 1)$ so $f_3 = 0$ so $x_1 = 1$ – a contradiction.

Case 8: $x_4 = 1$. then $f_2 = x_1 - 1$. Like in case 7 we get $x_1 \in (0, 1)$ and $x_1 = 1$ – a contradiction.

Therefore the only NEs are fully mixed. This means that $x \in \text{NE}(G_4)$ satisfies $f_i(x) = 0$ for every i . We solve this now.

Let $F = \{f_1, f_2, f_3, f_4\}$. It is a set of polynomials in the polynomial ring $\mathbb{Q}[x_1, x_2, x_3, x_4]$. We begin by finding the Gröbner basis $G = \{g_1, g_2, g_3, g_4\}$ of the ideal $\langle F \rangle$, with the monomial ordering generated lexicographically by $x_1 \succ x_2 \succ x_3 \succ x_4$, given by

$$\begin{aligned} g_1 &= 5x_4^6 - 44x_4^5 + 143x_4^4 - 163x_4^3 + 85x_4^2 - 21x_4 + 2, \\ g_2 &= 28x_3 - 1465x_4^5 + 12302x_4^4 - 36947x_4^3 + 32897x_4^2 - 11699x_4 + 1447, \\ g_3 &= 4x_2 + 255x_4^5 - 2094x_4^4 + 6053x_4^3 - 4687x_4^2 + 1401x_4 - 149, \\ g_4 &= 7x_1 + 5x_4^5 - 39x_4^4 + 104x_4^3 - 59x_4^2 + 26x_4 - 9. \end{aligned}$$

These were found using Mathematica (?).

Now we compute the solutions of $\forall i \ g_i = 0$ to find those of $\forall i \ f_i = 0$. First, we approximate the 6 roots of g_1 numerically:

$$\begin{aligned} r_1 &\approx 0.3200651976, \ r_2 \approx 0.4231894254, \\ r_3 &\approx 0.4135410939 - 0.09991760306i, \ r_4 \approx 0.4135410939 + 0.09991760306i, \\ r_5 &\approx 3.614831595 - 1.802444362i, \ r_6 \approx 3.614831595 + 1.802444362i. \end{aligned}$$

$x_4 \in (0, 1)$ so $x_4 \in \{r_1, r_2\}$. If $x_4 = r_2$ then for every $r'_2 \in [0.42, 0.43]$, $g_2(r'_2, x_3) = 0$ iff $x_3 > 1$ – a contradiction to $x_3 < 1$. Therefore $x_4 = r_1$.²

²See ?? for how to approximate the roots of g_1 using exact arithmetic without relying on the problematic floating point arithmetic.

Using the equalities $g_2 = g_3 = g_4 = 0$ we arrive at the approximations

$$\begin{aligned}x_1 &\approx 0.529270752820, \quad x_2 \approx 0.846414728986, \\x_3 &\approx 0.523440476515, \quad x_4 \approx 0.320065197645.\end{aligned}$$

We now show that x_1, x_2, x_3, x_4 are irrational. First, for each of the 3 monomial orderings generated lexicographically by $x_2 \succ x_3 \succ x_4 \succ x_1$, $x_3 \succ x_4 \succ x_1 \succ x_2$ and $x_4 \succ x_1 \succ x_2 \succ x_3$, we find corresponding Gröbner bases, in addition to the one that we found initially. From them we collect the univariate polynomials

$$\begin{aligned}P_1(y) &= 7y^6 - 42y^5 + 89y^4 - 83y^3 + 40y^2 - 10y + 1, \\P_2(y) &= 4y^6 - 27y^5 + 70y^4 - 79y^3 + 45y^2 - 13y + 1, \\P_3(y) &= 140y^6 - 511y^5 + 701y^4 - 454y^3 + 141y^2 - 19y + 1, \\P_4(y) &= 5y^6 - 44y^5 + 143y^4 - 163y^3 + 85y^2 - 21y + 2.\end{aligned}$$

which we further know satisfy $P_i(x_i) = 0$. These exist by the elimination theorem for Gröbner bases (?).

Now, we use Murty's criterion (?) to show that P_1, P_2, P_3, P_4 are all irreducible over \mathbb{Q} .

1. P_1 is irreducible because $P_1(18) = 167595301$ is prime and the H value of P_1 is $\frac{89}{7} \approx 12.714$.
2. P_2 is irreducible because $P_2(28) = 1504207909$ is prime and the H value of P_2 is $\frac{79}{4} = 19.75$.
3. P_3 is irreducible because $P_3(11) = 175397399$ is prime and the H value of P_3 is $\frac{701}{140} \approx 5.007$.
4. P_4 is irreducible because $P_4(39) = 13945135583$ is prime and the H value of P_4 is $\frac{163}{5} = 32.6$.

As a result, P_i is the minimal polynomial of x_i (up to normalizing the leading coefficient of P_i), so x_i is irrational for every i . Now, it can be verified that the Galois group of P_i is S_6 for every $i \in \{1, 2, 3, 4\}$ (we used Magma (?)). Therefore x_1, x_2, x_3, x_4 are all irrational. This concludes the proof. \square

5. Five players

Define the $2 \times 2 \times 2 \times 2 \times 2$ game G_5 :

0	0	0	0	0	1	1	0	0	0	1
0	0	0	0	0	1	1	1	0	0	0
0	0	0	1	0	0	1	0	0	1	0
0	0	0	1	1	0	0	0	1	0	0
0	0	1	0	0	0	1	0	0	0	2
0	0	1	0	1	0	0	1	1	0	0
0	0	1	1	0	0	1	1	0	1	0
0	0	1	1	1	0	0	1	0	0	1
0	1	0	0	0	0	1	0	0	0	0
0	1	0	0	1	0	0	0	1	1	1
0	1	0	1	0	0	1	1	0	1	1
0	1	0	1	1	0	0	1	1	0	1
0	1	1	0	0	0	1	1	1	0	0
0	1	1	0	1	0	1	0	1	0	1
0	1	1	1	0	0	1	0	2	1	1
0	1	1	1	1	0	0	0	1	0	1
1	0	0	0	0	0	0	0	0	0	1
1	0	0	0	1	0	0	1	0	1	0
1	0	0	1	0	0	0	0	0	1	0
1	0	0	1	1	0	1	0	1	0	0
1	0	1	0	0	0	0	0	0	0	2
1	0	1	0	1	0	1	1	1	0	0
1	0	1	1	0	0	1	1	0	1	0
1	0	1	1	1	0	0	1	0	0	1
1	1	0	0	0	0	0	0	0	0	0
1	1	0	0	1	0	1	0	1	1	1
1	1	0	1	0	0	0	1	0	1	1
1	1	0	1	1	0	1	1	1	0	1
1	1	1	0	0	0	0	1	1	0	0
1	1	1	0	1	0	1	0	0	1	1
1	1	1	1	0	0	0	0	1	1	0
1	1	1	1	1	0	1	0	1	0	1

Figure 3: The game G_5

Proposition 5.1. G_5 has a unique NE $((x_1, 1 - x_1), (x_2, 1 - x_2), (x_3, 1 -$

$x_3), (x_4, 1 - x_4), (x_5, 1 - x_5))$ given by

$$\begin{aligned} x_1 &\approx 0.350370422, \quad x_2 \approx 0.646516479, \quad x_3 \approx 0.648711818, \\ x_4 &\approx 0.371748770, \quad x_5 \approx 0.368061687. \end{aligned}$$

The probabilities in the NE are all irrational and irradical.

Proof. First observe that $\text{PNE}(G_4) = \emptyset$ using this table of profitable deviations:

action profile	unsatisfied players
(0, 0, 0, 0, 0)	4
(0, 0, 0, 0, 1)	3, 5
(0, 0, 0, 1, 0)	2
(0, 0, 0, 1, 1)	1, 2
(0, 0, 1, 0, 0)	2, 4
(0, 0, 1, 0, 1)	1, 5
(0, 0, 1, 1, 0)	5
(0, 0, 1, 1, 1)	3
(0, 1, 0, 0, 0)	3, 4, 5
(0, 1, 0, 0, 1)	1, 2
(0, 1, 0, 1, 0)	3
(0, 1, 0, 1, 1)	1, 4
(0, 1, 1, 0, 0)	4, 5
(0, 1, 1, 0, 1)	2
(0, 1, 1, 1, 0)	2
(0, 1, 1, 1, 1)	1, 2
(1, 0, 0, 0, 0)	1, 4
(1, 0, 0, 0, 1)	1, 3, 5
(1, 0, 0, 1, 0)	1, 2
(1, 0, 0, 1, 1)	2, 4
(1, 0, 1, 0, 0)	1, 2, 4
(1, 0, 1, 0, 1)	5
(1, 0, 1, 1, 0)	5
(1, 0, 1, 1, 1)	3
(1, 1, 0, 0, 0)	1, 3, 4, 5
(1, 1, 0, 0, 1)	2
(1, 1, 0, 1, 0)	1, 3
(1, 1, 0, 1, 1)	4
(1, 1, 1, 0, 0)	1, 4, 5
(1, 1, 1, 0, 1)	2, 3
(1, 1, 1, 1, 0)	1, 2, 5
(1, 1, 1, 1, 1)	2, 4

Now the polynomials f_1, f_2, f_3, f_4, f_5 from Equation ?? are

$$\begin{aligned}
f_1(x_2, x_3, x_4, x_5) &= -5x_2x_3x_4x_5 + 4x_2x_3x_4 + 2x_2x_3x_5 - x_2x_3 + 3x_2x_4x_5 \\
&\quad - 2x_2x_4 - 2x_2x_5 + x_2 + x_3x_4x_5 - x_3x_4 - x_4x_5 + x_4 \\
&\quad + 2x_5 - 1, \\
f_2(x_1, x_3, x_4, x_5) &= x_3x_4x_5 + 2x_3x_4 - 2x_3 - 2x_4x_5 + 1, \\
f_3(x_1, x_2, x_4, x_5) &= -2x_1x_2x_4x_5 + x_1x_2x_4 + x_1x_2x_5 + 2x_1x_4x_5 - x_1x_4 - x_1x_5 \\
&\quad + 3x_2x_4x_5 - 3x_2x_4 + x_2 - x_4x_5 + x_4 - x_5, \\
f_4(x_1, x_2, x_3, x_5) &= 2x_1x_2x_3x_5 - 2x_1x_2x_3 - x_1x_2x_5 + x_1x_2 - x_1x_3x_5 + x_1x_3 \\
&\quad + x_1x_5 - x_1 - x_2x_3x_5 + x_2x_3 + x_2x_5 - x_2 - 2x_5 + 1, \\
f_5(x_1, x_2, x_3, x_4) &= -x_1x_2x_3x_4 + x_1x_2x_3 + x_1x_2x_4 - x_1x_2 + x_1x_3x_4 - x_1x_3 \\
&\quad - x_1x_4 + x_1 - x_2x_3x_4 + 3x_2x_4 - x_3x_4 + x_3 - 1.
\end{aligned}$$

We show using Lemma ?? that if $x_i \in \{0, 1\}$ for some i then $x \notin \text{NE}(G_5)$. The proof proceeds as follows. For each i in turn, we fix $x_i \in \{0, 1\}$ and consider the subgame between the 4 players obtained by restricting player i to play the pure action corresponding to x_i . We then perform a case analysis to find all profiles x_{-i} for these 4 players which are NEs in this subgame. The case analysis is similar to the case analysis in the proof of Proposition ?. We then show that for each such NE profile x_{-i} , $(2x_i - 1)f_i(x_{-i}) < 0$, meaning i is not playing a best response to x_{-i} in the full 5-player game G_5 . So for player i to play a best response in any NE of the full 5-player game player, i will need to choose $x_i \in (0, 1)$.

There are 10 cases in the proof (one for each $i \in [5]$ and $x_i \in \{0, 1\}$). Here is additional information regarding these 10 cases (here the approximate values were computed using Mathematica):

Case 1: $x_1 = 0$. Here the unique partial NE for players 2, 3, 4, 5 is given by

$$\begin{aligned}
x_2 &\approx 0.650518016106, \quad x_3 \approx 0.638238319763, \\
x_4 &\approx 0.402794248582, \quad x_5 \approx 0.433321011106.
\end{aligned}$$

And $f_1(x_{-1}) \approx 0.103778882911 > 0$. Therefore $x_1 \neq 0$ (observe that $|\nabla f_1| \leq 1000$ so the value of f_1 at the exact value of x_{-1} will still be > 0 when using 12-digit approximations).

Case 2: $x_1 = 1$. Here the unique partial NE for players 2, 3, 4, 5 is given by

$$\begin{aligned} x_2 &\approx 0.589697169563, \quad x_3 \approx 0.681923203912, \\ x_4 &\approx 0.338247172171, \quad x_5 \approx 0.218624870529. \end{aligned}$$

And $f_1(x_{-1}) \approx -0.245843038086 < 0$. Therefore $x_1 \in (0, 1)$.

Case 3: $x_2 = 0$. There is no partial NE where $x_1 \in (0, 1)$, so $x_2 \neq 0$.

Case 4: $x_2 = 1$. Here all partial NEs for players 1, 3, 4, 5 are given by

$$x_1 \in (0, 1), \quad x_3 = 1, \quad x_4 = 0, \quad x_5 = \frac{1}{2}.$$

And $f_2(x_{-2}) = -1 < 0$. Therefore $x_2 \in (0, 1)$.

Case 5: $x_3 = 0$. There is no partial NE where $x_1, x_2 \in (0, 1)$, so $x_3 \neq 0$.

Case 6: $x_3 = 1$. There is no partial NE where $x_1, x_2 \in (0, 1)$, so $x_3 \in (0, 1)$.

Case 7: $x_4 = 0$. There is no partial NE where $x_1, x_2, x_3 \in (0, 1)$, so $x_4 \neq 0$.

Case 8: $x_4 = 1$. There is no partial NE where $x_1, x_2, x_3 \in (0, 1)$, so $x_4 \in (0, 1)$.

Case 9: $x_5 = 0$. There is no partial NE where $x_1, x_2, x_3, x_4 \in (0, 1)$, so $x_5 \neq 0$.

Case 10: $x_5 = 1$. There is no partial NE where $x_1, x_2, x_3, x_4 \in (0, 1)$, so $x_5 \in (0, 1)$.

Therefore the only NEs are fully mixed. This means that $x \in \text{NE}(G_5)$ satisfies $f_i(x) = 0$ for every i . We solve this now.

Let $F = \{f_1, f_2, f_3, f_4, f_5\}$. It is a set of polynomials in the polynomial ring $\mathbb{Q}[x_1, x_2, x_3, x_4, x_5]$. We begin by finding the Gröbner basis $G = \{g_1, g_2, g_3, g_4, g_5\}$ of the ideal $\langle F \rangle$, with the monomial ordering generated lexicographically by $x_1 \succ x_2 \succ x_3 \succ x_4 \succ x_5$, given in full in ?? (using Mathematica).

We now compute the solutions of $\forall i \ g_i = 0$ to find those of $\forall i \ f_i = 0$. The proof proceeds as follows. We observe that like in the proof of Proposition ??, $g_1 \in \mathbb{Q}[x_5]$ and for $i \geq 2$, $g_i = a_i x_{6-i} + h_i(x_5)$ where $a_i \in \mathbb{Z}$ and $h_i \in \mathbb{Q}[x_5]$. So we first find the roots of g_1 , using a similar argument to the one in the proof

of Proposition ?? and ??, and find 5 roots of g_5 in $(0, 1)$, meaning 5 possible values for x_5 . Then for each of the latter four we show using g_2, \dots, g_5 that one of x_1, x_2, x_3, x_4 falls outside $(0, 1)$, implying a contradiction. In conclusion, there is a unique possible value for x_5 and by the structure of g_2, \dots, g_5 a unique possible value for each of x_1, x_2, x_3, x_4 . This solution is given in the approximated decimal form:

$$\begin{aligned} x_1 &\approx 0.3503704221, \quad x_2 \approx 0.6465164785, \quad x_3 \approx 0.6487118183, \\ x_4 &\approx 0.3717487703, \quad x_5 \approx 0.3680616872. \end{aligned}$$

To see that x_1, x_2, x_3, x_4, x_5 are irrational, as before, we find their minimal polynomials (up to scaling):

$$\begin{aligned} P_1(y) = & 576y^{26} + 7360y^{25} + 4496y^{24} + 250816y^{23} + 355564y^{22} \\ & - 7898280y^{21} + 7658244y^{20} + 75548910y^{19} - 142993432y^{18} \\ & - 189077220y^{17} + 334832314y^{16} + 474658874y^{15} + 1696776519y^{14} \\ & - 7562033350y^{13} + 1512397109y^{12} + 14656660866y^{11} \\ & - 10470778075y^{10} - 3782764895y^9 - 665928467y^8 + 10050992181y^7 \\ & - 5447529632y^6 - 1586217407y^5 + 562969676y^4 + 1493433257y^3 \\ & - 79516172y^2 - 522298360y + 132583232, \end{aligned}$$

$$\begin{aligned} P_2(y) = & 2057728y^{26} - 47527232y^{25} + 374368064y^{24} - 482543600y^{23} \\ & - 11253727536y^{22} + 90191808584y^{21} - 352161035060y^{20} \\ & + 829456047656y^{19} - 1120735934348y^{18} + 286387243298y^{17} \\ & + 2393470340090y^{16} - 6252418753985y^{15} + 9238553888534y^{14} \\ & - 9587641717941y^{13} + 7391393209142y^{12} - 4277517478697y^{11} \\ & + 1815337752171y^{10} - 515212602860y^9 + 59595129244y^8 \\ & + 25514336227y^7 - 17322169528y^6 + 5510430025y^5 - 1139164516y^4 \\ & + 160112123y^3 - 14874613y^2 + 828146y - 20988, \end{aligned}$$

$$\begin{aligned}
P_3(y) = & 5828y^{26} - 80590y^{25} + 471147y^{24} - 1473516y^{23} \\
& + 1995893y^{22} + 3280961y^{21} - 21791522y^{20} + 51425278y^{19} \\
& - 93080861y^{18} + 203283288y^{17} - 444991348y^{16} + 713613468y^{15} \\
& - 837466118y^{14} + 925602099y^{13} - 1210417319y^{12} + 1552957912y^{11} \\
& - 1585613560y^{10} + 1241271492y^9 - 772369636y^8 + 401799920y^7 \\
& - 180281904y^6 + 69151344y^5 - 21721184y^4 + 5260240y^3 \\
& - 909264y^2 + 99200y - 5120,
\end{aligned}$$

$$\begin{aligned}
P_4(y) = & 367535904y^{26} - 4069658144y^{25} + 15863607440y^{24} + 4840925000y^{23} \\
& - 332049318712y^{22} + 1798491222250y^{21} - 5880537317282y^{20} \\
& + 13925726681306y^{19} - 25570689763592y^{18} + 37716767814324y^{17} \\
& - 45641117387319y^{16} + 45922091951685y^{15} - 38746001553780y^{14} \\
& + 27551639828330y^{13} - 16547091697878y^{12} + 8389403526528y^{11} \\
& - 3578413101941y^{10} + 1274916123621y^9 - 374578292008y^8 \\
& + 88741617902y^7 - 16260708632y^6 + 2104604912y^5 - 141853168y^4 \\
& - 7308320y^3 + 2769472y^2 - 256128y + 8192,
\end{aligned}$$

$$\begin{aligned}
P_5(y) = & 25772032y^{26} - 399987456y^{25} + 2634374272y^{24} - 8416506944y^{23} \\
& + 2215910496y^{22} + 110722637568y^{21} - 612619393968y^{20} \\
& + 2025933659884y^{19} - 4916502391844y^{18} + 9383743371458y^{17} \\
& - 14566831934444y^{16} + 18743281388994y^{15} - 20217905397986y^{14} \\
& + 18405783324440y^{13} - 14192403999687y^{12} + 9280251573089y^{11} \\
& - 5141465321619y^{10} + 2406410117193y^9 - 946553896669y^8 \\
& + 310404824295y^7 - 83872244335y^6 + 18358684229y^5 \\
& - 3175419088y^4 + 417903630y^3 - 39340244y^2 + 2360724y - 67892.
\end{aligned}$$

These can all be verified to be irreducible using Mathematica. Then, the Galois groups of each of these 5 polynomials can be verified to be all S_{26} which is unsolvable (using e.g., Magma). Therefore x_1, x_2, x_3, x_4, x_5 are all irrational. This concludes the proof. \square

6. Six players and up

Here we use the constructions of Section ?? and Section ?? to construct games G_n for every $n \geq 6$ with the same property as G_4 and G_5 . This

completes the picture in terms of the (general worst case) inexpressibility of NEs with radicals.

First, we use the notion of *product of games*:

Definition 6.1. *Given are games H^1, H^2 such that H^j has n_j players, the pure action profiles A^j and the payoff functions $u_1^j, \dots, u_{n_j}^j$. Define the product game $H^1 \times H^2$ as follows:*

- *there are $n_1 + n_2$ players,*
- *the set of pure action profiles is $A^1 \times A^2$,*
- *the payoff functions are $u_i(a^1, a^2) = u_i^1(a^1)$ if $i \in [n_1]$ and $u_i^2(a^2)$ if $i \in \{n_1 + 1, \dots, n_1 + n_2\}$.*

If $H_1 = H_2$, also write $H_1 \times H_2 = H_1^{\times 2}$, and generalize the notation to a product of q copies of H_1 by $H_1^{\times q}$.

This operation is associative, therefore $H_1^{\times q}$ is well defined. Here is a property of $H^1 \times H^2$ that follows by definition:

Lemma 6.2. $\text{NE}(H^1 \times H^2) = \text{NE}(H^1) \times \text{NE}(H^2)$.

Using this lemma, for every $n = 4n_1 + 5n_2$ where $n_1, n_2 \in \mathbb{N}$ are not both 0, the game $G_4^{\times n_1} \times G_5^{\times n_2}$ has a unique NE and its probabilities are all irrational (see Proposition ?? below). However, these constructions do not cover the cases $n = 6, 7, 11$. To fill these holes, we now show how to add 2 players to each of G_4 and G_5 while maintaining their desired properties (a unique NE, whose probabilities are irrational).

Define the $2 \times 2 \times 2$ game H_3 :

$$\left[\begin{array}{ccc|ccc} 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \end{array} \right]$$

Figure 4: The game H_3

Here is the main property of interest of this game:

Proposition 6.3. *Let $x = ((x_1, 1 - x_1), (x_2, 1 - x_2), (x_3, 1 - x_3))$ such that $x_1 \in (0, 1)$. Then $x \in \text{NE}(H_3)$ iff $x_1 = x_2 = x_3$.*

Proof. First, the polynomials f_1, f_2, f_3 from Equation ?? of H_3 are

$$f_1(x_2, x_3) = 0, \quad f_2(x_1, x_3) = x_1 - x_3, \quad f_3(x_1, x_2) = x_2 - x_1.$$

By definition, for every $x_1 \in (0, 1)$ setting $x_1 = x_2 = x_3$ means that $f_1(x) = f_2(x) = f_3(x) = 0$, meaning that $x \in \text{NE}(H_3)$. Now we show that if $x_1 \in (0, 1)$ and $x \in \text{NE}(H_3)$ then $x_1 = x_2 = x_3$.

We use Lemma ?. $f_1 = 0$ so x_1 can take on any value, regardless of the values of x_2, x_3 . Assume that $x_1 \in (0, 1)$. Suppose for contradiction that $x_2 = 0$. Then $f_3 = -x_1 < 0$, so $x_3 = 0$ so $f_2 = x_1 > 0$ so $x_2 = 1 -$ a contradiction. Similarly, if $x_2 = 1$, then $f_3 = 1 - x_1 > 0$ so $x_3 = 1$ so $f_2 = x_1 - 1 < 0$ so $x_2 = 0$ – again a contradiction. Therefore $x_2 \in (0, 1)$ so $f_2 = 0$ so $x_3 = x_1$ so $f_3 = 0$ so $x_1 = x_2 = x_3$. \square

Next we define a variation of the product game with a small overlap in the players.

Definition 6.4. *Let G be a game such that*

- *there are n players,*
- *each player has 2 pure actions,*
- *the payoffs are given by functions u_1, \dots, u_n ,*
- *there is a unique NE, and this NE is fully mixed.*

Define the game $G \circ H_3$ as follows:

- *there are $n + 2$ players,*
- *each player has 2 pure actions,*
- *the payoff functions are given by $u_i(a, a_{n+1}, a_{n+2}) = u_i(a)$ for $i \in [n]$, and for $i \in \{n + 1, n + 2\}$ the payoff is defined as $u_i(a, a_{n+1}, a_{n+2}) = u_{i-(n-1)}^{H_3}(a_n, a_{n+1}, a_{n+2})$ (as defined in Figure ??).*

The most important thing about this operation $\circ H_3$ is that it is compatible with G_4 and G_5 . The key property of $G \circ H_3$ is as follows:

Lemma 6.5. *If $\text{NE}(G) = \{((x_1, 1 - x_1), \dots, (x_n, 1 - x_n))\}$ and $x_n \in (0, 1)$ then*

$$\text{NE}(G \circ H_3) = \{((x_1, 1 - x_1), \dots, (x_n, 1 - x_n), (x_n, 1 - x_n), (x_n, 1 - x_n))\}.$$

Proof. By definition. Observe that a NE of $G \circ H_3$, restricted to players $n, n + 1, n + 2$, is a NE to H_3 , because the payoff of the first player in H_3 (meaning player n in $G \circ H_3$) is constant. Since player n plays $(x_n, 1 - x_n)$ in the only NE of G , player n will need to play it in any NE of $G \circ H_3$. Therefore we can invoke Proposition ?? here to get the desired result. \square

Now the games G_n for $n \geq 6$ follow. Let $n \geq 6$. Write $n = 4q + r$ where $1 \leq q \in \mathbb{N}$ and $r \in \{0, 1, 2, 3\}$. Define

$$G_n = \begin{cases} G_4^{\times q} & r = 0, \\ G_4^{\times q-1} \times G_5 & r = 1, \\ G_4^{\times q-1} \times (G_4 \circ H_3) & r = 2, \\ G_4^{\times q-1} \times (G_5 \circ H_3) & r = 3. \end{cases} \quad (3)$$

Proposition 6.6. *For every $n \geq 6$, the game G_n has a unique NE, and the probabilities in this NE are all irrational.*

Proof. Let $n = 4q + r$. By Proposition ??, Proposition ?? and Lemma ?? we obtain that $G_4 \circ H_3$ and $G_5 \circ H_3$ both have a unique NE, and that NE is irrational in both games. Now, for any $n \geq 6$, G_n is a product of games which all have a unique NE, and that NE is irrational in all those games. Therefore each G_n has a unique NE, and its probabilities are all irrational. \square

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Appendix A. Further details in the proof of Proposition ??

In the proof of Proposition ?? we use the approximations of the roots of g_1 to determine the unique value of x_4 . To circumvent possible floating

point issues, we can use the Sturm sequence of g_1 to determine using exact arithmetic that (1) there are exactly 2 real roots, and (2) one is in the interval $(0.3, 0.4]$ and one in the interval $(0.4, 0.5]$ (?). The Sturm sequence of g_1 can be found using Mathematica:

$$\begin{aligned}
p_0 &= g_1 = 5x_4^6 - 44x_4^5 + 143x_4^4 - 163x_4^3 + 85x_4^2 - 21x_4 + 2, \\
p_1 &= g_1' = 30x_4^5 - 220x_4^4 + 572x_4^3 - 489x_4^2 + 170x_4 - 21, \\
p_2 &= -\text{Remainder}(p_0, p_1) = \frac{55}{9}x_4^4 - \frac{5249}{90}x_4^3 + \frac{943}{15}x_4^2 - \frac{433}{18}x_4 + \frac{47}{15}, \\
p_3 &= -\text{Remainder}(p_1, p_2) = -\frac{27110403}{30250}x_4^3 + \frac{15927363}{15125}x_4^2 - \frac{2514591}{6050}x_4 \\
&\quad + \frac{831852}{15125}, \\
p_4 &= -\text{Remainder}(p_2, p_3) = \frac{4813052861375}{81663772313601}x_4^2 - \frac{1017283844500}{27221257437867}x_4 \\
&\quad + \frac{417696219875}{81663772313601}, \\
p_5 &= -\text{Remainder}(p_3, p_4) = \frac{46666740312733677523326}{1531601841083641320125}x_4 \\
&\quad - \frac{19799411241381912287163}{1531601841083641320125}, \\
p_6 &= -\text{Remainder}(p_4, p_5) = \frac{3506083202136869448018665125}{26667695965024814331677001308676}.
\end{aligned}$$

where $\text{Remainder}(f, g)$ is the remainder of the polynomial division of f by g .

For every $a \in \mathbb{R} \cup \{-\infty, \infty\}$, let $V(a)$ be the number of sign changes in the sequence $(p_0, p_1, p_2, p_3, p_4, p_5, p_6)$ (and the sign of something like $p_0(\infty)$ is defined to be +). Sturm's theorem states that for every $-\infty \leq a < b \leq \infty$, the number of real roots in $(a, b]$ is equal to $V(a) - V(b)$.

So we have $V(-\infty) - V(\infty) = 4 - 2 = 2$, so there are 2 real roots. Note

that

$$\begin{aligned}
p_0\left(\frac{3}{10}\right) &= \frac{161}{40000}, \\
p_1\left(\frac{3}{10}\right) &= -\frac{2751}{10000}, \\
p_2\left(\frac{3}{10}\right) &= \frac{371}{7500}, \\
p_3\left(\frac{3}{10}\right) &= \frac{26761959}{30250000}, \\
p_4\left(\frac{3}{10}\right) &= -\frac{258737930605}{326655089254404}, \\
p_5\left(\frac{3}{10}\right) &= -\frac{28996945737809045150826}{7658009205418206600625}, \\
p_6\left(\frac{3}{10}\right) &= \frac{3506083202136869448018665125}{26667695965024814331677001308676},
\end{aligned}$$

and

$$\begin{aligned}
p_0\left(\frac{4}{10}\right) &= -\frac{4}{3125}, \\
p_1\left(\frac{4}{10}\right) &= \frac{27}{625}, \\
p_2\left(\frac{4}{10}\right) &= -\frac{4}{625}, \\
p_3\left(\frac{4}{10}\right) &= -\frac{236727}{1890625}, \\
p_4\left(\frac{4}{10}\right) &= -\frac{32955935705}{81663772313601}, \\
p_5\left(\frac{4}{10}\right) &= -\frac{5663575581442206389163}{7658009205418206600625}, \\
p_6\left(\frac{4}{10}\right) &= \frac{3506083202136869448018665125}{26667695965024814331677001308676},
\end{aligned}$$

and

$$\begin{aligned}
p_0\left(\frac{5}{10}\right) &= \frac{1}{64}, \\
p_1\left(\frac{5}{10}\right) &= \frac{7}{16}, \\
p_2\left(\frac{5}{10}\right) &= -\frac{31}{360}, \\
p_3\left(\frac{5}{10}\right) &= -\frac{383139}{242000}, \\
p_4\left(\frac{5}{10}\right) &= \frac{380134673875}{326655089254404}, \\
p_5\left(\frac{5}{10}\right) &= \frac{28271671319879411796}{12252814728669130561}, \\
p_6\left(\frac{5}{10}\right) &= \frac{3506083202136869448018665125}{26667695965024814331677001308676}.
\end{aligned}$$

So $V(0.3) - V(0.4) = 4 - 3 = 1$ and $V(0.4) - V(0.5) = 3 - 2 = 1$, so $(0.3, 0.4]$ and $(0.4, 0.5]$ each contain exactly one of the real roots. More specifically, $r_1 \in (0.3, 0.4]$ and $r_2 \in (0.4, 0.5]$.

From here the argument in the proof of Proposition ?? that x_4 cannot be r_2 holds and the rest of the proof goes through.

As for r_1 , the method above can be further used to verify that $r_1 \in (0.320065197644, 0.320065197645]$ using exact arithmetic, justifying the approximation for $x_4 = r_1$ given in the proposition statement.

Appendix B. The Gröbner basis in Proposition ??

Here is the Gröbner basis G used in Proposition ?. It is $\{g_1, g_2, g_3, g_4, g_5\}$ given as follows:

$$\begin{aligned} g_1 = & 25772032x_5^{26} - 399987456x_5^{25} + 2634374272x_5^{24} - 8416506944x_5^{23} \\ & + 2215910496x_5^{22} + 110722637568x_5^{21} - 612619393968x_5^{20} \\ & + 2025933659884x_5^{19} - 4916502391844x_5^{18} + 9383743371458x_5^{17} \\ & - 14566831934444x_5^{16} + 18743281388994x_5^{15} - 20217905397986x_5^{14} \\ & + 18405783324440x_5^{13} - 14192403999687x_5^{12} + 9280251573089x_5^{11} \\ & - 5141465321619x_5^{10} + 2406410117193x_5^9 - 946553896669x_5^8 \\ & + 310404824295x_5^7 - 83872244335x_5^6 + 18358684229x_5^5 - 3175419088x_5^4 \\ & + 417903630x_5^3 - 39340244x_5^2 + 2360724x_5 - 67892, \end{aligned}$$

$$\begin{aligned} g_2 = & -5229757014096166506106706852927054776656409436526688311934217666969182934601238598093813122993784633757481615360x_5^{25} \\ & + 77417348381807438139298712751795524837772256296658899465798071561263108113263168115187337389190007008413770618112x_5^{24} \\ & - 479003561865428975555899714436006958766026300658008374819291240222777323612217905003850065127110597318088478641792x_5^{23} \\ & + 1363516171631648731401937260410517759345140344024728313437486577525612859813544026537686042034050828432977603879232x_5^{22} \\ & + 533587922587820729708044335466442589479278858927142642702831828582442863250633819428985449118623257898991410037152x_5^{21} \\ & - 2210014743880930808861819847452779162469225691207064470932200568612066220139113476396459639356379722280870005175520x_5^{20} \\ & + 108454900964535452643280673745481810479068278931480402551733219020654394658331661285136003467695044177290332008077552x_5^{19} \\ & - 333073782919957459027278804921869980726997241882391079940534925956228969796725710815019469532991976447383544674526084x_5^{18} \\ & + 757624233673338301359440552712904560489502717271112905470949925790532905236334410472737236149940199890686132929621704x_5^{17} \\ & - 1357384404662525914311341693592403279296839308873698777339142031034705771862285501894438660803278932926999068794546582x_5^{16} \\ & + 1974990208789200599466586861637319326312916081007911794448182406682496223056735984172167459256026946263774303269130502x_5^{15} \\ & - 2374297193220515593964694404791585125075717016662628050023610781028665713690408237409030760913488268277385743966777032x_5^{14} \\ & + 2382354902550652961280745151068702383273085454800813608807850901474781766931357892102898412048158076148480110682671270x_5^{13} \\ & - 2006547338020399904478992449241580128008830980255928702874144615524529032076799032362502252260795483084720768354418906x_5^{12} \\ & + 142230371259443866763168836036071116173856475208264246894179873753516371223301768725492398493859422183929726651437959x_5^{11} \\ & - 848591809572807446526737307028116087664196830302677461687679329565088044403583949846081471077208029772845733783978556x_5^{10} \\ & + 425256712357840429041584855108259194843259926786949852822199851920502756754006632376292656084188179152265324768387541x_5^9 \\ & - 178190185696798219573888134723758197101657622267716428247687462654805037079392982664482126949337362000059358903070826x_5^8 \\ & + 61968796201222942179577897911545319747966347541812098462585270467890006274705019861192475346963305894981804268851065x_5^7 \\ & - 17686914368321420604395590031619931467353119221112708310058990785917431994157571306893482096694260412715975874006756x_5^6 \\ & + 4075259952587509321223128820658917157429499391726104350291091494649891620567629499432147402492820100061680066097185x_5^5 \\ & - 739721454661762561829150779036766386028574319580498394474729242346540997892430899746745978354273573458080278629282x_5^4 \\ & + 101889556627033780067706963405180770731046840551112096527796778096348672520153243935935166739545857450254458665078x_5^3 \\ & - 10014961581059325297534083799423222232193173182415278003356345107715054775890753680315781533615428759531894840428x_5^2 \\ & + 626198151376926556516994281746232986281114459167528107589000512836306110760157779105130081244956865692258077872x_5 \\ & + 14562790811532232105030294780908703577127741453867317051625330584746436330267698312201792323393195475584x_4 \\ & - 1873039573765845395448383944548962695520731214885148618902226778682417235942672754625522902251804284010201836, \end{aligned}$$

$$\begin{aligned}
g_3 = & 2334544751155071882026722780061349241642599026068676574946240154619529918561800486748917602409834307249946624x_5^{25} \\
& - 34826386272732604074820628503893799177194447217375096775483672338860118986103889174063050911927141610810379520x_5^{24} \\
& + 217633964950393711549186629464978437165300442561564386772405822217081697486573874407123129145535922950706091648x_5^{23} \\
& - 631004665463527900669117543672267152557837342795954164588602147543470354961220923336218777672617507047396540736x_5^{22} \\
& - 181187120952388236998386793367411756303633247262581447028534281445658265062871643279630089023039796374318343072x_5^{21} \\
& + 9925416763570921108480922821290611359261790440483739235088028728524852578911877914757663730293796516424536982112x_5^{20} \\
& - 49511025362638119989000055444640688180713101494312128922185056130200524969006756798221621831744228717747164071536x_5^{19} \\
& + 153606450244873998637400182694657307098185015597200929534563721055218708062776648371930100226217096327071566507812x_5^{18} \\
& - 352426714217615075454859339166993063141216123999988812644511442641474817930073511713523770604296026296813438396088x_5^{17} \\
& + 636541832234522264020461181180548709485760714008828971077784400214808511794157519019430471954449244547200377722358x_5^{16} \\
& - 933498979637348381777015122675301004004870810348162357840262569632788867775411763510678948565334319576360783673102x_5^{15} \\
& + 1131075337234367536418256190781035607571454240605206256230753284711414951414891505602860229050574408650155226782728x_5^{14} \\
& - 1143888525546125381293381264403258806739680192797154817504244626489887282710627211132657772617738453443292075022774x_5^{13} \\
& + 971114741160926999426755391516911794308562783195876583953124144282681326795227451967634738501132063379159602305106x_5^{12} \\
& - 693868399603588601921565117253658200432981408294055991494009684604764816981147591464337021509223488254637528974415x_5^{11} \\
& + 417307080155928782025068977202423092573461396691977026044607405101837819435313315500107795471975427014871318266976x_5^{10} \\
& - 210801786743647026314962428699955929372925777966515764581497272997599561283615798792354612095177686041986524450961x_5^9 \\
& + 89033363860194520670813956986147529620606515277585298169282801928853396835629522791071359703663275943842956580202x_5^8 \\
& - 31207237132427508982273789207629895352893344450292521048842395054212958515460240818669350564989492535868999725081x_5^7 \\
& + 8976459942331587279270618598121344735704429300848888815877070229254977551136189717073239186833151007865116037992x_5^6 \\
& - 2084149071494599066636840904363798736708176439312551310459408028002687362733692900701646271408658267467743099101x_5^5 \\
& + 381158308756241245183778096001699736234731027675341960830001210231765721434485154435182600335913406210638306866x_5^4 \\
& - 52889407110880592240935035564652467638860334385675951440377263860227391443955038430565826220210872540006884446x_5^3 \\
& + 5236307109205295690183209426599989107958678066276665974271201847161057962269146837250130645744715913371810412x_5^2 \\
& - 329731014995559194481139244161781866865235013681345347340507275899274691789786673963774969773751450999559968x_5 \\
& + 369474411292823674535444822098526709270851631110901155469041811943454903727724876142148905734577408x_3 \\
& + 9931369572758577248575536574765581760354618753802907709675771557134507504523669268186354299397125834523196,
\end{aligned}$$

$$\begin{aligned}
g_4 = & -3558500754024337082770180156322829498597886283821905562909626877933627437862004077143293724200219009927888896x_5^{25} \\
& + 53459533959080117233238158244911856852132196679058294422752524765582301160471501982607167430974092115094435584x_5^{24} \\
& - 337159203512668656410477904168737985969382049375082645168114655955720623507549883202440822289987871942813580672x_5^{23} \\
& + 994379773429930614855305353972967341994717800292461875312328526647809247215078010238428657356468774616503704256x_5^{22} \\
& + 189245805202958885164717457856217010089940221854661319802127336154003800666496337272960302229722564154334014816x_5^{21} \\
& - 15197125154505736973813689155185565411338769583659301109177897329934627331339207419412635177208389035828827853280x_5^{20} \\
& + 77035200685565232608193856879005729648408468927829648172264573836344951657512635824410473666817586520056592466384x_5^{19} \\
& - 241401429737158706348114011840171158855788925159564563866972792216940735637246580890337940955337976062511379408780x_5^{18} \\
& + 558630449888944438077465724771239908940298471522896091656782046424166365437557890923179651315268131895457834737840x_5^{17} \\
& - 1017224595597228136804728426231636929621054017646148023601819708031196666257616755854341239641312660582786669412250x_5^{16} \\
& + 1503817045296227892476950853512837221625294649034200024746679125649511207123301404686450154802628965251836163156638x_5^{15} \\
& - 1836897428293610264177371346711383956146293025290376833853869340646192152519382428015203859429983377866327285534440x_5^{14} \\
& + 1872995781108315067497312460676257927626533427837197474160535503357940379854809628626707176037342237011018857587474x_5^{13} \\
& - 1603373170906444757398030676224373499811872254724020941343536500445962959353227034242254754665909779095643373418578x_5^{12} \\
& + 1155288040878502319572608128049976975067745119574642103393020077529068588347437379307725141482112392993781461014925x_5^{11} \\
& - 700697376259091308302109508302940157504814446394086929947963510799016007546849283458571446545182585269186009462998x_5^{10} \\
& + 356934874693265870665365221583700170939530876165407760810376569278622878215127192877970425179784984014582661641105x_5^9 \\
& - 152001087875514448869480222043342566350774439463857717641593193787967622803143243410107426653140398261285409508742x_5^8 \\
& + 53706296084516238933941939388111267070100298333608164123229065510826751235963393953259745031831145831662722511511x_5^7 \\
& - 15566872495760296770636202098010682330153678839979729723749724038484459853264945820724345917034624145600833219006x_5^6 \\
& + 3640403667504735955909101560636890647144799421188362661421370193637130168818945991278667273963692398047273666453x_5^5 \\
& - 670173581881674007325938393635716663699469008474575049226005766285927286886676689504925496451118751189529451670x_5^4 \\
& + 93535104453294586615470233339118267735846441345555250692303574178503068296553438102544432800191989801053615838x_5^3 \\
& - 9305270756368072254559737977820811715941760923436884753321555175707328683305744660190066221202432676682105108x_5^2 \\
& + 588069610334459746429373597750058332341244517659467031174397469825967982566556498337804906924843284145942504x_5 \\
& + 5095805281351607233898259232880569267178981710856112472144845244406570095429556031988777094079363968x_2 \\
& - 17749305818938603275929612340121722447988753251997841330848673207984991309809988341490716053109866637501340,
\end{aligned}$$

$$\begin{aligned}
g_5 = & 9270243108944911855913423923940118860724549686336022978726911310961722458203883119715431043049984039600128x_5^{25} \\
& - 138175875266337145223434566252355090972050168617333616931556655001369528702606497111283046584897048123496192x_5^{24} \\
& + 862492337870286319208107141713916178120914949268939539624311844622498318270220300752453215856315985897987456x_5^{23} \\
& - 2495102461650519533506960742919876844791681341902817276675711486882107127757546342958357009723667467045595840x_5^{22} \\
& - 749636554146947543936421259610110568924854870730577569236289936020650163568330306595235407217786364530222432x_5^{21} \\
& + 39402580158298384618280673190243773364515744583137648398647346524048246636434435392106680856462315406779353056x_5^{20} \\
& - 196122144147700348424157453329717858016119647923907735554534428793476055982854327245096135183281374007838006224x_5^{19} \\
& + 607567407220738066950645872823562629356481100952338581837849090942575790518184039147606780264543360404788805388x_5^{18} \\
& - 1392034559331843275751880128686367355159492660859513756685269123763970049996344394913220744816481337236164688752x_5^{17} \\
& + 2510558512785233508346113941902852818446593588622592969503358098040942372525809496185640497523813014968232095706x_5^{16} \\
& - 3675774626200611204927149730882154016887117306216041250003754929446407301397549595440548848847486181249067607390x_5^{15} \\
& + 4445471379414986634335124565974011524961495924338229912795322982376585158552336715200597575424938671909289449640x_5^{14} \\
& - 4486118641224030633093582292243068588732437612076418746519037181965028002339068594596667317742675531315640964002x_5^{13} \\
& + 379887430855571159848287364010160216595125000958154051962643647863636596526569190224129406287877541961423998546x_5^{12} \\
& - 2706192488036664027366752238840352044709794830418822829301155231922766756880774252399289499491242862301112642493x_5^{11} \\
& + 1621786402882167377760236931048484324618716111562222154462408173075352980667480821275949711332943322996801085942x_5^{10} \\
& - 815797787394998843496442733011612610801250048164205077534210923183660530362409166766004966247838070909131717133x_5^9 \\
& + 342842754795609541453395389548504455515865006985724325769495898905537672061053294325586795901382310703074484590x_5^8 \\
& - 119464200488048451770969329661905353520359137813236139439207882676912887567274947822290425167328332819674138255x_5^7 \\
& + 34124598338008014536089722419271075905275461000608984295483205503028442881007977804213586157616567502762395862x_5^6 \\
& - 7858371647506645013235958461172093630568069977126879574130925982176847618427475770739613780276848717890340861x_5^5 \\
& + 1423375266711743719442643469712816450858096992327301545160933803057177533347744363470751224596704554150864598x_5^4 \\
& - 195278103772230475236462201852070232583953616027553465198362842948766419805745572080063898831820910760883150x_5^3 \\
& + 1907722911030273239358991830671587449405158270858559469272326336390718508460877526710123484137952808852276x_5^2 \\
& - 1182613115373742448605260562696400057836772355234957546845962252089854311499956589035815582233191712586248x_5 \\
& + 22822715865720062256822260809119245287191259217569053122320779895326763303740236022841002163082496x_1 \\
& + 34970112057581607639480208646156766876639044746724958622577550768034487124422136034942339429080145084988.
\end{aligned}$$

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