Lecture 6
String Matching

Bo Tang @ SUSTech, Fall 2022

Our Roadmap

- String Concepts
- String Searching Problem
  - Brute Force Solution
  - Rabin-Karp
  - Finite State Automata
  - Knuth-Morris-Pratt

String Definition

- String:
  - Sequence of characters over some alphabet
  - Binary {0,1}: S1 = “1000010101010101001010101”
  - DNA {ACGT}: S2 = “ACGTACGTACGTTCGA”
  - English Characters {a…z, A..Z}: S3 = “Hello World”
- Applications
  - Word processors
  - Virus scanning
  - Text retrieval
  - Natural language processing
  - Web search engine

String Operators

- append: append to string
- assign: assign content to string
- insert: insert to string
- erase: erase characters from string
- replace: replace portion of string
- swap: swap string values
- find: find the specific char in the string
- Give string s=“SUSTechCS203”, how many sub string it has?

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Why String Searching?

- **Applications in Computational Biology**
  - DNA sequence is a long word (or text) over a 4-letter alphabet
  - GTTTGAGTGGTCAGTCTTTTCGTTTCGACGGAGCCC.....
  - Find a Specific pattern W
- **Finding patterns in documents formed using a large alphabet**
  - Word processing
  - Web searching
  - Desktop search (Google, MSN)
- **Matching strings of bytes containing**
  - Graphical data
  - Machine code
- **grep in unix**
  - grep searches for lines matching a pattern.

String Searching

Search Text

a	s	s	u	s	u	s	t	c	s	c

Search Pattern

s	u	s	t	c

Successful Search

a	s	s	u	s	u	s	t	c	s	c

- Parameter
  - n: # of characters in text
  - m: # of characters in pattern
  - Typically, n >> m
    - e.g., n = 1 Billion, m = 100

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Brute Force

- Brute force
  - Check for pattern starting at every text position
- **Algorithm:** BruteForce(T, P):

```
n ← len(T), m ← len(P)
for i ← 0 to n-m-1
        for j ← 0 to m-1
                if P[j] != T[i+j] then
                        break;
        if j = m-1
            pattern occurs with shift i
```

- Time complexity?

Analysis of Brute Force

- Analysis of brute force
  - Running time depends on pattern and text
  - Can be slow when strings repeat themselves
  - Worst case: mn comparisions
  - Too slow when m and n are large

Search Pattern

a	a	a	a	b

Search Text

a	a	a	a	a	a	a	a	a	a	a	a	a	a	a	b
a	a	a	a	b											
	a	a	a	a	b										
		a	a	a	a	b									
			a	a	a	a	b								

......

Can we do better?

- How to avoid re-computation?
  - Pre-analyze search pattern
  - Example: suppose the first 4 chars of pattern are all a's
    - If t[0..3] matches p[0..3] then t[1..3] matches p[0..2]
    - No need to check i=1, j=0,1,2
    - Saves 3 comparisons
  - Need better ideas in general

Search Pattern

a	a	a	a	b

Search Text

a	a	a	a	a	a	a	a	a	a	a	a	a	a	a	b
a	a	a	a	b											
	a	a	a	a	b										

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Rabin-Karp Algorithm

- Given search text T and search pattern P as follows:

Pattern			
1	3	5	9

Search Text												
2	4	6	8	0	1	2	1	3	5	9	7	2
							1	3	5	9		

- Any idea?

2468	4680	6801	8012	121	1213	2135	1359	3597	5972
							1359		

Rabin-Karp Algorithm

- General idea
  - Convert search pattern to a number p
  - Convert search text to an array of numbers t[0],…,t[n-m-1]
  - Compare p with t[i], for each i in [0,n-m-1]
  - if p=t[i], pattern p occurs
- Example
  - p = 1359
  - Array t is:

2468	4680	6801	8012	121	1213	2135	1359	3597	5972

  - t[7] = p → T[7,8,9,10]=P[0,1,2,3]

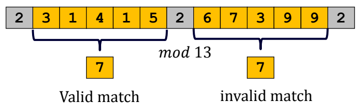
Rabin-Karp Algorithm

- How to convert size-m characters to a number?
  - E.g., the alphabet Σ = {a,…,z,A,…,Z}
  - Solution: radix-d (d=|Σ|) Horner's rule
  - p = P[m-1]+d(P[m-2]+d(P[m-3]+…+d(P[1]+dP[0])))
- When m is large, p may be too large to work
  - Modulo a proper prime number q
  - p = P[m-1]+d(P[m-2]+d(P[m-3]+…+d(P[1]+dP[0]))) mod q
- Compute t[0],t[1],…,t[n-m-1] in time O(n-m)
  - Compute t[i+1] by using t[i] in O(1) time
  - t[i+1] = d(t[i]-d^{m-1}T[i])+T[i+m]
  - t[i+1] = ((t[i]-hT[i])+T[i+m]) mod q, where h≡d^{m-1} (mod q)
  - t[0] → t[1] → t[2] → t[3] → … → t[n-m-1] in O(n-m)

Rabin-Karp Algorithm

- Correctness analysis
  - $p \not\equiv t[i]\ (mod\ q)$ we have $p \neq t[i]$, thus, P[0,..m-1] != T[i,i+m-1]
  - $p \equiv t[i]\ (mod\ q)$, it does not imply $p = t[i]$ (**spurious hit**)
- Example: search P:

3	1	4	1	5



- Additional test to check
  - P[0,…,m-1] = T[i, i+m-1]

Rabin-Karp Algorithm

- **Algorithm:** Rabin-Karp(T, P, d, q):

```
n ← len(T), m ← len(P)
h ← d^(m-1) (mod q), p ← 0, t0 ← 0
for j ← 0 to m-1
      p ← (dp + P[j]) mod q,
      t_0 ← (dt_0 + T[j]) mod q,
for i ← 0 to n-m
      if p != t_i then
            t_i+1 ← (d(t_i-T[i]h)+T[i+m]) mod q
      else
            If P[0,..m-1]=T[i,i+m-1]
                  pattern occurs with shift i
            Else
                  t_i+1 ← (d(t_i-T[i]h)+T[i+m]) mod q
```


Analysis of Rabin-Karp Alg.

- **Algorithm:** Rabin-Karp(T, P, d, q):

```
Overall Cost: O(mn)
```

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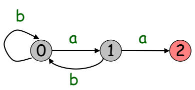
Midterm Exam (tentative)

- **Time: 12 Nov. 16:30-18:30**
- **Venue: To be decided**
- **Scope: Lecture 1 to 6**

Finite State Automata

- A finite State automaton is defined by:
  - *Q, a set of states*
  - $q_0 \in Q$ *, the start state*
  - $A \subseteq Q$ *, the accepting states*
  - Σ*, the input alphabet*
  - δ*, the transition function, from* $Q \times \Sigma$ *to* Q

	0	1
a	1	2
b	0	0

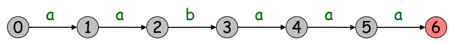


FSA idea for String Matching

- Start in state q_0
- Perform a transition from q_0 to q_1 if next character of T = P[1]
- State q_i means first i characters of P match.
- Transition from q_i to q_{i+1} if the next character of T = P[i+1]

Search Pattern					
a	a	b	a	a	a

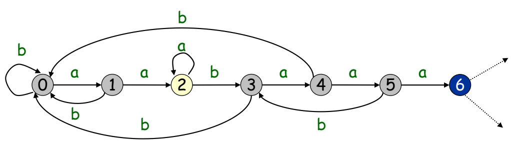
	0	1	2	3	4	5
a	1	2	?	4	5	6
b	?	?	3	?	?	?



- How to fill these ???
  - Reset to q_0? Why not?

FSA construction

- FSA construction
  - FSA builds itself
- Example. Build FSA for `aabaaabb`
  - State 6. P[0..5]=aabaaa
  - assume you know state for p[1..5] = abaaa X = 2
  - if next char is b (match): go forward 6 + 1 = 7
  - if next char is a (mismatch): go to state for abaaaa X + 'a' = 2
  - update X to state for p[1..6] = abaaab X + 'b' = 3



FSA construction

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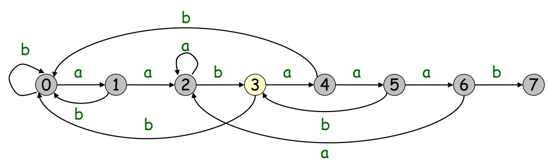
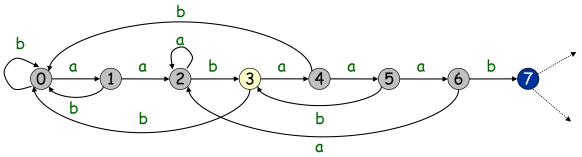
- Example. Build FSA for `aabaaabb`



FSA construction

- FSA construction
  - FSA builds itself
- Example. Build FSA for `aabaaabb`
  - State 7. p[0..6]=aabaaab
  - assume you know state for p[1..6] = abaaab X = 3
  - if next char is b (match): go forward 7 + 1 = 8
  - if next char is a (mismatch): go to state for abaaaba X + 'a' = 4
  - update X to state for p[1..7] = abaaabb X + 'b' = 0



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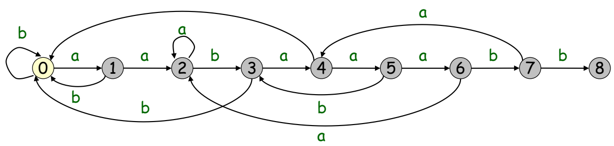


FSA construction

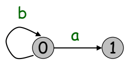
- FSA construction
  - FSA builds itself
- Crucial Insight
  - To compute transitions for state n of FSA, suffices to have:
    - FSA for state 0 to n-1
    - State X that FSA ends up in with input p[1..n-1]
  - To compute state X' that FSA ends up in with input p[1..n], it suffices to have
    - FSA for states 0 to n-1
    - State X that FSA ends up in with input p[1..n-1]

FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

j	pattern[1..j]	X

a	
b	

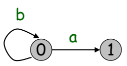


FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

j	pattern[1..j]							X
0								0

	0
a	1
b	0

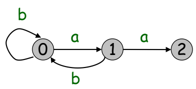


FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

j	pattern[1..j]							X
0								0
1	a							1

	0	1
a	1	2
b	0	0

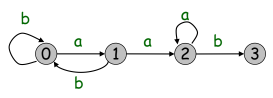


FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

	0	1	2
a	1	2	2
b	0	0	3

j	pattern[1..j]								X
0									0
1	a								1
2	a	b							0

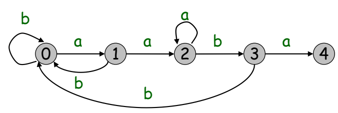


FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

	0	1	2	3
a	1	2	2	4
b	0	0	3	0

j	pattern[1..j]								X
0									0
1	a								1
2	a	b							0
3	a	b	a						1

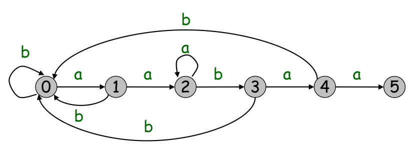


FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

	0	1	2	3	4
a	1	2	2	4	5
b	0	0	3	0	0

j	pattern[1..j]								X
0									0
1	a								1
2	a	b							0
3	a	b	a						1
4	a	b	a	a					2

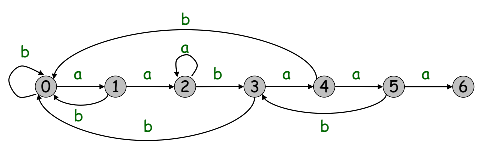


FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

	0	1	2	3	4	5
a	1	2	2	4	5	6
b	0	0	3	0	0	3

j	pattern[1..j]								X
0									0
1	a								1
2	a	b							0
3	a	b	a						1
4	a	b	a	a					2
5	a	b	a	a	a				2

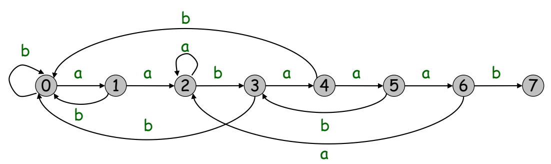


FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

	0	1	2	3	4	5	6
a	1	2	2	4	5	6	2
b	0	0	3	0	0	3	7

j	pattern[1..j]								X
0									0
1	a								1
2	a	b							0
3	a	b	a						1
4	a	b	a	a					2
5	a	b	a	a	a				2
6	a	b	a	a	a	b			3

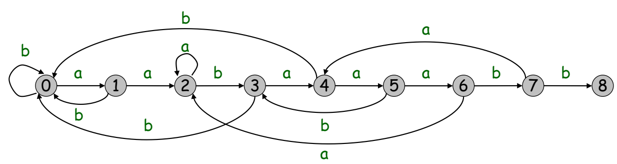


FSA construction

Search Pattern							
a	a	b	a	a	a	b	b

	0	1	2	3	4	5	6	7
a	1	2	2	4	5	6	2	4
b	0	0	3	0	0	3	7	8

j	pattern[1..j]								X
0									0
1	a								1
2	a	b							0
3	a	b	a						1
4	a	b	a	a					2
5	a	b	a	a	a				2
6	a	b	a	a	a	b			3
7	a	b	a	a	a	b	b		0



Transition function

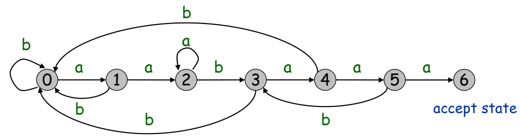
- **Algorithm:** Transition(P, Σ):

```
m ← len(P)
X ← 0
Initialize δ(0,a) for each a ∈ Σ
for j ← 1 to m-1
      for each character a ∈ Σ
            if P[j+1] = a then   // char match
                  δ(j,a) ← j + 1
            else                 // char mismatch
                  δ(j,a) ← δ(X,a)
      X ← δ(X,P[j+1])
return δ
```


Finite State Automata (FSA)

- FSA-matching algorithm.
  - Use knowledge of how search pattern repeats itself.
  - ➡ Build FSA from pattern.
  - Run FSA on text.

Search Pattern					
a	a	b	a	a	a



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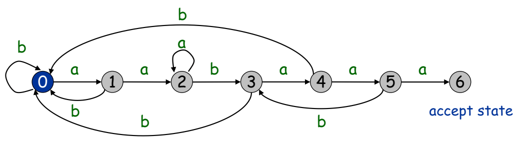
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



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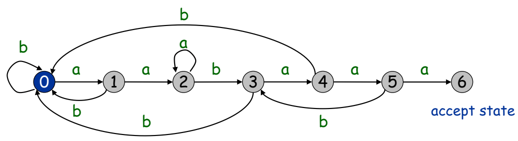
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



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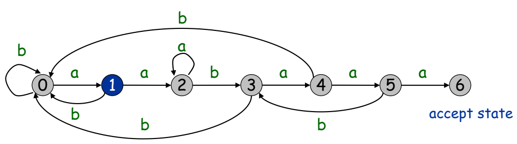
Search Pattern					
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



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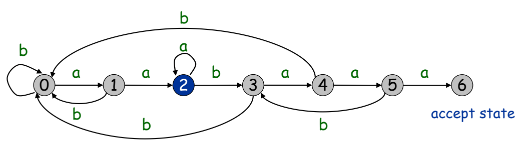
Search Pattern					
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



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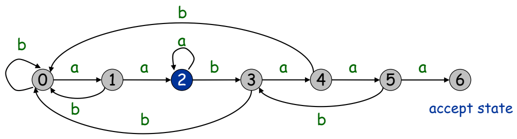
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



Finite State Automata (FSA)

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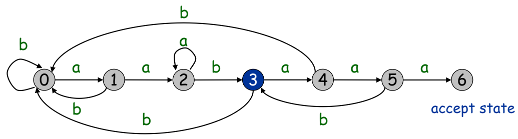
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



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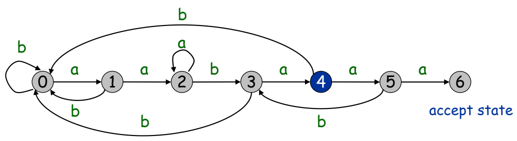
Search Pattern					
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



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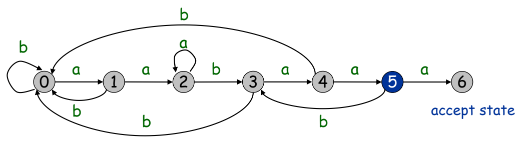
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



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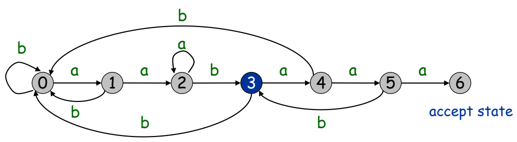
Search Pattern					
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



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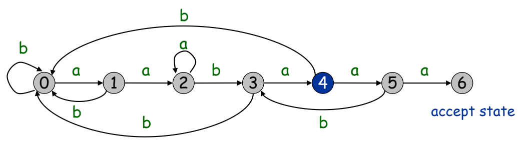
  - Run FSA on text.

Search Pattern					
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b

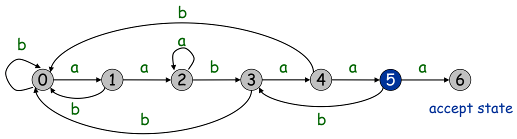


Finite State Automata (FSA)

- FSA-matching algorithm.
  - Use knowledge of how search pattern repeats itself.
  - Build FSA from pattern.
  - Run FSA on text.

Search Pattern					
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b



Finite State Automata (FSA)

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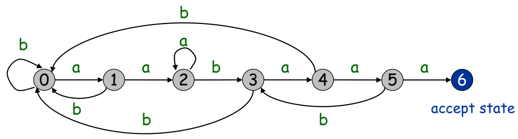
Search Pattern					
a	a	b	a	a	a

Search Text										
a	a	a	b	a	a	b	a	a	a	b

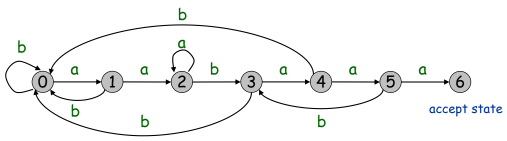


Finite State Automata (FSA)

- FSA used in KMP has special property
  - If match, go to next state
  - Only need to keep track of where to go upon character mismatch.
    - go to state next[j] if character mismatches in state j

Search Pattern					
a	a	b	a	a	a

	0	1	2	3	4	5
a	1	2	2	4	5	6
b	0	0	3	0	0	3
next	0	0	2	0	0	3



FSA algorithm

- **Algorithm:** FSA(T, P):

```
n ← len(T), m ← len(P)
δ ← Transition(P, Σ)
q ← 0 // q is the state of the FSA.
for i ← 1 to n
      q ← δ(q,T[i])
      if q = m
            pattern occurs with shift i – m
```


Analysis of FSA

- **Algorithm:** FSA(T, P):

Overall Cost: O(|Σ|m+n)

Our Roadmap

- String Concepts
- String Searching Problem
  - Brute Force Solution
  - Rabin-Karp
  - Finite State Automata
  - Knuth-Morris-Pratt

History of KMP

- Inspired by the theorem of Cook that says O(m+n) algorithm should be possible
- Discovered in 1976 independently by two groups
- Knuth-Pratt
- Morris was hacker trying to build an editor
- Resolved theoretical and practical problem
  - Surprise when it was discovered
  - In hindsight, seems like right algorithm

String

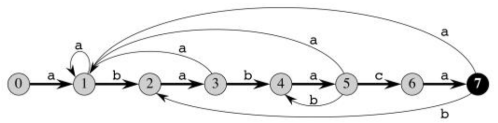
- **String**: “HelloCS203”
- **Substring**: a substring of s string S is a string S’ that occurs in S, e.g., P[2,…,4] = “ell”
- **Prefix (P[1,…])**: a prefix of a string S is a substring of S that occurs at the beginning of S, e.g., P[1,…,1] = “H” (note that P[1]=‘H’), P[1,…,2] = “He”, P[1,…,5] = “Hello”, we denote prefix as: **P[1,…]**
- **Suffix**: a suffix of a string S is a substring of S that occurs at the end of S, e.g., P[10,…,10]=“3”, P[8,…,10]=“203”, P[6,…,10] = “CS203”, we denote suffix as: **P[…,m]**

Finite State Automata

- P =“ababaca”
- Transition function table

State	0	1	2	3	4	5	6	7
a	1	1	3	1	5	1	7	1
b	0	2	0	4	0	4	0	2
c	0	0	0	0	0	6	0	0
P	a	b	a	b	a	c	a	

- State transition graph



Finite State Automata

- P =“ababaca” and T =“abababacaba”

i	1	2	3	4	5	6	7	8	9	10	11
T	a	b	a	b	a	b	a	c	a	b	a
1	a	b	a	b	a	c	a				
2			a	b	a	b	a	c	a		
3									a	b	

- After **failure**: at i=6, ‘c’ was expected, but not found in T[6], FSA transition to state δ(5,b)=4, it means pattern prefix P[1..4] =“abab” has matched the text suffix T[2..6] = “abab”

	0	1	2	3	4	5	6	7
a	1	1	3	1	5	1	7	1
b	0	2	0	4	0	4	0	2
c	0	0	0	0	0	6	0	0

Finite State Automata

- In general, the FSA is constructed so that the state number tells us how much of a prefix of P has been matched.
- FSA transition function:
  - 1) Find the longest prefix of P is also a suffix of T[…,i], denote as k, i.e., *P[1,…,k]=T[i-k+1,…,i]*
  - 2) Read the next character at “*k+1*” (i.e., T[i+1]), there are two kinds of transitions:
    - P[k+1] = T[i+1], it is matched, continues.
    - Otherwise, it is mismatched, go to δ(k,T[i+1])

Prefix Function

- Consider the first step of FSA transition function:
  - Find the longest prefix of P is also a suffix of T[…i], denote as *k*, i.e., *P[1,…,k]=T[i-k+1,…,i]*
- Suppose it is mismatched at *“P[k+1]”*, it means:
  - *P[k+1] != T[i+1]* then,
  - we should find the longest prefix of *P[1,…,k]* is also a suffix of *T[i-k+1, …, i]*.
- **Prefix function (next array in general),** given *P*[1..*m*], the prefix function π for *P* is π : {1, 2 …, *m*} -> {0, 1, …, *m-1*} such that:

  π[i]=max{k, k<i and P[1,..,k]= P[i-k+1,…,i] }

Prefix Function

- **Prefix function,** given *P*, the prefix function π for *P* is π : {1, 2 ..., m} -> {0, 1, ..., m-1} such that:

  π[q]=max{k, k<q and P[1,..,k]= P[q-k+1,...,q] }

- Example: P ="ababaca"

i	1	2	3	4	5	6	7
P[i]	a	b	a	b	a	c	a
π[i]	0	0	1	2	3	0	1

Compute next array

- **Algorithm:** NextArray(P):

```
m ← len(P)
Let π[1,…,m] be a new array
π[1] = 0, k ← 0
for q = 2 to m
      while k > 0 and P[k+1] != P[q]
            k ← π[k]
      if P[k+1] = P[q]
            k ← k + 1
      π[q] ← k
return π
```


KMP algorithm

- **Algorithm:** KMP(T, P):

```
n ← len(T), m ← len(P)
π ← NextArray(P)
q ← 0
for i = 1 to n
      while q > 0 and P[q+1] != T[i]
            q ← π[q]
      if (P[q+1] = T[i])
            q ← q + 1
      if q == m
            print "Pattern occurs with shift" i-m
            q ← π[q]
```


Analysis of KMP

- **Algorithm:** KMP(T, P):

```
Overall Cost: O(m+n)
```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Thank You!