

# Condorcet Domains

Clay Thomas  
claytont@cs.princeton.edu

Corey Sinnamon  
sinnamon@cs.princeton.edu

June 17, 2019

## 1 Introduction

Arrow’s impossibility theorem says voting doesn’t work out like you want it to. Condorcet domains work around this.

**Theorem 1.** *Let  $\mathcal{C} \subseteq S_n$  be some set of linear orders. The following are equivalent:*

1. *Simple majority voting over any profile of voters is always acyclic.*
2. *Simple majority voting over any odd-size profile of voters always yields a total order.*
3. *Every triple of outcomes  $i, j, k \in [n]$  is acyclic, meaning that  $\mathcal{C}$  restricted to  $\{i, j, k\}$  does not contain the orders  $ijk$ ,  $jki$ , and  $kij$ .*
4. *Among every triple  $i, j, k \in [n]$  of distinct outcomes in  $[n]$ , one of them is either never first, never last, or never in the middle.*

The above theorem is not too hard, but it’s very standard so we do not prove it here. Some intuition/a guide to the proof is the following: majority voting being acyclic is equivalent to it being *transitive*, and being transitive is simply a property of triples of outcomes ( $a > b$  and  $b > c$  implies  $a > c$ ).

**Definition 2.** *A Condorcet domain  $\mathcal{C} \subseteq S_n$  is any set of orders satisfying any of the above conditions.*

## 2 Preliminaries

If  $\mathcal{C}$  is a Condorcet domain, theorem 1 part 4 tells us that every triple  $i, j, k$  of outcomes falls into some “value restriction case”: one of  $i$ ,  $j$ , or  $k$  is either never first, never second, or never third among  $\{i, j, k\}$ . In those cases, if  $u$  is the “restricted” outcome among  $\{i, j, k\}$ , we write  $uN_1$ ,  $uN_2$ , or  $uN_3$ , respectively. If the triple  $\{i, j, k\}$  is not clear from the context, we will write  $uN_i v w$  for  $\{v, w\} = \{i, j, k\} \setminus \{u\}$ .

**Definition 3.** *A value restriction (VR) casting on  $[n]$  is a function  $\psi$  from triples of outcomes  $(i, j, k) \in \binom{[n]}{3}$  to one of the nine “value restriction cases”  $uN_v$  for  $u \in \{i, j, k\}$  and  $v \in \{1, 2, 3\}$ .*

*The corresponding domain  $\mathcal{D}(\psi)$  of some value restriction is the collection of all orders  $\sigma \in S_n$  which satisfy  $\psi$  on every triple. That is,  $\sigma \in \mathcal{D}(\psi)$  if for all  $i, j, k \in [n]$  with  $\psi(i, j, k) = uN_v$ ,  $\sigma|_{i,j,k}$  does not have  $u$  in position  $v$ .*

CDs, even closed CDs, are not in general defined by the VR castings they satisfy<sup>1</sup>. However, maximal CDs are always the corresponding domain of a VR casting:

**Proposition 4.** *If  $\mathcal{C}$  is a maximal CD, then  $\mathcal{C} = \mathcal{D}(\psi)$  for some  $\psi$ .*

*Proof.* By theorem 1 part 4, every CD is contained in  $\mathcal{D}(\psi)$  for some  $\psi$ . If  $\mathcal{C} \subsetneq \mathcal{D}(\psi)$ , then  $\mathcal{C}$  is not maximal.  $\square$

**Example 5.** *When  $n = 3$ , there are three maximal CDs up to relabeling. They are important enough to give a name ((AND A PICTURE IN THE FINAL VERSION)).*

**Proposition 6.** *For any  $\psi$ , the domain  $\mathcal{D}(\psi)$  is median-closed.*

*Proof.* Observe that medians are preserved under taking restrictions, that is, for any  $S \subseteq [n]$ ,  $\text{med}(p, q, r)|_S = \text{med}(p|_S, q|_S, r|_S)$ . For any triple  $T = \{i, j, k\}$ , we have  $\mathcal{D}|_T \subseteq \mathcal{D}(\psi(\{i, j, k\}))$ , which turn out to be a *representative* domain on three alternatives. ((THIS IS JUMPING A BIT AHEAD)). Thus,  $\text{med}(p|_T, q|_T, r|_T)$  will always lie in  $\mathcal{D}|_T$  when  $p, q, r \in \mathcal{D}$ , so  $\text{med}(p, q, r)$  will satisfy  $\psi$  on triple  $i, j, k$  ((A BIT INFORMAL HERE)).  $\square$

**Definition 7.** *We distinguish the two natural orders  $\alpha = 12 \dots n$  and  $\omega = n(n-1) \dots 1$ .*

*The reversed order of any  $\sigma \in S_n$  is given by  $\text{rev}(\sigma) = \sigma(n)\sigma(n-1) \dots \sigma(1)$*

*A domain  $\mathcal{C}$  is maximal width if  $\mathcal{C}$  contains a reversed pair.*

*A domain  $\mathcal{C}$  is normal if  $\alpha \in \mathcal{C}$  and  $\omega \in \mathcal{C}$ .*

We typically assume without loss of generality that any maximal width domain is in fact normal.

### 3 Single Crossing Domains

**Theorem 8.** *Let  $\mathcal{D} \subseteq S_n$  be any set of orders. The following are equivalent:*

1. *There exists an ordering  $\succ$  on  $\mathcal{D}$  such that, for all pairs of outcomes  $i, j \in [n]$ , all of the orderings  $\succ$  such that  $i \succ j$  are on the same side of all of the orderings  $\succ'$  such that  $j \succ' i$*

### 4 Normal Domains and Distributive Lattices

Maximal width (without loss of generality, normal) CDs have a significantly nicer structure than general domains. For one intuitive reason why this might be the case, observe that if you assume  $\alpha, \omega \in \mathcal{C}$ , then for any triple  $i < j < k$ , you can only lie in one of the following four cases:

$$\begin{array}{ll} jN_1ik & iN_2jk \\ jN_3ik & kN_2ij \end{array}$$

instead of the naive 9 possible cases.

**Definition 9.** *The inversion set  $\text{inv}(\sigma)$  of an ordering  $\sigma \in S_n$  is the collection of (unordered) pairs  $(ij)$  such that  $i < j$  yet  $j \succ i$ .*

---

<sup>1</sup> For example, single crossing domains  $\mathcal{C}$  are closed but in general much smaller than  $\mathcal{D}(\psi(\mathcal{C}))$ .

So, for example,  $\text{inv}(\alpha) = \emptyset$  and  $\text{inv}(\omega) = \binom{[n]}{2}$ . In the language of SYMMETRIC CONDORCET DOMAINS BY DANILOV ET AL, this next proposition relates normal CDs to sublattices of the Bruhat lattice on  $S_n$ .

**Proposition 10.** *Let  $\mathcal{D}$  be a normal CD, closed under taking medians. Then  $\mathcal{D}$  forms a distributive lattice, with join and meet operations given by*

$$p \wedge p' := \text{med}(\alpha, p, p') \qquad p \vee p' := \text{med}(\omega, p, p')$$

Moreover,  $p \leq p'$  in this lattice if and only if  $\text{inv}(p) \subseteq \text{inv}(p')$ , and we have

$$\text{inv}(p \wedge p') = \text{inv}(p) \cap \text{inv}(p') \qquad \text{inv}(p \vee p') = \text{inv}(p) \cup \text{inv}(p') \qquad (*)$$

*Proof.* The lattice operations  $\wedge$  and  $\vee$  are well-defined because  $\mathcal{D}$  is closed under taking medians. It's not too hard to verify that these operations satisfy the algebraic conditions to form a lattice. However, we use a less direct approach: essentially everything follows from the formulas (\*). We prove the first one (the second is similar). If  $(ij) \in \text{inv}(p) \cap \text{inv}(p')$ , then  $j$  is ranked higher than  $i$  in both  $p$  and  $p'$ . Thus,  $j$  is again ranked higher in  $p \wedge p'$ , so  $(ij) \in \text{inv}(p \wedge p')$ . On the other hand, if  $(ij) \notin \text{inv}(p) \cap \text{inv}(p')$ , then one of  $p$  or  $p'$  must rank  $i$  before  $j$  (as does  $\alpha$ ) and thus  $p \wedge p'$  does as well.

Now, the existence of  $p \wedge p'$  and  $p \vee p'$  satisfying (\*) can be used to easily show that  $\mathcal{D}$  forms a distributive lattice where  $p \leq p'$  precisely when  $\text{inv}(p) \subseteq \text{inv}(p')$ .  $\square$

Now, the first thing you should do when you find a distributive lattice is look to Birkhoff's representation theorem and ask yourself: what are the join-irreducible elements?

**Definition 11.** *Given a normal CD  $\mathcal{D}$  and any pair of outcomes  $i < j$ , define*

$$\mathcal{D}_{ij} := \{\succ \in \mathcal{D} : j \succ i\} \qquad p_{ij} := \bigwedge \mathcal{D}_{ij}$$

That is,  $\mathcal{D}_{ij}$  is the set of all preferences which invert pair  $(ij)$  and  $p_{ij}$  is the lowest preference which inverts  $(ij)$ .

**Proposition 12.** *Let  $\mathcal{D}$  be a median-closed Condorcet domain.*

1. *The join irreducible elements of  $\mathcal{D}$  are precisely the collection  $\{p_{ij} : (ij) \in \binom{[n]}{2}\}$ .*
2. *Any  $p \in \mathcal{D}$  can be expressed as the join of elements of  $\{p_{ij}\}$  as follows:*

$$p = \bigvee_{(ij) \in \text{inv}(p)} p_{ij}$$

3. *For any  $i < j$ , a preference  $p$  has  $(ij) \in \text{inv}(p)$  (i.e.  $p$  ranks  $j$  higher than  $i$ ) if and only if  $p \geq p_{ij}$  in the lattice.*

*Proof.* First, we show that each  $p_{ij}$  is join-irreducible. Let  $p_{ij} = p_1 \vee p_2$ , so that  $\text{inv}(p_{ij}) = \text{inv}(p_1) \cup \text{inv}(p_2)$ . At least one of  $p_1$  or  $p_2$  must invert  $i, j$ , say  $(ij) \in \text{inv}(p_1)$ . Then  $p_1 \in \mathcal{D}_{ij}$ , so  $p_{ij} \leq p_1$ . For any lattice, if  $u = p_1 \vee p_2$ , then  $p_1 \leq u$ . Thus, we have  $p_{ij} = p_1$ , and  $p_{ij}$  is join-irreducible.

To show  $\{p_{ij}\}_{ij}$  is actually all irreducible elements in  $\mathcal{D}$ , it will suffice to show 2. For, if  $p \notin \{p_{ij}\}_{ij}$ , then formula 2 gives a way to write  $p$  as the join of elements other than  $p$ .

We now prove 2. Certainly  $(ij) \in \text{inv}(p_{ij})$ , so

$$\text{inv}(p) \subseteq \bigcup_{(ij) \in \text{inv}(p)} \text{inv}(p_{ij}) = \text{inv} \left( \bigvee_{(ij) \in \text{inv}(p)} p_{ij} \right) \implies p \leq \bigvee_{(ij) \in \text{inv}(p)} p_{ij}$$

On the other hand, for each  $(ij) \in \text{inv}(p)$ , we have  $p \in \mathcal{D}_{ij}$ , so actually  $p_{ij} \leq p$ . Thus  $\bigvee_{(ij) \in \text{inv}(p)} p_{ij} \leq p$ .

To prove 3, notice that if  $(ij) \in \text{inv}(p)$ , then  $p \in \mathcal{D}_{ij}$ , so  $p \geq p_{ij}$ . On the other hand, if  $p \geq p_{ij}$ , then  $(ij) \in \text{inv}(p)$ .  $\square$

Finally, we arrive at the following:

**Definition 13.** For any median-closed normal domain  $\mathcal{D}$ , define the inversion poset  $\mathcal{P}(\mathcal{D})$  as follows: the elements of  $\mathcal{P}(\mathcal{D})$  are the preimages in  $\binom{[n]}{2}$  of the map  $(ij) \mapsto p_{ij}$  (that is, the collection of sets of pairs  $\{(kl) \mid p_{kl} = p_{ij}\}$  for some fixed  $(ij)$ ). The order relations in  $\mathcal{P}(\mathcal{D})$  are as follows: if  $S \ni (ij)$  and  $T \ni (kl)$ , then  $S \leq T$  if and only if  $p_{ij} \leq p_{kl}$  in  $\mathcal{D}$ .

**Theorem 14.** For any closed normal CD  $\mathcal{D}$ , the distributive lattice on  $\mathcal{D}$  is isomorphic to the lattice of downward-closed sets of  $\mathcal{P}(\mathcal{D})$ . Moreover, the correspondance  $\mathcal{D} \rightarrow 2^{\mathcal{P}(\mathcal{D})}$  is given by  $p \mapsto \{S : S \subseteq \text{inv}(p)\}$ .

*Proof.* TODO  $\square$

**Example 15.** 1. A single crossing

2. A symmetric

3. Fishburn's?

## 5 Normal Peak/Pit Domains

**Theorem 16.** Let  $\mathcal{D} = \mathcal{D}(\psi)$  be a closed normal peak/pit CD given by a peak/pit VR casting  $\psi$ .  $\mathcal{D}$  is maximal if and only if the map  $(ij) \mapsto p_{ij}$  is injective (that is, every pair in  $\binom{[n]}{2}$  appears as its own element in  $\mathcal{P}(\mathcal{D})$ ).

*Proof.* Hopefully the following outline pans out.

( $\Leftarrow$ ) Consider any maximal chain from  $\alpha$  to  $\omega$  in the poset  $\mathcal{D}$ . By theorem 14, this corresponds to an ordering of the set of inversions  $\binom{[n]}{2}$ . This is actually a maximal single crossing domain. Observe that maximal single crossing domains have full projection onto every triple. Thus, I should write a “full projection lemma” and be done with this direction.

( $\Rightarrow$ )  $\square$

The following lemma tells us there is a canonical way to make a normal peak pit domain into a distributive lattice.

**Proposition 17.** If  $\mathcal{C}$  is a peak/pit domain, then  $\mathcal{C}$  can contain at most one reversed pair.

*Proof.* Suppose  $\mathcal{C}$  contains  $\alpha$  and  $\omega$  and a some preference  $p, \text{rev}(p) \in \mathcal{C}$  with  $p \neq \alpha, \omega$ . There must be some triple  $i < j < k$  which are neither in increasing order nor in reversed order in  $p$ . Up to symmetry, this leaves two cases:  $p|_{\{i,j,k\}} = jik$  and  $\text{rev}(p)|_{\{i,j,k\}} = kij$ , or  $p|_{\{i,j,k\}} = ikj$  and  $\text{rev}(p)|_{\{i,j,k\}} = jki$ . In either case, each of  $i, j, k$  are sometimes the top and sometimes the last choice in  $\mathcal{C}|_{\{i,j,k\}}$ . Thus, this triple can satisfy no peak/pit restriction in  $\mathcal{C}$ .  $\square$

VR castings are especially useful when studying restricted classes of Condorcet domains. In special cases, some simple observations about which VR cases are possible allow us to reduce the complexity of the set of VR castings considerably.

**Definition 18.** *A normal peak/pit VR casting  $\psi$  is ...*

$$\forall i < j < k : \quad \begin{cases} \text{either} & \psi(\{i, j, k\}) = jN_1ik \\ \text{or} & \psi(\{i, j, k\}) = jN_3ik \end{cases}$$

**Definition 19.** *Given a normal peak/pit VR casting, define the inversion graph  $R(\psi)$  with vertices  $\binom{[n]}{2}$  and, for each  $i < j < k$ , add edges*

$$\begin{array}{ll} \text{if } \psi(\{i, j, k\}) = jN_1ik \text{ then} & (ij) \longrightarrow (ik) \longrightarrow (jk) \\ \text{if } \psi(\{i, j, k\}) = jN_3ik \text{ then} & (jk) \longrightarrow (ik) \longrightarrow (ij) \end{array}$$

The above definition looks a bit strange, but  $\overline{R}(\psi)$  is actually defined very naturally such that the following theorem is essentially obvious:

**Proposition 20.** *An order  $p \in S_n$  is in  $\mathcal{D}(\psi)$  if and only if  $\text{inv}(p)$  is a closed set (i.e. a set closed under following directed edges) in the inversion graph  $R(\psi)$ .*

*Proof.* TODO □

We thus obtain the following correspondence with the previous section:

**Proposition 21.** *For domain given by a normal peak/pit casting  $\psi$ , the inversion poset  $\mathcal{P}(\mathcal{D}(\psi))$  is the transitive closure of the relation given by  $p_1 \geq p_2$  whenever  $p_1 \longrightarrow p_2$  is an edge in  $R(\psi)$ .*

## 6 Bad News – Counterexamples

The property of maximality gives a lot of structure to Condorcet domains. However, this structure is complicated.

**Example 22.** *There are maximal Condorcet domains without full projection onto some triple. Moreover, maximality is not preserved under restrictions.*