How to Make JOINs in ClickHouse go Brrr



9 Dec 2024 Robert Schulze ClickHouse Mysterious black box, not part of the presentation



Example SELECT Query

```
SELECT
    customer.name
FROM
    usage JOIN customers ON
        users.user_id =
customers.user_id
GROUP BY
    usage.customer_id
WHERE
    users.support_agreement = 'GOLD'
ORDER BY
    customers.name
```

Example SELECT Query

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```

Relational Operator

- 5 Project
- 2 Join



- (3) Aggregate
- (1) Filter
- (4) Sort

Why are JOINs important?

De-normalized Schema

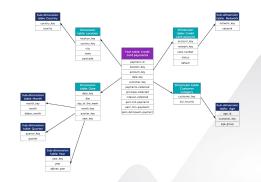
Table DesignSingle fact table

Characteristics Fast SELECTs, redundancy

Normalized Schema

Star or snowflake schema

Lower storage costs, fewer ETL steps, requires fast joins

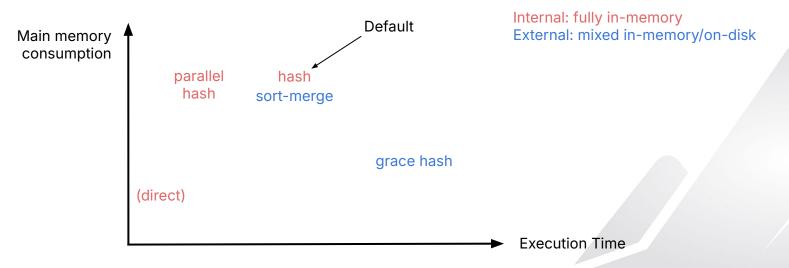


Status Quo

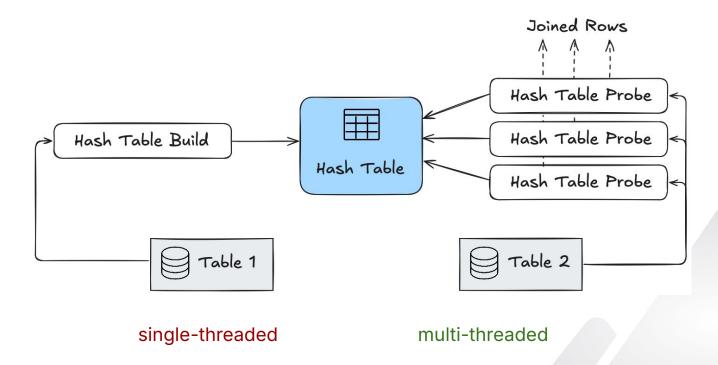
Support for all major join types:

```
o equi join tab_a.col = tab_b.col
o non-equi join tab_a.col > tab_b.col v24.12
```

Support for all major join algorithms

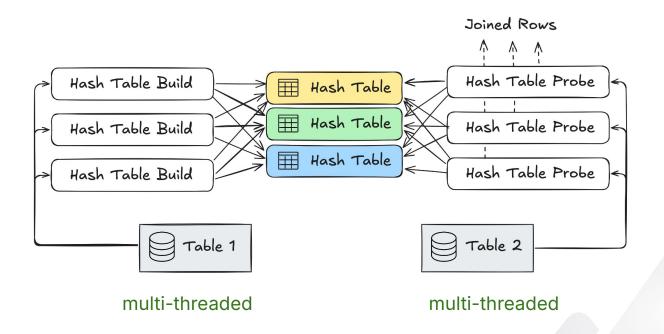


Hash Join



Optimization 1: Parallel Hash Join by Default





Up to 30% faster join queries

Optimization 2: Automatically Select the Build Side Table

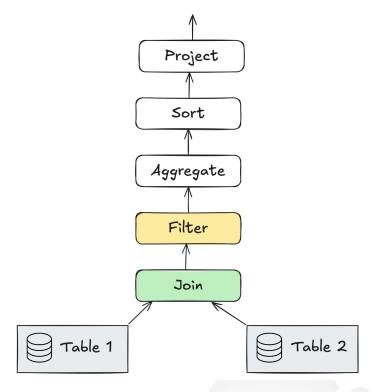
v24.12

- Hash tables are ideally as small as possible for L1/L2/Lx cache locality
- Previously: Hash table always built from right table in FROM clause of SELECT query
- Now: Hash table build from the smaller of both tables

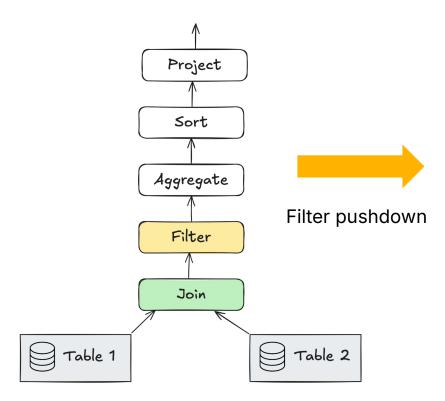
Up to 40% faster join queries

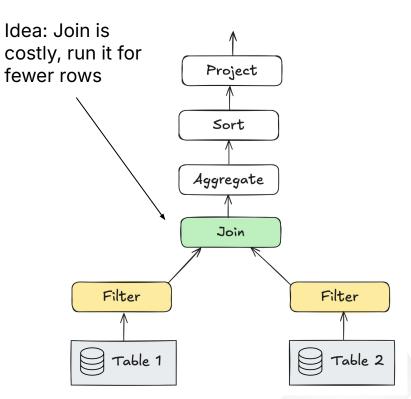
Filter Pushdown

```
SELECT
    customer.name
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    usage JOIN customers ON
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GROUP BY
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ORDER BY
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```



Filter Pushdown





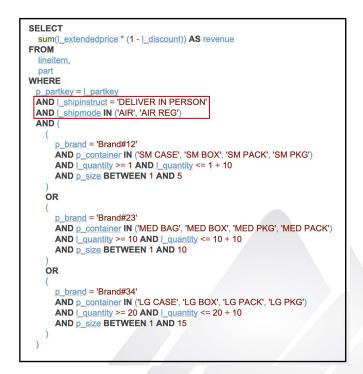
Optimization 3: Pushdown Filters Aggressively

v24.12

```
SELECT
 sum(| extendedprice * (1 - | discount)) AS revenue
FROM
 lineitem,
 part
WHERE
   p partkey = I partkey
   AND p_brand = 'Brand#12'
   AND p. container in ('SM CASE', 'SM BOX', 'SM PACK', 'SM PKG')
   AND | quantity >= 1 AND | quantity <= 1 + 10
   AND p size BETWEEN 1 AND 5
   AND | shipmode in ('AIR', 'AIR REG')
   AND | shipinstruct = 'DELIVER IN PERSON'
 OR
   p partkey = I partkey
   AND p brand = 'Brand#23'
   AND p container in ('MED BAG', 'MED BOX', 'MED PKG', 'MED PACK')
   AND | quantity >= 10 AND | quantity <= 10 + 10
   AND p size BETWEEN 1 AND 10
   AND I shipmode in ('AIR', 'AIR REG')
   AND | shipinstruct = 'DELIVER IN PERSON'
 OR
   p partkey = I partkey
   AND p brand = 'Brand#34'
   AND p container in ('LG CASE', 'LG BOX', 'LG PACK', 'LG PKG')
   AND | quantity >= 20 AND | quantity <= 20 + 10
   AND p size BETWEEN 1 AND 15
   AND I shipmode in ('AIR', 'AIR REG')
   AND | shipinstruct = 'DELIVER IN PERSON'
```

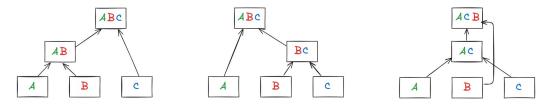


Applying De Morgan's Law enables filter pushdown

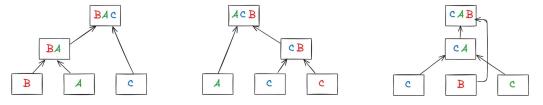


The rabbit hole goes deeper ...

Observation 1: JOINs are associative



Observation 2: JOINs are commutative (switch build & probe sides)



 Observation 3: As more tables are joined, the number of possible join orders explodes

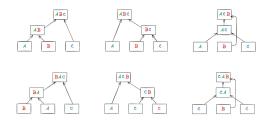
$$NumJoins(N) = \frac{(2N-2)!}{(N-1)!}$$

$$NumJoins(3) = 12$$

Work in Progress: Join Reordering

- v25.x
- Performance is mostly influenced by the JOIN order
- To find a good JOIN order, we need ...

Optimization Algorithm



Enumerate exhaustively vs. sub-set

Cost Model

Input: join order



Output: estimated costs

$$C = \begin{cases} |R| & \text{if R is base table} \\ |R| + C(|S|) + C(|T|) & \text{if R is a join between S and T} \end{cases}$$

Statistics about Base Tables

Cost models need statistics about values in a column:

- how many distinct values?
- top-10 most frequent values, etc.

Summary

- First JOIN optimizations will be in v24.11 / v24.12
- More foundational optimizations for JOINs (join reordering) are work-in-progress and expected in v25.x.

