# appendix 1

#### 1. The generation process of rK in MN-SC mechanism

The generation process of rK in MN-SC mechanism is as follows:

## (1)Initialization stage

Distributor  $KGC_C$ ,  $NP_i = \{NP_1, NP_2, \dots, NP_n\}$ , master key is  $r_k$ . The distributor  $KGC_C$  randomly selects n points  $K_1, K_2, \dots, K_n$ , on the elliptic curve  $E_p(a, b)$ , and makes its vertical coordinate  $r_{k_1}, r_{k_2}, \dots, r_{k_n}$  as the subkey of  $NP_i$ , G is the base point selected on  $E_p(a, b)$ .  $KGC_C$  distributes its subkeys to  $NP_i$ ,  $i \in [1, n]$  over a secure channel.  $KGC_C$  calculates  $G'_i = k_i G$ ,  $i \in [1, n]$ . Then, parameters (E, G, n, H(x), G') is disclosed.

### (2)Key distribution

- a. Distributor  $KGC_C$  randomly selects a point Q on  $E_p(a, b)$  and a t-1 degree polynomial  $f(x) = a_0 + \sum_{i=1}^{t-1} a_i x^i$ , where  $f(0) = a_0 = r_K$  and Q is the public parametlates and f(x) is retained by the distributor.
- b. Distributor  $KGC_C$  calculates  $f(i) = r_{k_i}$ ,  $A_l = a_lG$ ,  $1 \le l \le t 1$ ,  $D_i = (i, f(i)) K_iQ$ . The parameters  $A_l$  and  $D_i$  are disclosed.
- c. Distributor  $KGC_C$  calculates  $F_i = H(K_iQ)$ , which is used to verify the authenticity of the subkeys interacted by the  $NP_i$ .

#### (3) Verification of subkeys

After receiving the subkey distributed by  $KGC_C$ ,  $NP_i(1 \le i \le n)$  first verifies the authenticity of the subkey. Then it calculates  $C_i = K_iQ$  and  $(i, r_{k_i}) = (D_i + C_i)$ . If the equation  $r_{k_i}G = \sum_{k=0}^{t-1} A_l i_l^i$  holds, then the subkey is true. If the equation  $F_i = H(C_i)$  holds, it means that the identity between  $NP_i$  is true.

### (4) Master key reconstruction:

When the number l of participants satisfies  $l \ge t$ , allow the participants be  $NP_{i_1}, NP_{i_2}, \dots NP_{i_l}$ , and the master key  $r_K$  can be reconstructed by Lagrange interpolation. For the subkeys  $r_{k_i}, i \in [1,t]$  of t different  $NP_i$ , a unique Lagrange interpolation polynomial  $L(x) = \sum_{i=0}^{t-1} r_{k_i} l_i(x)$  can be constructed by  $(x_i, r_{k_i})$ , where  $l_i(x) = \prod_{0 \le j \le t-1} \frac{x-x_i}{x_i-x_j}, 0 \le i \le t-1$ . Because L(x) is a polynomial of degree t-1 and corresponds to different values, there is  $L(0)=f(0)=r_k$ . Get the master key and send  $r_k$  to  $NP_1$ .