



CS-3002 Information Security

Lecture # 2: Introduction to Cryptography and Classical Cryptography

Prof. Dr. Sufian Hameed

Department of Computer Science

FAST-NUCES



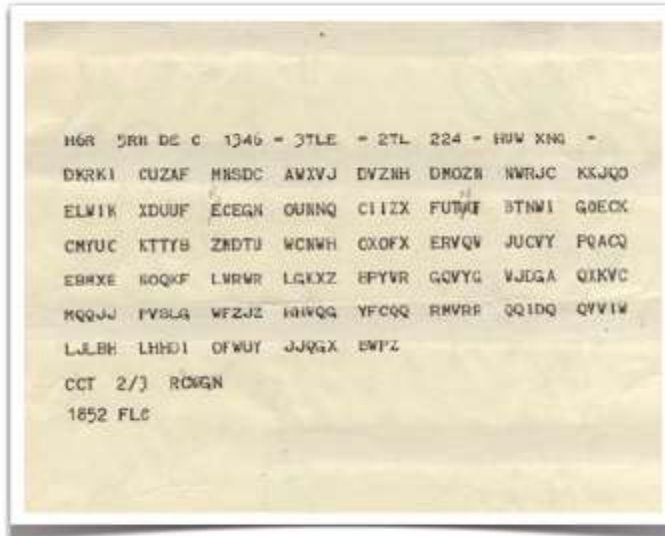
Cryptography

- » **Cryptography** (*kryptos*: secret; *graphein*: writing)
= art and science of keeping information secure
⇒ protection of confidentiality and integrity
- » **Cryptanalysis** = study of attacks against cryptography
- » **Cryptology** = cryptography and cryptanalysis
- » **Steganography** (*steganos*: covered; *graphein*: writing)
= art and science of hiding information
⇒ deniability and unobservability of communication



Examples

Cryptography



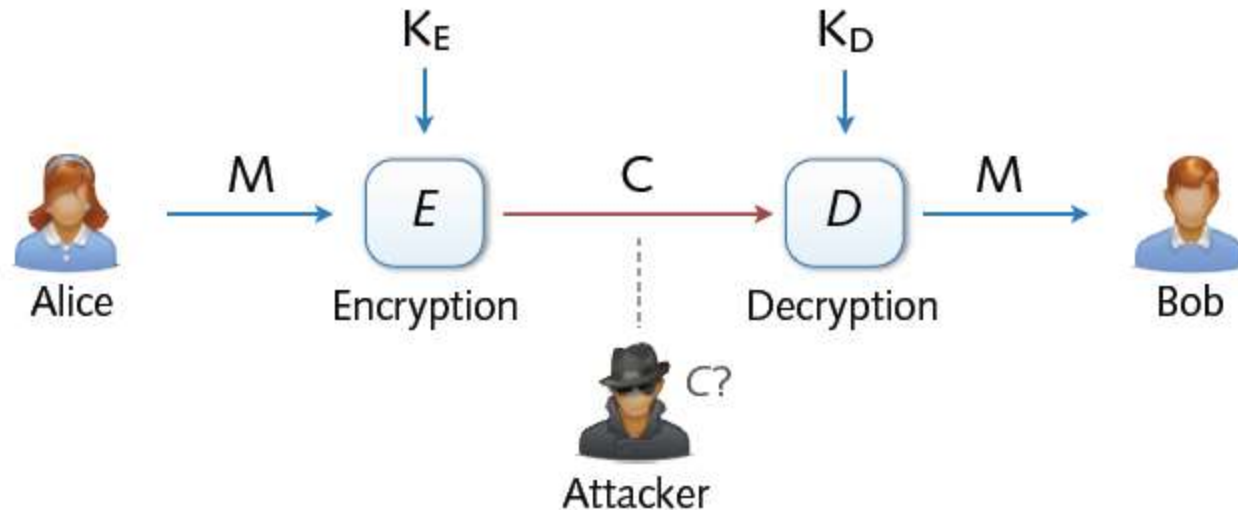
Message encrypted using
the Enigma during WW2

Steganography



Message of 500 bytes
hidden within image

Cryptosystem



» Cryptographic system for en/decrypting messages

- » M = plaintext message C = ciphertext message
- » K_E = encryption key K_D = decryption key

Attacks against Cryptosystems

- 1.) **Cipher text-only:** Attacker possesses a string y of the cipher text
- 2.) **Known plaintext:** Attacker possesses a string x of the plaintext and the corresponding cipher text y . The problem now is to find out the key which produces y from x
- 3.) **Chosen plaintext:** Attacker has access to the encryption machinery. Hence he can chose a plaintext string x and construct the corresponding cipher text string y .
- 4.) **Chosen cipher text:** Attacker has access to the decryption machinery. Hence, he can chose a cipher text string y and construct the corresponding plaintext string x .



Security of Keys

» Kerckhoffs's Principle

- » Cryptosystem is known, security depends on key only
- » Contrasting concept: "security by obscurity"

» Keyspace defined over bits of key

- » n -bit key \mapsto size of keyspace 2^n
- » Time of brute-force attacks grows exponentially in n

Bits of key	10^9 checks per second	Cluster of 100,000 nodes
16	0.07 milliseconds	0.000 milliseconds
32	4.29 seconds	0.004 milliseconds
64	585 years	5 hours
128	10^{22} years	10^{16} years

considered secure



Cryptography is everywhere

Secure communication:

- web traffic: HTTPS
- wireless traffic: 802.11i WPA2 (and WEP), GSM, Bluetooth

Encrypting files on disk:

- EFS (Encrypting File System)
- TrueCrypt (open-source disk encryption software)

Content protection

- DVD --- Content Scramble System (**CSS**) is a Digital Rights Management (DRM) and encryption system employed on almost all commercially produced **DVD-Video**
 - Easy to break
- Blu-Ray --- Advance Access Content System (AACS)

User authentication

... and much much more



Things to remember

Cryptography is:

- A tremendous tool
- The basis for many security mechanisms

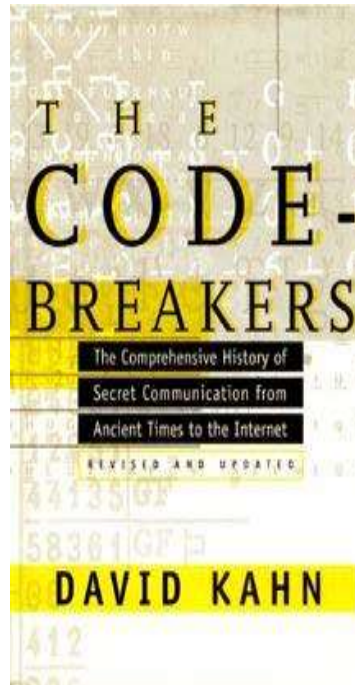
Cryptography is not:

- The solution to all security problems
 - Software bugs
 - Social engineering attacks
- Reliable unless implemented and used properly
 - Wired Equivalent Privacy (WEP -- good example on how not to use cryptography)
- ***Something you should try to invent yourself***
 - many examples of broken ad-hoc designs
 - Proprietary ciphers, once re-engineered are easily broken



History

David Kahn, “The code breakers” (1996)

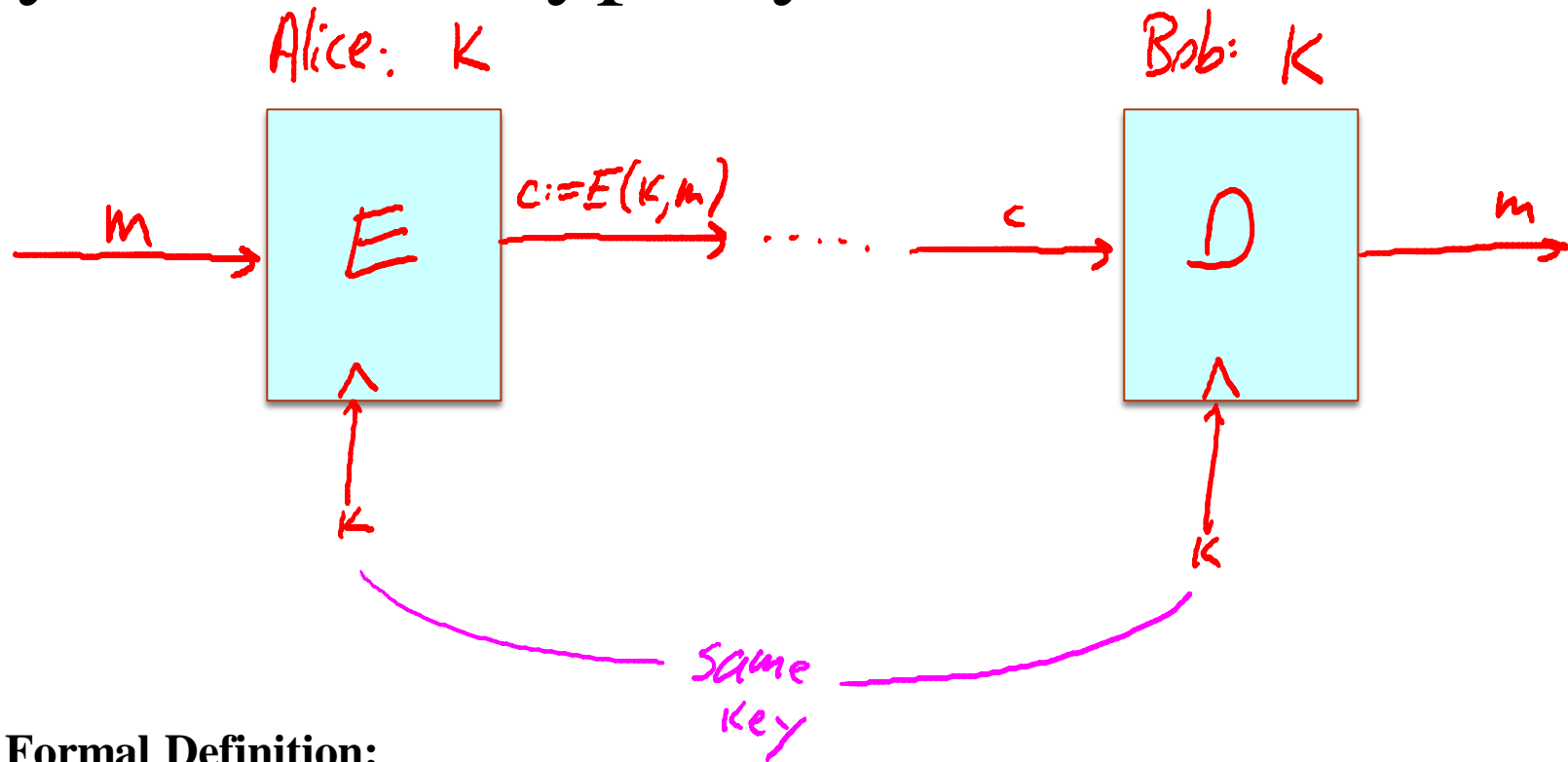


Historical Cryptosystems

- **Monoalphabetic cipher:** Each alphabetic character is mapped onto a unique alphabetic character
 - Examples: Shift Cipher, Substitution Cipher, Affine Cipher
- **Polyalphabetic cipher:** Each alphabetic character is mapped onto various alphabetic characters
 - Examples: Vigenere Cipher, Hill Cipher, Permutation Cipher



Symmetric Cryptosystems



Formal Definition:

Cryptosystem is defined over (K, M, C) and a pair of “efficient” algorithms (E, D) s.t.

$$\forall m \in M, k \in K \text{ and } c \in C : E(k, m) = c, D(k, E(k, m)) = m$$

Efficient means run in polynomial time



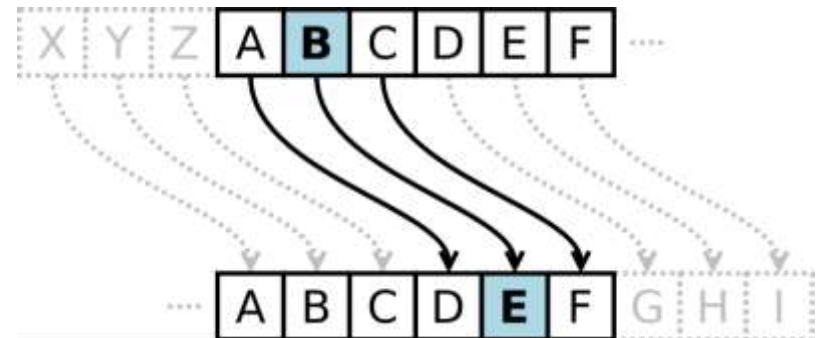
Shift Cipher

- Cipher in which each letter in the plaintext is replaced by a letter some fixed number of positions down the alphabet. Example includes Ceasar cipher, ROT13
- **Ceasar Cipher**
 - Each letter is replaced with a fixed shift of 3 letters

Example of Ceasar cipher using left rotation of 3 places

Plain: ABCDEFGHIJKLMNOPQRSTUVWXYZ

Cipher: DEFGHIJKLMNOPQRSTUVWXYZABC



Shift Cipher (ROT13)

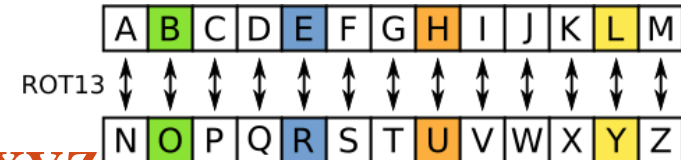
- **ROT13**

- Each letter is replaced with a fixed shift of 13 letters

The transformation can be done as follows

Plain: ABCDEFGHIJKLMNOPQRSTUVWXYZ

Cipher: NOPQRSTUVWXYZABCDEFGHIJKLM



Modular arithmetic representation:

- Encryption of a letter x by a shift n can be described mathematically as

$$E_n(x) = (x+n) \bmod 26$$

- Decryption is performed in a similar manner

$$D_n(x) = (x-n) \bmod 26$$

Key space is ridiculously small, very easy to break

Substitution Cipher

Idea: use a permutation over the set of characters as key to get a more flexible scheme as in the shift cipher

- Keyspace significantly larger
- Character frequencies are preserved

• a	• b	• c	• d	• e	• f	• g	• h	• i	• j	• k	• l	• m
• F	• G	• N	• E	• A	• T	• X	• Z	• O	• I	• Q	• B	• H
• n	• o	• p	• q	• r	• s	• t	• u	• v	• w	• x	• y	• z
• K	• Y	• W	• V	• C	• P	• J	• L	• S	• D	• M	• U	• R

computerscience → NYHWLJACPN OAKNA



What is the size of key space in the substitution cipher assuming 26 letters?

$$|\mathcal{K}| = 26$$

$$|\mathcal{K}| = 26! \quad (26 \text{ factorial})$$

$$|\mathcal{K}| = 2^{26}$$

$$|\mathcal{K}| = 26^2$$



$$26! \approx 2^{88}$$



Breaking Monoalphabetic Ciphers

Monoalphabetic ciphers preserve the frequency of alphabetic characters, pairs, etc.

→ *Identify alphabetic characters due to their frequency*

Method to decipher natural languages:

1. Determine frequency of alphabetic characters of the cipher text
2. Identify alphabetic characters according to their frequency: *e, n, i, s, r, a, t*
(in Germany: *e, n, r, i, s, t, u, d, a, g, l, o, ...*)
3. Determine frequency of pairs
4. Identify e.g. *th he*
5. Look at identified text, re-substitute, guess, ...



Breaking Monoalphabetic Ciphers

letter	probability
a	.082
b	.015
c	.028
d	.043
e	.127
f	.022
g	.020
h	.061
i	.070
j	.002
k	.008
l	.040
m	.024

letter	probability
n	.067
o	.075
p	.019
q	.001
r	.060
s	.063
t	.091
u	.028
v	.010
w	.023
x	.001
y	.020
z	.001

Partition into five groups:

1. **E**, having probability about 0.12
2. **T,A,O,I,N,S,H,R**, each having probabilities between 0.06 and 0.09
3. **D,L**, each having probabilities around 0.04
4. **C,U,M,W,F,G,Y,P,B**, each having probabilities between 0.015 and 0.028
5. **V,K,J,X,Q,Z**, each having probabilities less than 0.01

Digram frequencies

th	.0315	in	.0169
he	.0251	er	.0154
an	.0172	re	.0148

Vigenere Cipher

Popular polyalphabetic substitution cipher

- Known as “le chiffre indéchiffrable” (‘the indecipherable cipher’);-)
- Combination of simple substitution ciphers
- Rotations determined by a word (key)

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

Message	T	H	I	S	A	T	E	S	T
Running key	K	E	Y	K	E	Y	K	E	Y
	+10	+4	+24	+10	+4	+24	+10	+4	+24
Ciphertext	D	L	G	C	E	R	O	W	R

Polyalphabetic **Period**

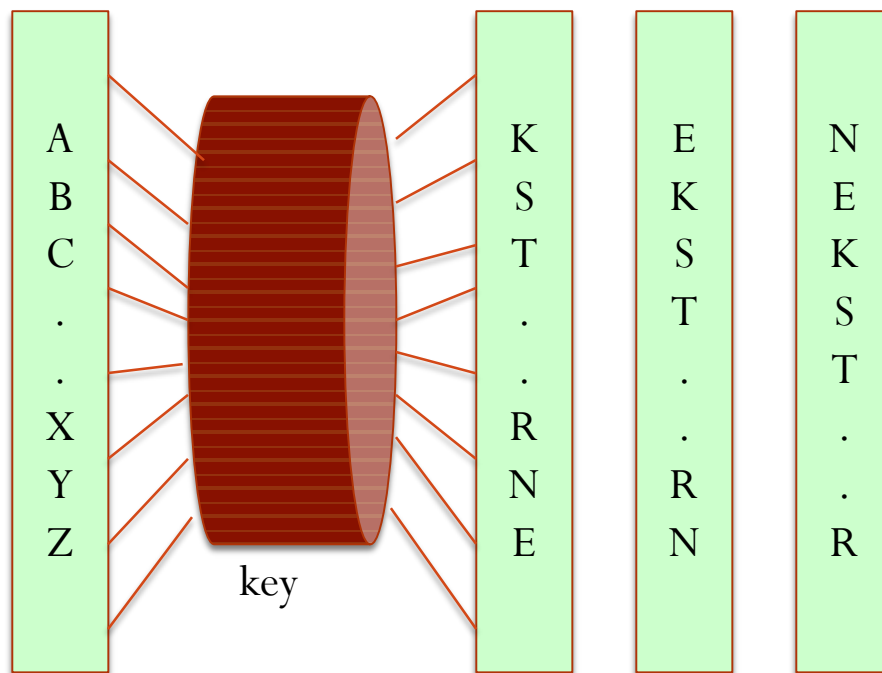
Breaking Vigenere Cipher

- Frequency analysis trivial if period can be guessed
- **Kasiski test**
 - Repeated words may, by chance, sometimes be encrypted using the same key letters, leading to repeated groups in the ciphertext
 - Consider the following encryption using the keyword ABCD
Key: ABCD ABCD ABCD ABCD ABCD ABCD ABCD
Plaintext: **CRYPTOIS SHOR TFOR CRYPTOGR APHY**
Ciphertext: **CSAS TPKV SIQU TGQU CSAS TPIU AQJB**
 - Repetitions of CSASTP is at a distance 16
 - Assuming that the repeated segments represent the same plaintext segments, this implies that the key is 16, 8, 4, 2, or 1 characters long



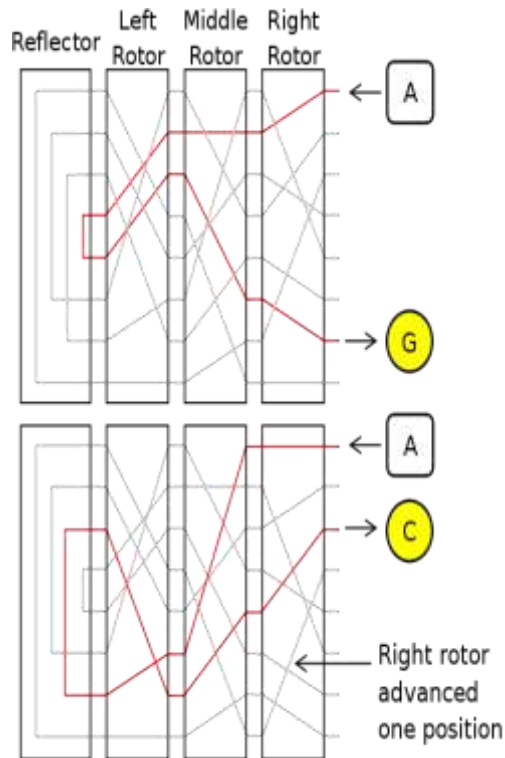
Rotor Machines (1870-1943)

- The Hebern Machine (single rotor)
 - Easily broken (CT only) using letter frequency, diagram frequency, trigram frequency



Rotor Machines (cont.)

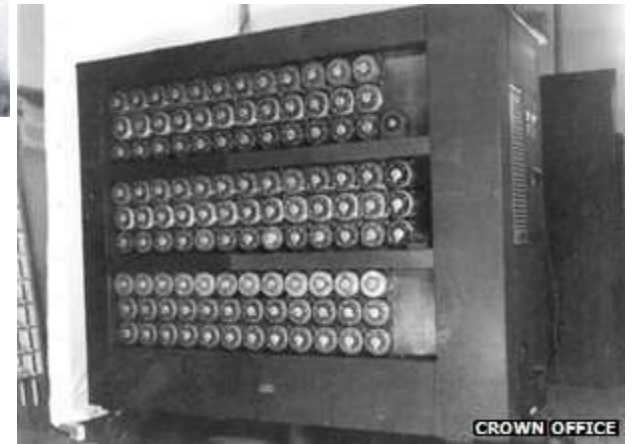
Most famous: the Enigma (3-5 rotors)



With 4 rotors keys = $26^4 = 2^{18}$ (actually 2^{36} due to optional plugboard)



Turing Bombe



FAST-NUCES

Must watch

“The Man Who Cracked Enigma”



FAST-NUCES

Acknowledgements

Material in this lecture are taken from the slides prepared by:

- Prof. Dr. Konrad Rieck (Uni-Göttingen)
- Prof. Dan Boneh (Stanford)

