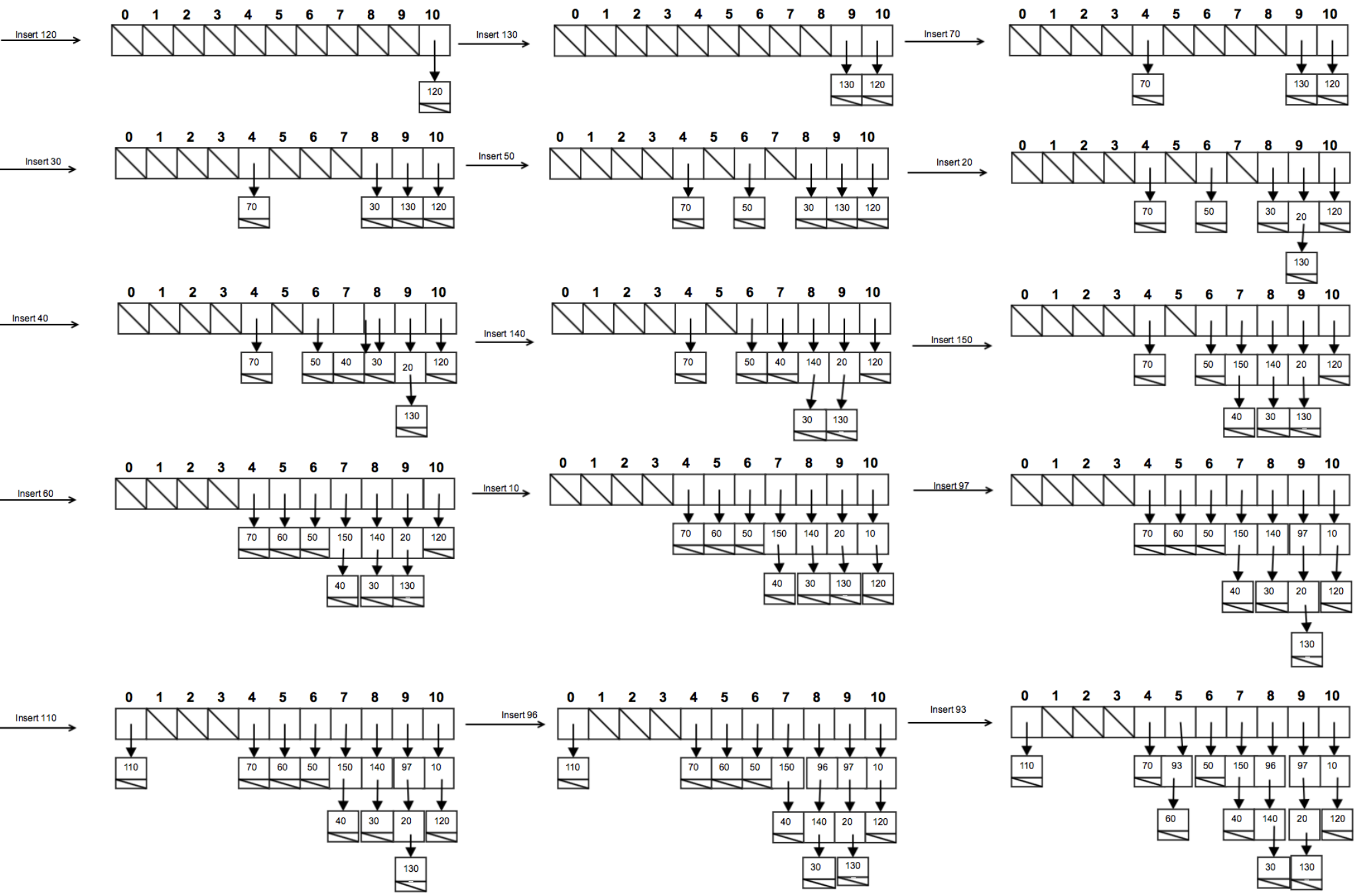
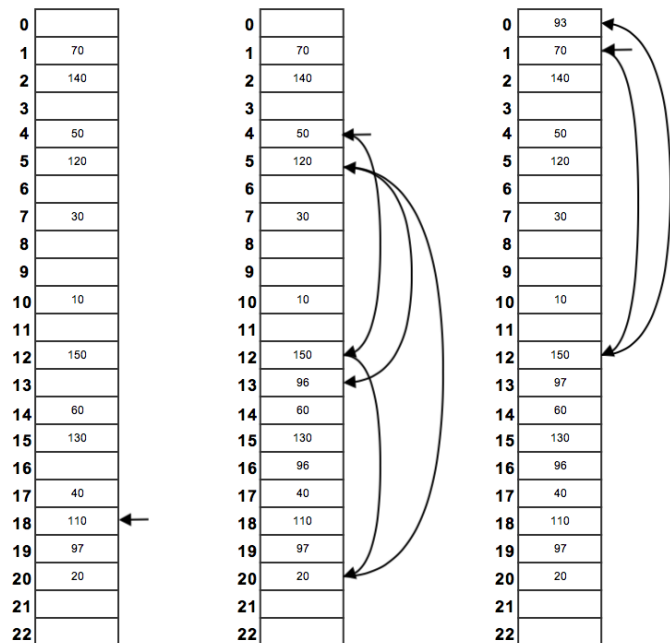
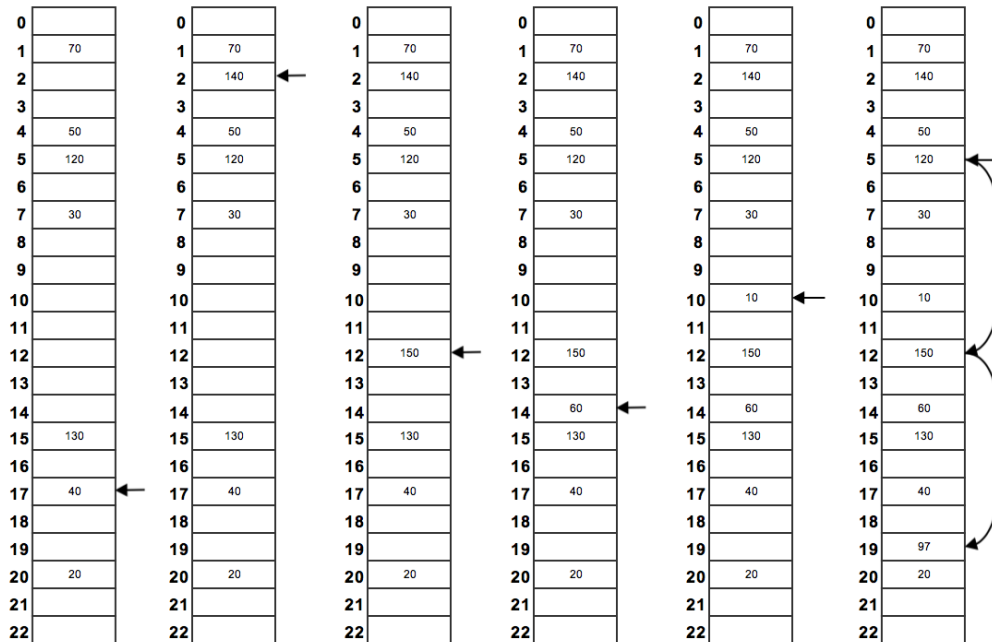
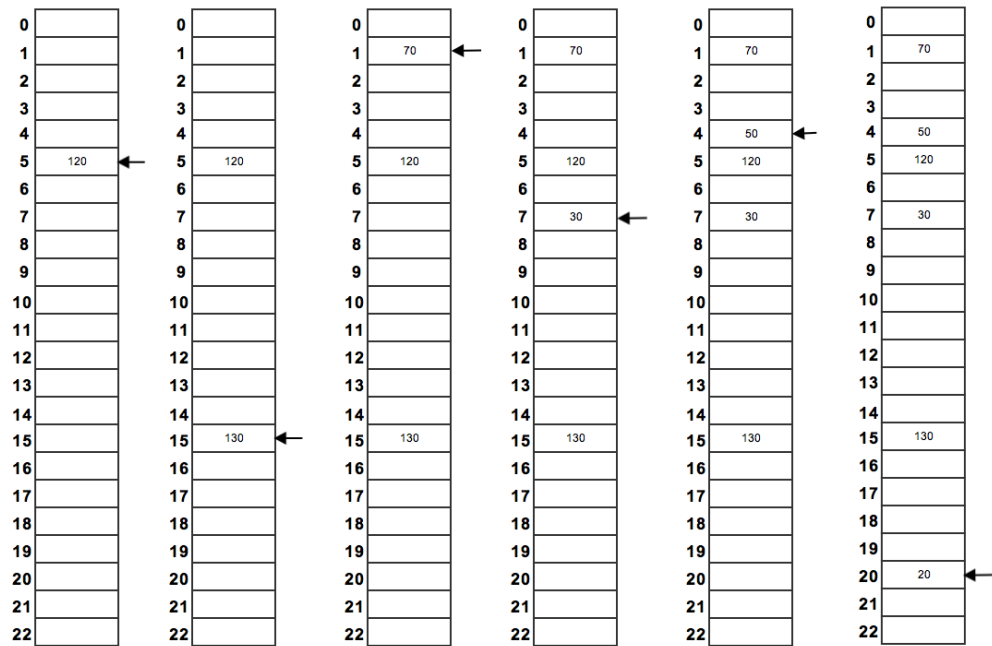


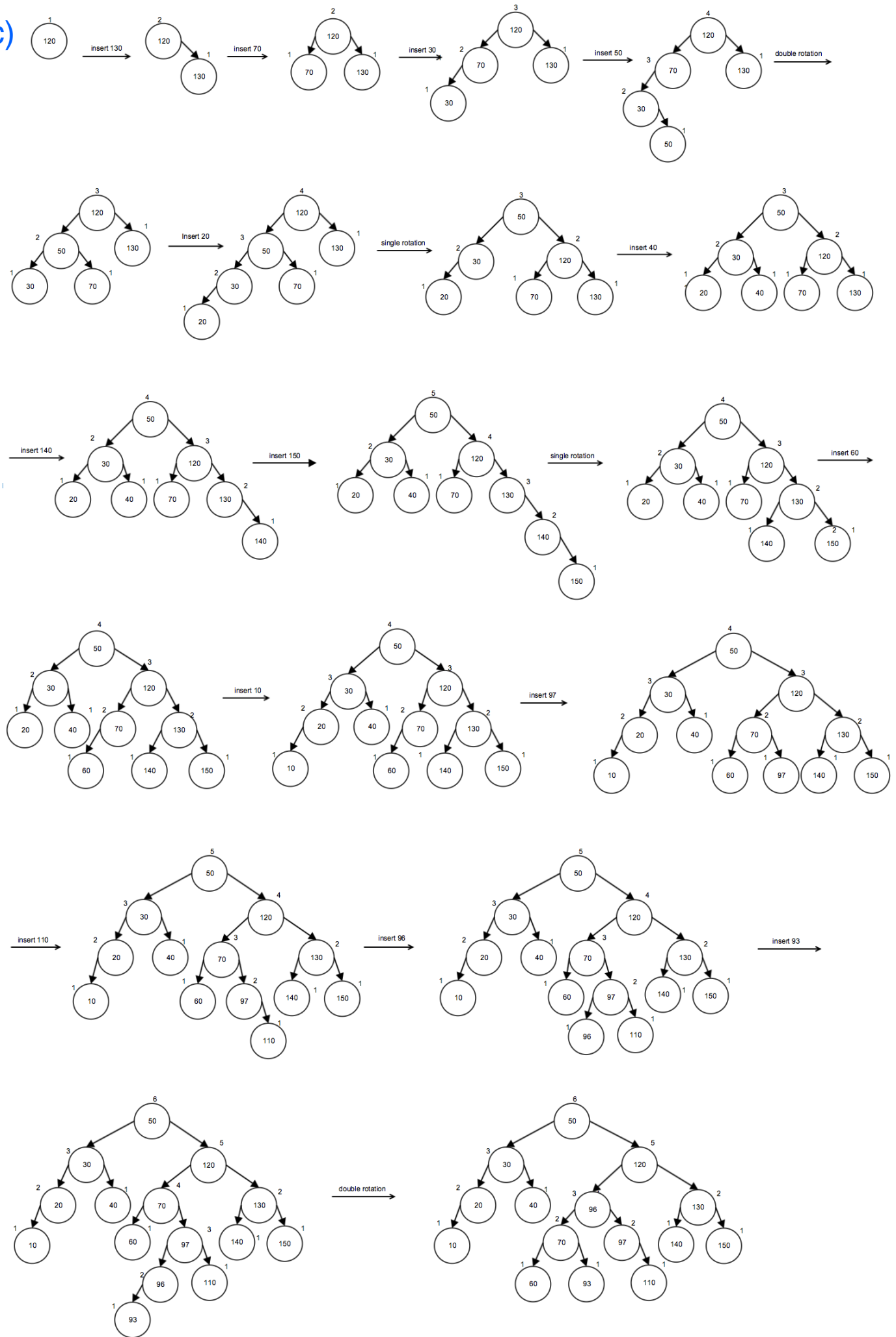
Question#1 (a)



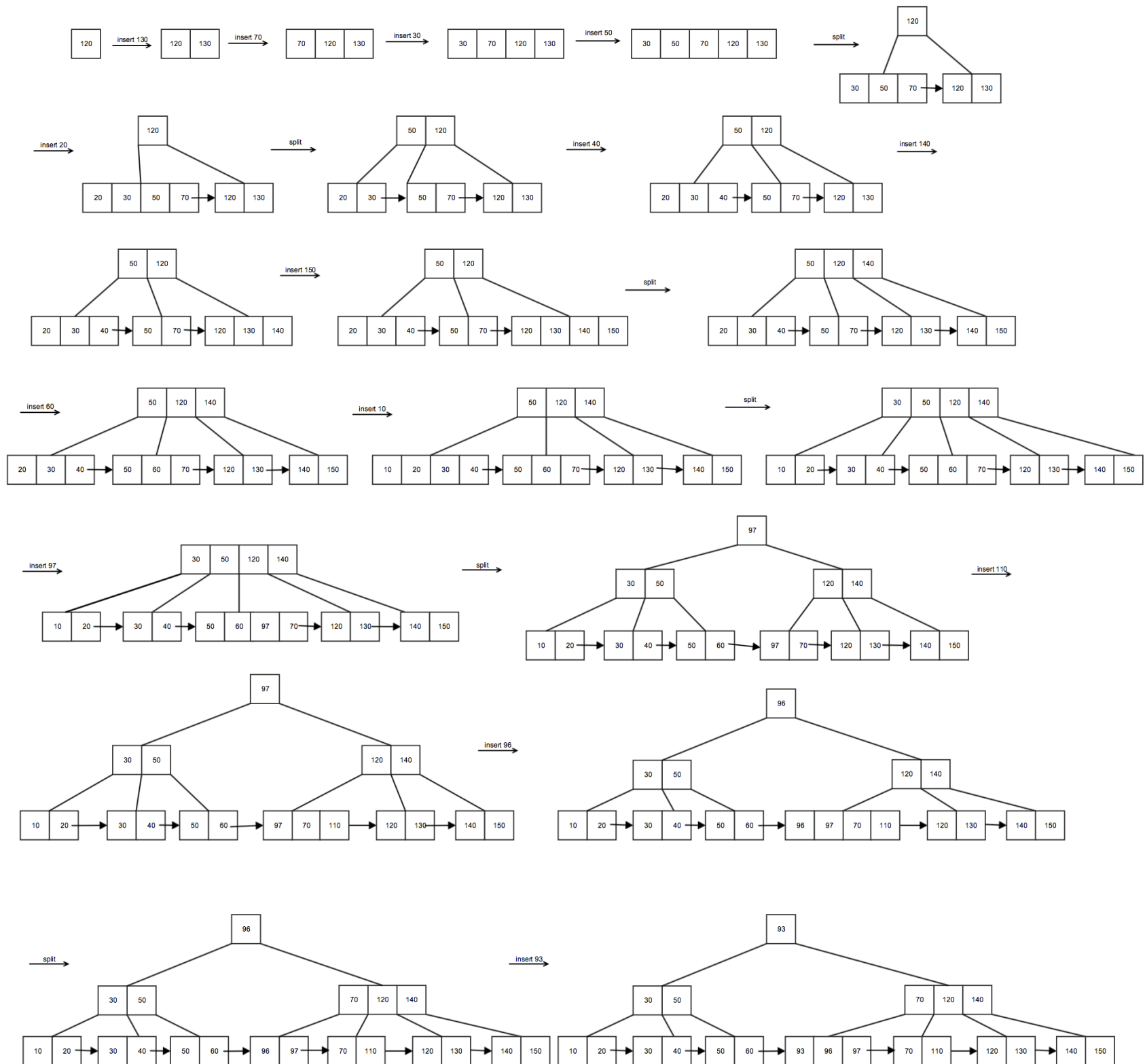
(b)



(c)



(d)



Q2:

Each computer can connect to 0-5 computers. However, let's say there is a computer that is not connected to any other computers, then the max number of connections among all six computers is 4, so

of connections $\in \{0, 1, 2, 3, 4\}$, 5 choices.

Similar things happen when we assume that there is a computer connecting to 1, 2, 3, 4, 5 respectively. There are always 5 possible choices for computer connections and we have 6 computers in total. According to the pigeonhole principle, two computers must have the same number of connections.

Q3: In an AVL tree of height h , at least one of the left and right subtrees is an AVL tree of height $h-1$. Since the height of left and right subtrees can differ by at most 1, then the other subtree must have height of $h-2$. Thus, we have the relation

$$N(h) = N(h-1) + N(h-2) + 1$$

Given the base case: $N(0) = 1$ and $N(1) = 2$

Base case:

- $F(3) = F_1 + F_2 = 2$
- $\therefore N(0) = F(3) - 1 = 2 - 1 = 1$, $N(h) = F_{h+3} - 1$ applies
- $F(4) = F_3 + F_2 = 2 + 1 = 3$
- $\therefore N(1) = F(4) - 1 = 3 - 1 = 2$, $N(h) = F_{h+3} - 1$ applies

Induction hypothesis:

Assume that the hypothesis $N(h) = F_{h+3} - 1$ is true for $h \geq 1$

Inductive step: assume $h = k+1$, $k \in \mathbb{Z}$

$$N(k+1) = N(k) + N(k-1) + 1$$

$$= F(k+3) - 1 + F(k+2) - 1 + 1$$

$$= F(k+4) - 1$$

replace $h = k+1$, so we have $N(h) = F_{h+3} - 1$

Thus, it is true for all h .