

Figure 1: Illustrations

Claim 1. For any pair of points $i = (i_1, i_2), j = (j_1, j_2) \in X$, where neither dominates the other, let $a = (a_1, a_2) \in \mathbb{R}_{\geq 0}^2$ be the preference weight vector that gives $i =_a j$, then for any preference weight vector $b = (b_1, b_2) \in \mathbb{R}_{\geq 0}^2$,

1. $\frac{a \times b}{\|a \times b\|} > 0$ if and only if $i >_b j$.
2. $\frac{a \times b}{\|a \times b\|} < 0$ if and only if $i <_b j$.

Proof. 1 and 2 are completely analogous, here we only prove 1.

- “ \implies ”: On the one hand, from $\frac{a \times b}{\|a \times b\|} > 0$ we have

$$a_1 b_2 - a_2 b_1 < 0 \implies a_1 < \frac{b_1}{b_2} a_2. \quad (1)$$

On the other hand, by $i =_a j$,

$$a_1 i_1 + a_2 i_2 = a_1 j_1 + a_2 j_2 \implies a_1(i_1 - j_1) = a_2(j_2 - i_2). \quad (2)$$

Plug (1) into (2), $a_2(j_2 - i_2) < \frac{b_1}{b_2} a_2(i_1 - j_1)$, this gives

$$b_1 i_1 + b_2 i_2 > b_1 j_1 + b_2 j_2 \implies i >_b j.$$

- “ \impliedby ”: From (2) we get

$$\frac{j_2 - i_2}{i_1 - j_1} = \frac{a_1}{a_2}. \quad (3)$$

By $i >_b j$ and (3),

$$b_1 i_1 + b_2 i_2 > b_1 j_1 + b_2 j_2 \implies \frac{b_1}{b_2} > \frac{j_2 - i_2}{i_1 - j_1} = \frac{a_1}{a_2} \implies a_1 b_2 - a_2 b_1 < 0 \implies \frac{a \times b}{\|a \times b\|} > 0.$$

□

Using Claim 1 we are able to show the following lemma.

Lemma 2. Let i, j, i', j' be two pairs of points in X satisfying

- neither of i, j dominates the other: $i_1 < j_1$ and $i_2 > j_2$,

- neither of i', j' dominates the other: $i'_1 < j'_1$ and $i'_2 > j'_2$,
- the slope of line ij is less than the slope of line $i'j'$: $\frac{j_2 - i_2}{j_1 - i_1} < \frac{j'_2 - i'_2}{j'_1 - i'_1}$.

Then for any preference weight vector $b \in \mathbb{R}_{\geq 0}^2$,

1. If $i >_b j$, $i' >_b j'$.
2. If $j' >_b i'$, $j >_b i$.

Proof. Again 1 and 2 are symmetric, we only prove 1. Let a, a' be the preference weight vectors that gives $i =_a j$ and $i' =_{a'} j'$ respectively. By definition of a' ,

$$a'_1 i'_1 + a'_2 i'_2 = a'_1 j'_1 + a'_2 j'_2 \implies \frac{a'_1}{a'_2} = \frac{j'_2 - i'_2}{i'_1 - j'_1}. \quad (4)$$

Since $i >_b j$,

$$b_1 i_1 + b_2 i_2 > b_1 j_1 + b_2 j_2 \implies \frac{b_1}{b_2} > \frac{j_2 - i_2}{i_1 - j_1}. \quad (5)$$

From the relation of slopes we have

$$\frac{j_2 - i_2}{i_1 - j_1} > \frac{j'_2 - i'_2}{i'_1 - j'_1}. \quad (6)$$

Combine (4)(5)(6) we get

$$\frac{b_1}{b_2} > \frac{j_2 - i_2}{i_1 - j_1} > \frac{j'_2 - i'_2}{i'_1 - j'_1} = \frac{a'_1}{a'_2} \implies a'_1 b_2 - a'_2 b_1 < 0 \implies \frac{a' \times b}{\|a' \times b\|} > 0.$$

By Proposition ??, this implies $i' >_b j'$. □

Let P_2 be a set of 2-dimensional preferences on $[n]$.

Theorem 3. $|P_2| \leq \binom{n}{2} + 1$.

Proof. We prove this upper bound algorithmically. Consider the following algorithm of generating a preference in P_2 given X :

Algorithm 1 Preference Generator for P_2

- 1: Relabel all points in X such that $y_1 \geq y_2 \geq \dots \geq y_n$
 - 2: **for all** $i, j \in X, i < j$ **do**
 - 3: Compute slope s_{ij}
 - 4: **if** $s_{ij} < 0$ **then**
 - 5: $L \leftarrow s_{ij}$ ▷ Suppose that $L[1]$ is the first element in the list L
 - 6: Sort L in ascending order
 - 7: Choose a pivot $p \in \{0, \dots, |L|\}$
 - 8: **if** $p \geq 1$ **then**
 - 9: Let $L[p] = s_{ij}$, set $i >_a j$
 - 10: **if** $p + 1 \leq |L|$ **then**
 - 11: Let $L[p + 1] = s_{i'j'}$, set $i' <_a j'$
 - 12: Output $>_a$
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Lemma 2 ensures that by setting $i <_a j$ and $i' >_a j'$ in steps 8-11, all pairs k, l before i, j in list L are set to $k <_a l$, and all pairs k', l' after i', j' in L are set to $k' <_a l'$. Moreover, there is a one-to-one correspondence between the actual preference, and choosing a pivot to flip from $<_a$ to $>_a$. Since $|L| \leq \binom{n}{2}$, and the number of pivots that we can choose from is $|L| + 1$, we get that $|P_2| \leq \binom{n}{2} + 1$.

The upper bound is tight. ADD the example here! □

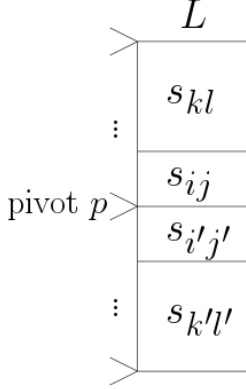


Figure 2: List L