

WEAKLY DISTANCE-REGULAR DIGRAPHS OF DIAMETER 2

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ABSTRACT. Weakly distance-regular digraphs is a directed version of distance-regular graphs. In this paper, we characterize all weakly distance-regular digraphs of diameter 2.

1. INTRODUCTION

As a natural generalization of distance-regular graphs (for the background of distance-regular graphs, refer to [? ?]) and distance-regular digraphs (see the definition in Section 2), Wang and Suzuki [?] proposed the concept of weakly distance-regular digraphs. Distance-regular graphs, distance-regular digraphs and weakly distance-regular digraphs can all be regarded as association schemes.

A d -class association scheme \mathfrak{X} is a pair $(X, \{R_i\}_{i=0}^d)$, where X is a finite set, and each R_i is a nonempty subset of $X \times X$ satisfying the following axioms (see [? ? ? ?] for a background of the theory of association schemes):

- (i) $R_0 = \{(x, x) \mid x \in X\}$ is the diagonal relation;
- (ii) $X \times X = R_0 \cup R_1 \cup \cdots \cup R_d$, $R_i \cap R_j = \emptyset$ ($i \neq j$);
- (iii) for each i , $R_i^\top = R_{i^*}$ for some $0 \leq i^* \leq d$, where $R_i^\top = \{(y, x) \mid (x, y) \in R_i\}$;
- (iv) for all i, j, h , the cardinality of the set

$$P_{i,j}(x, y) := R_i(x) \cap R_{j^*}(y)$$

is constant whenever $(x, y) \in R_h$, where $R(x) = \{y \mid (x, y) \in R\}$ for $R \subseteq X \times X$ and $x \in X$. This constant is denoted by $p_{i,j}^h$.

A d -class association scheme is also called an association scheme with d classes. The integers $p_{i,j}^h$ are called the *intersection numbers* of \mathfrak{X} . We say that \mathfrak{X} is *commutative* if $p_{i,j}^h = p_{j,i}^h$ for all $0 \leq i, j, h \leq d$. The subsets R_i are called the *relations* of \mathfrak{X} . For each i , the integer $k_i (= p_{i,i^*}^0)$ is called the *valency* of R_i . A relation R_i is called *symmetric* if $i = i^*$, and *non-symmetric* otherwise. An association scheme is called *symmetric* if all relations are symmetric, and *non-symmetric* otherwise.

In order to state our main result, we introduce some basic notations about weakly distance-regular digraphs. A *digraph* Γ is a pair $(V\Gamma, A\Gamma)$ where $V\Gamma$ is a finite nonempty set of vertices and $A\Gamma$ is a set of ordered pairs (*arcs*) (x, y) with distinct vertices x and y . For any arc $(x, y) \in A\Gamma$, if $A\Gamma$ also contains an arc (y, x) , then $\{(x, y), (y, x)\}$ can be viewed as an *edge*. We say that Γ is an *undirected graph* or a *graph* if $A\Gamma$ is a symmetric relation. For the digraphs Γ and Σ , the *lexicographic product* of Γ and Σ is the digraph with the vertex set $V\Gamma \times V\Sigma$ such that $((u_1, u_2), (v_1, v_2))$ is an arc if and only if either $(u_1, v_1) \in A\Gamma$, or $u_1 = v_1$ and $(u_2, v_2) \in A\Sigma$.

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A strongly connected digraph Γ is said to be *weakly distance-regular* if the configuration $\mathfrak{X}(\Gamma) = (V\Gamma, \{\Gamma_{\tilde{i}}\}_{\tilde{i} \in \delta(\Gamma)})$ is an association scheme. We call $\mathfrak{X}(\Gamma)$ the *attached scheme* of Γ . We say that Γ is *commutative* if $\mathfrak{X}(\Gamma)$ is commutative. For more information about weakly distance-regular digraphs, see [? ? ? ? ? ? ? ? ? ? ? ? ? ?].

Theorem 1.1. *Let $\mathfrak{X} = (X, \{R_l\}_{l=0}^d)$ be an association scheme. The digraph (X, A) is a weakly distance-regular digraph of diameter 2 with \mathfrak{X} as its attached scheme if and only if \mathfrak{X} has exactly one pair of non-symmetric relations, say $R_1^\top = R_2$, and one of the following holds:*

- Here, $i \in \{1, 2\}$ and $j \in \{3, 4\}$.

2. ASSOCIATION SCHEMES

In this section, we always assume that $\mathfrak{X} := (X, R)$ is an association scheme, where $R = \{R_i\}_{i=0}^d$. In the following, we present some basic concepts, notations and results concerning association schemes.

First we list basic properties of intersection numbers which are used frequently in the remainder of this paper.

Lemma 2.1. ([? , Proposition 5.1] and [? , Chapter II, Proposition 2.2]) *The following hold:*

- (i) $k_i k_j = \sum_{h=0}^d p_{i,j}^h k_h$;
- (ii) $p_{i,j}^h k_h = p_{h,j^*}^i k_i = p_{i^*,h}^j k_j$;
- (iii) $\sum_{j=0}^d p_{i,j}^h = k_i$;
- (iv) $\sum_{r=0}^d p_{e,l}^r p_{m,r}^h = \sum_{t=0}^d p_{m,e}^t p_{t,l}^h$.

The *adjacency matrix* A_i of R_i is the $|X| \times |X|$ matrix whose (x, y) -entry is 1 if $(x, y) \in R_i$, and 0 otherwise. By the *adjacency* or *Bose-Mesner algebra* \mathfrak{A} of \mathfrak{X} we mean the algebra generated by A_0, A_1, \dots, A_d over the complex field. Axioms ??-?? are equivalent to the following:

$$A_0 = I_{|X|}, \quad \sum_{i=0}^d A_i = J_{|X|}, \quad A_i^\top = A_{i^*}, \quad A_i A_j = \sum_{h=0}^d p_{i,j}^h A_h,$$

where $I_{|X|}$ and $J_{|X|}$ are the identity and all-one matrices of order $|X|$, respectively.

We say that \mathfrak{X} is a *P-polynomial* association scheme with respect to the ordering R_0, R_1, \dots, R_d , if there exist some complex coefficient polynomials $v_i(x)$ of degree i ($0 \leq i \leq d$) such that $A_i = v_i(A_1)$.

Distance-regular graphs are equivalent to symmetric association schemes which are *P-polynomial*. As a directed version of distance-regular graph, Damerell [?] introduced the concept of distance-regular digraphs. A digraph Γ of diameter d is said to be *distance-regular* if $(V\Gamma, \{\Gamma_i\}_{0 \leq i \leq d})$ is a non-symmetric association scheme. Notice that, for a distance-regular digraph Γ , $(V\Gamma, \{\Gamma_i\}_{0 \leq i \leq d})$ is a *P-polynomial* association scheme with respect to the ordering $\Gamma_0, \Gamma_1, \dots, \Gamma_d$. On the contrary, every *P-polynomial* non-symmetric association scheme arises from a distance-regular digraph in this way.

In [?], Damerell presented the following basic results concerning distance-regular digraphs.

Theorem 2.2. ([? , Theorems 2 and 4]) *Let Γ be a distance-regular digraph of diameter d and girth g . Then $d = g - 1$ (short type) or $d = g$ (long type). Moreover, if $d = g$, then Γ is a lexicographic product of a short distance-regular digraph of diameter $g - 1$ and an empty graph.*

Let E, F be two nonempty subsets of R . Define

$$EF := \{R_h \mid \sum_{R_i \in E, R_j \in F} p_{i,j}^h \neq 0\}.$$

We write $R_i R_j$ instead of $\{R_i\}\{R_j\}$, and R_i^2 instead of $\{R_i\}\{R_i\}$. If $R_{i^*} R_j \subseteq F$ for any $R_i, R_j \in F$, we say that F is *closed*. For $x \in X$, set $E(x) := \{y \in X \mid (x, y) \in \bigcup_{e \in E} e\}$.

Theorem 2.3. *Let $x \in X$ and F be a closed subset of R . Then $(F(x), \{f \cap ([F(x)] \times [F(x)])\}_{f \in F})$ is an association scheme.*

Proof. This follows from [? , Theorem 1.5.1 (ii)]. □

Let $x \in X$ and F be a closed subset of R . Then $(F(x), \{f \cap ([F(x)] \times [F(x)])\}_{f \in F})$ is called a *subscheme* of \mathfrak{X} . Let $X/F = \{F(x) \mid x \in X\}$, $R_i^F = \{(F(x), F(y)) \mid y \in FR_i F(x)\}$ for $0 \leq i \leq d$ and $R \parallel F = \{R_i^F \mid R_i \in R\}$. By [?, Theorem 4.1.3 (i)], the pair $(X/F, R \parallel F)$ forms an association scheme, called the *quotient scheme* of \mathfrak{X} over F .

Now we recall the definition of a wreath product of two association schemes. This notion appears in the monograph [?] by Weisfeiler, in a more general context of coherent configurations. Let $\mathfrak{Y} = (Y, \{S_j\}_{j=0}^e)$ be an association scheme. Then the *wreath product* of \mathfrak{X} and \mathfrak{Y} denoted by $\mathfrak{X} \wr \mathfrak{Y}$ is defined on the set $X \times Y$ with the relations W_0, W_1, \dots, W_{d+e} satisfying the following:

$$W_0 = \{((x, y), (x, y)) \mid (x, y) \in X \times Y\};$$

$$W_k = \{((x_1, y), (x_2, y)) \mid (x_1, x_2) \in R_k, y \in Y\}, \text{ for } 1 \leq k \leq d;$$

$$W_k = \{((x_1, y_1), (x_2, y_2)) \mid x_1, x_2 \in X, (y_1, y_2) \in S_{k-d}\}, \text{ for } d+1 \leq k \leq d+e.$$

Then, it is easy to see that $\mathfrak{X} \wr \mathfrak{Y} = (X \times Y, \{W_k\}_{k=0}^{d+e})$ is an association scheme.

As a generalization of the wreath product, the wedge product (generalized wreath product) of schemes has been defined by Muzychuk [?] as a generalization of the wedge product of Schur rings. By [?, Theorem 2.3], we have the following equivalent definition.

Let $K \subseteq F$ be two closed subsets of R such that:

- (a) $\sum_{R_i \in K} A_i A_j = \sum_{R_i \in K} k_i A_j = \sum_{R_i \in K} A_j A_i$, for every $R_j \in R \setminus F$;
- (b) $K R_i = R_i K$ for all $0 \leq i \leq d$.

Then \mathfrak{X} is the *wedge product* of subschemes $(F(x), \{f \cap ([F(x)] \times [F(x)])\}_{f \in F})$ for each $x \in X$ and the quotient scheme $(X/K, R \parallel K)$.

Let F be a nonempty closed subset of R . For a digraph Γ with X as its vertex set, the *quotient digraph* of Γ over F , denoted by Γ/F , is defined as the digraph with the vertex set X/F in which $(F(x), F(y))$ is an arc whenever there is an arc in Γ from $F(x)$ to $F(y)$.

Let K be a nonempty subset of R . We define $\langle K \rangle$ to be the smallest closed subset containing K . If $K = \{R_{i_1}, R_{i_2}, \dots, R_{i_l}\}$ for $1 \leq l \leq d+1$, we write $\langle R_{i_1}, R_{i_2}, \dots, R_{i_l} \rangle$ instead of $\langle \{R_{i_1}, R_{i_2}, \dots, R_{i_l}\} \rangle$. We close this section with a result concerning weakly distance-regular digraphs.

Lemma 2.4. *Let \mathfrak{X} be a commutative association scheme such that $R_1^\top = R_2$ and $\langle R_a \rangle = \{R_0, R_a\}$ with $3 \leq a \leq d$. Suppose $p_{1,1}^a = k_1$. Assume $\Delta_l = (X, R_1 \cup R_l)$ for $1 \leq l \leq d$. The following hold:*

- (i) Δ_1 is the lexicographic product of $\Delta_1/\langle R_a \rangle$ and the empty graph \overline{K}_{k_a+1} ;
- (ii) Δ_a is the lexicographic product of $\Delta_a/\langle R_a \rangle$ and the complete graph K_{k_a+1} .

Moreover, if Δ_l is a weakly distance-regular digraph with \mathfrak{X} as its attached scheme for some $l \in \{1, a\}$, then $\Delta_l/\langle R_a \rangle$ is also weakly distance-regular.

Proof. Let $l \in \{1, a\}$ and $\Lambda_l = \Delta_l/\langle R_a \rangle$. We claim that $(\langle R_a \rangle(x), \langle R_a \rangle(y))$ is an arc in Λ_l if and only if (x', y') is an arc in Δ_l for all $x' \in \langle R_a \rangle(x)$ and $y' \in \langle R_a \rangle(y)$. The sufficiency is immediate. Now suppose that $(\langle R_a \rangle(x), \langle R_a \rangle(y))$ is an arc in Λ_l . Then $y \notin \langle R_a \rangle(x)$. Without loss of generality, we may assume that (x, y) is an arc in Δ_l . Since $A \Delta_l = R_1 \cup R_l$ with $l \in \{1, a\}$, we have $(x, y) \in R_1$. Let $y' \in \langle R_a \rangle(y)$. If $y = y'$, then $(x, y') \in R_1$; if $y \neq y'$, then $(y, y') \in R_a$ since $\langle R_a \rangle = \{R_0, R_a\}$, which implies $x \in R_{1^*}(y) = P_{1^*,1}(y, y')$ by the commutativity of \mathfrak{X} . Then $(x, y') \in R_1$. Let $x' \in \langle R_a \rangle(x)$. If $x' = x$, then $(x', y') \in R_1$; if $x' \neq x$, then $(x, x') \in R_a$ since

$\langle R_a \rangle = \{R_0, R_a\}$, which implies $y' \in R_1(x) = P_{1,1^*}(x, x')$, and so $(x', y') \in R_1$. Since $x' \in \langle R_a \rangle(x)$ and $y' \in \langle R_a \rangle(y)$ were arbitrary, the necessity is valid. Thus, our claim is valid.

Since $\langle R_a \rangle = \{R_0, R_a\}$ and $A\Delta_l = R_1 \cup R_l$, from the claim, Δ_1 is the lexicographic product of Λ_1 and the empty graph \overline{K}_{k_a+1} , and Δ_a is the lexicographic product of Λ_a and the complete graph K_{k_a+1} . Thus, ?? and ?? are both valid.

Now suppose that Δ_l is a weakly distance-regular digraph with \mathfrak{X} as its attached scheme for some $l \in \{1, a\}$. Let $x_0, y \in V\Delta_l$ and $y \notin \langle R_a \rangle(x_0)$. Pick a shortest path $(x_0, x_1, \dots, x_{m-1}, y)$ in Δ_l with $m > 0$. Since $\langle R_a \rangle = \{R_0, R_a\}$ and $A\Delta_l = R_1 \cup R_l$, we have $\partial_{\Delta_l}(\langle R_a \rangle(x_0), \langle R_a \rangle(y)) \leq \partial_{\Delta_l}(x_0, y)$.

Pick a shortest path $(\langle R_a \rangle(x_0), \langle R_a \rangle(x_1), \dots, \langle R_a \rangle(x_{n-1}), \langle R_a \rangle(y))$ in Λ_l with $n > 0$. By the claim, $(x_0, x_1, \dots, x_{n-1}, y)$ is a path in Δ_l . It follows that $\partial_{\Delta_l}(x_0, y) \leq \partial_{\Lambda_l}(\langle R_a \rangle(x_0), \langle R_a \rangle(y))$. Thus, $\partial_{\Delta_l}(x_0, y) = \partial_{\Lambda_l}(\langle R_a \rangle(x_0), \langle R_a \rangle(y))$.

Let $\tilde{i}, \tilde{j}, \tilde{h} \in \tilde{\partial}(\Delta_l) \setminus \{(0, 0)\}$ and $(\langle R_a \rangle(u), \langle R_a \rangle(v)) \in (\Delta_l)_{\tilde{h}}$. Since $\partial_{\Delta_l}(x, y) = \partial_{\Lambda_l}(\langle R_a \rangle(x), \langle R_a \rangle(y))$ for all $x, y \in X$ and $y \notin \langle R_a \rangle(x)$, we have $\tilde{i}, \tilde{j} \in \tilde{\partial}(\Delta_l)$ and $(u, v) \in (\Delta_l)_{\tilde{h}}$. Since \mathfrak{X} is the attached scheme of Δ_l , we may assume $(\Delta_l)_{\tilde{i}} = R_i$, $(\Delta_l)_{\tilde{j}} = R_j$ and $(u, v) \in R_h$ for some $i, j, h \in \{1, 2, \dots, a-1, a+1, \dots, d\}$. Since $\partial_{\Delta_l}(x, y) = \partial_{\Lambda_l}(\langle R_a \rangle(x), \langle R_a \rangle(y))$ for all $x, y \in X$ and $y \notin \langle R_a \rangle(x)$, one gets

$$\begin{aligned} R_i(u) \cap R_{j^*}(v) &= (\Delta_l)_{\tilde{i}}(u) \cap (\Delta_l)_{\tilde{j}^*}(v) \\ &= \cup_{w \in (\Delta_l)_{\tilde{i}}(u) \cap (\Delta_l)_{\tilde{j}^*}(v)} \langle R_a \rangle(w) \\ &= \cup_{\langle R_a \rangle(w) \in (\Delta_l)_{\tilde{i}}(\langle R_a \rangle(u)) \cap (\Delta_l)_{\tilde{j}^*}(\langle R_a \rangle(v))} \langle R_a \rangle(w). \end{aligned}$$

Since $|\langle R_a \rangle(w)| = k_a + 1$ for all $w \in X$, we have

$$|(\Delta_l)_{\tilde{i}}(\langle R_a \rangle(u)) \cap (\Delta_l)_{\tilde{j}^*}(\langle R_a \rangle(v))| = \frac{p_{i,j}^h}{k_a + 1}.$$

Thus, Λ_l is also a weakly distance-regular digraph. \square

3. PROOF OF THEOREM ??

In this section, we always assume that $\Gamma := (X, A\Gamma)$ is a weakly distance-regular digraph and $\mathfrak{X} := (X, \{R_i\}_{i=0}^d)$ is an association scheme with $R_1^\top = R_2$ and $R_h^\top = R_h$ for $3 \leq h \leq d$. Before giving a proof of Theorem ??, we need some auxiliary facts and lemmas.

Fact 3.1. *If \mathfrak{X} is the attached scheme of Γ , then $R_i \subseteq A\Gamma$ or $R_i \cap A\Gamma = \emptyset$ for all $1 \leq i \leq d$.*

Fact 3.2. *The digraph $(X, (A\Gamma)^\top)$ is weakly distance regular, and has the same attached scheme with Γ .*

In the remainder of this paper, we may assume that Γ has diameter 2. Since

$$\{(0, 0), (1, 2), (2, 1)\} \subseteq \tilde{\partial}(\Gamma) \subseteq \{(0, 0), (1, 2), (2, 1), (1, 1), (2, 2)\}, \quad (1)$$

its attached scheme has at most 4 classes. By a theorem of Higman [?], association schemes with at most four classes are commutative.

3.1. The attached scheme of Γ has 2 classes. In this subsection, we consider the case that the attached scheme of Γ has 2 classes.

Lemma 3.3. *If $d = 2$, then (X, R_i) is a weakly distance-regular digraph of diameter 2 with \mathfrak{X} as the attached scheme for $i \in \{1, 2\}$.*

Proof. Let $\Delta = (X, R_1)$. Since $A\Delta = R_1$ and R_1 is non-symmetric, we have $R_1 = \Delta_{(1,a)}$ with $a > 1$, which implies that $R_2 = \Delta_{(a,1)}$. Since $d = 2$ and $2 = 1^*$, by setting $(i, j) = (1, 1)$ and $(i, j) = (1, 1^*)$ in Lemma ?? ?? respectively, we have $k_1 k_1 = p_{1,1}^1 k_1 + p_{1,1}^{1^*} k_{1^*}$ and $k_1 k_{1^*} = k_1 + p_{1,1^*}^1 k_1 + p_{1,1^*}^{1^*} k_{1^*}$. The fact $k_1 = k_{1^*}$ implies $k_1 = p_{1,1}^1 + p_{1,1}^{1^*} = 1 + p_{1,1^*}^1 + p_{1,1^*}^{1^*}$. In view of Lemma ?? ??, one gets $p_{1,1}^1 = p_{1,1^*}^{1^*}$, and so $p_{1,1}^{1^*} = 1 + p_{1,1^*}^1 \geq 1$. This implies $R_2 = R_{1^*} \in R_1^2$, and so $a = 2$. Thus, Δ is a weakly distance-regular digraph of diameter 2 with the attached scheme

$$(X, \{\Delta_{(0,0)}, \Delta_{(1,2)}, \Delta_{(2,1)}\}).$$

The desired result follows from Fact ??.

□

Lemma 3.4. *Suppose that \mathfrak{X} is the attached scheme of Γ . Then $d = 2$, and $\Gamma = (X, R_i)$ for some $i \in \{1, 2\}$.*

Proof. Since the attached scheme of Γ has 2 classes, we have $d = 2$. Then it is immediately from Fact ?? and Lemma ??.

□

3.2. The attached scheme of Γ has 3 classes. In this subsection, we consider the case that the attached scheme of Γ has 3 classes.

Lemma 3.5. *If $d = 3$, then $p_{1,1}^{1^*} \neq 0$ or $p_{1,1}^3 \neq 0$.*

Proof. Suppose, to the contrary that $p_{1,1}^{1^*} = 0$ and $p_{1,1}^3 = 0$. Then $R_1^2 = \{R_1\}$. By setting $(i, j) = (1, 1)$ in Lemma ?? ??, we obtain $k_1^2 = p_{1,1}^1 k_1$, and so $p_{1,1}^1 = k_1$. Then $p_{1,1}^1 + p_{1,0}^1 = k_1 + 1$, contrary to Lemma ?? ??. □

Lemma 3.6. *If $d = 3$, then $(X, R_i \cup R_3)$ is a weakly distance-regular digraph of diameter 2 with \mathfrak{X} as the attached scheme for $i \in \{1, 2\}$.*

Proof. According to Fact ??, it suffices to prove that $\Delta = (X, R_1 \cup R_3)$ is a weakly distance-regular digraph of diameter 2.

Since R_1 is non-symmetric and R_3 is symmetric, we have $R_1 = \Delta_{(1,a)}$ with $a > 1$ and $R_3 = \Delta_{(1,1)}$.

Suppose $a > 2$. Since $A\Delta = R_1 \cup R_3$, one has $p_{1,1}^{1^*} = p_{1,3}^{1^*} = 0$. By Lemma ?? ??, we get $p_{1,1}^3 = p_{1,3}^{1^*} = 0$, contrary to Lemma ??.

Note that $a = 2$, and so $R_1 = \Delta_{(1,2)}$. It follows that Δ is a weakly distance-regular digraph of diameter 2 with the attached scheme $(X, \{\Delta_{(0,0)}, \Delta_{(1,2)}, \Delta_{(2,1)}, \Delta_{(1,1)}\})$. Thus, the desired result follows. □

Lemma 3.7. *Suppose $d = 3$. If \mathfrak{X} is neither a P -polynomial association scheme nor a wedge product of at least two 2-class association schemes with the same parameters and a 1-class association scheme, then (X, R_i) is a weakly distance-regular digraph of diameter 2 with \mathfrak{X} as the attached scheme for $i \in \{1, 2\}$.*

Proof. Let $\Delta = (X, R_1)$. Since R_1 is non-symmetric and R_3 is symmetric, we have $R_1 = \Delta_{(1,a)}$ with $a > 1$ and $R_3 = \Delta_{(b,b)}$ with $b > 1$. If $a = b = 2$, then Δ is a weakly distance-regular digraph of diameter 2 with the attached scheme $(X, \{\Delta_{(0,0)}, \Delta_{(1,2)}, \Delta_{(2,1)}, \Delta_{(2,2)}\})$, which implies that the desired result follows from Fact ??. Therefore, it suffices to show that $a = b = 2$.

Suppose, to the contrary that $a > 2$ or $b > 2$. Since $p_{1,1}^{1^*} \neq 0$ or $p_{1,1}^3 \neq 0$ from Lemma ??, we have $a > 2$ and $b = 2$, or $a = 2$ and $b > 2$. Note that the induced subdigraph of Δ on the vertex set $\langle R_1 \rangle(x)$ is strongly connected for all $x \in X$. If $a > 2$ and $b = 2$, then $a = 3$ and $(X, \{\Delta_i\}_{i=0}^3)$ is a non-symmetric association scheme

since $\tilde{\partial}(\Delta) = \{(0, 0), (1, a), (a, 1), (2, 2)\}$, which imply that Δ is a distance-regular digraph of girth 4, contrary to the fact that \mathfrak{X} is not P -polynomial.

Note that $a = 2$ and $b > 2$. The fact $R_3 = \Delta_{(b,b)}$ implies $p_{1,1}^3 = 0$. Since $d = 3$ and $2 = 1^*$, by setting $(i, j) = (1, 1)$ in Lemma ?? ??, we have $k_1^2 = p_{1,1}^1 k_1 + p_{1,1}^{1*} k_{1^*}$, and so $k_1 = p_{1,1}^1 + p_{1,1}^{1*}$.

In view of Lemma ?? ??, we obtain $p_{1,1}^1 = p_{1,1}^{1*} = p_{1,1}^{1*}$. By setting $(i, j) = (1, 1^*)$ in Lemma ?? ??, we get $k_1^2 = k_1 k_{1^*} = k_1 + 2p_{1,1}^1 k_1 + p_{1,1}^{3*} k_3$. By substituting $k_1 = p_{1,1}^1 + p_{1,1}^{1*}$ and $p_{1,1}^{3*} k_3 = p_{3,1}^1 k_1$ from Lemma ?? ?? into the above equation respectively, one obtains

$$(p_{1,1}^{1*})^2 - p_{1,1}^1 - p_{1,1}^{1*} - (p_{1,1}^1)^2 - p_{1,1}^{3*} k_3 = 0 \quad (2)$$

and $p_{1,3}^1 = k_1 - 2p_{1,1}^1 - 1$.

By setting $(h, m, e, l) = (1, 1, 1, 1^*)$ in Lemma ?? ??, we have

$$k_1 + p_{1,1}^{1*} p_{1,1}^{1*} + p_{1,1}^{3*} p_{1,3}^1 = p_{1,1}^{1*} p_{1,1}^{1*}.$$

Since $p_{1,1}^1 = p_{1,1}^{1*} = p_{1,1}^{1*}$ and $p_{1,1}^{1*} = p_{1,1}^{1*}$ from Lemma ?? ??, we obtain

$$k_1 + (p_{1,1}^1)^2 + p_{1,1}^{3*} p_{1,3}^1 = (p_{1,1}^{1*})^2.$$

Substituting $k_1 = p_{1,1}^1 + p_{1,1}^{1*}$ and $p_{1,3}^1 = k_1 - 2p_{1,1}^1 - 1$ into the above equation, one has

$$(p_{1,1}^1 - p_{1,1}^{1*} + 1)(p_{1,1}^1 - p_{1,1}^{3*} + p_{1,1}^{1*}) = 0.$$

Case 1. $p_{1,1}^{3*} = p_{1,1}^1 + p_{1,1}^{1*}$.

Since $k_1 = p_{1,1}^1 + p_{1,1}^{1*}$, we have $p_{1,1}^{3*} = k_1$. Pick $(x, y), (y, z) \in R_3$. Suppose $(x, z) \notin R_0 \cup R_3$. Then $(x, z) \in R_1 \cup R_{1^*}$. Since the proofs are similar, we may assume $(x, z) \in R_1$. Since $p_{1,1}^{3*} = k_1$, we have $z \in R_1(x) = P_{1,1^*}(x, y)$, contrary to the fact that $(y, z) \in R_3$. Then $(x, z) \in R_0 \cup R_3$. Since $(x, y), (y, z) \in R_3$ were arbitrary, one has $R_3^2 \subseteq \{R_0, R_3\}$. By induction, we get $R_3^h \subseteq \{R_0, R_3\}$ for $h > 0$, which implies $\langle R_3 \rangle = \{R_0, R_3\}$.

Let $\Lambda = \Delta / \langle R_3 \rangle$. In view of Lemma ?? ??, Δ is the lexicographic product of Λ and the empty graph \overline{K}_{k_3+1} . Since $R_3 = \Delta_{(b,b)}$, we get $\tilde{\partial}(\Lambda) = \tilde{\partial}(\Delta) \setminus \{(b, b)\} = \{(0, 0), (1, 2), (2, 1)\}$. By Theorem ?? and Lemma ??, Λ is a short distance-regular digraph of girth 3. In view of the definition of the lexicographic product, one gets $b = 3$. Since $(X, \{\Delta_i\}_{i=0}^3)$ is a non-symmetric association scheme, from Theorem ?? again, Δ is a long distance-regular digraph of girth 3, contrary to the fact that \mathfrak{X} is not P -polynomial.

Case 2. $p_{1,1}^{1*} = p_{1,1}^1 + 1$.

Substituting $p_{1,1}^{1*} = p_{1,1}^1 + 1$ into (??), one obtains $-p_{1,1}^{3*} k_3 = 0$. It follows that $p_{1,1}^{3*} = 0$, and so $R_1 R_{1^*} \subseteq \{R_0, R_1, R_{1^*}\}$. Since $p_{1,1}^3 = 0$, we obtain $R_1^2 \subseteq \{R_1, R_{1^*}\}$. By induction, we get $R_1^h \subseteq \{R_0, R_1, R_{1^*}\}$ for $h > 0$, and so $\langle R_1 \rangle = \{R_0, R_1, R_{1^*}\}$. Then $p_{l,1}^3 = p_{l,1}^{3*} = 0$ for $l \in \{0, 1, 2\}$. By Lemma ?? ??, we have $k_l = p_{l,3}^3$ for $l \in \{0, 1, 2\}$, which implies $A_l A_3 = k_l A_3$ from Lemma ?? ??. Thus, $\sum_{l=0}^2 A_l A_3 = \sum_{l=0}^2 k_l A_3$.

In view of Theorem ??, $(\langle R_1 \rangle(x), \{R_i \cap ([\langle R_1 \rangle(x)] \times [\langle R_1 \rangle(x)])\}_{i=0}^2)$ is a subscheme of \mathfrak{X} for $x \in X$. By the definition of the quotient scheme, $(X / \langle R_1 \rangle, \{R_i\}_{i=0}^3 // \langle R_1 \rangle)$ is a 1-class association scheme. By the commutativity of \mathfrak{X} , \mathfrak{X} is a wedge product of subschemes $(\langle R_1 \rangle(x), \{R_i \cap ([\langle R_1 \rangle(x)] \times [\langle R_1 \rangle(x)])\}_{i=0}^2)$ for each $x \in X$ and a 1-class association scheme, a contradiction. \square

Lemma 3.8. *Suppose that \mathfrak{X} is the attached scheme of Γ . Then $d = 3$, and exactly one of the following holds:*

- (i) $A\Gamma = R_i \cup R_3$;
- (ii) $A\Gamma = R_i$, and \mathfrak{X} is neither a P -polynomial association scheme nor a wedge product of at least two 2-class association schemes with the same parameters and a 1-class association scheme.

Here, $i \in \{1, 2\}$.

Proof. Since the attached scheme of Γ has 3 classes, we have $d = 3$. Then $\tilde{\partial}(\Gamma) = \{(0, 0), (1, 2), (2, 1), (2, 2)\}$ or $\tilde{\partial}(\Gamma) = \{(0, 0), (1, 2), (2, 1), (1, 1)\}$. It follows from Fact ?? that $A\Gamma = R_i$ or $A\Gamma = R_i \cup R_3$ for $i \in \{1, 2\}$. If $A\Gamma = R_i \cup R_3$ for some $i \in \{1, 2\}$, from Lemma ??, then ?? holds.

Suppose $A\Gamma = R_i$ for some $i \in \{1, 2\}$. Note that R_1 and R_2 are the only pair of non-symmetric relations. If \mathfrak{X} is a P -polynomial association scheme, from Theorem ??, then Γ is a short distance-regular digraph of girth 4 or a long distance-regular digraph of girth 3, which implies that Γ has diameter 3, a contradiction. If \mathfrak{X} is a wedge product of at least two 2-class association schemes with the same parameters and a 1-class association scheme, then $\langle R_1 \rangle \neq \{R_i\}_{i=0}^3$, which implies that Γ is not strongly connected, a contradiction. By Lemma ??, ?? is valid.

This completes the proof of this lemma. \square

3.3. The attached scheme of Γ has 4 classes. In this subsection, we consider the case that the attached scheme of Γ has 4 classes.

Lemma 3.9. *Suppose $d = 4$ and $j \in \{3, 4\}$. If \mathfrak{X} is neither a wreath product of a 1-class association scheme and a P -polynomial association scheme nor a wedge product of subschemes $(\langle R_1, R_j \rangle(x), \{R_h \cap ([\langle R_1, R_j \rangle(x)] \times [\langle R_1, R_j \rangle(x)])\}_{h \in \{0, 1, 2, j\}})$ for each $x \in X$ and a 1-class association scheme, then $(X, R_i \cup R_j)$ is a weakly distance-regular digraph of diameter 2 with \mathfrak{X} as the attached scheme for $i \in \{1, 2\}$.*

Proof. Let $\Delta = (X, R_1 \cup R_j)$. Since the proofs are similar, we may assume that $j = 3$. Since R_1 is non-symmetric and R_3, R_4 are symmetric, we have $R_1 = \Delta_{(1,a)}$ with $a > 1$, $R_3 = \Delta_{(1,1)}$ and $R_4 = \Delta_{(b,b)}$ with $b > 1$. If $a = b = 2$, then Δ is a weakly distance-regular digraph of diameter 2 with the attached scheme $(X, \{\Delta_{(0,0)}, \Delta_{(1,2)}, \Delta_{(2,1)}, \Delta_{(1,1)}, \Delta_{(2,2)}\})$, which implies that the desired result is valid from Fact ??. Suppose, to the contrary that $a > 2$ or $b > 2$.

Case 1. $a > 2$.

Since $A\Delta = R_1 \cup R_3$, one has $p_{1,1}^{1*} = p_{1,3}^{1*} = p_{3,3}^{1*} = 0$. In view of Lemma ?? ??, we have $p_{1,1}^{3*} = p_{1,3}^{1*} = 0$. By setting $(i, j) = (1, 1)$ in Lemma ?? ??, we obtain

$$k_1^2 = p_{1,1}^1 k_1 + p_{1,1}^4 k_4. \quad (3)$$

If $b > 2$, then $p_{1,1}^4 = 0$, and so $k_1^2 = p_{1,1}^1 k_1$. This implies that $p_{1,1}^1 = k_1$, and so $p_{1,1}^1 + p_{1,0}^1 = k_1 + 1$, contrary to Lemma ?? ??. Therefore, $b = 2$. Note that the induced subdigraph of Δ on the vertex set $\langle R_1, R_3 \rangle(x)$ is strongly connected for all $x \in X$. Since $\tilde{\partial}(\Delta) = \{(0, 0), (1, a), (a, 1), (1, 1), (2, 2)\}$, one gets $a = 3$. Then Δ is a weakly distance-regular digraph with the attached scheme

$$(X, \{\Delta_{(0,0)}, \Delta_{(1,3)}, \Delta_{(3,1)}, \Delta_{(1,1)}, \Delta_{(2,2)}\}).$$

By setting $(i, j) = (1, 1^*)$ in Lemma ?? ??, we obtain $k_1^2 = k_1 k_{1^*} = k_1 + p_{1,1^*}^1 k_1 + p_{1,1^*}^{1^*} k_{1^*} + p_{1,1^*}^3 k_3 + p_{1,1^*}^4 k_4$. Lemma ?? ?? implies $p_{1,1}^1 = p_{1,1^*}^{1^*}$. By (??), we have

$$p_{1,1}^4 k_4 = k_1^2 - p_{1,1}^1 k_1 = k_1 + p_{1,1^*}^1 k_1 + p_{1,1^*}^3 k_3 + p_{1,1^*}^4 k_4, \quad (4)$$

which implies $p_{1,1}^4 \neq 0$.

Since $p_{1,1}^{1^*} = p_{1,1}^3 = 0$, by setting $(h, m, e, l) = (3, 1, 1, 1^*)$ in Lemma ?? ??, we get

$$p_{1,1^*}^{1^*} p_{1,1^*}^3 + p_{1,1^*}^3 p_{1,3}^3 + p_{1,1^*}^4 p_{1,4}^3 = p_{1,1}^1 p_{1,1^*}^3 + p_{1,1}^4 p_{4,1^*}^3. \quad (5)$$

In view of Lemma ?? ??, one has $p_{1,3}^3 = p_{3,3}^{1^*} = 0$. By substituting $p_{1,3}^3 = 0$, $p_{1,4}^3 = p_{4,1^*}^3$ and $p_{1,1}^1 = p_{1,1^*}^{1^*}$ from Lemma ?? ?? into (??), one obtains $p_{1,1^*}^4 p_{1,4}^3 = p_{1,1}^4 p_{1,4}^3$. If $p_{1,4}^3 \neq 0$, then $p_{1,1^*}^4 = p_{1,1}^4$, which implies $k_1 + p_{1,1^*}^1 k_1 + p_{1,1^*}^3 k_3 = 0$ from (??), a contradiction. Thus $p_{1,4}^3 = 0$.

Since $p_{1,1}^{1^*} = p_{1,1}^3 = 0$, by setting $(h, m, e, l) = (3, 1, 1, 3)$ in Lemma ?? ??, we get

$$p_{1,3}^{1^*} p_{1,1^*}^3 + p_{1,3}^3 p_{1,3}^3 + p_{1,3}^4 p_{1,4}^3 = p_{1,1}^1 p_{1,3}^3 + p_{1,1}^4 p_{4,3}^3.$$

By substituting $p_{1,3}^{1^*} = p_{1,3}^3 = p_{1,4}^3 = 0$ into the above equation, we obtain $p_{1,1}^4 p_{4,3}^3 = 0$. Since $p_{1,1}^4 \neq 0$, one obtains $p_{4,3}^3 = 0$. In view of Lemma ?? ??, we have $p_{3,3}^4 = p_{4,3}^3 = 0$. Since $p_{3,3}^{1^*} = p_{3,3}^3 = 0$, one gets $R_3^2 \subseteq \{R_0, R_3\}$. By induction, we get $R_3^h \subseteq \{R_0, R_3\}$ for $h > 0$, which implies $\langle R_3 \rangle = \{R_0, R_3\}$. Since $p_{1,1}^3 = p_{1,3}^3 = p_{1,4}^3 = 0$, we have $p_{1,1^*}^3 = k_1$ by Lemma ?? ??.

Let $\Lambda = \Delta / \langle R_3 \rangle$. Note that Δ is a weakly distance-regular digraph of $\tilde{\partial}(\Delta) = \{(0, 0), (1, 3), (3, 1), (1, 1), (2, 2)\}$. By Lemma ?? ??, Δ is a lexicographic product of Λ and a complete graph, where Λ is a weakly distance-regular digraph. Since $R_3 = \Delta_{(1,1)}$, from the definition of the lexicographic product, one gets $\tilde{\partial}(\Lambda) = \{(0, 0), (1, 3), (3, 1), (2, 2)\}$. By Theorem ??, Λ is a short distance-regular digraph of girth 4. It follows that \mathfrak{X} is a wreath product of a 1-class association scheme and a P -polynomial association scheme, a contradiction.

Case 2. $a = 2$.

Note that $b > 2$. The fact $R_4 = \Delta_{(b,b)}$ implies $p_{1,1}^4 = p_{1,3}^4 = p_{3,3}^4 = 0$. By setting $(h, m, e, l) = (4, 1, 3, 3)$ in Lemma ?? ??, we have

$$p_{3,3}^1 p_{1,1}^4 + p_{3,3}^{1^*} p_{1,1^*}^4 + p_{3,3}^3 p_{1,3}^4 + p_{3,3}^4 p_{1,4}^4 = p_{1,3}^1 p_{1,3}^4 + p_{1,3}^{1^*} p_{1^*,3}^4 + p_{1,3}^3 p_{3,3}^4 + p_{1,3}^4 p_{4,3}^4,$$

and so $p_{3,3}^1 p_{1,1^*}^4 = p_{1,3}^{1^*} p_{1^*,3}^4$. In view of Lemma ?? ??, one has $p_{1^*,3}^4 = p_{1,3}^4 = 0$, which implies $p_{3,3}^1 p_{1,1^*}^4 = 0$.

Case 2.1. $p_{1,1^*}^4 = 0$.

Since $p_{1,1}^4 = p_{1,3}^4 = p_{3,3}^4 = 0$, we obtain

$$\{R_1, R_3\}^2 = R_1^2 \cup R_1 R_3 \cup R_3^2 \subseteq \{R_0, R_1, R_1^*, R_3\}.$$

Since $p_{1,1^*}^4 = 0$, by induction, we get $(\{R_1, R_3\})^h \subseteq \{R_0, R_1, R_1^*, R_3\}$ for $h > 0$, which implies $\langle R_1, R_3 \rangle = \{R_0, R_1, R_1^*, R_3\}$. Then $p_{l,1}^4 = p_{l,1^*}^4 = p_{l,3}^4 = 0$ for $l \in \{0, 1, 2, 3\}$. By Lemma ?? ??, we have $k_l = p_{l,4}^4$ for $l \in \{0, 1, 2, 3\}$, which implies $A_l A_4 = k_l A_4$ from Lemma ?? ???. Thus, $\sum_{l=0}^3 A_l A_4 = \sum_{l=0}^3 k_l A_4$.

By Theorem ??, $(\langle R_1, R_3 \rangle(x), \{R_h \cap ([\langle R_1, R_3 \rangle(x)] \times [\langle R_1, R_3 \rangle(x)])\}_{h=0}^3)$ is a subscheme of \mathfrak{X} for every $x \in X$. In view of the definition of the quotient scheme, $(X / \langle R_1, R_3 \rangle, \{R_i\}_{i=0}^4 // \langle R_1, R_3 \rangle)$ is a 1-class association scheme. By the commutativity of \mathfrak{X} , \mathfrak{X} is a wedge product of subschemes $(\langle R_1, R_3 \rangle(x), \{R_h \cap ([\langle R_1, R_3 \rangle(x)] \times [\langle R_1, R_3 \rangle(x)])\}_{h=0}^3)$ for each $x \in X$ and a 1-class association scheme, a contradiction.

Case 2.2. $p_{1,1^*}^4 \neq 0$.

Let $(x, y) \in R_4$. Since $p_{1,1^*}^4 \neq 0$, there exists $z \in P_{1,1^*}(x, y)$. The fact $R_1 = \Delta_{(1,2)}$ implies $b = \partial_\Delta(x, y) \leq \partial_\Delta(x, z) + \partial_\Delta(z, y) = 3$, and so $R_4 = \Delta_{(3,3)}$. It follows that Δ is a weakly distance-regular digraph with $\tilde{\partial}(\Delta) = \{(0, 0), (1, 2), (2, 1), (1, 1), (3, 3)\}$.

Since $p_{3,3}^1 p_{1,1^*}^4 = 0$, one has $p_{3,3}^1 = 0$. In view of Lemma ?? ??, we obtain $p_{3,3}^1 = 0$. Since $p_{3,3}^4 = 0$, one gets $R_3^2 \subseteq \{R_0, R_3\}$. By induction, we get $R_3^h \subseteq \{R_0, R_3\}$ for $h > 0$, which implies $\langle R_3 \rangle = \{R_0, R_3\}$. By setting $(h, m, e, l) = (4, 1, 1, 3)$ in Lemma ?? ??, we have

$$p_{1,3}^1 p_{1,1}^4 + p_{1,3}^1 p_{1,1^*}^4 + p_{1,3}^3 p_{1,3}^4 + p_{1,3}^4 p_{1,4}^4 = p_{1,1}^1 p_{1,3}^4 + p_{1,1}^1 p_{1^*,3}^4 + p_{1,1}^3 p_{3,3}^4 + p_{1,1}^4 p_{4,3}^4.$$

By substituting $p_{1,1}^4 = p_{1,3}^4 = p_{1^*,3}^4 = p_{3,3}^4 = 0$ into the above equation, we obtain $p_{1,3}^1 p_{1,1^*}^4 = 0$, which implies $p_{1,3}^1 = 0$. In view of Lemma ?? ??, we have $p_{1,1}^3 = p_{1,3}^1 = 0$, $p_{1,3}^3 = p_{3,3}^1 = 0$ and $p_{1,4}^3 = p_{4,3}^1 = 0$. By setting $(i, h) = (1, 3)$ in Lemma ?? ??, one obtains $p_{1,1^*}^3 = k_1$.

Let $\Lambda = \Delta / \langle R_3 \rangle$. Note that Δ is a weakly distance-regular digraph with $\tilde{\partial}(\Delta) = \{(0, 0), (1, 2), (2, 1), (1, 1), (3, 3)\}$. By Lemma ?? ??, Δ is a lexicographic product of Λ and a complete graph, where Λ is weakly distance-regular. Since $R_3 = \Delta_{(1,1)}$, from the definition of the lexicographic product, we have $\tilde{\partial}(\Lambda) = \{(0, 0), (1, 2), (2, 1), (3, 3)\}$. By Theorem ??, Λ is a long distance-regular digraph of girth 3. Thus, \mathfrak{X} is a wreath product of a 1-class association scheme and a P -polynomial association scheme, a contradiction. \square

Lemma 3.10. *Suppose that \mathfrak{X} is the attached scheme of Γ . Then $d = 4$, $A\Gamma = R_i \cup R_j$ for some $i \in \{1, 2\}$ and $j \in \{3, 4\}$, and \mathfrak{X} is neither*

- (i) *a wreath product of a 1-class association scheme and a P -polynomial association scheme, nor*
- (ii) *a wedge product of subschemes*

$$(\langle R_1, R_j \rangle(x), \{R_h \cap ([\langle R_1, R_j \rangle(x)] \times [\langle R_1, R_j \rangle(x)])\}_{h \in \{0,1,2,j\}})$$

for each $x \in X$ and a 1-class association scheme.

Proof. Since the attached scheme of Γ has 4 classes, we have $d = 4$. Then $\tilde{\partial}(\Gamma) = \{(0, 0), (1, 2), (2, 1), (1, 1), (2, 2)\}$. It follows from Fact ?? that $\Gamma = (X, R_i \cup R_j)$ for some $i \in \{1, 2\}$ and $j \in \{3, 4\}$.

If \mathfrak{X} is a wreath product of a 1-class association scheme and a P -polynomial association scheme, by Theorem ??, then Γ is a lexicographic product of a short distance-regular digraph of girth 4 or a long distance-regular digraph of girth 3, and a complete graph, which implies that Γ has diameter more than 2, a contradiction. If \mathfrak{X} is a wedge product of subschemes $(\langle R_1, R_j \rangle(x), \{R_h \cap ([\langle R_1, R_j \rangle(x)] \times [\langle R_1, R_j \rangle(x)])\}_{h \in \{0,1,2,j\}})$ for each $x \in X$ and a 1-class association scheme, then $\langle R_1, R_j \rangle \neq \{R_h\}_{h=0}^4$, which implies that Γ is not strongly connected, a contradiction.

The desired result follows from Lemma ??.

\square

Now, we are ready to prove Theorem ??.

Proof of Theorem ??. We first prove the sufficiency. By Lemmas ??, ??, ?? and ??, (X, A) is a weakly distance-regular digraph of diameter 2 with \mathfrak{X} as its attached scheme, as desired.

We next prove the necessity. In view of (??), the attached scheme of (X, A) has exactly one pair of non-symmetric relations, and has at most 4 classes. Without loss

of generality, we may assume $R_1^\top = R_2$. Therefore, by Lemmas ??, ?? and ??, the desired result follows. \square

DATA AVAILABILITY STATEMENT

No data was used for the research described in the article.

CONFLICT OF INTEREST

The authors declare that they have no conflict of interest.

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