Introduction to Apache Cassandra

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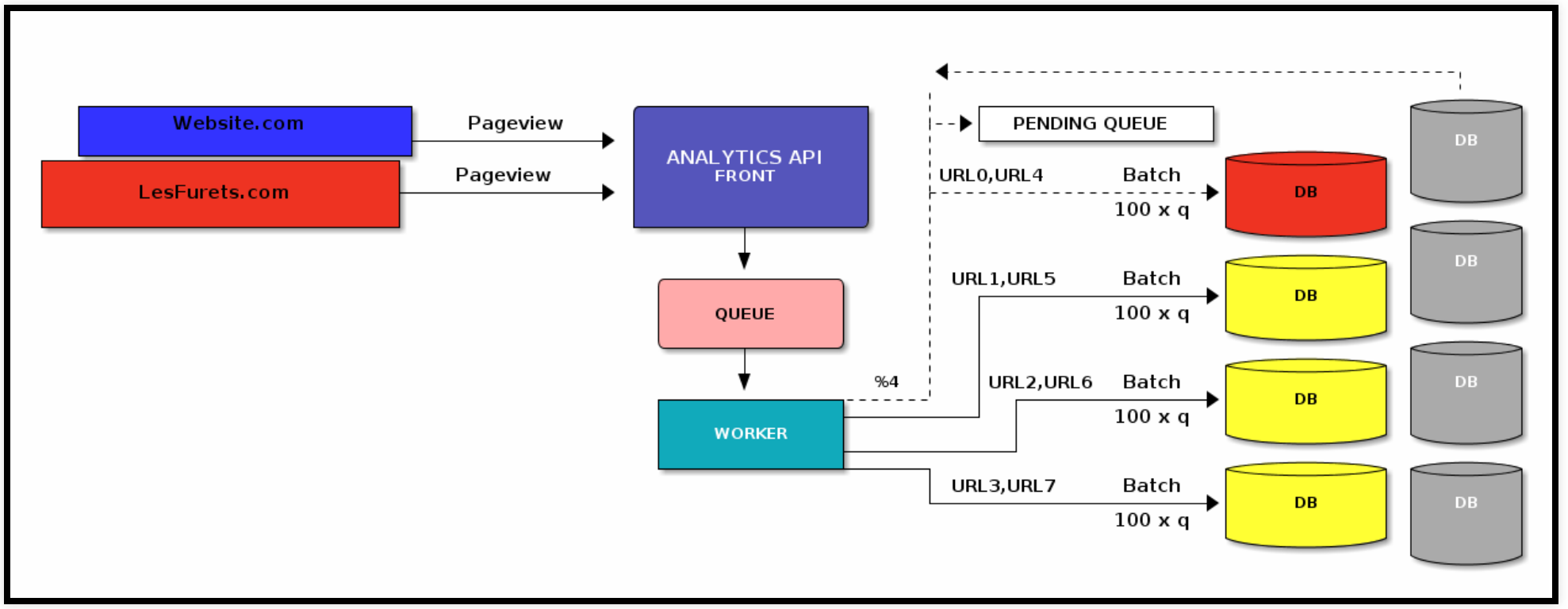
Plan

- **Motivation**
- Apache Cassandra
  - Partitioning and replication
  - Consistency
- Practice: Tune consistency in Apache Cassandra

Motivation



- do I build new features for customers?
- or just dealing with reading/writting the data?

What went wrong?

- **A single server cannot take the load ⇒** solution / **complexity**
  - *Better database*
    - easy to add/remove nodes (**scalling**)
    - transparent data distribution (**auto-sharding**)
    - handle failures (**auto-replication**)

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Apache Cassandra

- started @Facebook inspired by *BigTable* model and *Amazon DynamoDB*
- 2008 Open Source Project
- **Datastax**: commercial offering Datastax Enterprise
  - monitoring(OpsCenter) automating repairs backup…
  - other features: search, analytics, graph, encryption
- 2010 Top Level Apache Project
  - Datastax biggest committer
- 2016 Datastax/ASF disagreements(#Staxit)

Apache Cassandra

Open source, distributed database *designed to handle large amounts of data across many commodity servers, providing* ***high availability*** *with* ***no single point of failure****. It offers robust support for clusters spanning* ***multiple datacenters****, with* ***asynchronous masterless replication*** *allowing* ***low latency*** *operations for all clients.*

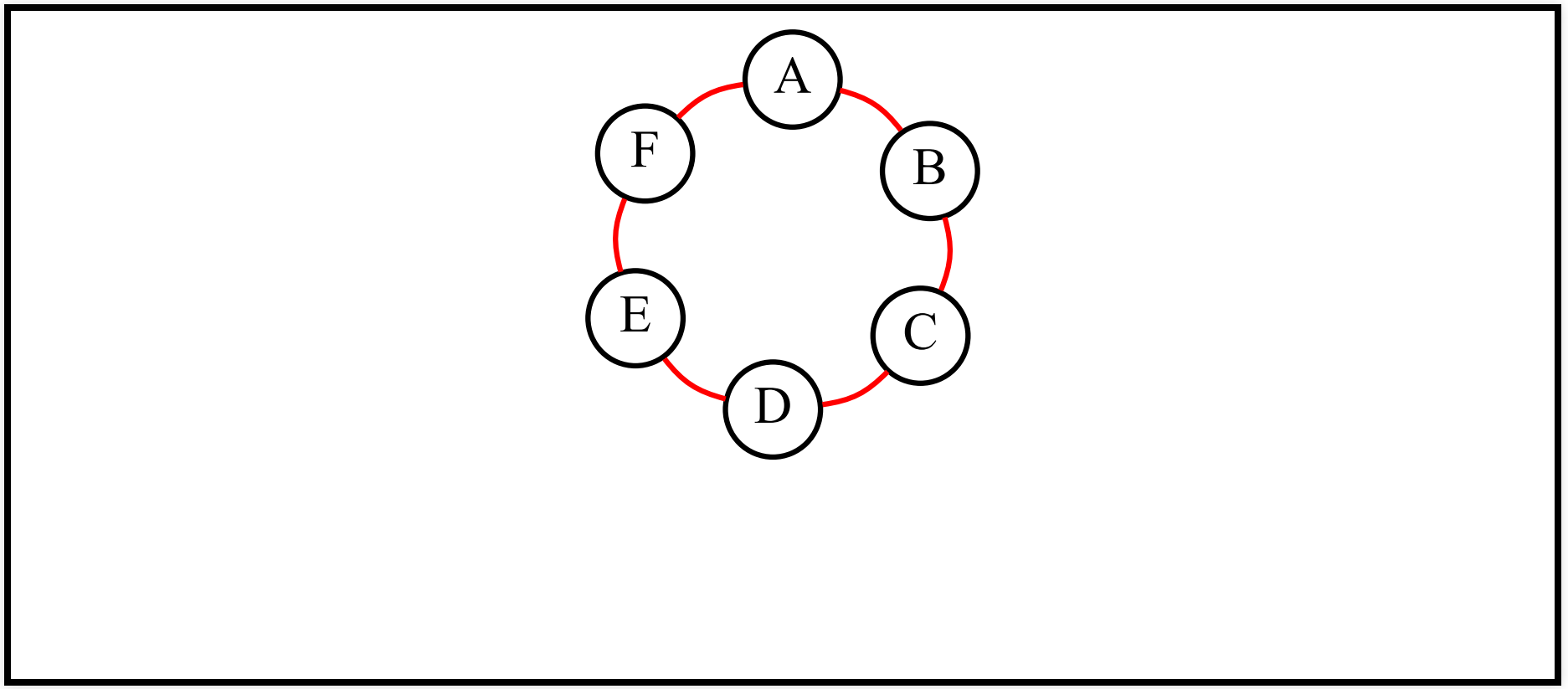
Cassandra

- column oriented NoSQL database
- distributed (data, query)
- resilient (no SPOF)
  - we can query any node ⇒ coordinator to dispatch and gather the results

Cassandra terminology

RDBMS	**Cassandra**
Schema (set of tables)	***Keyspace***
Table	***Table****/column family*
Row	Row
Database server	***Node***
Master/Slave	***Cluster***: *a set of nodes groupped in one or more datacenters (can span physical locations)*

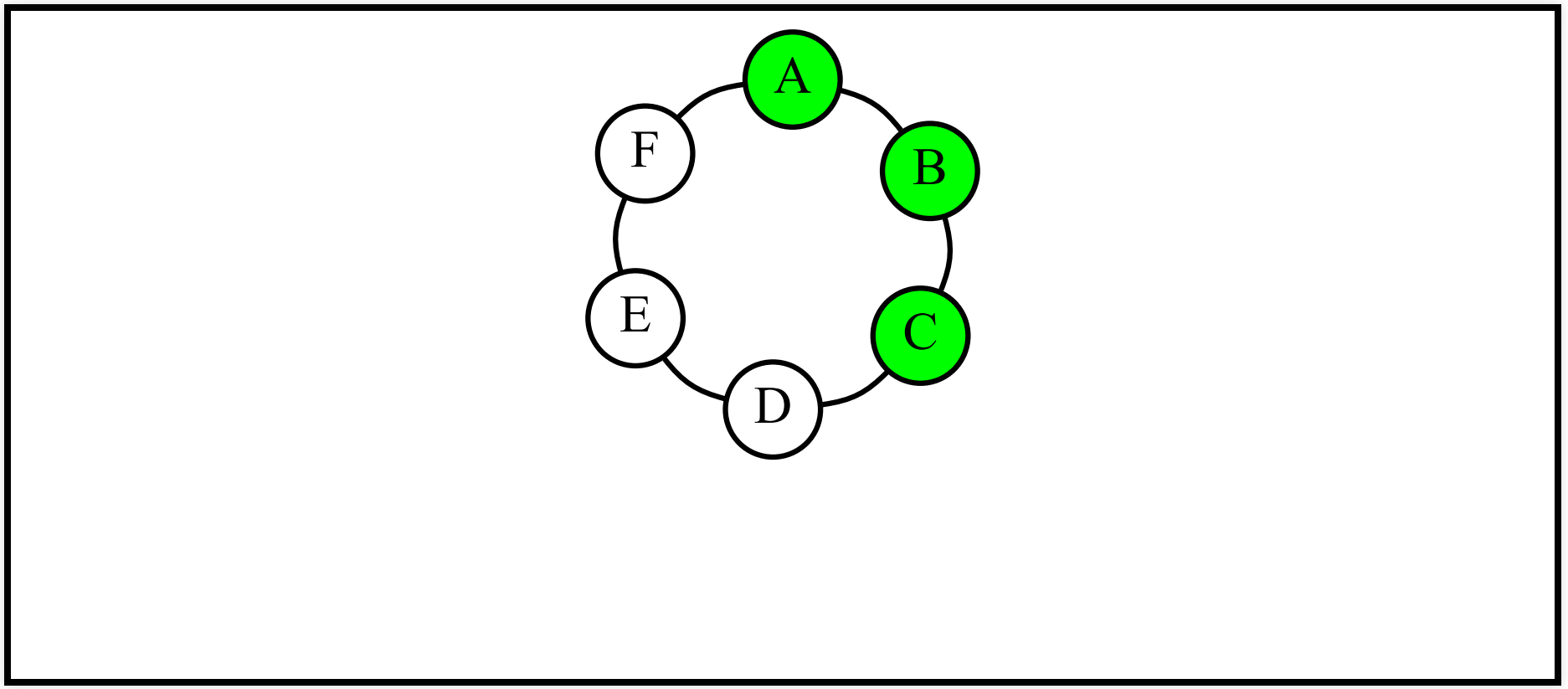
Cassandra ring



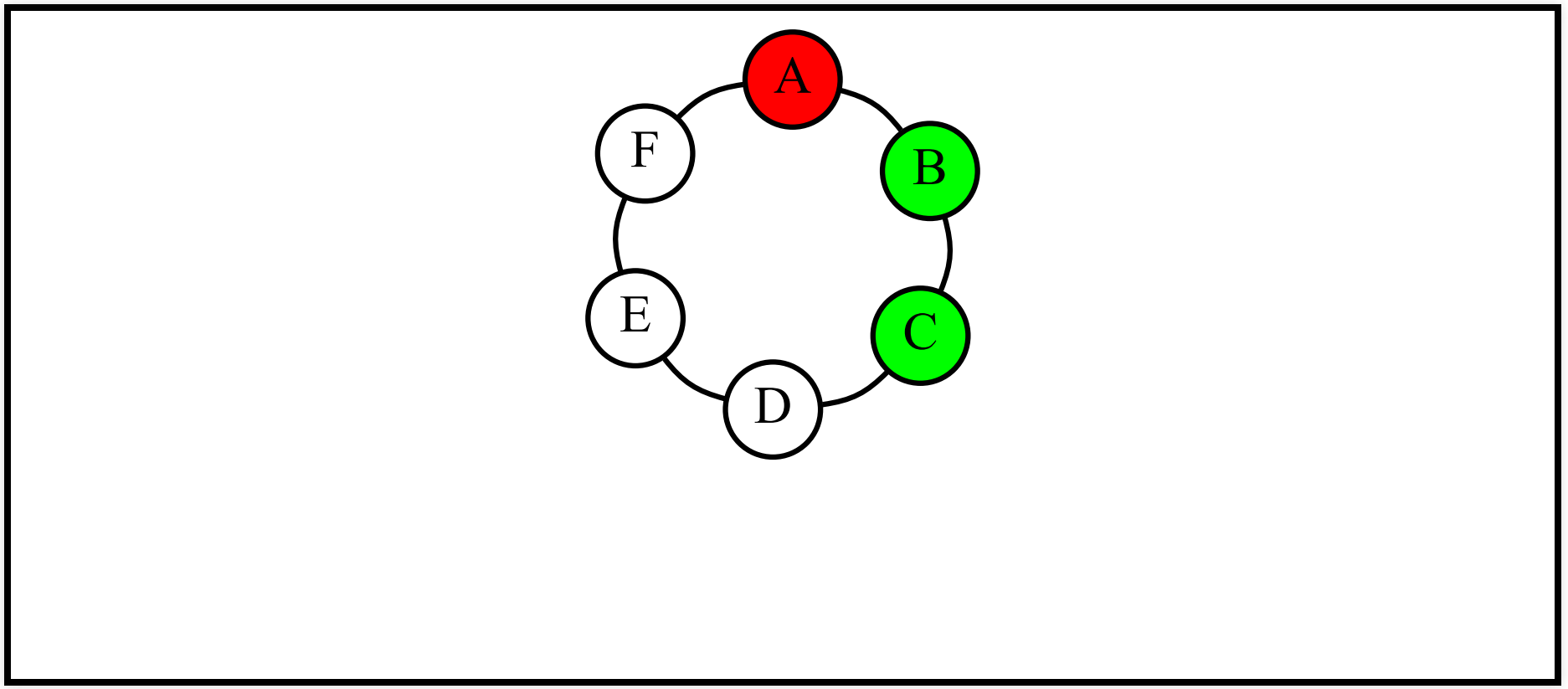
Gossip

- peer-to-peer communication protocol
  - discover and share ***location*** and ***state information*** about nodes
  - persist gossip info locally to use when a node restarts
- **seed nodes** ⇒ bootstrapping the gossip process for new nodes joining the cluster

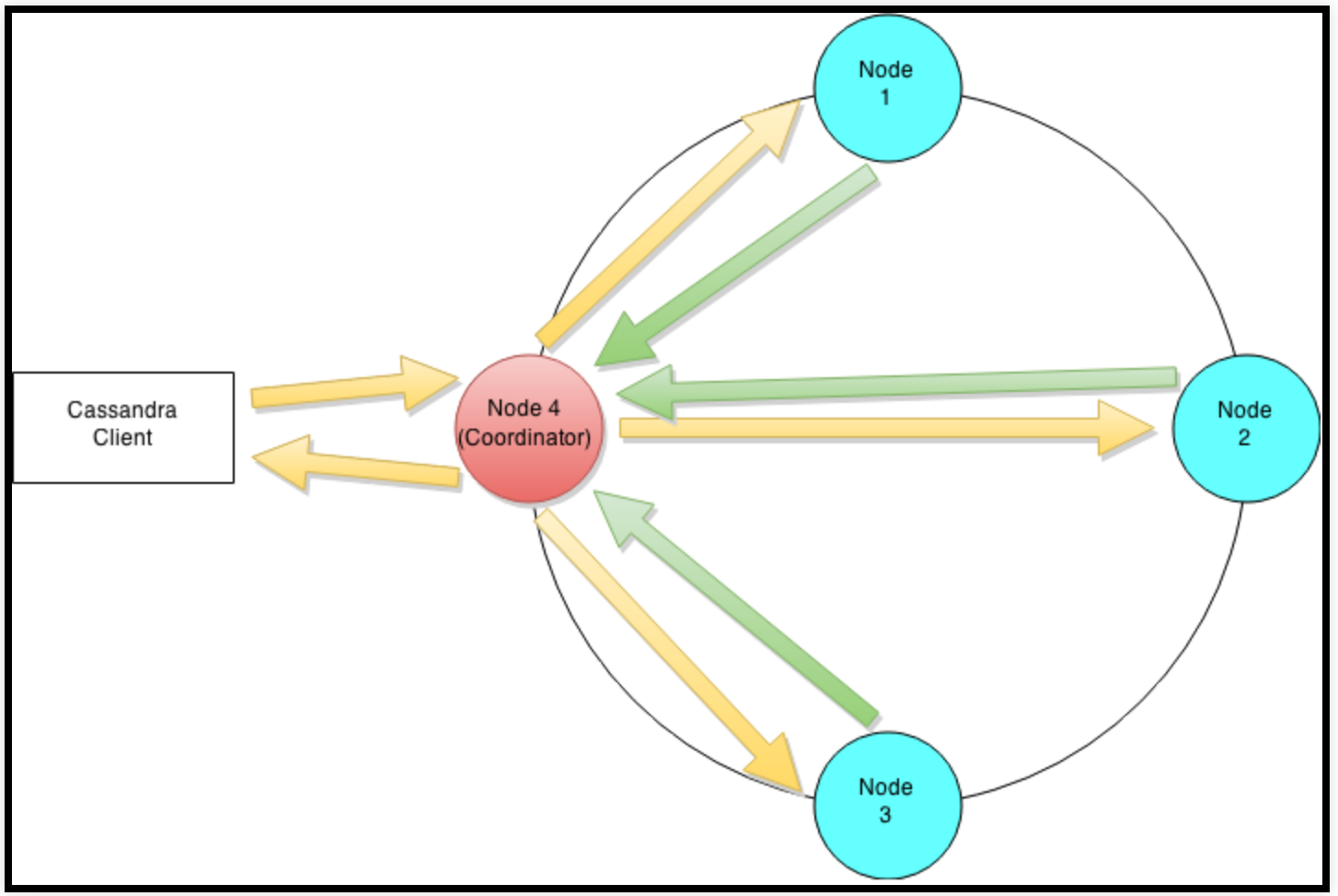
Cassandra replicas



Cassandra node failure



Query



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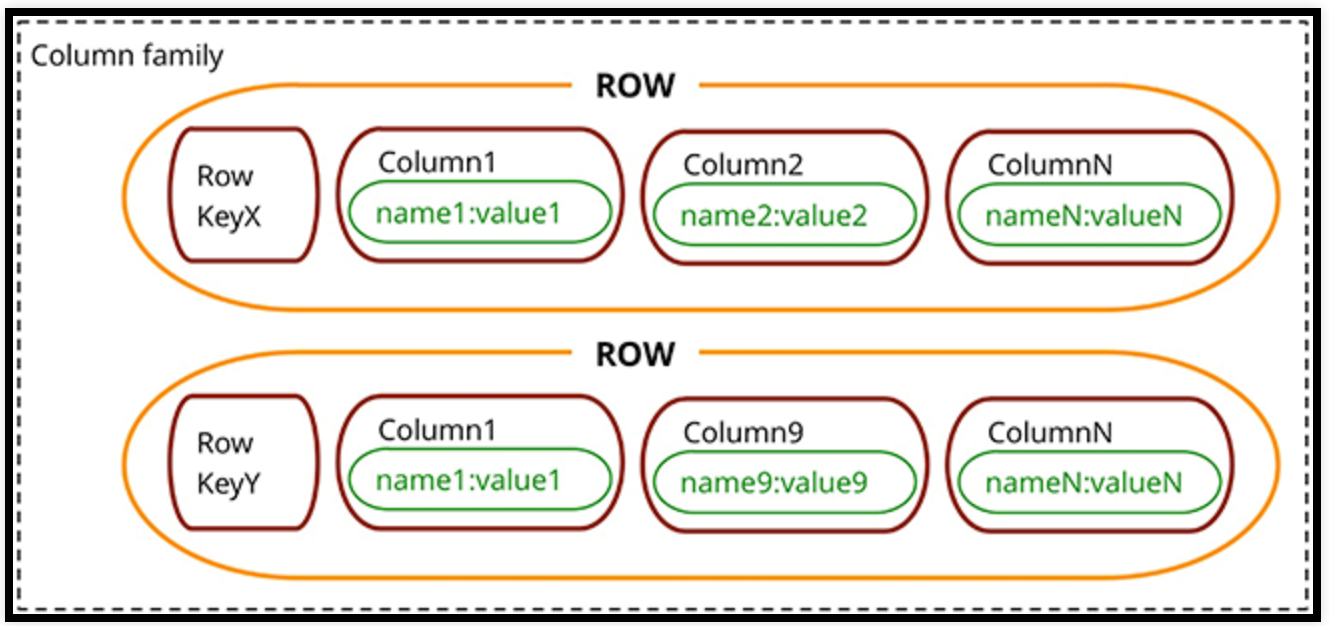
Cassandra data model

- Stores data in **tables/column families** as rows that have many columns associated to a row key
- Free format: two rows can have different columns
- **Rows are the atomic agregate** ⇒ distributed and replicated

"Map<RowKey, SortedMap<ColumnKey, ColumnValue>>"



Data partitioning

- C* = single logical database spread across a cluster of nodes
- ***How to divide data evenly around its cluster of nodes?***

- **Problems**:
  - distribute data **efficiently**, **evenly**
    - determining a node on which a specific piece of data should reside on
    - minimise the data movements when nodes join or leave the cluster
- Algorithm of **Consistent Hashing**

Mapping data to nodes

- **Problem**: map **k** entries to **n** physical nodes
- *Naive hashing (NodeID = hash(key) % n)* ⇒ remap a large number of keys when nodes join/leave the cluster
- ***Consistent hashing***: only k/n keys need to be remapped on average

Consistent Hashing

- **Idea** :
  - use a part of the data as a partition key
  - compute a hash value for each
  - The range of values from a consistent hashing algorithm is a fixed circular space which can be visualised as a ring.

Partitioner

- **hash function** that derives a token from the primary key of a row
- determines which node will receive the ***first replica***
- RandomPartitioner, **Murmur3Partitioner**, ByteOrdered

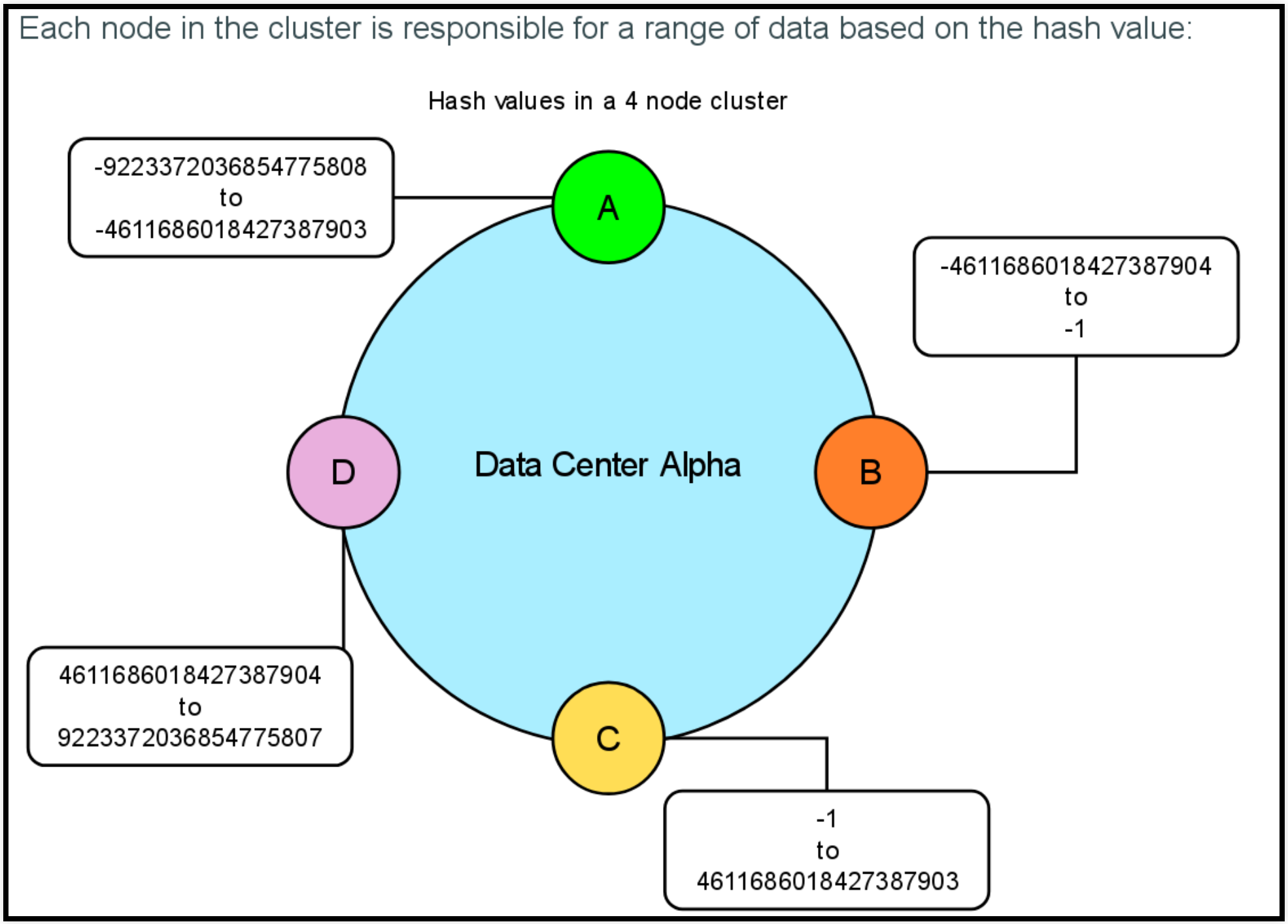
Murmur3Partitionner

name	age	car	gender
jim	36	camaro	M
carol	37	bmw	F
johnny	12		M
suzy	10		F

Cassandra assigns a hash value to each partition key:

Partition key	Murmur3 hash value
jim	-2245462676723223822
carol	7723358927203680754
johnny	-6723372854036780875
suzy	1168604627387940318

Consistent Hashing: mapping



Consistent Hashing: mapping

Node	Start range	End range	Partition key	Hash value
A	-9223372036854775808	-4611686018427387903	johnny	-6723372854036780875
B	-4611686018427387904	-1	jim	-2245462676723223822
C	0	4611686018427387903	suzy	116860462738794031 8
D	4611686018427387904	9223372036854775807	carol	772335892720368075 4

Data Replication

- create copies of the data, thus avoiding a single point of failure.
- **Replication Factor (RF)** = # of replica for each data
  - set at the *Keyspace* level

How data is replicated: ***Snitches***

- determine the *proximity of nodes within a ring*
  - *Dynamic snitch* ⇒ monitors the performance of reads from the replicas and chooses the best replica
  - ***SimpleSnitch*** ⇒ place the copy to the next available node (clockwise)
  - *RackInferringSnitch* ⇒ place copies of rows to different racks (same DC)

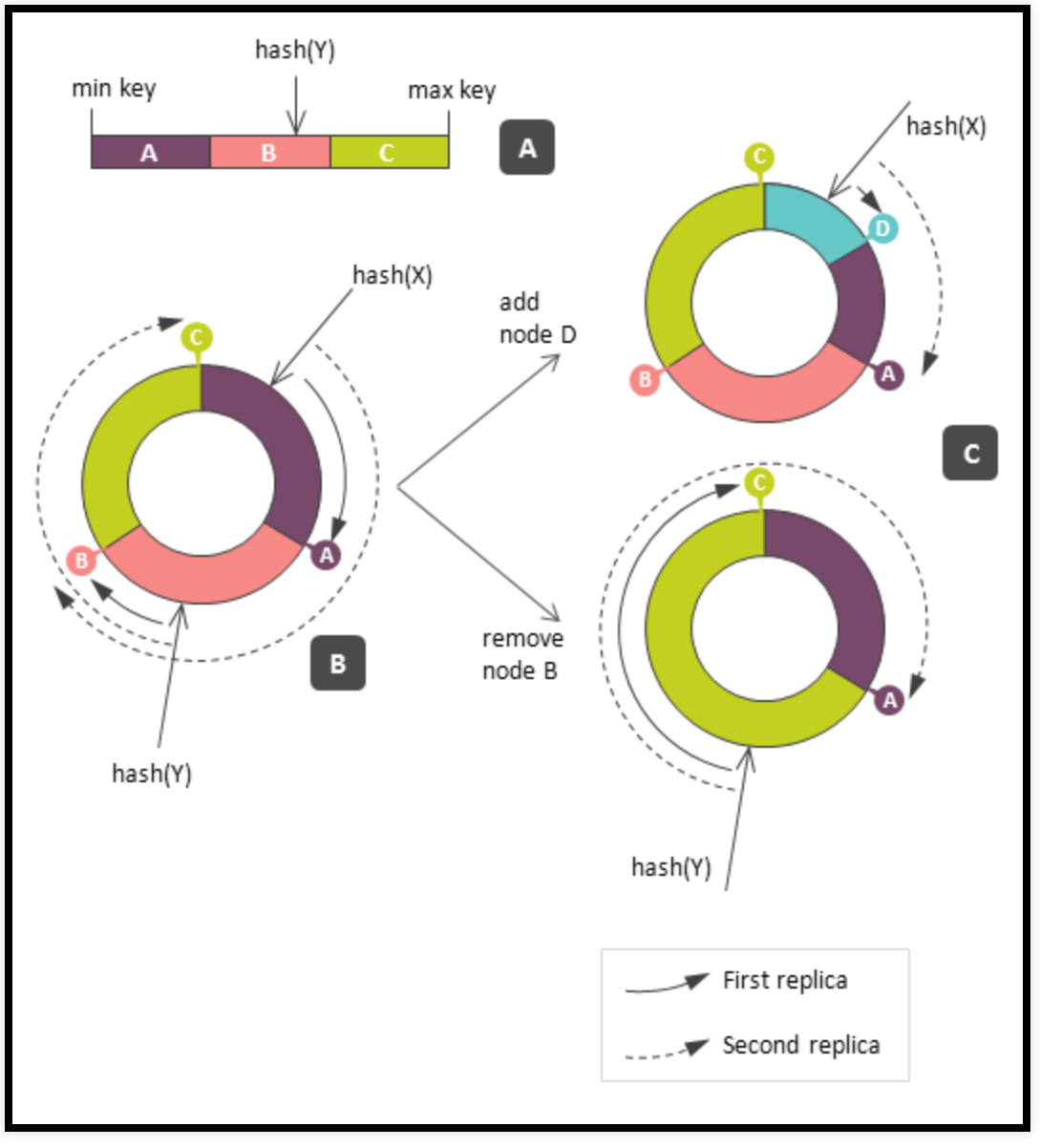
Configured in the cassandra.yaml file.

How data is replicated: ***Replication strategies***

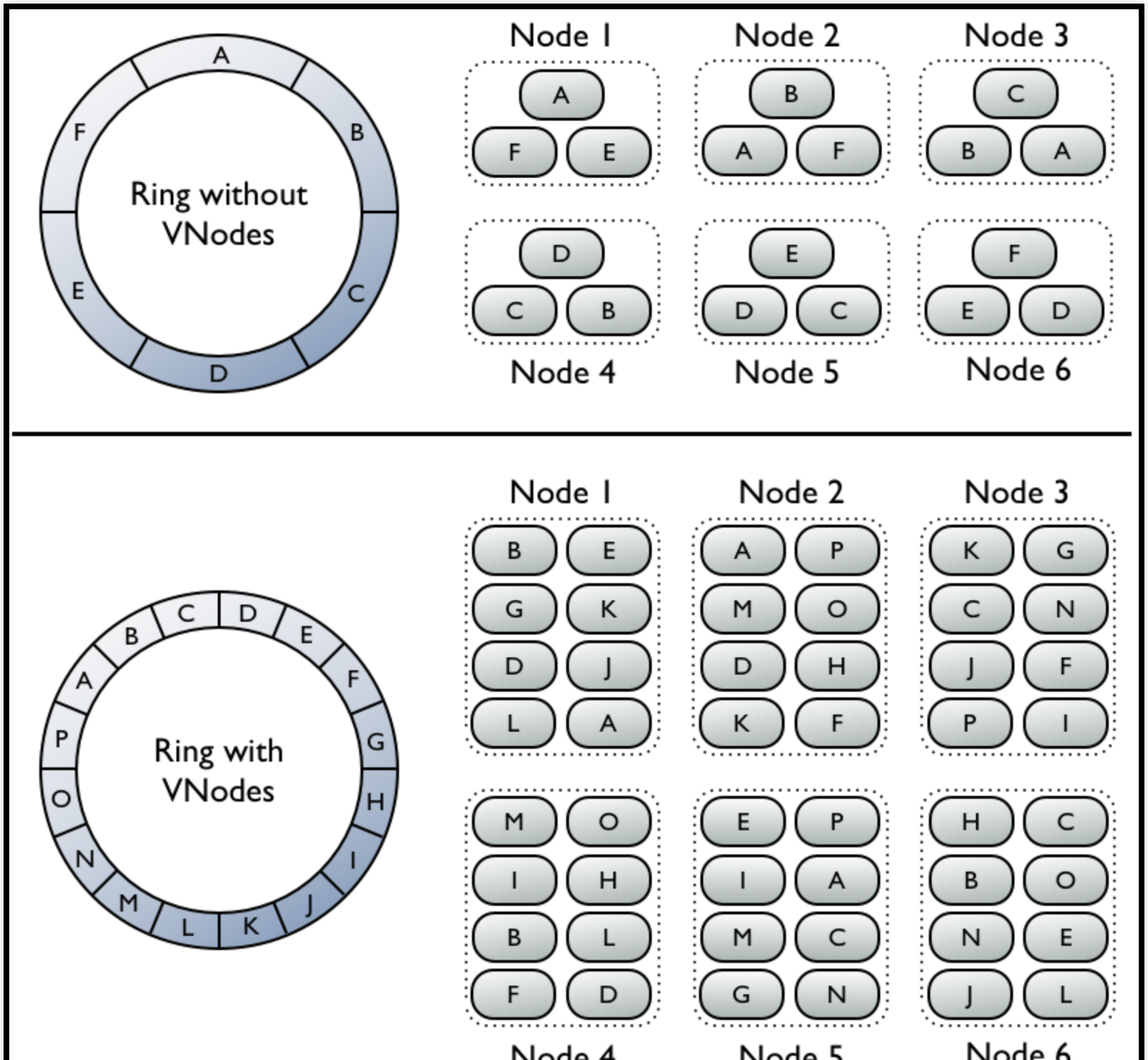
- use proximity information provided by snitches to determine locality of a copy
  - **SimpleStrategy**: use only for a single data center and one rack
  - **NetworkTopologyStrategy**: specifies how many replicas you want in each DC
- defined at ***keyspace*** level

```
CREATE KEYSPACE temperature
WITH replication =
 {'class': 'SimpleTopologyStrategy', 'replication_factor':'2'};
```

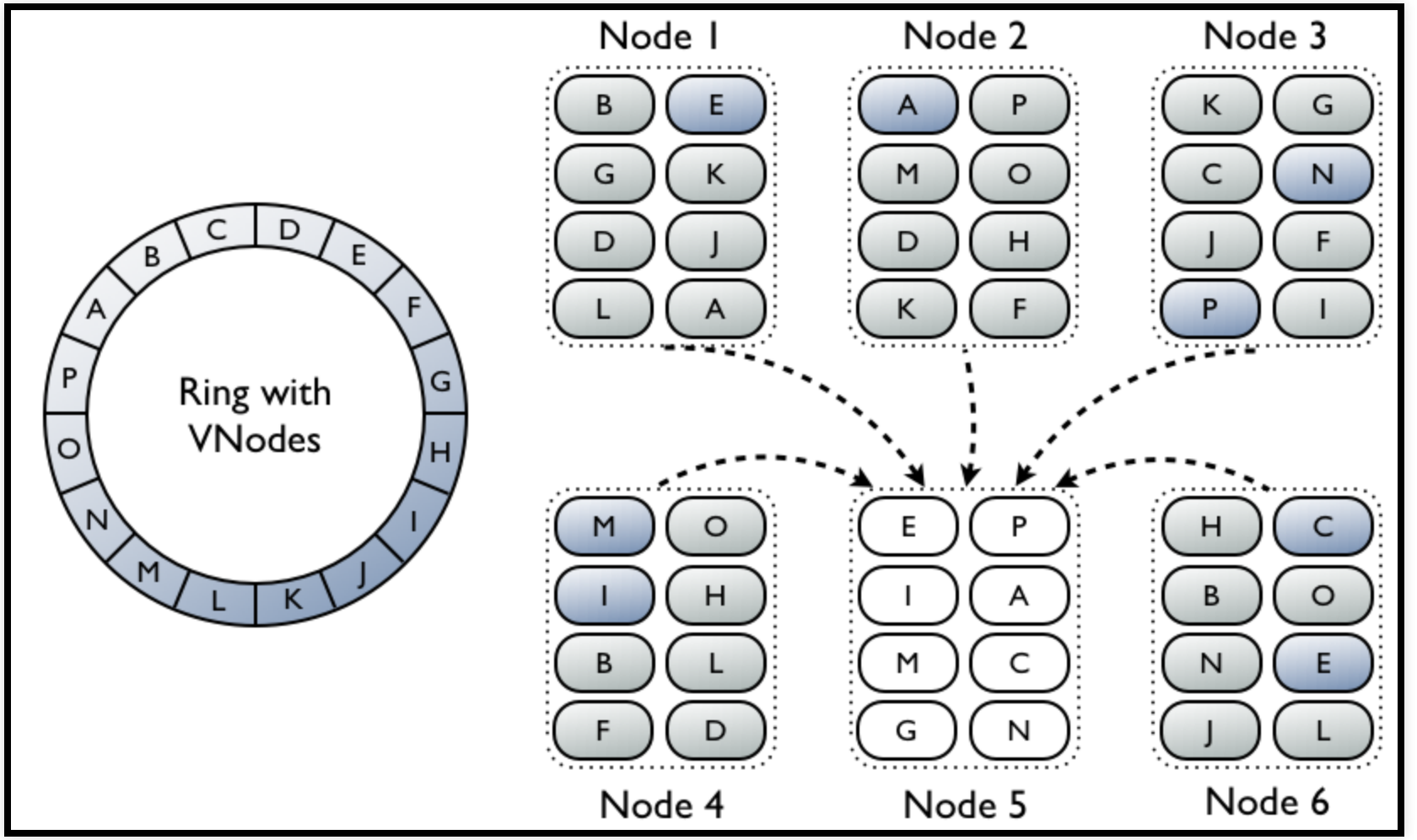
Example: RF2 + SimpleSnitch + SimpleReplicationStrategy



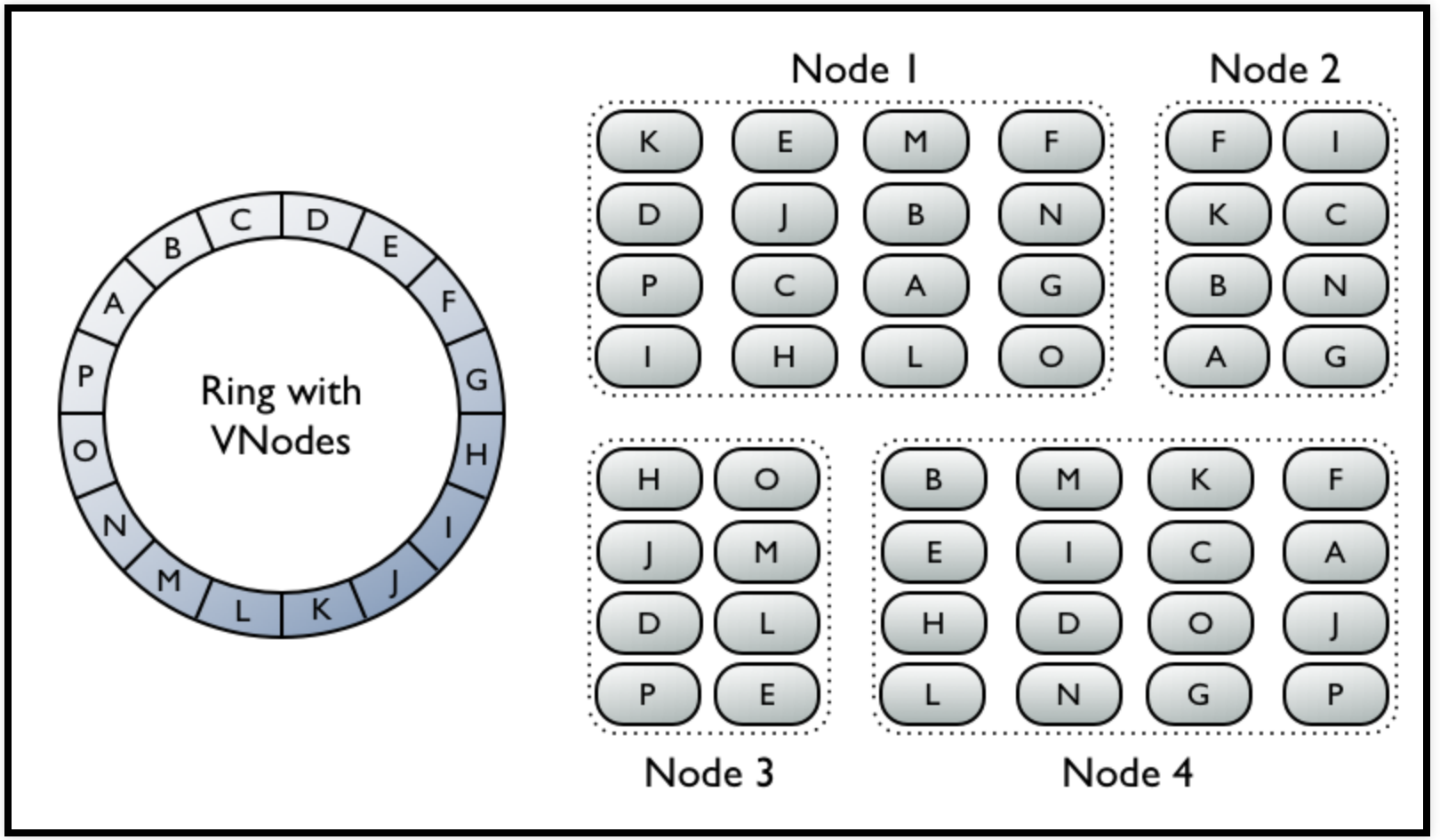
Virtual nodes (VNodes)



VNodes: remapping keys



VNodes and heteregenous clusters

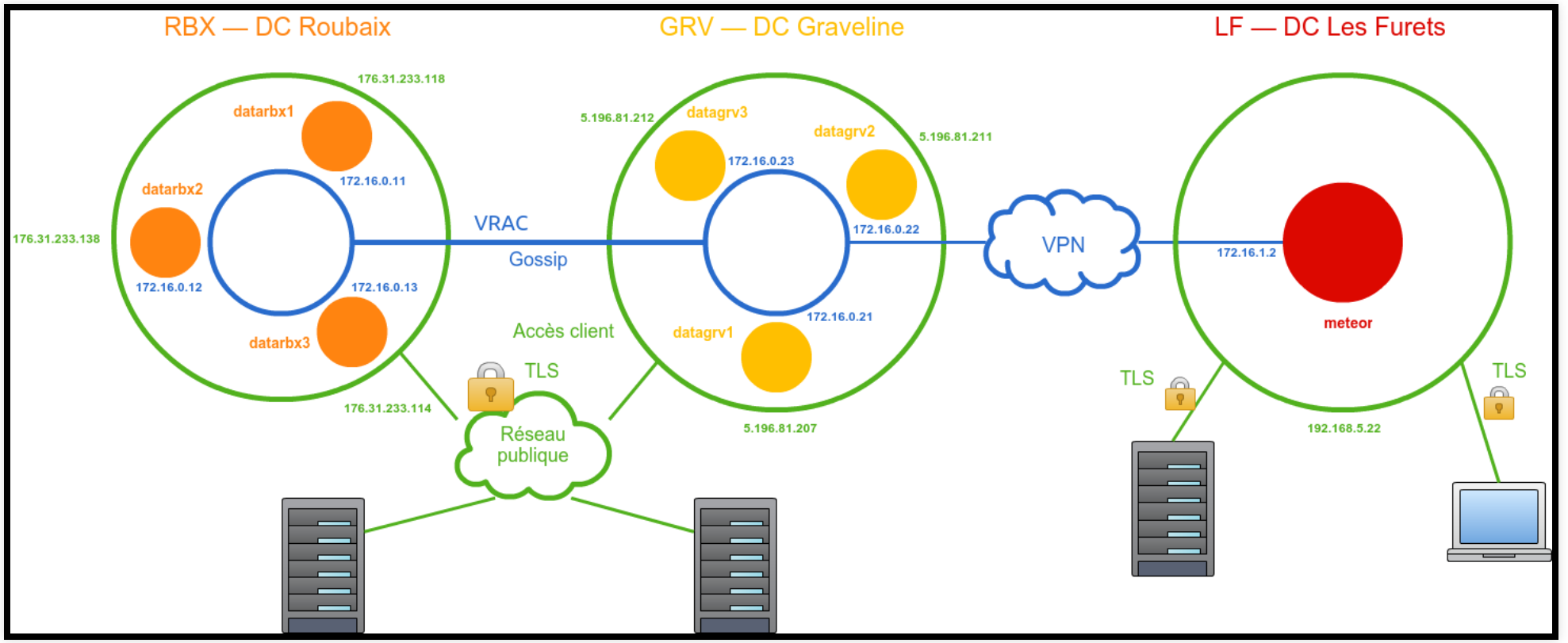


Token allocation

- no VNODES ⇒ manual (initial-token="-29334…" dans cassandra.yaml)
  - need to be modified at each topology change
- VNODES (num_tokens)
  - random slot allocation (< 3.0)
  - smart (3.0+)

Example: NetworkTopologyStrategy

```
CREATE KEYSPACE lesfurets
WITH replication =
{'class': 'NetworkTopologyStrategy', 'RBX': 2,'GRV':2,'LF':1};
```



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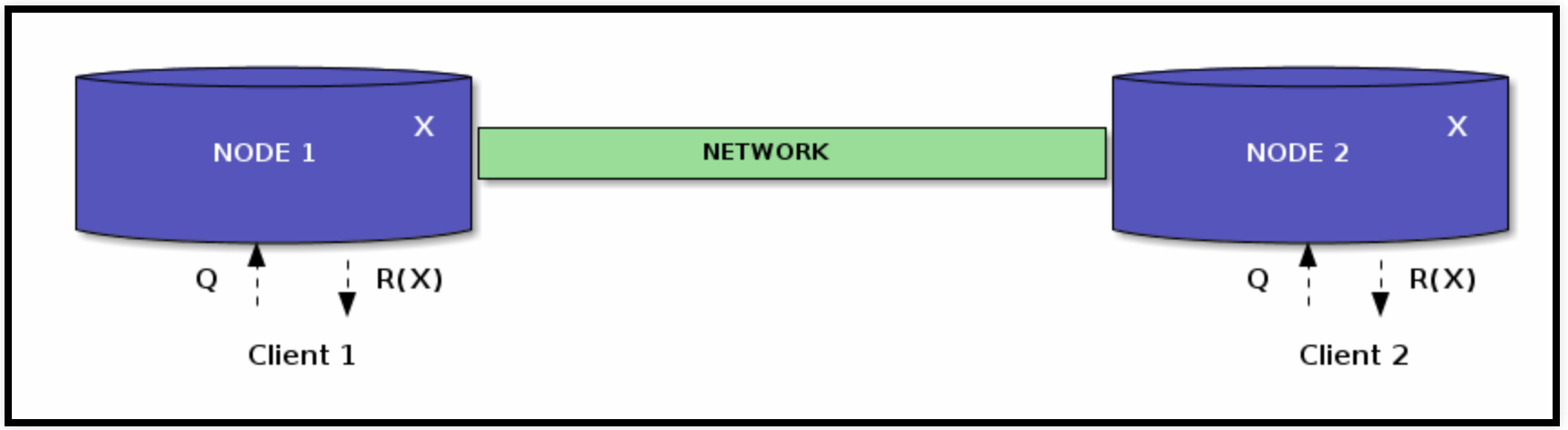
Properties of distributed systems

- **Consistency**: read is guaranteed to return the most recent write for a given client.
- **Availability**: non-failing node will return a reasonable response within a reasonable amount of time (no error or timeout)
- **Partition Tolerance**: the system will continue to function when network partitions occur.

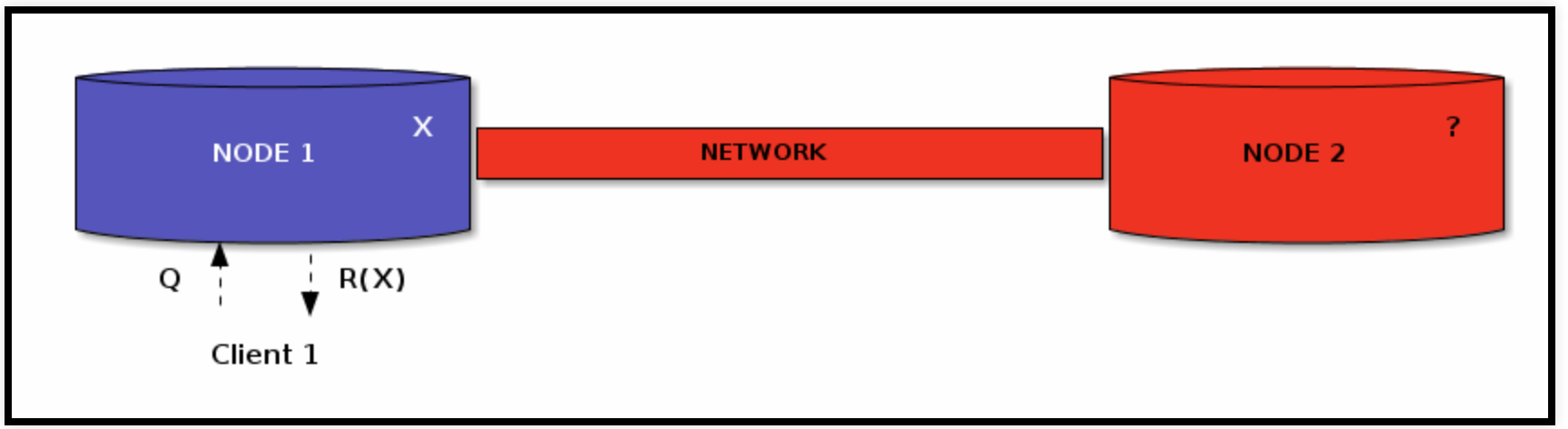
Consistency

- *a read returns the most **recent** write*
- *eventually consistent* : guarantee that the system will evolve in a consistent state
  - provided there are no new updates, all nodes/replicas will eventually return the last updated value (~DNS)



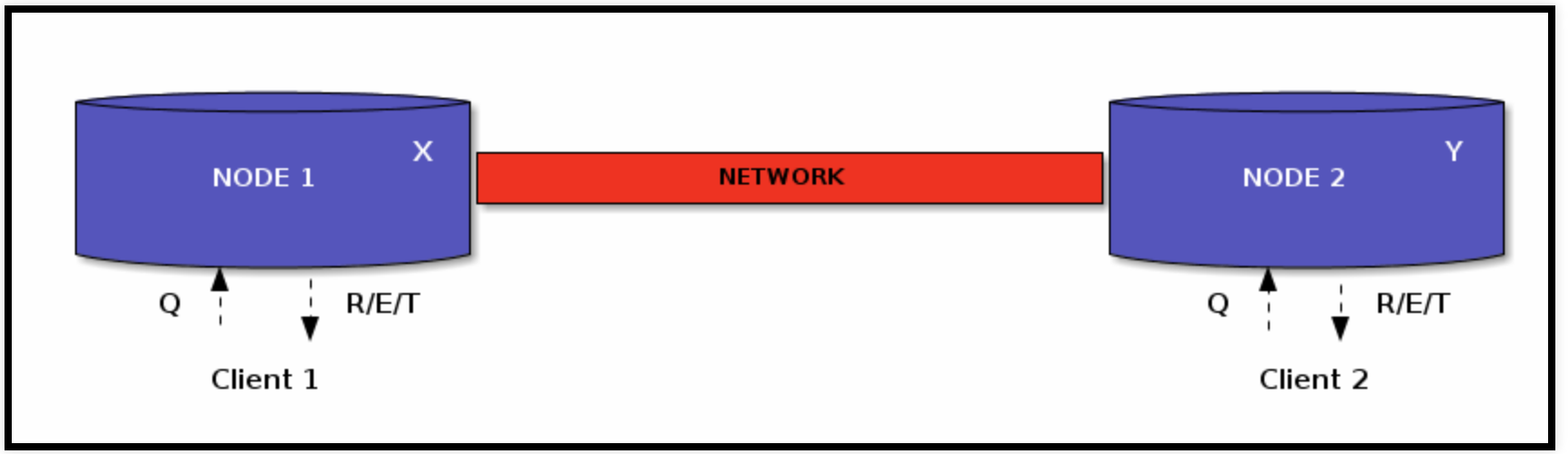
Availability

- *a non-failing node will return a reasonable response (no error or timeout)*



Partition tolerance

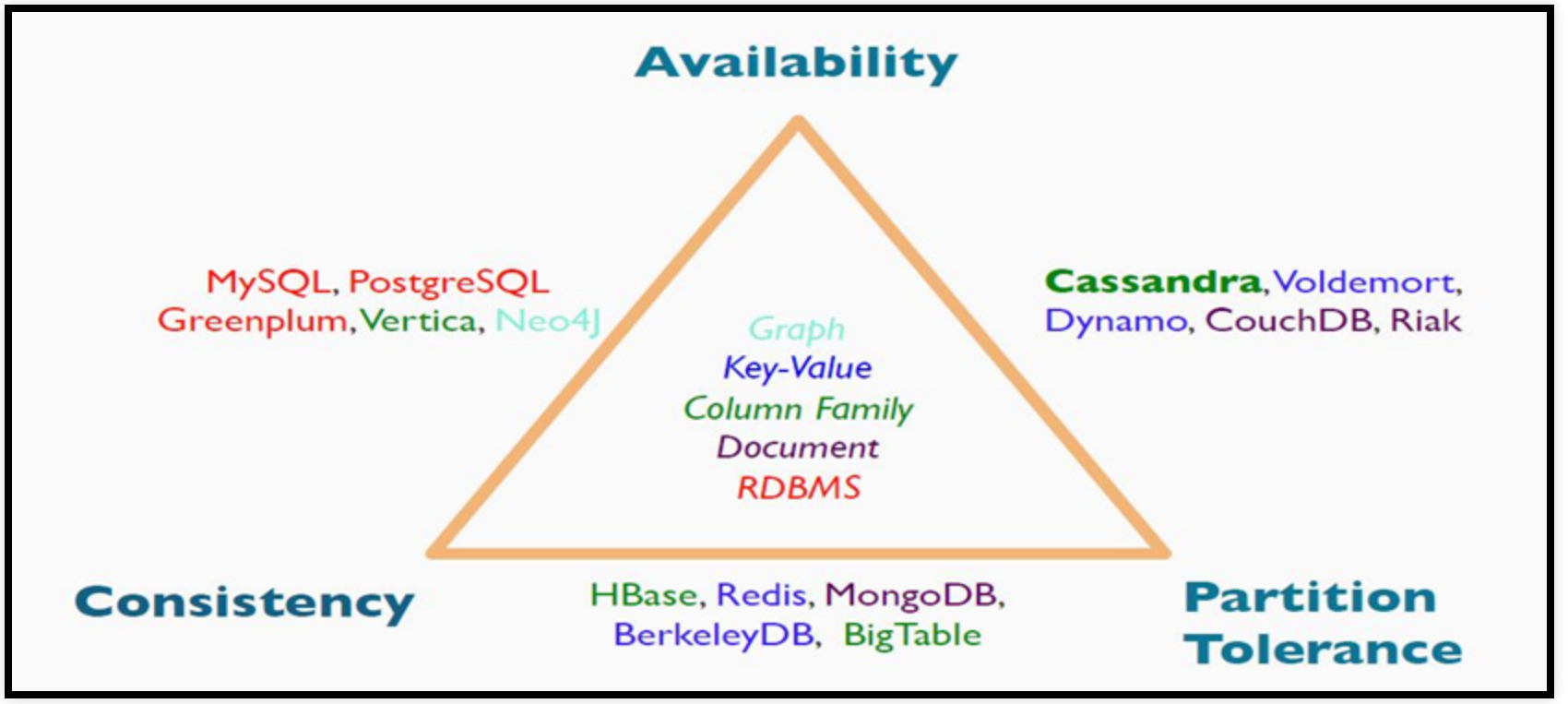
- *ability to **function (return a response, error, timeout)** when network partitions occur*



Partition tolerance

- network **is** unreliable
- **you can choose** how to handle errors
  - return an old value
  - wait and eventually timeout, or return an error at once
- in practice: choice between AP and CP systems

CAP theorem


Availability
MySQL, PostgreSQL
Greenplum, Vertica, Neo4J
Graph
Key-Value
Column Family
Document
RDBMS
Cassandra, Voldemort,
Dynamo, CouchDB, Riak
Consistency
HBase, Redis, MongoDB,
BerkeleyDB, BigTable
Partition
Tolerance

Cassandra consistency

- **AP** system
  - *eventually consistent*
    - without updates the system will converge to a consistent state due to *repairs*
  - *tunable consistency* :
    - Users can determine the consistency level by tuning it during **read** and **write** operations.

Consistency Level (CL)

- mandatory **protocol-level** parameter for each query (read/write),
- #replicas in a cluster that must acknowledge the **read** / **write**
  - **write consistency - R**: #replicas on which the write *must succeed** before returning an acknowledgment to the client application.
  - **read consistency - W**: #replicas that must *respond to a read* request before returning data to the client application
- default level: *ONE*
- most used: *ONE*, *QUORUM*, *ALL*, *ANY* … (*LOCAL_ONE*, *LOCAL_QUORUM…*)

Durability

- guarantees that writes, once completed, will survive permanently
  - appending writes to a commitlog first
    - default: flushed to disk every *commitlog_sync_period_in_ms*
    - batch mode ⇒ sync before ACK the write

Which consistency level ?

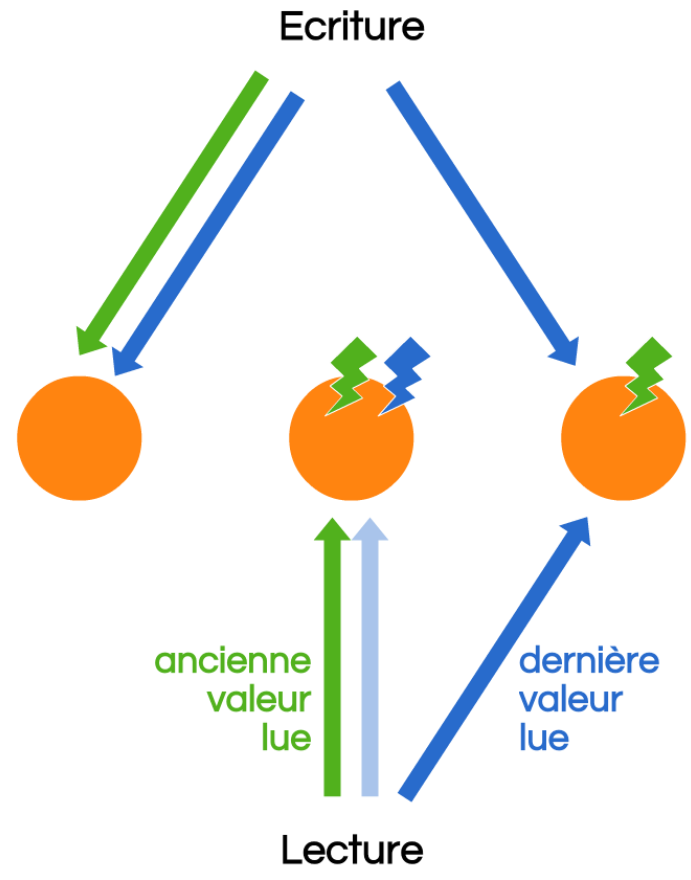
W : Niveau de cohérence en écriture (nb de nœuds)
R : Niveau de cohérence en lecture (nb de nœuds)
RF : Facteur de réplication

Cohérence immédiate si :

W + R > RF

W : ONE
R : ONE
Cohérence : à terme
Disponibilité : 2 nœuds en panne

W : QUORUM
R : QUORUM
Cohérence : immédiate
Disponibilité : 1 nœuds en panne



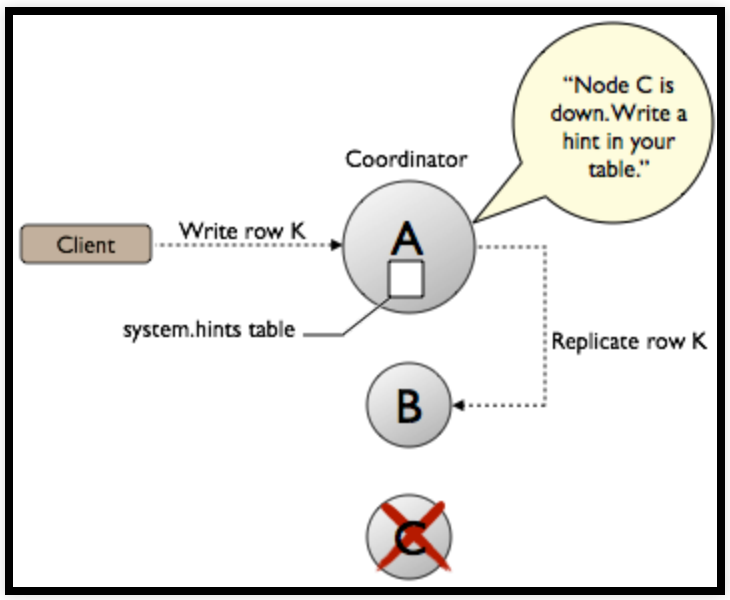
nœud en panne pendant l'écriture

How to achieve consistency

1. hinted handoff
2. read repairs
3. anti-entropy repair (nodetool repair)

Hinted handoff

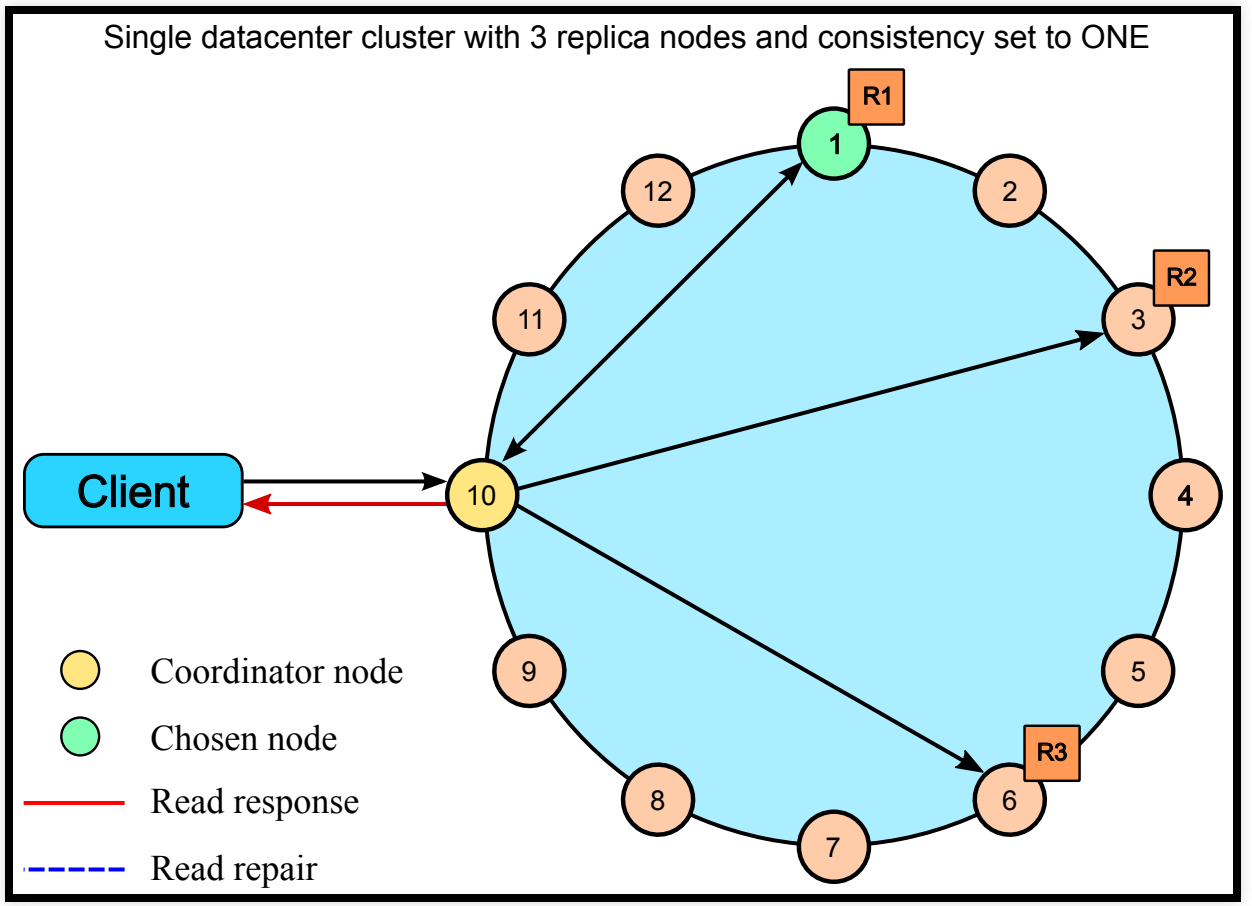
- ONE/QUORUM vs ANY (any node may ACK even if not a replica)
- if one/more replica(s) are down ⇒ **hinted handoff**



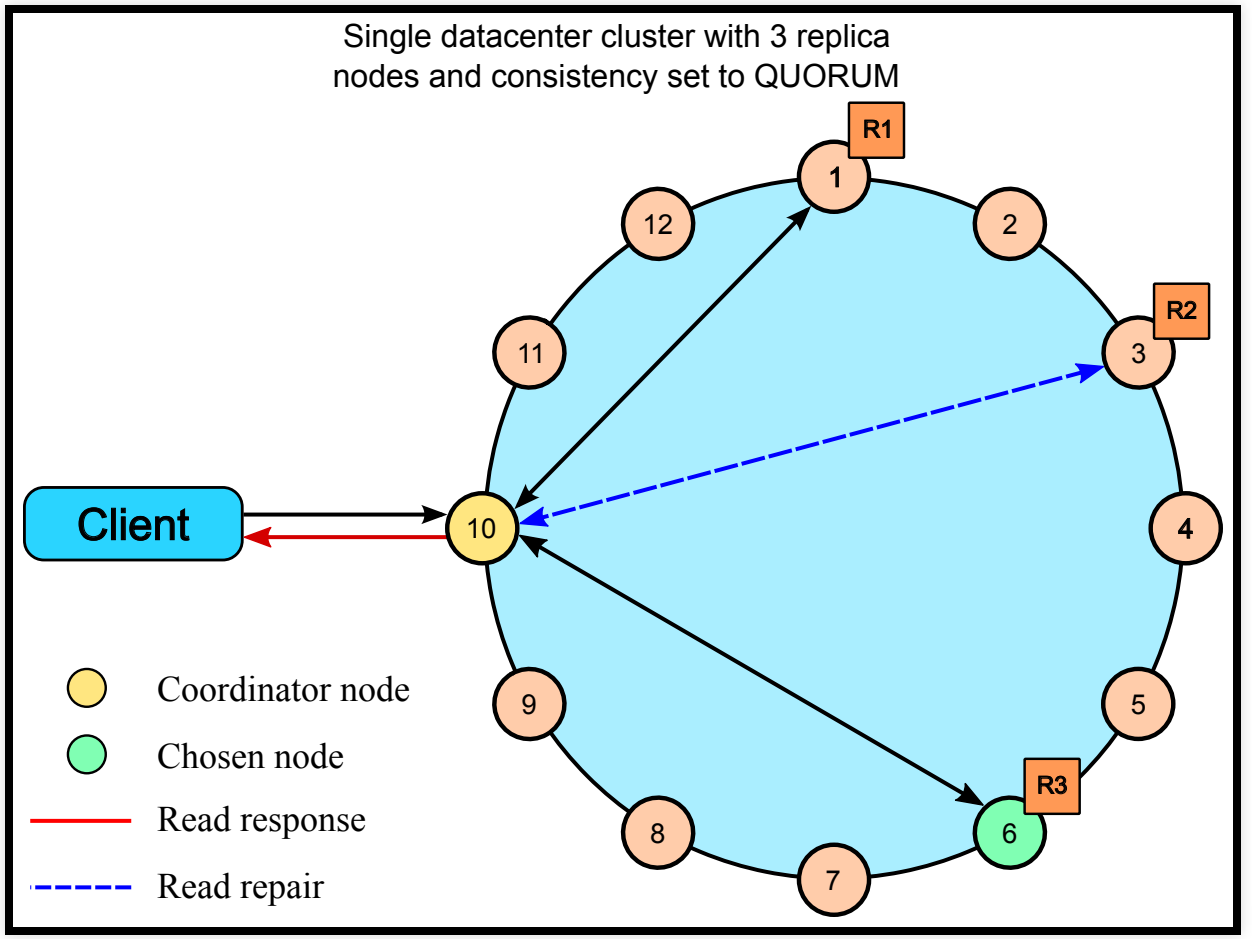
Read repairs

- Goal: **detect and fix inconsistencies during reads**
- **CL** = *ONE* ⇒ no data is repaired because no comparison takes place (unless **read_repair_chance** >0)
- **CL** > *ONE* ⇒ repair participating replica nodes in the foreground before the data is returned to the client.
  - C* sends a digest request to each replica not directly involved in the read
  - Cassandra compares all replicas and writes the most recent version to any replica node that does not have it.

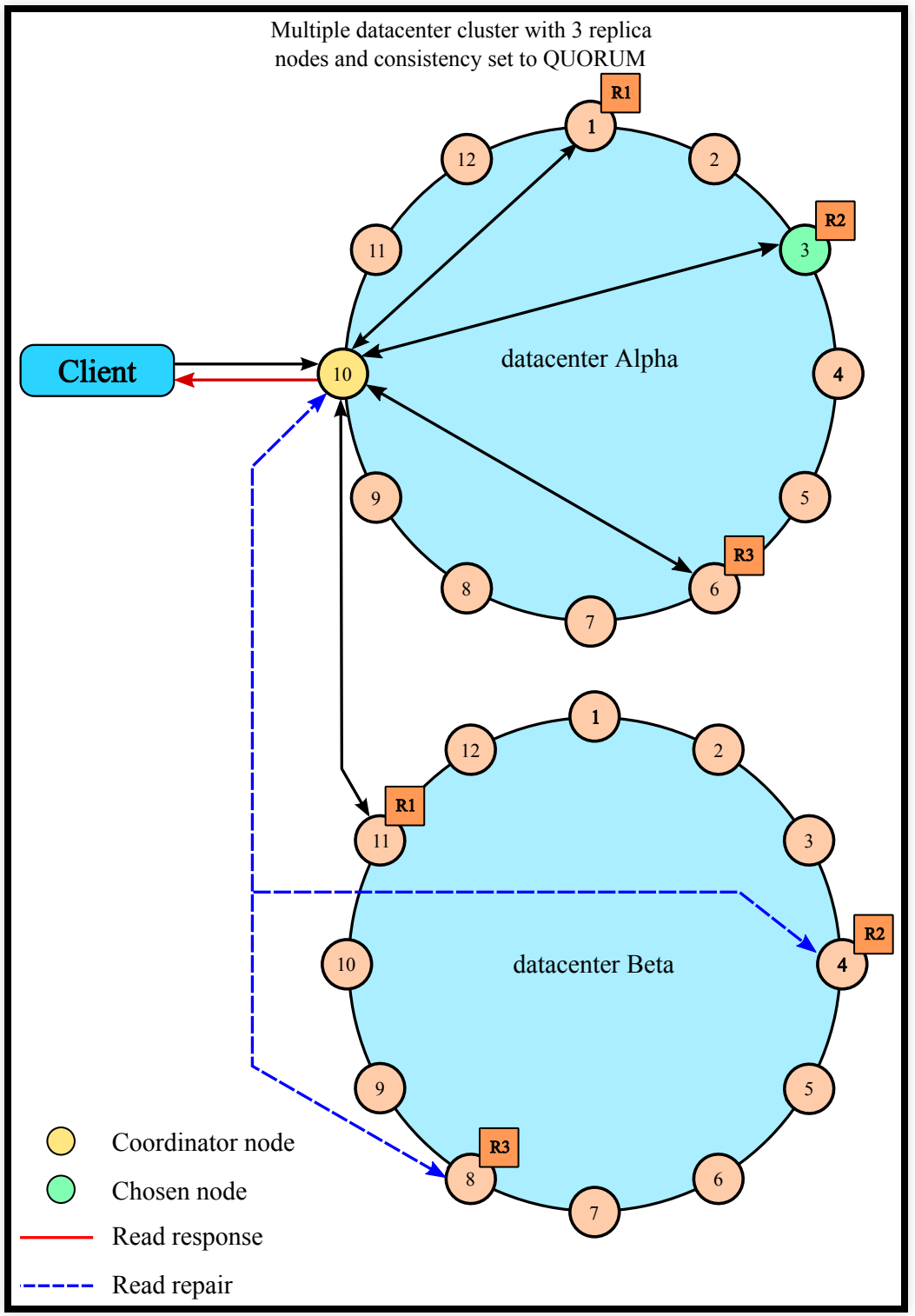
Read repairs ONE



Read repairs QUORUM



Read repairs QUORUM DC



Anti-entropy repair

- for each token range, read and synchronize the rows
- to insure the consistency this tool must be run regularly !
- not automatised (only on Datastax…)

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Practice: Tune consistency in Apache Cassandra

1. create local test clusters
2. explore configuration options and consistency properties

Cassandra cluster manager

- create multi-node cassandra clusters on the local machine
- great for quickly setting up clusters for development and testing

```
$ccm create test -v 2.0.5 -n 3 -s (1)

$ccm node1 stop (2)

$ccm node1 cqlsh (3)
```

Nodetool

- a command line interface for managing a cluster
  - explore, debug, performance test
  - maintenance operations, repairs

```
$ccm node1 nodetool status mykeyspace (1)

Datacenter: datacenter1
=======================
Status=Up/Down
|/ State=Normal/Leaving/Joining/Moving
--  Address    Load       Tokens  Owns    Host ID
UN  127.0.0.1  47.66 KB   1       33.3%   aaa1b7c1-6049-4a08-ad3e-3697a0e30e
UN  127.0.0.2  47.67 KB   1       33.3%   1848c369-4306-4874-afdf-5c1e95b873
UN  127.0.0.3  47.67 KB   1       33.3%   49578bf1-728f-438d-b1c1-d8dd644b6f
```

CQLSh

- standard CQL client

```
[bigdata@bigdata ~]$ ccm node2 cqlsh (1)
Connected to test at 127.0.0.2:9160.
[cqlsh 4.1.1 | Cassandra 2.0.5-SNAPSHOT | CQL spec 3.1.1 | Thrift protocol 1
Use HELP for help.
cqlsh> SELECT * FROM system.schema_keyspaces ; (2)
```

Ressources:

Datastax documentation

https://dzone.com/articles/introduction-apache-cassandras

https://highlyscalable.wordpress.com/2012/09/18/distributed-algorithms-in-nosql-databases/