

File: lbriaghat2  
21K-3153

```

10
#define array_size 10
int main() {
    int array1[array_size]    // creating arrays
    int array2[array_size]
    int array3result[array_size] //
}

for (int i = 0; i < array_size; i++) {
    array1[i] = i
    array2[i] = 2 * i
    #pragma omp parallel for
    for (int i = 0; i < array_size; i++)
        result[i] = array1[i] + array2[i];
    printf("Array 1:")
    for // print array1
    printf("Array 2:")
    // print array2
    # print result array
    return 0
}

```

after creating 3 arrays, initializing array1 and 2 with values we use "parallel for" to add the values into our 3rd array. "#pragma omp parallel" creates threads and "for" distributes workload in the for loop between threads.

Since we have used parallel for, we need to declare a parallel block as it is implicitly understood that threads will execute for loop because of "parallel for" pragma.

Static scheduling of reduction:

```
#define array_size 10
int main() {
    int array[array_size]
    int sum = 0
```

```
for(int i=0; i<array_size; i++)
    array[i] = i+1;
```

```
#pragma omp parallel for schedule(static) reduction(+:sum)
for (int i=0; i<array_size; i++)
    sum += array[i];
```

```
#include <iostream>
// print array
print f("/n %d", sum);
return 0
}
```

Here, we are initializing an array and summing its elements in the sum variable.

"#pragma omp parallel" to create threads. "for" to specify that they will run in a for the "for loop"

schedule (static) ensures that iterations will be performed by threads cyclically. For example, thread 0 has iteration 1, thread 1 has iteration 2 and so on. Assume 4 threads, thread 0 will do iteration 4, thread 1 will do iteration 5 and so on.

The reduction clause makes the sum variable private to each thread, which prevents the variable being changed and ensures consistency by something else.



## ③ dynamic scheduling reduction

#define size 10

int main()

int array1 [size]

int array2 [size]

int result [size]

for (i &lt; size)

~~array1~~ array1[i] = i

array2[i] = i \* 2

#pragma omp for parallel for schedule(dynamic)

for (int i = 0; i &lt; size; i++)

result[i] = array1[i] + array2[i];

printf("Array1")

// print array 1

printf("Array2")

// print array 2

printf("result")

// print result

return 0

}

~~Two~~ 3 arrays created,  
two initialized~~Two~~ For Each element  
of the 2 arrays is  
summed into 3rd array.~~omp~~ pragma omp  
parallel to create threads  
for to specify workload  
of for loop.Dynamic scheduling  
makes it so that  
when one thread finishes  
it's task, it communicates  
with other threads  
to find an iteration which  
is not being worked on,  
and executes that iteration.

However it is expensive.

```

4) int main()
   int main() {
       int array[size]
       int sum = 0

       for (i < size)
           array[i] = i + 1

       #pragma omp parallel shared(sum)
       {
           int local_sum = 0
           #pragma omp for
           for (i < size)
               local_sum += array[i]

           #pragma omp critical
           {
               sum += local_sum
           } // end of critical
       } // end of parallel

       printf("sum")
       return 0
   }

```

threads. "initialized array and ~~are~~ summing it. "omp parallel" to create threads. "shared (sum)" shares the global variable sum ~~between~~ threads.

every thread sums their part of the array in ~~the~~ local\_sum variable.

~~to ensure~~ one thread at a time then accesses critical section and update the shared sum variable.





## ⑤ Collapse clause

```
#pragma omp for collapse(2) printe (i,k,j)
for (k=kl, kc=ku, k+=ks)
    for (j=jl; j<=ju, jr=js)
        for (i=il; i<=iu, it=is)
            bar (a[i][j],k)
```

J and k are part of the loop construct, so they are collapsed into 1 loop. i is not part of the loop construct this is executed in every iteration of collapsed J and k

## ⑥ #pragma omp sections

```
#pragma omp section
{ task() }
#pragma omp section
{ taskB() }
}
```

~~now~~ this allows  
for non iterative  
parallel tasks.

## ⑦ Vector add and multiply using sections

```
int main()
{
    int vectorA[size], vectorB[size], resultAdd[size], resultMul[size]
    for (i=0; i<size; i++)
        vectorA[i] = i*1
        vectorB[i] = i*2
}
```

```
#pragma omp parallel sections
{
    #pragma omp section
    {
```

```
        for (i=0; i<size; i++)
            resultAdd[i] = vectorA[i] + vectorB[i]
    }
}
```

day / date:

```
#pragma omp section
{
    for (i < size)
        result[i] = vector A[i] + vector B[i]
}
```

```
Print result of array
Print result of array
return 0
}
```

- this splits ~~parallel~~ sections into different sections to be parallelized by different threads



day / date:

⑧ print message using sync clause

```
int main() {  
    omp_set_num_threads(4);  
    #pragma omp parallel  
    {  
        printf("%d", or get_thread_num());  
        #pragma omp barrier  
  
        #pragma omp single  
        {  
            printf("barrier complete")  
        }  
  
        #pragma printf("%d's print second message", omp_get_thread_num()  
    )  
}
```

omp parallel to create threads.

omp barrier to wait for each thread to print 1<sup>st</sup> message

omp single for only 1 thread to print message

~~omp~~ then all threads print second message

① ~~Serially~~ ~~is that~~ ② generating prime via serial

```

int is_prime (int num)
{
    if (num <= 1) return 0
    if (num == 2) return 1

    if (num % 2 == 0) return 0
    for (i = 3, i * i <= num, i += 2)
        if (num % i == 0)
            return 0
}

```

```

return 1
}

int main ()
{
    int start = 2
    int end = 100

    #regions are parallel for
    for (i = start, i <= end; i++)
        if (is_prime(i))
            #pragma omp critical
            printf ("%d is prime prime", i)
    }

    return 0;
}

```

omp parallel for creates threads that run is\_prime check.  
 If is\_prime true, one thread at a time can access  
 critical section and print the prime number.