

# ||||· ClickHouse - Lightning Fast Analytics for Everyone

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## Open source **column-oriented distributed OLAP database**

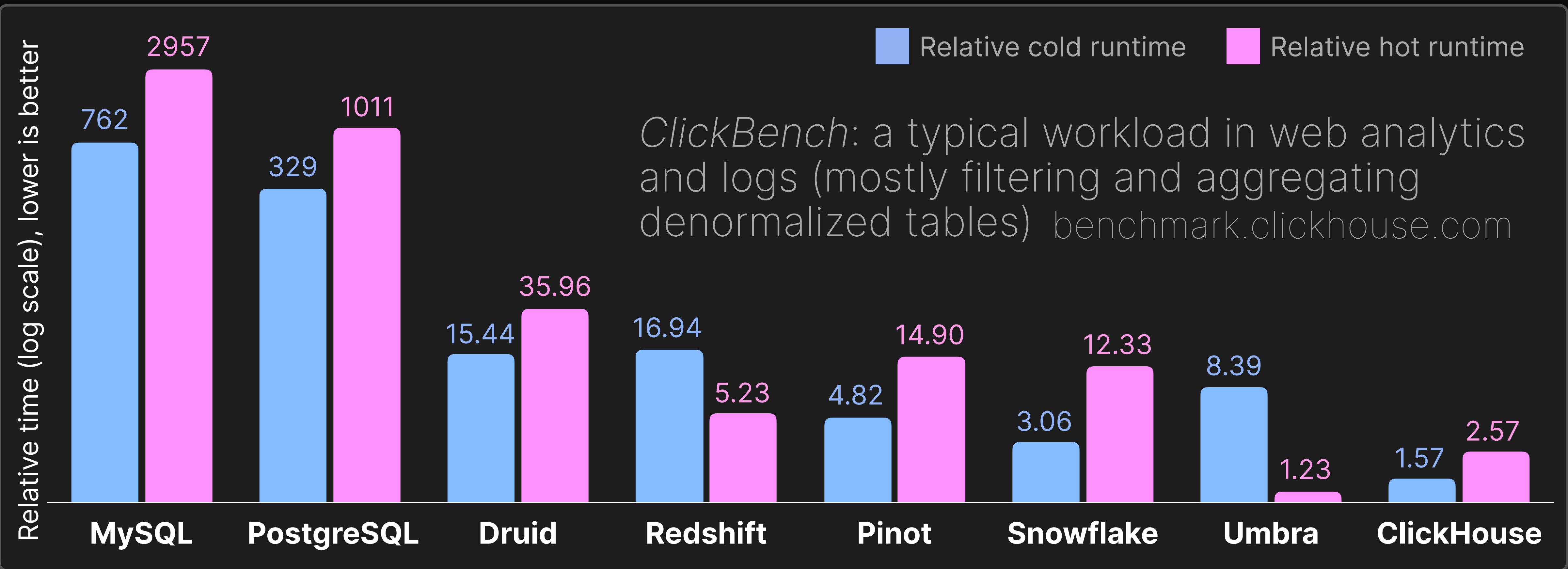
Developed since 2009, built in C++  
OSS (Apache 2.0) since 2016

Best for filter and aggregation queries  
Optimized for append-only workloads

Replication  
Sharding  
Eventually consistent

Business intelligence  
Logs, events, traces  
Real-time analytics

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Total relative cold and hot runtimes for sequentially executing all ClickBench queries in databases frequently used for analytics.

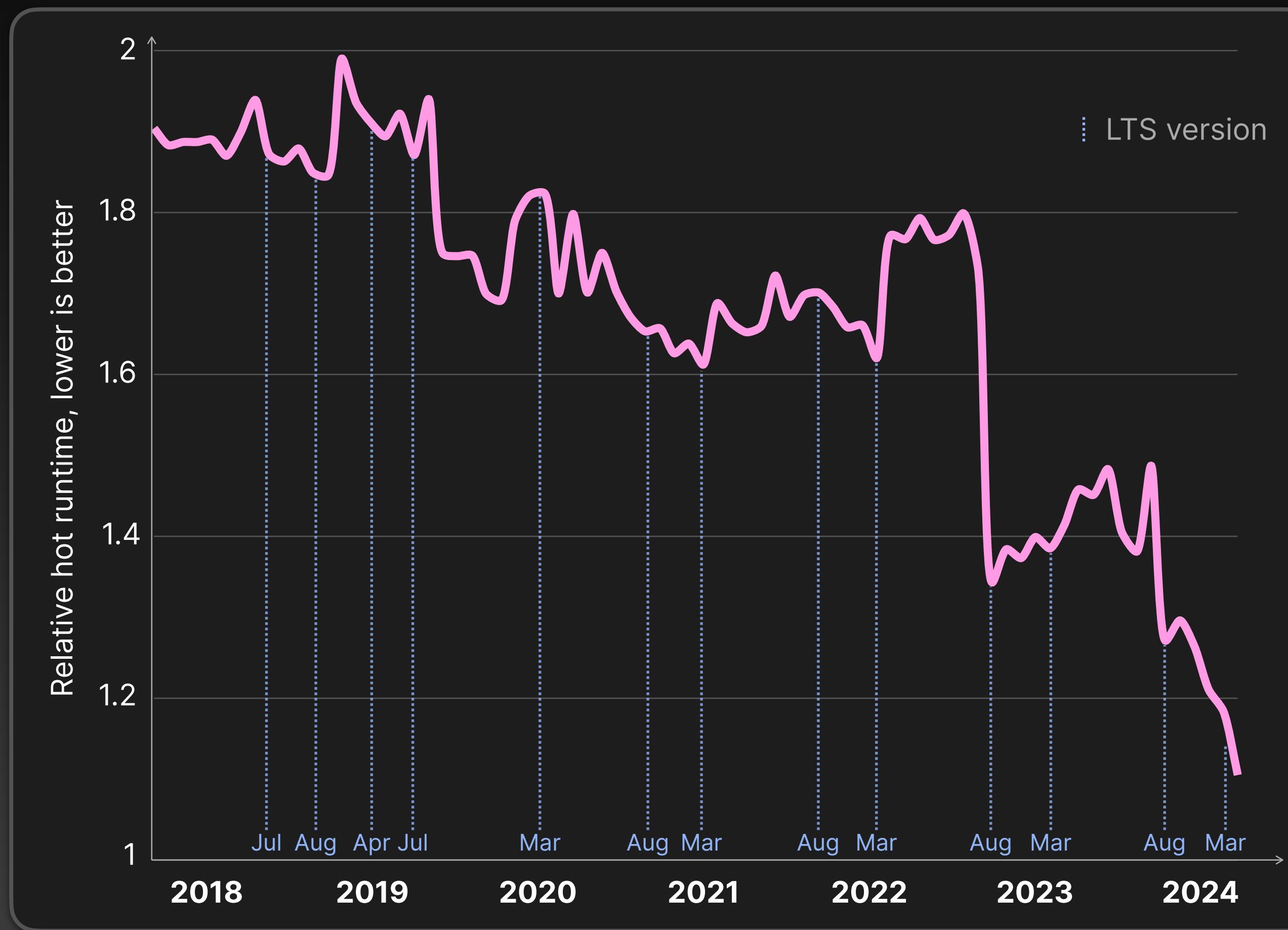
Measurements taken on a single-node AWS EC2 c6a.4xlarge instance with 16 vCPUs, 32 GB RAM, and 5000 IOPS / 1000 MiB/s disk.

Comparable systems were used for Redshift (ra3.4xlarge, 12 vCPUs, 96 GB RAM) and Snowflake (warehouse size S: 2×8 vCPUs, 2×16 GB RAM).

ClickHouse has the best query performance amongst production-grade analytics databases.

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Performance improvements by  $1.72 \times$  since 2018

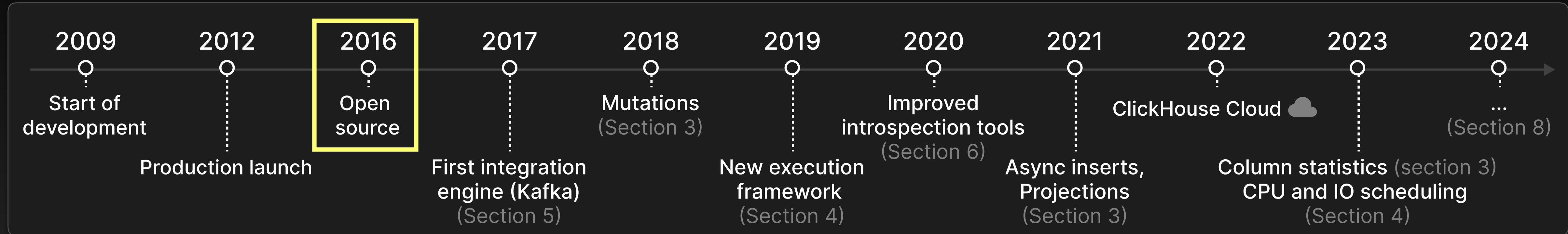


- *VersionBench* benchmark is run when a new release is published to check its performance and identify regressions.
- Combination of four benchmarks:

	# Queries	# Rows
ClickBench	42	100 million
MgBench	15	200 million
Star Schema Benchmark (denormalized schema)	13	600 million
NYC Taxi Rides Benchmark	4	3.4 billion

Query performance is a top priority and continuously improved.

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The most popular analytics database on GitHub  
**(36k ★, 2k+ contributors)**.

[github.com/ClickHouse/ClickHouse](https://github.com/ClickHouse/ClickHouse)

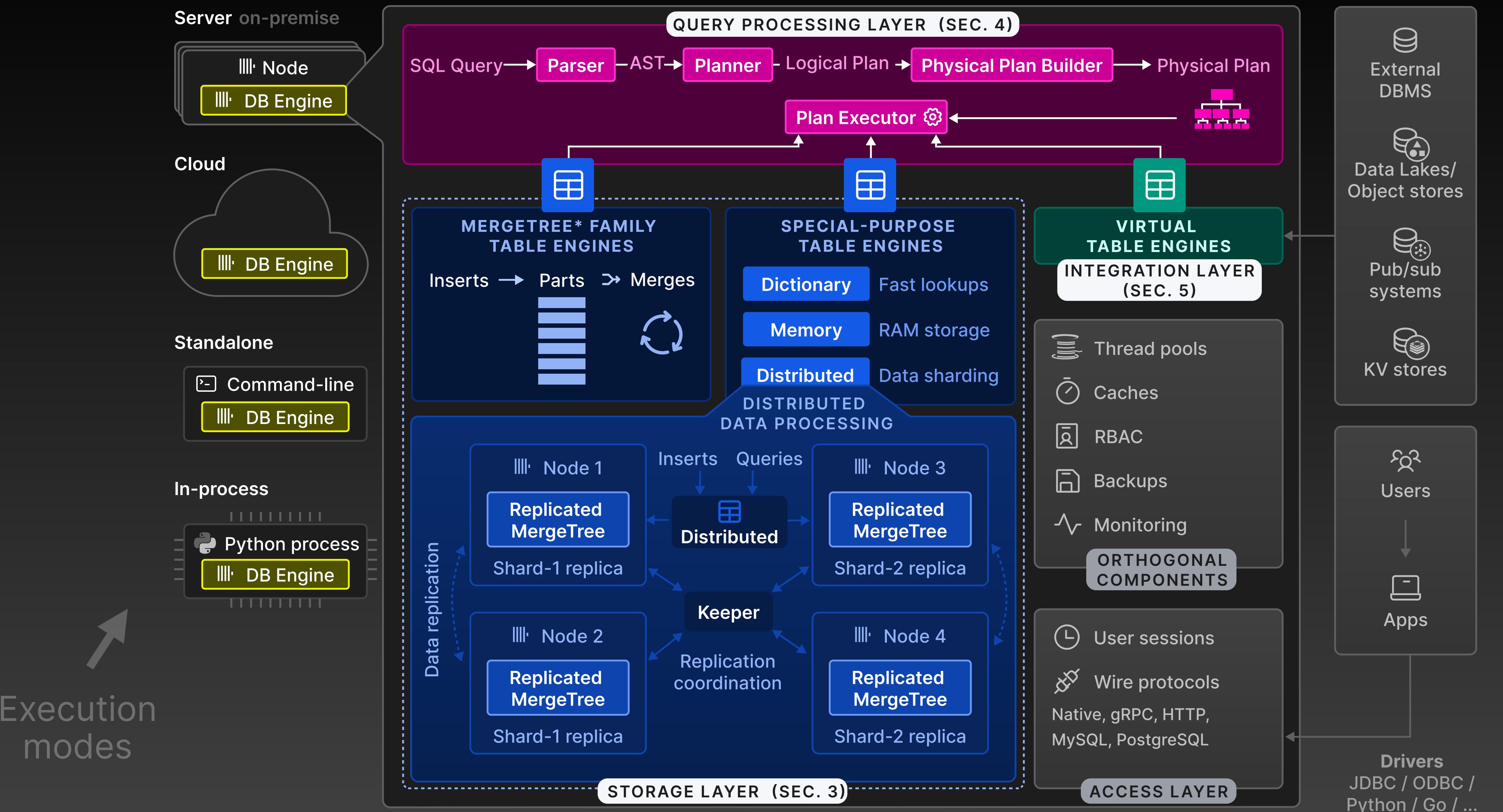
**Runs on anything** from Raspberry Pi to clusters with hundreds of nodes, largest known cluster is 4000 servers.

**Used by hundreds of companies globally for production workloads.**

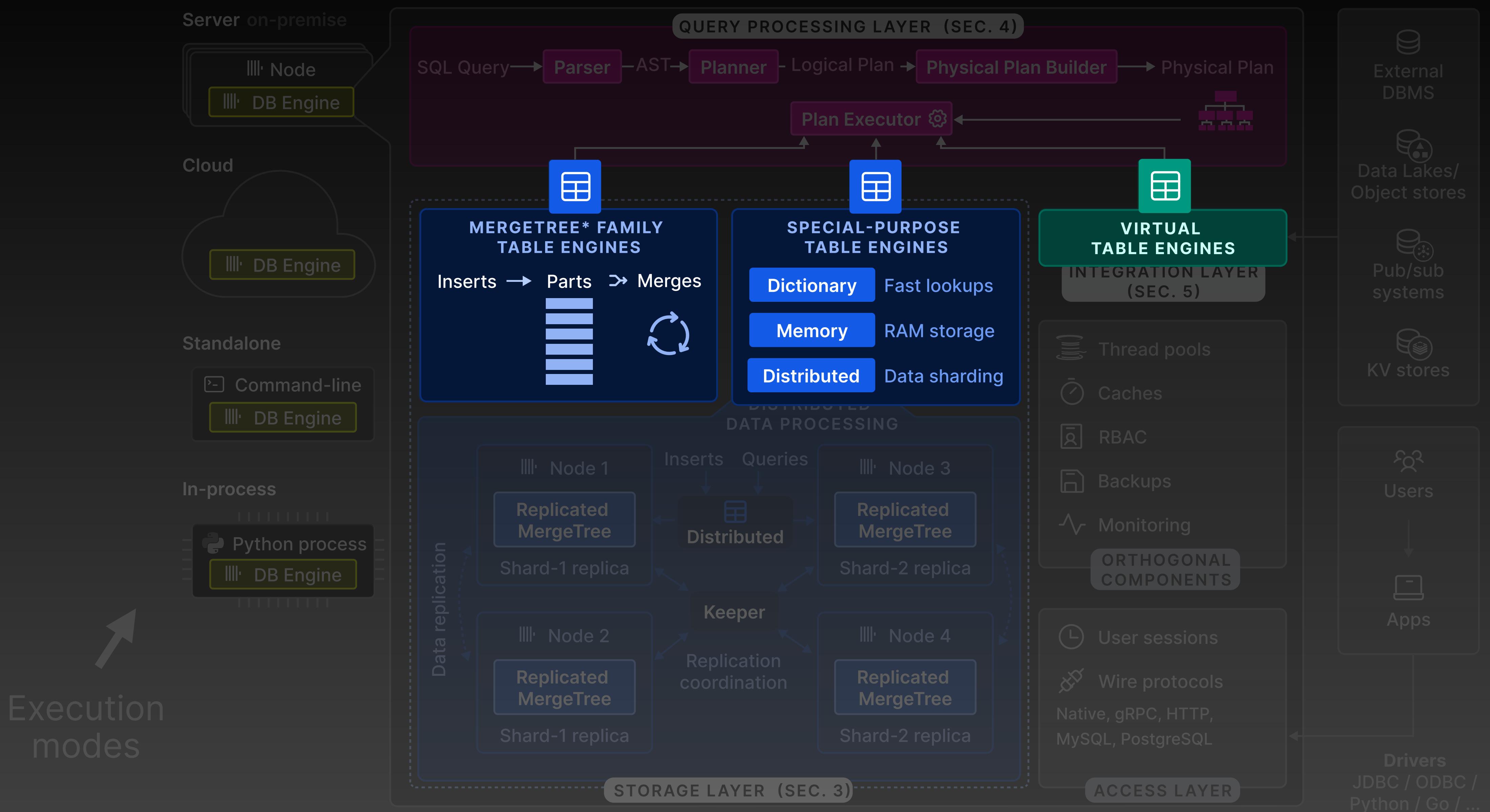
[clickhouse.com/docs/en/faq/general/who-is-using-clickhouse](https://clickhouse.com/docs/en/faq/general/who-is-using-clickhouse)

ClickHouse is trusted by 50%+ of Fortune's Global Top 2000 companies.

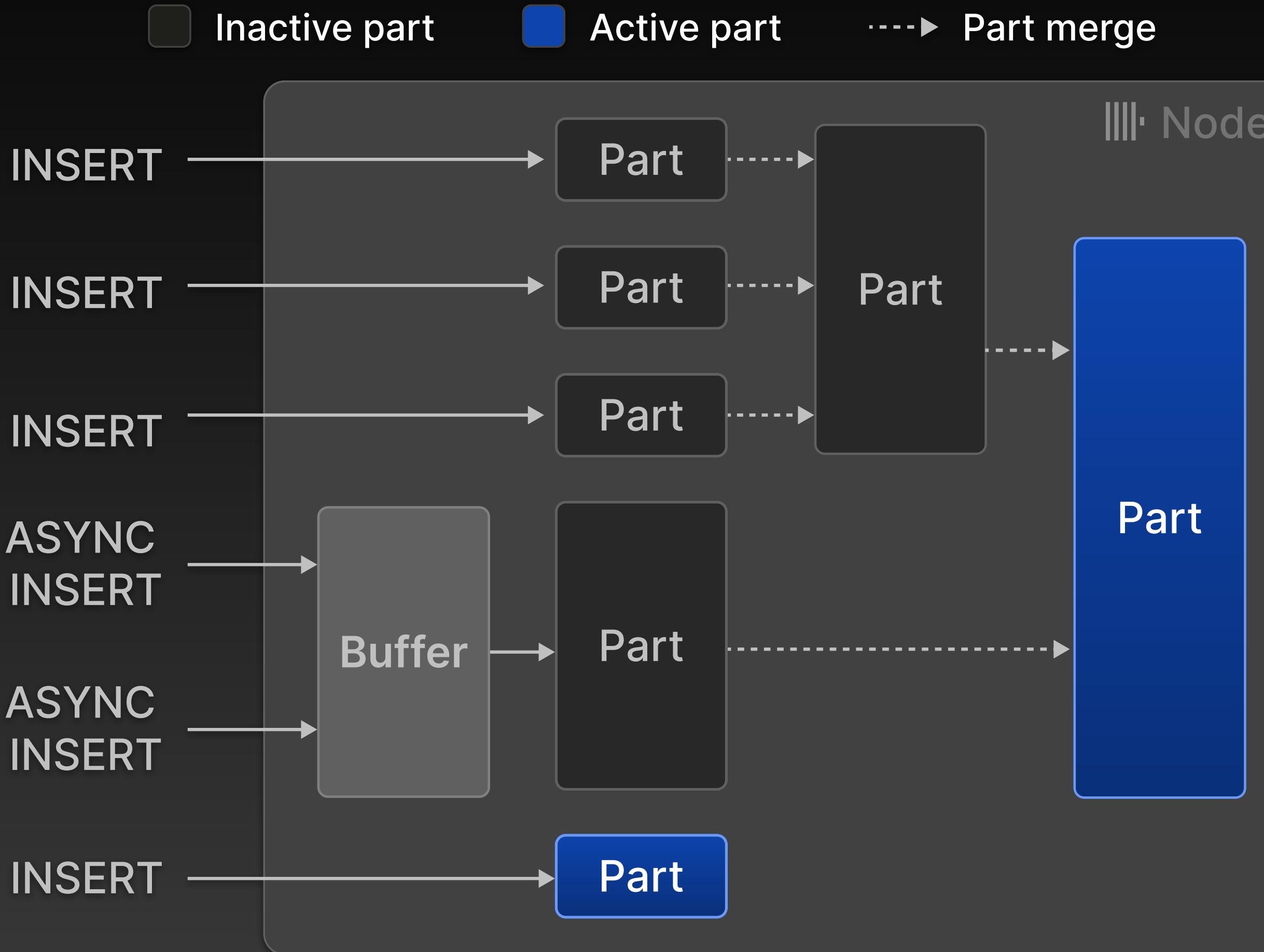
# Architecture



# Table engines encapsulate the location and format of table data



# An LSM-Tree Inspired Storage Layer



- Like LSM trees, INSERTs create sorted and immutable parts.
- Parts are continuously merged by a background job.
- Unlike LSM trees, all parts are equal (i.e., no levels or notion of recency).
- INSERTs can be synchronous or asynchronous.

Ingestion rates are only limited by the speed of disk.

# Column Layout and Compression

Row	EventTime	RegionID	URL
0	2023-10-19 17:03:05.154	EMEA	https://...
.	.		.
8,191	2023-10-19 17:03:07.490		https://...
8,192	2023-10-19 17:03:07.492		https://...
.	.		.
16,383	2023-10-19 17:03:09.838		https://...
.	.		.
.	.		.
.	.		.
Compressed block			
:	:	:	:

- Local (per-part) sorting defined by primary key:

```
CREATE TABLE page_hits
(
    EventTime Date     CODEC(Delta, ZSTD),
    RegionId String  CODEC(LZ4),
    URL       String  CODEC(AES),
    PRIMARY KEY (EventTime)
)
```

- Parts are further divided into *granules*  $g_0, g_1, \dots$
- Consecutive granules in a column form *blocks* which are encoded:
  - Generic codecs: LZ4, ZSTD, DEFLATE, ...
  - Logical codecs: Delta, GCD, ...
  - Specialized Codecs: Gorilla(FP), AES,...
- Codecs can be combined: CODEC(Delta, ZSTD)

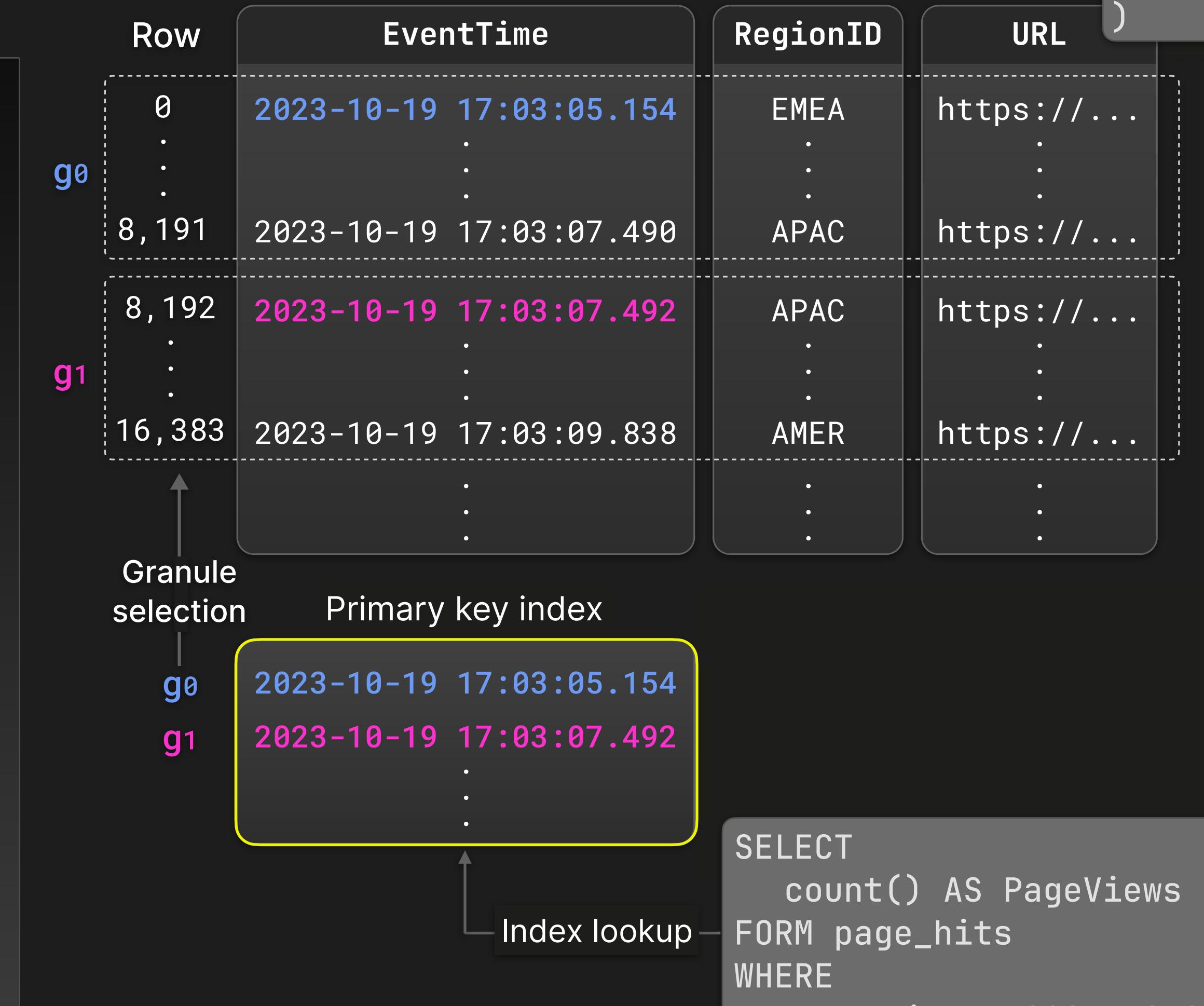
High compression rates are in many use cases critical for cost efficiency and performance.

# Data Pruning

## Primary key indexes

## Table projections

## Skipping indexes



```
CREATE TABLE page_hits
(
    EventTime Date,
    RegionId String,
    URL       String,
    PRIMARY KEY (EventTime)
```

- Define the local part sorting (clustered index).
- Also create a mapping from primary key column values to granules.
- The mapping is small enough to remain in DRAM at all times.

```
SELECT
    count() AS PageViews
FROM page_hits
WHERE
    EventTime >= '2023-12-09'
```

Quickly find granules containing rows that match a predicate on a prefix of the PK columns.

# Data Pruning

Primary key indexes

Table projections

Skipping indexes

EventTime	RegionID	URL
2023-10-19 17:03:05.154	EMEA	https:// ...
2023-10-19 17:03:05.462	APAC	https:// ...
2023-10-19 17:03:05.875	AMER	https:// ...
2023-10-19 17:03:06.104	AMER	https:// ...
2023-10-19 17:03:07.550	APAC	https:// ...

```
ALTER TABLE page_hits ADD PROJECTION proj (
    SELECT *
    ORDER BY RegionID
);
ALTER TABLE page_hits MATERIALIZE PROJECTION proj;
```

EventTime	RegionID	URL
2023-10-19 17:03:05.875	AMER	https:// ...
2023-10-19 17:03:07.550	AMER	https:// ...
2023-10-19 17:03:06.104	APAC	https:// ...
2023-10-19 17:03:05.462	APAC	https:// ...
2023-10-19 17:03:05.154	EMEA	https:// ...

- Alternative table versions sorted by different primary keys.
- Works at the granularity of parts.
- Speed up queries on columns different than primary key columns.

```
SELECT
    count() AS PageViews
FROM page_hits
WHERE
    RegionID = 'AMER'
```

Powerful but increase space consumption and insert/merge overhead.

# Data Pruning

## Primary key indexes

## Table projections

## Skipping indexes

- Store small amounts of metadata at the level of granules or multiple granules which allows to skip data during scans.

- Skipping index types:
  - Min/Max values
  - Unique values
  - Bloom filter
  - ...

```
ALTER TABLE T  
ADD INDEX idx_minmax (Clicks) TYPE minmax;  
ALTER TABLE T MATERIALIZE INDEX idx_minmax;
```

```
SELECT *  
FROM T  
WHERE  
Clicks BETWEEN 15 AND 30
```

Clicks	min/max index
25	min: 7
8	max: 25
7	
25	
25	
18	
20	
22	
19	
17	
8	
6	
6	
13	
5	

Some match → Load and Scan block

All match → SKIP load

None match → SKIP load

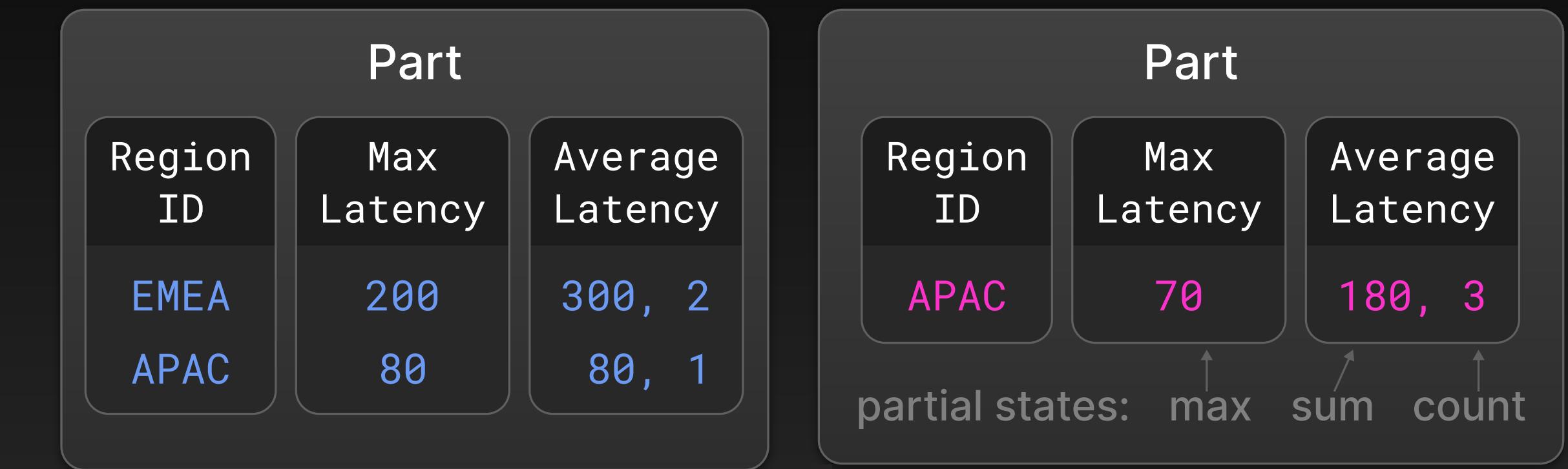
Skipping indexes are a light-weight alternative to projections.

# Merge-time Data Transformation

Merges optionally perform additional data transformations or maintenance.

## Replacing merges

Retain the most recently inserted version of the same rows in multiple input parts.



## Aggregating merges

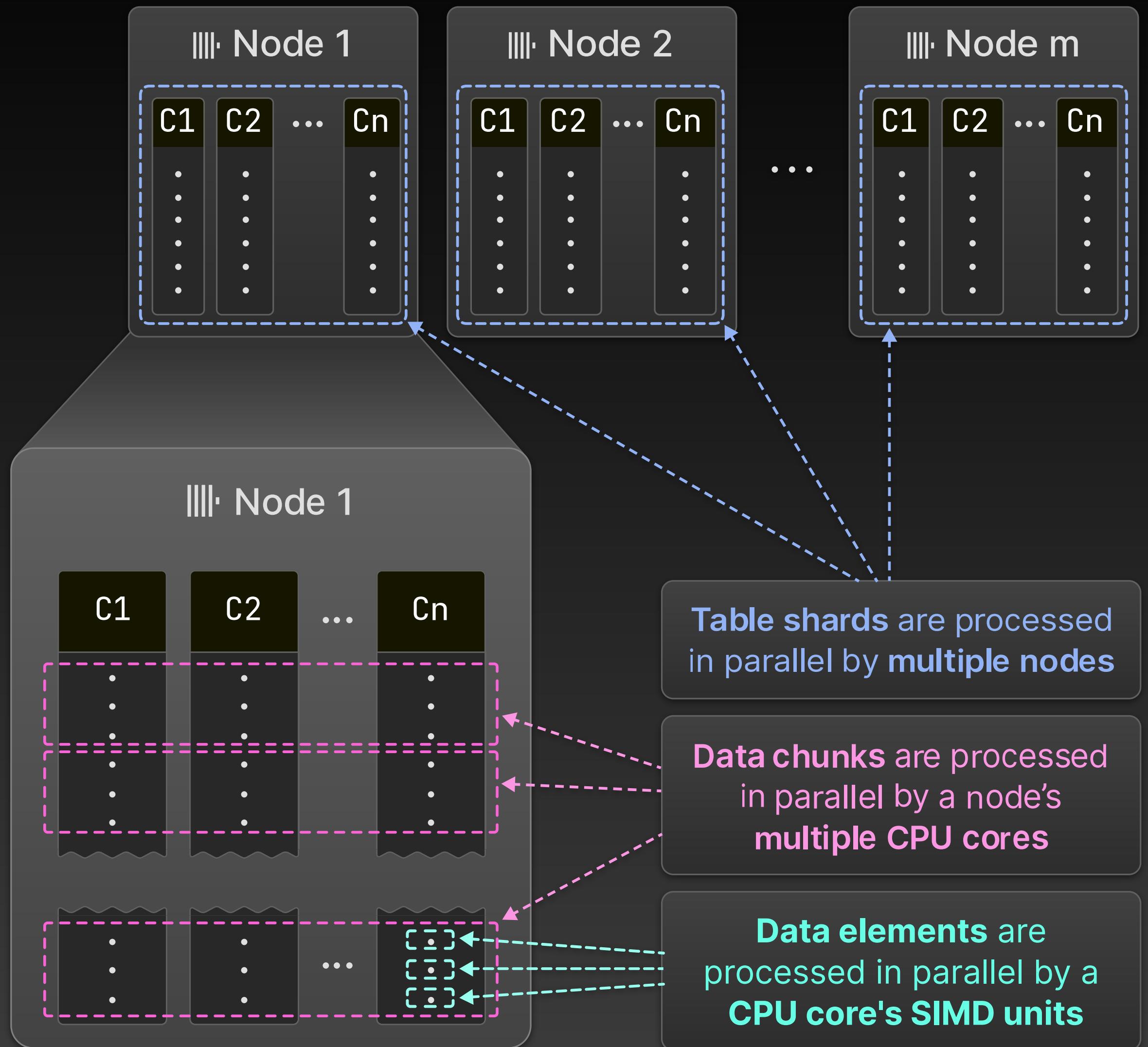
Combine aggregation states into new aggregation states.

## TTL (time-to-live) merges

Compress, move, or, delete rows or parts.

Data transformations don't compromise the performance of parallel INSERTs and SELECTs.

# State-of-the-art Vectorized Query Execution Engine



Query execution utilizes all server and cluster resources.

# Parallelization Across SIMD ALUs

- Based on compiler auto-vectorization or manually written intrinsics.
- SQL expressions are compiled into *compute kernels*.
- The fastest kernel is selected at runtime based on the system capabilities (cpuid).

```
SELECT col1 + col2  
FROM T
```



```
if (isArchSupported(TargetArch::AVX512))  
    implAVX512BW(in1, in2);  
else if (isArchSupported(TargetArch::AVX2))  
    implAVX2(in1, in2, out);  
else if (isArchSupported(TargetArch::SSE42))  
    implSSE42(in1, in2, out);  
else  
    implGeneric(in1, in2, out);
```

Dispatch code based on cpuid

```
MULTITARGET_FUNCTION_AVX512F_AVX2_SSE42(  
    MULTITARGET_FUNCTION_HEADER(),  
    impl,  
    MULTITARGET_FUNCTION_BODY()  
    const double * in1, const double * in2  
    double * out, size_t num_elements)  
{  
    for (size_t i = 0; i < (sz & ~0x7); i += 8)  
    {  
        const __m512d _in1 = _mm512_load_pd(&in1[i]);  
        const __m512d _in2 = _mm512_load_pd(&in2[i]);  
        const __m512d _out = _mm512_add_pd(_in1, _in2);  
        out[i] = (double*)&_out;  
    } /* tail handling */  
})
```

AVX-512 kernel, manually vectorized

```
MULTITARGET_FUNCTION_AVX2_SSE42(  
    MULTITARGET_FUNCTION_HEADER(),  
    impl,  
    MULTITARGET_FUNCTION_BODY()  
    const double * in1, const double * in2  
    double * out, size_t num_elements)  
{  
    for (size_t i = 0; i < num_elements; ++i)  
        *out[i] = *in1[i] + *in2[i];  
})
```

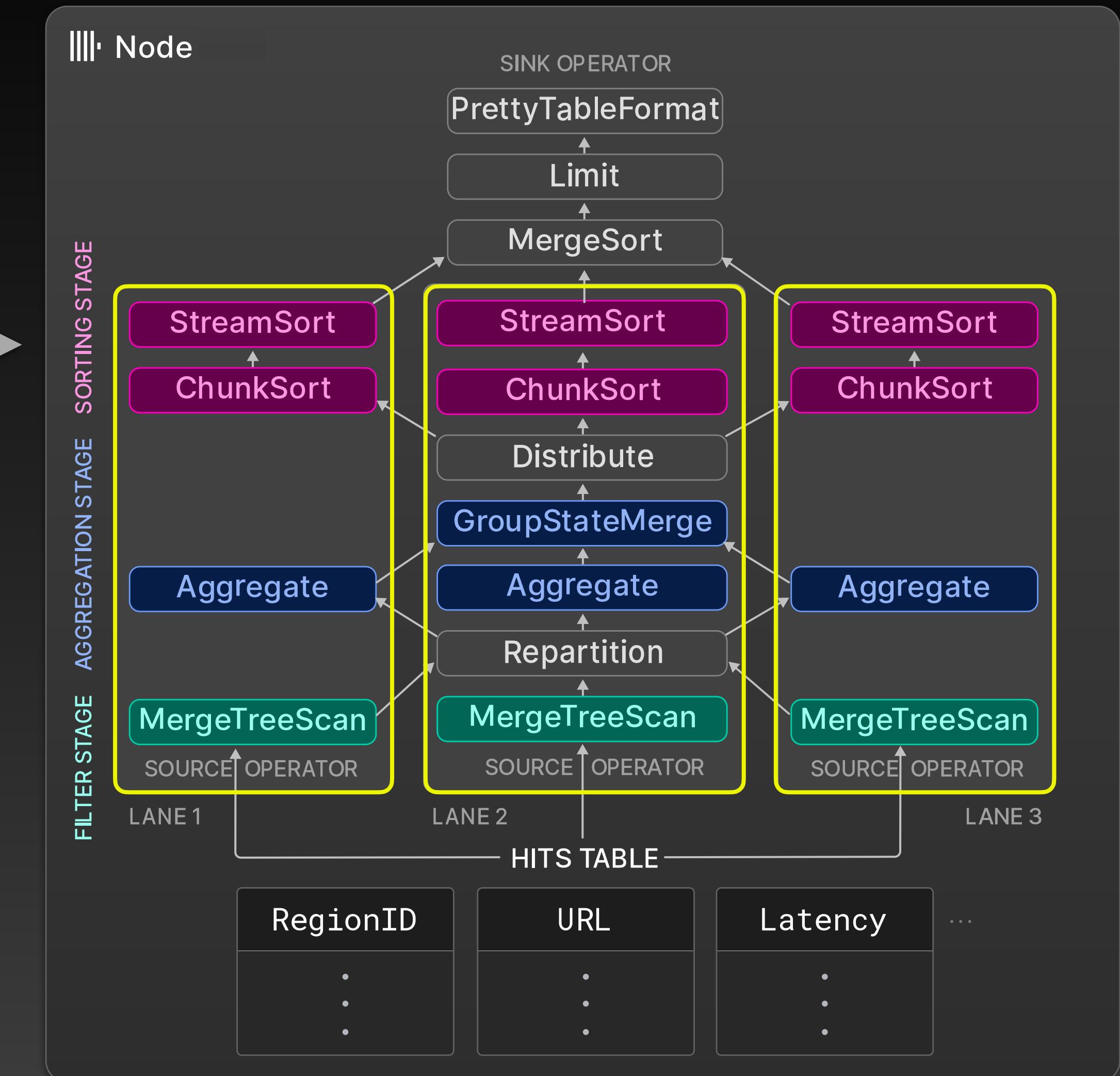
AVX2 kernel, compiler auto-vectorized

Remain compatible with legacy hardware while utilizing modern hardware fully.

# Parallelization Across CPU Cores

```
SELECT RegionID, avg(Latency) AS AvgLatency  
FROM hits  
filter WHERE URL = 'https://clickhouse.com'  
aggregation GROUP BY RegionID  
sort ORDER BY AvgLatency DESC  
LIMIT 3
```

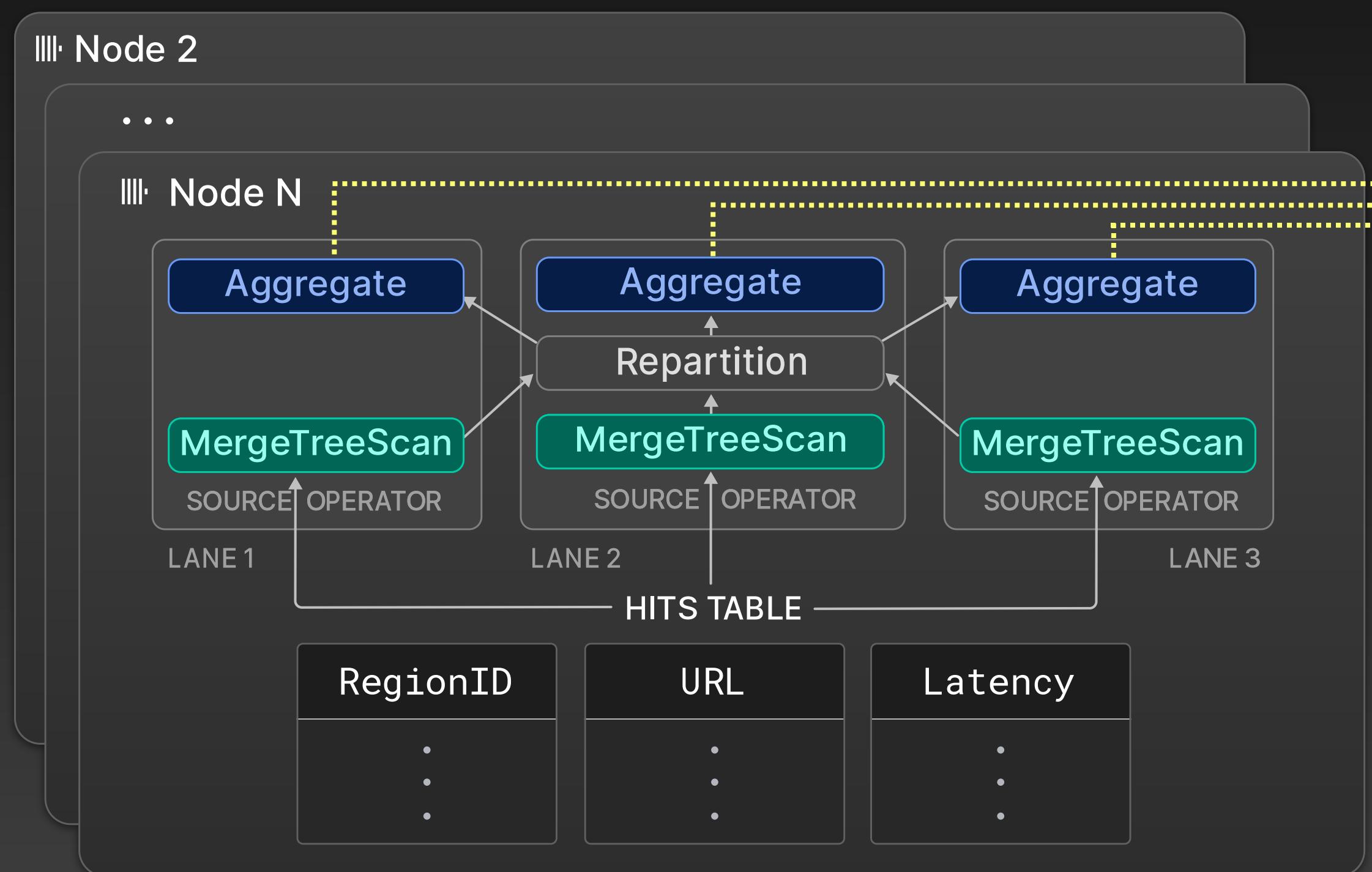
- Execution plan gets unfolded into N lanes (typically 1 lane per CPU core).
- Lanes decompose the data to be processed into non-overlapping ranges.
- Exchange operators (Repartition, Distribute) ensure lanes remain balanced.



Enables vertical scaling by adding more CPU cores.

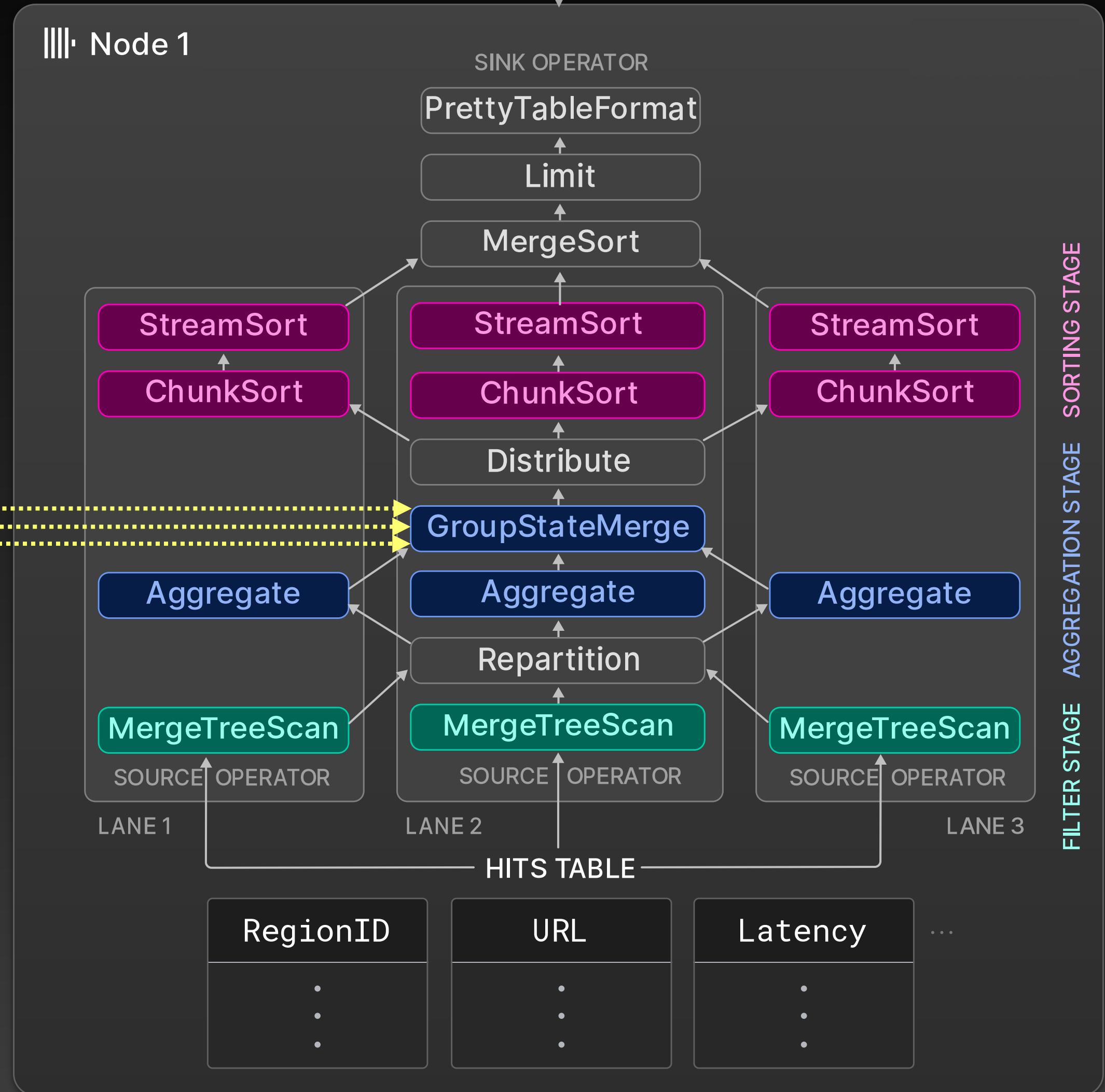
# Parallelization Across Cluster Nodes

- For sharded tables, the initiator node pushes as much work as possible to the other nodes.
- Results from remote nodes are integrated into different points of the initiator query plan.



```
SELECT RegionID, avg(Latency) AS AvgLatency  
FROM hits  
WHERE URL = 'https://clickhouse.com'  
GROUP BY RegionID  
ORDER BY AvgLatency DESC  
LIMIT 3
```

filter  
aggregation  
sort

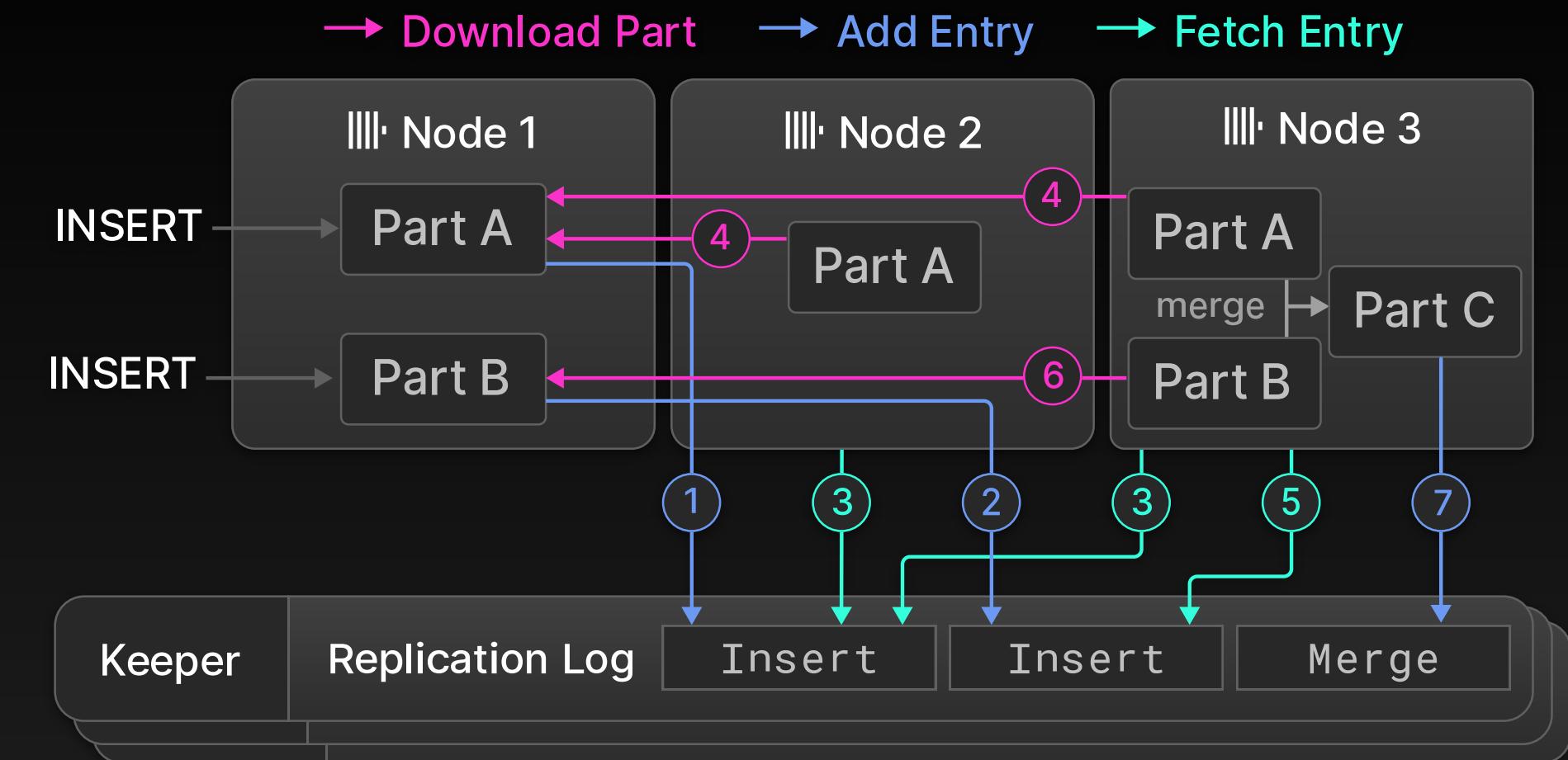


Enables horizontal scaling by adding more cluster nodes.

# What else is described in the paper?

## Additional storage layer details

- Updates and deletes, idempotent inserts
- Data replication
- ACID compliance

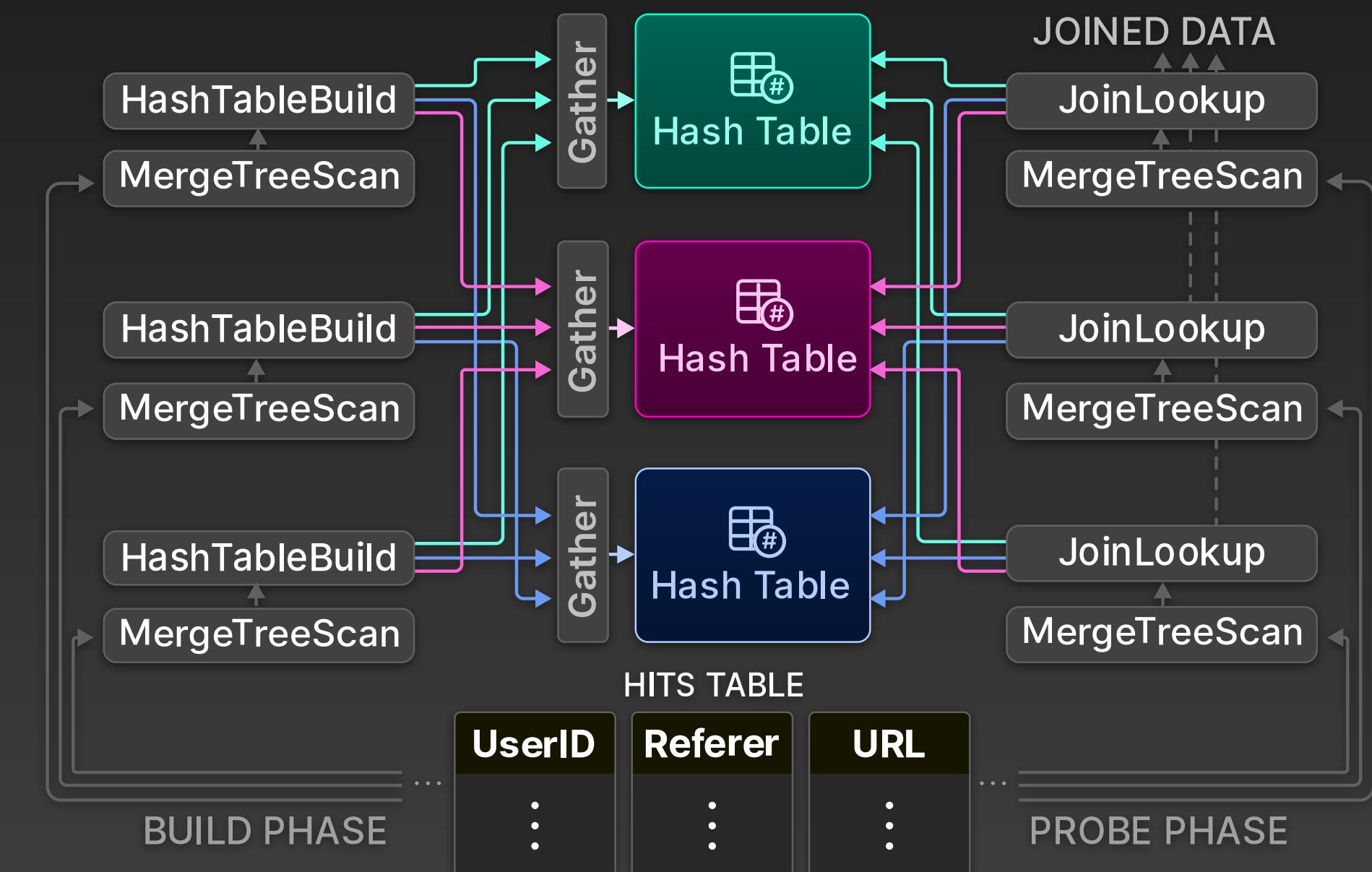


## Low-level query optimizations

- JIT query compilation based on LLVM
- Hash table framework for aggregations and joins
- Parallel join execution

## Integration layer

- Native support for 90+ file formats and 50+ integrations with external systems



**Come and join us on GitHub in our mission  
to build the world's fastest analytics database**



[github.com/ClickHouse/ClickHouse](https://github.com/ClickHouse/ClickHouse)

# Backup slides

# Benchmarks on Normalized Tables - TPC-H

	Q1	Q2	Q3	Q4	Q5	Q6	Q7-Q9	Q10	Q11	Q12	Q13	Q14	Q15	Q16	Q17	Q18	Q19	Q20- Q22
·	1.86		4.13		7.01	0.39		3.59	0.83	1.53		1.00	1.04	0.48		2.18		
*·	2.20		2.10		1.90	0.23		4.30	1.30	0.88		0.65	0.77	1.90		3.40		

The results of eleven queries are excluded:

- Queries containing **correlated subqueries**  
(not supported as of ClickHouse v24.6)
- Queries requiring **extended plan-level optimizations for joins**  
(missing as of ClickHouse v24.6)

Hot runtimes of the TPC-H queries based on the parallel hash join algorithm described in Section 4.4. The fastest of five runs was recorded.

Measurements taken on a single-node AWS EC2 c6i.16xlarge instance with 64 vCPUs, 128 GB RAM, and 5000 IOPS / 1000 MiB/s disk.

Comparable size was used for Snowflake (warehouse size L, 8×8 vCPUs, 8×16 GB RAM)

- Queries over normalized tables are an emerging use case for ClickHouse
- Automatic subquery decorrelation and better plan optimizer support for joins are planned for 2024.

<https://github.com/ClickHouse/ClickHouse/issues/58392>

# One more thing...

## There's a lot more to uncover

- Powerful SQL dialect with higher-order functions and lambda functions.
- 150+ built-in aggregate functions plus aggregate function combinator.
- 1300+ regular functions (mathematics, geo, machine learning, time series, etc.)
- Parallelized window functions and joins.
- JSON, maps and arrays plus 80+ array functions.