# **CCNP Route Master Study Guide**

QUIZLET FLASHCARDS - https://quizlet.com/\_4ckrl4

# 1.0 Network Principles

# 1.1 Identify Cisco Express Forwarding concepts

- GENERAL CONCEPT
  - Considered as switching within routers
  - Enables faster packet transfer from one interface to another
  - This is due embedding relevant information in hardware
- PROCESS-BASED SWITCHING
  - Eats up CPU resources to switch packets
  - Known as "IP Input" in the CPU table
- FAST SWITCHING
  - Store data in "Fast Switching Cache" for faster lookups
  - CPU would have more breathing room
  - Packet would first need to be analyzed by the CPU before entering the cache
  - Does not track the routing table so the cached entry aging out may take a bit
- CISCO EXPRESS FORWARDING
  - Pre-populated cache, downloads information from arp/etc as soon as available
  - Enabled by default on routers and multilayer switches
  - CEF cannot be disabled on switches due to specialized hardware
  - Some packets on switches cannot be CEF-switched:
    - ARP Requests (Glean)
    - Packets requiring response from router CPU such as MTU too large
    - Routing Protocol Traffic
    - CDP or LLDP
    - Packets needing encryption

## 1.1.a FIB

- A.K.A. "Forwarding Information Base"
  - Shadow copy of the routing table
  - AD, Tags, Metrics, and more are not included in the FIB table
- VERIFICATION
  - show ip cef [detail]
  - show ip cef <ip address><mask> detail

# 1.1.b Adjacency Table

- GENERAL CONCEPT

- Pre-Populated with L2 tables such as:
  - ARP Table
  - Frame-Relay Map Table
- If FIB does not have a specific IP address...
  - An ARP request is sent out so that it can be seen in the FIB as 'attached'
  - This means there was a successful lookup

#### - ADJACENCY TYPES

- Glean
  - Need to wait for an ARP, so the request is punted to the CPU
  - Not enough L2 information to forward the L3 packet
  - If reply does not come back, entry stays as Glean until it ages out
- Null
  - Valid packet but needs to be dropped
  - Could happen because route has null0
  - Processed in hardware
- Drop
  - Valid packet but needs to be dropped
  - Could happen because no route or something is wrong with packet
- Discard
  - Valid packet but needs to be dropped
  - Could happen because we have a security policy, such as an ACL
- Punt
  - Similar to Glean, packets will be forwarded to the CPU
  - Punt is for packets that are destined for our processor to begin with

#### VERIFICATION

- show adjacency <intf type/number> [summary | detail]
- show adjacency vlan <vlan-id> detail

# 1.2 Explain general network challenges

# 1.2.a Unicast Flooding

- GENERAL CONCEPT
  - Occurs when no entry in L2 table (CAM)
    - Packet is flooded out all interfaces in the Vlan
    - Only exception is the received interface
  - To reduce flooding, make the ARP entry expire earlier than the CAM entry
    - The ARP entry would repopulate before CAM entry expires
  - To block unicast flooding on a per interface basis:
    - Router(config-if)# switchport block unicast

# 1.2.b Out-of-order packets

- GENERAL CONCEPT
  - Packets arriving in a different order than they were sent

- Could be caused by asymmetric routing or packet loss
- TCP attempts to alleviate this issue via the use of sequence numbers
  - Allows receiver to reorder the packets
- If packets arrive out of order, the receiver may send *duplicate ACKs*
- This could trigger the "Fast Retransmit" algorithm
  - The sender will assume packet loss, and retransmit sequences
  - The sender will reduce window size
  - This leads to reduced TCP throughput

## 1.2.c Asymmetric routing

- GENERAL CONCEPT
  - When the return path is different from the original path of a packet
  - Not normally a problem unless there are mechanisms to drop the data
    - Examples: NAT, Firewalls, Unicast RPF

# 1.3 Describe IP operations

## 1.3.a ICMP Unreachable and Redirects

- ICMP UNREACHABLE (TYPE 3)
  - Disabled by default for security reasons
  - Feedback mechanism to notify the host about an unreachable destination
  - Up to the sender to act upon that information from router
  - Six failure feedback codes:
    - **Code-0**: "Network Unreachable", no route to the network in question
    - Code-1: "Host Unreachable", no host exists on that subnet
    - **Code-2**: "Protocol Unreachable", unknown protocol destined for router
    - Code-3: "Port Unreachable"
    - Code-4: "Fragmentation required and DF bit set"
    - Code-5: "Source Route Failed"
- ICMP REDIRECTS (TYPE 5)
  - A packet destined for a specific gateway which is not the preferred gateway is then redirected to the preferred gateway
  - The sender is then sent an ICMP Redirect message
    - Indicates that the message has been redirected
    - Suggests to use the preferred gateway going forward

## 1.3.b IPv4 and IPv6 fragmentation

- GENERAL CONCEPT
  - Payload of an IP packet is greater than the MTU of the link
    - Must be fragmented, otherwise it will be dropped
    - Routers will fragment packets in IPv4, unless the DF bit is set
  - Only source host fragments IPv6 packets using Extension Header

- Two choices for IPv6 packets to avoid fragmentation:
  - Use **Path MTU Discovery** (PMD) to find the lowest MTU in path
  - Use the minimum MTU size (1280 Bytes)

## 1.3.c TTL

- GENERAL CONCEPT
  - Originally intended for routers to decrement the seconds
    - Based upon how long each packet was queue
    - This feature was never implemented
  - Instead, TTL is used as a glorified hop count system
    - Each router decrements the TTL by 1
    - Upon reaching 0, the packet is dropped
  - Traceroute uses this concept to determine hops
    - Increments TTL for each packet until destination is reached
    - Each hop will report back "ICMP Time Exceeded"
    - Final destination will report back with "ICMP Destination Unreachable"

# 1.4 Explain TCP operations

## 1.4.a IPv4 and IPv6 (P)MTU

- GENERAL CONCEPT
  - MTU A.K.A. "Maximum Transmission Unit"
  - The maximum amount of bytes supported in the payload of the transmission
  - Default MTU is set as 1500 bytes
  - Different technologies can add more bytes to headers requires MTU adjustment
- CONFIGURATION STEPS
  - Router(config-if)# mtu <bytes>
- VERIFICATION
  - show system mtu
- PATH MTU DISCOVERY CONCEPT
  - PMD can be used to avoid fragmentation
  - IPv6 source sends out probes to discover lowest MTU in path
  - Packet Too Big ICMPv6 message will be sent back to host, containing MTU
  - IPv4 will have DF A.K.A. "Don't Fragment" bit set
  - Any router that drops the packet in the path will send back ICMP Type 3 Code 4

## 1.4.b MSS

- GENERAL CONCEPT
  - MSS A.K.A. "Maximum Segment Size"
  - The amount of data a host will accept in a single TCP/IP datagram
  - MSS > MTU + protocol overhead = fragmentation at an IP level
    - Should attempt to avoid fragmentation when using tunneling techniques

- Typically, MSS is calculated at 1460 bytes
  - Allocates 20 bytes respectively for TCP and IP header overhead
- The MSS of a host is sent in TCP SYN
  - Each host uses the lowest of the two values

#### CONFIGURATION STEPS

- Router(config-if)# ip tcp adjust-mss <max-segment-size>
- Router(config-if)# ip mtu <mtu-size>

#### VERIFICATION

- show ip interface <intf type/number>

## 1.4.c Latency

## - GENERAL CONCEPT

- Often defined by RTT A.K.A "Round Trip Timer"
- Length of time it takes to receive a response back
- TCP Latency has an inverse relationship with throughput
  - More latency, less throughput

## 1.4.d Windowing

#### GENERAL CONCEPT

- Allows single acknowledgment of multiple TCP segments
- Window Size specifies how many bytes may be sent before an ACK is required
  - **Example**: Window Size 1000, ACK 1001 = successful, increase size
  - **Example**: Window Size 1000, ACK 900 = unsuccessful, decrease size
- Can be adjusted based upon host requirements

## 1.4.e Bandwidth-delay product

## - GENERAL CONCEPT

- Amount of data that can be at transit at any given point of time
- Calculated by multiplying the link's bandwidth in bits and Round Trip Delay Time
- Networks with large bandwidth-delay product are known as LFN
  - LFN A.K.A. "Long Fat Network"
  - **Example**: Satellite Link high bandwidth but huge delays
- TCP Window Scaling is used to alleviate the issue

# 1.4.f Global synchronization

#### - GENERAL CONCEPT

- Refers to how TCP streams gradually increase window sizes until drops occur
- This causes TCP streams to shrink at once, then repeat the process
- Ends up with a sawtooth-like graph for bandwidth utilization on the link
- One way to prevent this issue is by discarding random TCP streams
  - A.K.A. "Random Early Detection" queuing
  - Only some streams will back off, allowing more efficient link utilization

# 1.5 Describe UDP operations

## 1.5.a Starvation

#### - GENERAL CONCEPT

- Occurs when TCP and UDP streams occupy the same queue
- When congestion occurs...
  - TCP reacts by reducing Window Size, reducing bandwidth utilization
  - UDP reacts by eating up the newly available bandwidth
- This results in **UDP Dominance** and **TCP Starvation**
- Solution is the create separate queues for TCP and UDP, or use QoS

## 1.5.b Latency

## - GENERAL CONCEPT

- UDP does not suffer from throughput related issues as TCP
  - UDP does not expect to receive ACKs for the data
  - UDP is not affected by packet loss or latency
- UDP latency can affect application performance, such as VoIP
- UDP is generally used for real time applications
  - These applications can be sensitive to latency and jitter
  - Jitter is the variation of latency over time
- These applications may attempt to alleviate these issues by using buffers

# 1.6 Recognize proposed changes to the network

# 1.6.a Changes to routing protocol parameters

#### - GENERAL CONCEPT

- Metrics can be changed to provide added flexibility to routing protocols
- Other types of changes involve redistribution or added routes

## 1.6.b Migrate parts of the network to IPv6

## - GENERAL CONCEPT

- Migration techniques involve "Dual Stack" and "Tunneling"
- Tunneling is encapsulating IPv6 data into an IPv4 packet
- Dual Stack involves running both IPv4 and IPv6 on the device

#### TUNNELING CONCEPT

- Router or host encapsulates the IPv6 packet inside an IPv4 packet
- Results in fewer routers needing any IPv6 configuration at all
- Two main categories of tunnels are: point-to-point and multipoint
- Protocol number of the IP Packet will be 41 (IPv6) which means encapsulation

#### - POINT-TO-POINT IPv6 TUNNEL CONCEPT

- Two devices sit at the ends of the tunnel
- These point-to-point tunnels work like virtual point-to-point serial links
- Each router configures a type of virtual interface called tunnel interface
- Point-to-Point tunnels work best when IPv6 occurs regularly

#### POINT-TO-MULTIPOINT IPv6 TUNNEL CONCEPT

- Allows router to use a tunnel interface to send packets to multiple destinations
- Additional logic so that the sending router knows which remote router to send to
- Point-to-Multipoint works best when IPv6 traffic occurs infrequently
- Does not support IPv6 IGPs, thus requiring statics or BGP

#### OVERVIEW OF IPv6 TUNNELING OPTIONS

- Manually Configured Tunnel
  - Point-to-Point tunnel that is configured manually
  - Acts like a virtual point-to-point link, supporting IPv6 IGPs
  - Good for more permanent tunnels
  - Slightly less overhead than GRE

#### GRE Tunnel

- Point-to-POint tunnel that is configured manually
- Same advantages as manually configured tunnels
- Can support other Layer 3 protocols over the same tunnel

#### 6to4 Tunnel

- Multipoint tunnel that is formed dynamically
- Require less configuration than all other types when adding new site
- Supports global unicasts, with some extra configuration
- Uses second and third quartets to store IPv4 address

#### - ISATAP Tunnel

- Multipoint tunnel is dynamically formed
- Easily supports global unicast addresses for all prefixes
- Uses seventh and eighth quartets to store IPv4 address

#### - IPv4-Compatible Tunnel

- Does not scale well for large networks, not recommended
- Address format is ::IPv4Address, automatically created

#### - IPv6 Rapid Deployment (6rd) Tunnel

- An extension of the 6to4 feature
- Allows an ISP to offer IPv6 service to customers over its IPv4 network

#### - Teredo Tunnel

- Known as shipworm
- Two dual-stacked devices can speak even with IPv4 NAT in between
- Uses UDP port 3544 to communicate with Teredo servers
- These servers are used as dispatchers between clients and relays
- Although proposed by Microsoft, there is a version for Linux
- Defined in RFC 4380
- Uses 2001::/32 as the prefix, with the format being:
  - 32-bits: 2001::/32
  - 32-bits: Server Public IPv4 Address

- 16-bits: Flags
- 16-bits: Obfuscated Client UDP Port
- 32-bits: Obfuscated Client Public IPv4 Address

#### MANUALLY CONFIGURED TUNNELS

- Many similarities between GRE:
  - Both create virtual point-to-point links between two IPv4 routers
  - IPv6 IGP routing protocols can be run over these virtual links
  - Difference is that MCT encapsulates packets without an additional header

#### Configuration Steps:

- Router(config)# interface loopback {number}
- Router(config-if)# ip address {ipv4-address} {subnet}
- Router(config)# interface tunnel {number}
- Router(config-if)# tunnel source {interface | ipv4-address}
- Router(config-if)# tunnel destination {ipv4-address}
- Must match tunnel source on the other router!
- Router(config-if)# tunnel mode ipv6ip
- Router(config-if)# ipv6 address {ipv6-address/prefix}

#### - Verification:

- show interface tunnel
- show ipv6 interface brief

#### - GENERIC ROUTING ENCAPSULATION TUNNELS

- Only one difference between MCT and GRE: tunnel mode
- Configuration Steps:
  - Router(config-if)# tunnel mode gre ip
- IOS Default MTU Settings:

GRE: 1476MCT: 1480

### - IPv6 AUTOMATIC 6TO4 TUNNELS

- Addressing Option-1:
  - Use the 2002 prefix, reserved for 6to4 protocol
  - Example: 2002:0101:0101:xxxx::/64 -> 1.1.1.1
  - Embeds the IPv6 information within the IPv6 address

#### Addressing Option-2:

- Provide global prefix to the IPv6 cloud segments
- IPv4 address cannot be embedded in this scenario.
- This in return requires routing to be specific with next-hop

#### Configuration Steps:

- Router(config)# interface tunnel {number}
- Router(config-if)# tunnel source {interface | ipv4-address}
- Derive the /48 prefix used for allocating local IPv6 subnets
- Router(config-if)# ipv6 address 2002:0101:0101::1/64
- Router(config-if)# tunnel mode ipv6ip 6to4
- Router(config)# ipv6 route <remote global prefix> tunnel 0 <next-hop>
- The next hop command is only required with option-2 addressing

#### Verification:

- show ipv6 route
- show ipv6 interface tunnel

#### ISATAP TUNNELS

- Enables reachability for IPv6 Hosts not directly connected to an IPv6 router
- ISATAP A.K.A. "Intra-Site Tunnel Addressing Protocol"
- Designed for transporting IPv6 packets within a site
- 64-Bit Interface Identifier (after the /64 prefix):
  - First 32 bits contain the value 0000:5EFE to indicate ISATAP address
  - Remaining 32 bits encode the IPv4 address
  - ipv6prefix:0000:5efe:ipv4address
- Allows hosts to tunnel IPv6 packets through an IPv4 domain
- This requires the use of a DNS server
- Configuration Steps:
  - Router(config)# ipv6 unicast-routing
  - Router(config)# interface tunnel {number}
  - Router(config-if)# ipv6 add 2001:1111:2222:aaaa::/64 eui-64
  - Router(config-if)# no ipv6 nd suppress-ra
  - Allows a router to send router advertisements
  - Router(config-if)# tunnel source {interface | ipv4-address}
  - Router(config-if)# tunnel mode ipv6ip isatap
- IPv6 RAPID DEPLOYMENT (6RD) TUNNELS
  - Main Differences Between 6rd and 6to4:
    - Does not require 2002::/16 prefix and can be from service provider
    - Not all 32 bits of the IPv4 destination address need to be carried
    - The IPv4 destination address is obtained from a combination of:
      - Bits in the payload header
      - The information configured in the router
    - Defined in RFC 5569

## 1.6.c Routing protocol migration

- GENERAL CONCEPT
  - Routing protocol migration can be done by setting up boundaries
  - Modifying AD can ensure no routing issues

# 2.0 Layer 2 Technologies

# 2.1 Configure and verify PPP

- GENERAL CONCEPT
  - PPP A.K.A. "Point-to-Point Protocol"
  - Operates in the LLC sub-layer of the data link layer in OSI
  - Server generally known as **Access Concentrator**
- PPP STAGES
  - **Phase-1**: Active Discovery (only in PPPoE)
  - **Phase-2**: Link Control Protocol (LCP)
  - **Phase-3**: Authentication (CHAP or PAP)

- **Phase-4**: Network Control Protocol (NCP)
- ACTIVE DISCOVERY (PPPoE)
  - → PADI (PPPoE Active Discovery Initialization)
    - Broadcast from client to AC
    - "Are there any PPPoE Servers out there? My unique Host-ID is xx-xx"
  - ← PADO (PPPoE Active Discovery Offer)
    - Unicast from AC to client
    - "Yes, I'm here xx-xx. My unique AC ID is yy.yy"
  - → PADR (PPPoE Active Discovery Request)
    - Unicast from client to AC
    - "Thanks for the info! Can I have a Session-ID please?"
  - ← PADS (PPPoE Active Discovery Session-Confirmation)
    - Unicast from AC to client
    - "Yes. let's use Session-ID 0x02"

#### - LINK CONTROL PROTOCOL (LCP) OPTIONS

- Callback Option
  - Reinitiate dial session on the lower rate side
  - Must be configured on both ends of the session
- Multilink Option
  - If frames are larger than MTU, will initiate fragmentation
  - MRRU A.K.A. "Maximum Received Reconstructable Unit"
    - Specifies maximum amount of bytes reconstructable
    - Example: MTU is 1500 bytes, MRRU is 3000 bytes
    - Once byte limit is reached, packet will be dropped
  - MRU A.K.A. "Maximum Receivable Unit"
    - Specifies the MTU size, do not confuse with MRRU
- Authentication Option
  - Specifies the type of authentication if any are to be done
- Magic Number Option
  - Loop prevention system
  - If the same magic number is sent and received, then assumes loop

#### - LINK CONTROL PROTOCOL (LCP) CONTROL MESSAGES

- Configuration-Request
  - Lists all link specific options a sender wishes to implement
- Configuration-Reject
  - When a receiver does not support a particular feature and offers no alternatives
- Configuration-NAK (Negative Acknowledgement)
  - Receiver does not support a particular feature and offers an alternative
- Configuration-Acknowledgement
  - Acknowledging all LCP options in the most recent received Config-Req

#### LCP CONFIGURATION STEPS

- Router(config-if)# encapsulation ppp
- Router(config-if)# ppp multilink
- NETWORK CONTROL PROTOCOL (NCP)

- Negotiate what Layer 3 Protocol to use:
  - IP: IPCPIPX: IPXCP
  - CDP: CDPCP
- DHCP could also be used but not really necessary as IPCP is easier

#### NCP CONFIGURATION STEPS

- Specifically, this is for IPCP in case DHCP is not preferred
- Router(config-if)# ip address negotiated
- Router(config-if)# peer default ip address <x.x.x.x | dhcp | dhcp-pool | pool>

## 2.1.a Authentication (PAP, CHAP)

#### PAP GENERAL CONCEPT

- PAP A.K.A. "Password Authentication Protocol"
- Sends clear text username and password for authentication
- Two-way handshake and by default sends hostname as username
- Once PPP authentication is done, PPP session will stay up and established

#### - PAP CONFIGURATION STEPS

- Client One Way
- Router(config-if)# encapsulation ppp
- Router(config-if)# ppp authentication pap callin
- Router(config-if)# ppp pap sent-username Chris password Cisco
- Client Two Way
- Router(config)# username Sally password Server
- Router(config-if)# encapsulation ppp
- Router(config-if)# ppp authentication pap callin
- Router(config-if)# ppp pap sent-username Chris password Cisco

#### - CHAP GENERAL CONCEPT

- CHAP A.K.A. "Challenge Handshake Authentication Protocol"
- Three-way handshake and by default sends hostname as username
- A CHAP challenge is sent from the server
  - Combination of random numbers
  - Password never sent across the link
  - Challenge + Username + Password = MD5 HASH
- Client sends MD5 hash across link to server and server verifies validity
- Once PPP authentication is done, PPP session will stay up and established

#### CHAP CONFIGURATION STEPS

- Client One Way
- Router(config-if)# encapsulation ppp
- Router(config-if)# ppp authentication chap callin
- Router(config-if)# ppp chap hostname Chris
- Router(config-if)# ppp chap password Cisco
- Alternative Client
- Router(config)# username Sally password Cisco
- Router(config-if)# encapsulation ppp
- Router(config-if)# ppp authentication chap callin
- Client Two Way
- Router(config)# username Sally password Cisco
- Router(config-if)# encapsulation ppp

- Router(config-if)# ppp authentication chap
- Router(config-if)# ppp chap hostname Chris
- Router(config-if)# ppp chap password Cisco

#### VERIFICATION

- debug ppp negotiations
- debug ppp authentication
- **show interface serial** <number>
- show users

## 2.1.b PPPoE (client side only)

#### - GENERAL CONCEPT

- PPPoE A.K.A. "PPP Over Ethernet"
- Allows single DSL connection to support an entire LAN of PPP clients
- BBA A.K.A. "BroadBand Access"

#### PPPoE CONFIGURATION STEPS

- Server Side
- Router(config)# hostname Router
- Router(config)# username client password cisco
- Router(config)# bba-group pppoe PPPOE-GROUP
- Router(config-bba-group)# virtual-template 1
- Router(config)# interface virtual-template 1
- Router(config-if)# ip unnumbered loopback 0
- Router(config-if)# peer default ip address pool MY-POOL
- Router(config-if)# ppp authentication chap
- Router(config)# ip local pool MY-POOL 1.2.1.2 1.2.1.254
- Router(config)# interface FastEthernet0/0
- Router(config-if)# pppoe enable group PPPOE-GROUP
- Client Side
- Router(config)# hostname Router
- Router(config)# interface Dialer 7
- Router(config-if)# ip address negotiated
- Router(config-if)# encapsulation ppp
- Router(config-if)# dialer pool 1
- Router(config-if)# ppp chap password cisco
- Router(config)# interface FastEthernet0/0
- Router(config-if)# pppoe-client dial-pool-number 1

### - VERIFICATION

- show pppoe session
- show pppoe summary
- **show interface virtual-access** <number>

#### PPPoE MTU CONSIDERATIONS

- PPP adds 8 bytes of overhead on Ethernet Frame
- 1508 Bytes + 14 Bytes = 1522 Bytes
  - Results in fragmentation
  - CPU-intensive process
- Server side not a huge issue
  - Virtual-Access interfaces defaults to 1492 MTU
- Client side it is an issue

- Dialer interface defaults to 1500 MTU
- Adjust-MSS intercepts TCP packets and adjusts MSS to avoid fragmentation
  - Default for Windows is 1460 otherwise

#### PPPoE MTU CONFIGURATION STEPS

- Client Side
- Router(config)# interface Dialer 7
- Router(config-if)# ip mtu 1492
- Router(config)# interface FastEthernet0/0
- Router(config-if)# ip tcp adjust-mss 1452

#### PPPoE VPDN CONFIGURATION STEPS

- Client Side
- Router(config)# vpdn enable
- Router(config)# vpdn group CBTNuggets request dialout pppoe
- Router(config)# vpdn group CBTNuggets ppp authentication cpap/chap/mschap>
- Router(config)# vpdn group CBTNuggets localname ninja
- Router(config)# vpdn username ninja password secretninja store-local
- Router(config)# interface fa0/0
- Router(config-if)# ip address pppoe setroute
- Router(config-if)# pppoe client vpdn group CBTNuggets

# 2.2 Explain Frame Relay

## 2.2.a Operations

#### - GENERAL CONCEPT

- Multipoint layer 2 technology
- Legacy technology typically used in service provider end
- ISP Frame Relay Switches monitor data usage
  - This is done by monitoring transmit and receive buffers
  - If congestion occurs, various bits in Frame-Relay header are modified

#### FRAME RELAY ADVANCED CONCEPTS

- PVC A.K.A. "Private Virtual Circuit"
  - Service provider carries multiple customer traffic in a single link
- DLCI A.K.A. "Data-Link Connection Identifier"
  - Works as a Layer 2 address in Frame Relay
  - 10-bit value (can be extended) that ranges from 0 to 1023
    - 0-15 and 1007-1023 are reserved
- LMI A.K.A. "Local Management Interface"
  - Works as a keepalive between Frame Relay switch and end device
  - Propagates DLCI information to hub and spoke
  - Standards of LMI: Cisco, Q933a, ANSI
  - Uses reserved DLCI number of 1023 or 0
- CIR A.K.A. "Committed Information Rate"
  - Amount of data rate that a service provider guarantees
  - Anything above CIR is considered as BURST
  - Closer CIR is to access rate (total bandwidth), the more expensive

- Inverse ARP
  - Dynamically maps the destination IP with corresponding local DLCI
  - Only one DLCI can be mapped with a Layer 3 address
  - Subinterfaces disable this functionality

#### FRAME RELAY CONGESTION

- DE A.K.A. "Discard Eligible"
  - If customer goes above CIR, their packets are marked with DE bit
  - If there is congestion and DE bit is set, packets can be dropped WAN
- FECN A.K.A. "Forward Explicit Congestion Notification"
  - If the transmit buffer is near full, WAN switch will set the FECN bit
  - Just informational, notifies switch that there is congestion going towards it
- BECN A.K.A. "Backward Explicit Congestion Notification"
  - Sent to notify router about congestion if frames are sent toward DLCI
  - Can have the router respond to BECNs by slowing down transmission

#### - LMI EXTENSION OPTIONS

- Virtual Status Messages
  - Provide communication between the network and user device
  - Prevents data from entering black holes when PVC fails
- Multicasting
  - Allows a sender to transmit a single frame to multiple recipients
- Global Addressing
  - Gives connection identifiers global rather than local significance
  - This allows the Frame Relay cloud to perform like one big LAN
- Simple Flow Control
  - Provides flow control mechanisms that applies to the entire interface
  - Intended for devices that cannot use the congestion notification bits

#### - VERIFICATION

- show frame-relay map

## 2.2.b Point-to-point

### CONFIGURATION STEPS

- Router(config-if)# no ip address
- Router(config-if)# encapsulation frame-relay
- Router(config-if)# no shut
- Router(config-if)# exit
- Router(config)# interface serial x/y.<subinterface number> point-to-point
- Router(config-sub-if)# ip address <address> <subnet mask>
- Router(config-sub-if)# frame-relay interface-dlci <dlci>

## 2.2.c Multipoint

#### - CONFIGURATION STEPS

- Physical Interface Configuration
- Router(config-if)# encapsulation frame-relay
- Router(config-if)# ip address <address> <subnet mask>
- Configure Static Mapping (required with subinterfaces)

# 3.0 Layer 3 Technologies

# 3.1 Identify, configure, and verify IPv4 addressing and subnetting

- 3.1.a Address types (Unicast, broadcast, multicast, and VLSM)
  - GENERAL CONCEPT
    - Unicast
      - Assigned to a single network interface located on a specific subnet
      - Used for **one-to-one** communications
    - Broadcast
      - Assigned to all network interfaces located on a subnet on the network
      - Used for **one-to-everyone** communications
    - Multicast
      - Assigned to one or more network interfaces located on various subnets
      - Used for **one-to-many** communications
    - VLSM
      - VLSM A.K.A. "Variable-Length Subnet Mask"
      - Allows divisions of an IP address space beyond classful system
      - Introduced alongside of CIDR to slow IPv4 address pool depletion

## 3.1.b ARP

- ARP A.K.A. "Address Resolution Protocol"
  - Communication protocol
  - Used to discover link layer address associated with a given IPv4 address
- 3.1.c DHCP relay and server
  - DHCP A.K.A. "Dynamic Host Configuration Protocol"
    - See **Section 6.5** for expanded details
  - DHCP RELAY AGENT
    - Router forwards DHCP discovery as unicast from client to DHCP server
    - This transaction is transparent to the client
  - DHCP RELAY CONFIGURATION
    - Router(config-if)# ip helper-address 192.168.23.3
  - DHCP SERVER CONFIGURATION
    - DHCP(config)# ip dhcp pool EXAMPLE
    - DHCP(dhcp-config)# network 192.168.12.0
    - DHCP(dhcp-config)# dns-server 208.67.222.222
    - DHCP(dhcp-config)# **default-router** 192.168.12.1
    - DHCP(config)# ip dhcp excluded-address 192.168.12.100
  - VERIFICATION

- show ip dhcp binding

## 3.1.d DHCP protocol operations

- D.O.R.A. A.K.A. "Discover, Offer, Request, Acknowledge"
  - See <u>Section 6.5</u> for expanded details
  - DHCP built on top of bootstrap protocol (bootp)

# 3.2 Identify IPv6 addressing and subnetting

## 3.2.a Unicast

- GENERAL CONCEPT
  - IPv6 String Format:
    - They are 128-bits in length and in hex
    - Represented in all lowercase
    - x:x:x:x:x:x:x
      - Each x is a 16-bit section
      - Each section referred to as a hextet
  - Helpful Notation Rules:
    - Rule 1: Omit Leading 0s
      - **Example**: 2001:0001:0010... -> 2001:1:10...
    - Rule 2: Omit All-0s Hextets
      - **Example**: fe80:0000:0000: ... :0001 -> fe80::1
- GUA A.K.A. "Global Unicast Address"
  - Global Routing Prefix: Prefix of the address assigned by provider
  - **Subnet ID**: Separate field for allocating subnets within customer site
  - Interface ID: Identifies the interface on a subnet, generally 64-bits
  - Several ways a device can be configured with GUA:
    - Manually configured
    - Stateless Address Autoconfiguration (SLAAC)
    - Stateful DHCPv6
- LINK-LOCAL UNICAST ADDRESS
  - All devices MUST have an IPv6 link-local address, fe80::/10
  - Not routable off the link
  - Only unique on the link
  - Only one link-local address per interface
  - Devices can use DAD A.K.A. "Duplicate Address Detection"
- LOOPBACK ADDRESSES
  - An IPv6 loopback address is ::1
  - Equivalent to IPv4 127.0.0.0/8
- UNSPECIFIED ADDRESSES
  - An all-0s address (::)
  - Used as a source address to indicate the absence of an address

- Cannot be assigned to an interface
- Cannot be used as destination address
- Router will never forward a packet that has an unspecified source address
- ULA A.K.A. "Unique Local Addresses"
  - Counterpart of IPv4 private addresses, fc00::/7
  - Allow sites to be combined or privately interconnected
  - NAT should **not** be used to translate between ULA and GUA
  - RFC 4193 defines a process whereby...
    - All locally assigned Global IDs can be generated dynamically
    - Pseudo random algorithm gives it a very high probability of being unique
  - **Site-Local Addresses** (the original ULA) have been deprecated (fec0::/10)
- IPv4 EMBEDDED ADDRESS
  - Used to aid the transition from IPv4 to IPv6
  - Features such as NAT64 are required to translate between the two address families
  - **Example Address**: ::ffff:192.168.10.10

## 3.2.b EUI-64

- GENERAL CONCEPT
  - Two options to generate the Interface ID of an IPv6 address:
    - EUI-64 process
    - Random 64-bit value (privacy extension)
  - The EUI-64 uses the Ethernet MAC address to generate the Interface ID
- CONVERSION STEPS
  - Step-1: Convert MAC to binary, then split in half
  - Step-2: FF:FE is inserted between the two halves
  - Step-3: Flip the seventh bit -> Universally/Locally bit A.K.A. "Local/Global bit"
  - Result: aaaa.aaaa.aaaa -> fe80::a8aa:aaff:feaa:aaaa

## 3.2.c ND, RS/RA

- Covered in depth in Section 6.5.a

# 3.2.d Autoconfig (SLAAC)

- Covered in depth in **Section 6.5.a** 

# 3.2.e DHCP relay and server

- Covered in depth in Section 6.5.a

## 3.2.f DHCP protocol operations

- Covered in depth in Section 6.5.a

# 3.3 Configure and verify static routing

- GENERAL CONCEPT
  - Static Routes are manually configured routes
- IPv4 CONFIGURATION STEPS
  - Router(config)# ip route <prefix> <mask> {interface | next-hop}
- IPv6 CONFIGURATION STEPS
  - Router(config)# ipv6 route <prefix>/<length> {interface | next-hop}
- VERIFICATION
  - show ip route
  - show ipv6 route

# 3.4 Configure and verify default routing

- GENERAL CONCEPT
  - Default Route is a route a router would use if there is no specific route available
- STATIC ROUTE CONFIGURATION STEPS
  - Router(config)# ip route 0.0.0.0 0.0.0.0 {interface | next-hop}
- EIGRP CONFIGURATION STEPS
  - Three Methods:
    - Option-1: By advertising a static default route with EIGRP
    - **Option-2**: By configuring a default network
    - **Option-3**: By using the summary-address command
  - Static Default Route Method:
    - Router(config)# ip route 0.0.0.0 0.0.0.0 null0
    - Router(config-router)# network 0.0.0.0 or
    - Router(config-router)# redistribute static
  - Default Network Method:
    - A **classful** network advertised into EIGRP and marked with a flag
    - On the router on which all traffic should be directed
      - Identify a classful network that can be advertised into EIGRP
      - Ensure the network is being advertised
    - Router(config)# ip default-network {network-number}
  - Summary-Address Method:
    - Identify specific interfaces for outgoing default route
    - ONLY the default route will be advertised
    - Default will be locally installed as a summary route
    - Router(config-if)# ip summary-address eigrp <ASN> 0.0.0.0 0.0.0.0
- OSPF CONFIGURATION STEPS
  - Two Methods:
    - Option-1: Using the default-information originate command
    - Option-2: Using stub areas
  - Default Information Originate Method:
    - Used on ASBRs to flood a default route into the entire OSPF domain
    - Only injects the default route as Type 2 using Type 5 LSA with metric 20
      - Default route must exist in the routing table already

- With the always parameter...
  - Default route is advertised no matter what
- Router(config-router)# default-information originate [always]

#### Stub Areas Method:

- With a stub area, ABR injects a default route into an area
- In parallel, ABR will not advertise external routes (5 LSA)
- There are several features of stub areas:
  - ABRs create a default route via Type 3 LSA
  - ABRs do not flood Type 5 LSA into the stub area
  - ABRs may not flood other Type 3 LSAs into the area
  - The default route has a metric of 1 unless otherwise configured
  - Router(config-router)# area <area-num> default-cost <cost>
- Cannot redistribute external routes into the stubby area
- Router(config-router)# area <area-num> stub

#### BGP CONFIGURATION STEPS

- Three Methods:
  - Option-1: Advertise via network command
  - Option-2: Redistribution
  - Option-3: Default Originate

#### - Advertise via Network Command Method:

- Inject default route into BGP only if...
  - Default route is currently present in the routing table
  - Router(config-router)# network 0.0.0.0
- Redistribution:
  - Injects default route into BGP only if...
    - Default route is currently present in the routing table
    - Learned by the specific protocol we are redistributing from
  - Router(config-router)# redistribute <protocol>
- Default Originate:
  - Two ways, either globally or per neighbor
  - Router(config-router)# default-information originate <- global
  - Router(config-router)# neighbor <ip-addr> default-originate <- neighbor</li>

#### VERIFICATION

show ip route

# 3.5 Evaluate routing protocol types

#### 3.5.a Distance vector

- GENERAL CONCEPT
  - The router has no idea or scope of the entire topology, just the relative costs

#### 3.5.b Link state

- GENERAL CONCEPT

- The router has the entire scope of the topology
- The router constructs a map based off the info

## 3.5.c Path vector

- GENERAL CONCEPT
  - The path of the route is maintained, but nothing else
  - **Example**: Which Autonomous Systems did the route traverse through

## 3.6 Describe administrative distance

Routing Protocol	Administrative Distance	
Connected	0	
Static	1	
EIGRP Summary	5	
eBGP	20	
EIGRP	90	
OSPF	110	
RIP	120	
External EIGRP	170	
iBGP	200	
NHRP	250	

- AD A.K.A. "Administrative Distance"
  - A value used to rank routes from most preferred to least preferred
  - Used to prevent potential loops by installing a route based on preferred protocol
- CONFIGURATION STEPS
  - Router(config-router)# distance <ad-num>
- VERIFICATION
  - show ip route

# 3.7 Troubleshoot passive interfaces

- GENERAL CONCEPT
  - Passive interfaces do not send hello messages
    - Prevents adjacencies on the interface
    - Prevents send/receive routing updates
- CONFIGURATION STEPS
  - Router(config-router)# passive-interface <interface> or
  - Router(config-router)# passive-interface default

#### VERIFICATION

- show ip ospf interface
- show ip protocols

# 3.8 Configure and verify VRF lite

- VRF A.K.A. "Virtual Routing and Forwarding"
  - Multiple routing tables
  - Separation maintained via VRF-to-Interface Allocation
  - Originally designed only for MPLS VPNs with BGP
  - By default, Routing Protocols are only active on the Global Routing Table

#### VRF-LITE INITIAL CONFIGURATION STEPS

- Router(config)# ip vrf Company-A
- Router(config-if)# ip vrf forwarding Company-A

#### GENERAL VERIFICATION

- show ip vrf <name>
- show ip vrf interfaces <name>
- **show ip interface** <type/name>
- show ip route vrf <name>
- show ip protocols vrf <name>

#### VRF-LITE STATIC ROUTE CONFIGURATION STEPS

Router(config)# ip route vrf <name> <subnet> <mask> {interface | next-hop}

#### VRF-LITE RIP CONFIGURATION STEPS

- Router(config)# router rip
- Router(config-router)# address-family ipv4 vrf Customer-A
- Router(config-router-af)# network 1.0.0.0

#### RIP VRF VERIFICATION

- show ip rip database vrf <name>

#### VRF-LITE EIGRP CONFIGURATION STEPS

- Router(config)# router eigrp 100
- Router(config-router)# address-family ipv4 vrf Customer-A autonomous-system 2
- Router(config-router-af)# network 1.1.1.0 0.0.0.255

#### EIGRP VRF VERIFICATION

- show ip eigrp vrf <name> topology
- show ip eigrp vrf <name> neighbors
  - show ip eigrp vrf <name> interfaces

#### VRF-LITE OSPF CONFIGURATION STEPS

- Router(config)# router ospf 1 vrf Customer-A
- Router(config-router)# network 1.1.1.0 0.0.0.255 area 1
- Router(config-router)# capability vrf-lite
- Router(config)# router ospf 2 vrf Customer-B ... Must use different IDs

#### OSPF VRF VERIFICATION

- show ip ospf 1 neighbors
- show ip ospf 1 database

# 3.9 Configure and verify filtering with any protocol

- ACL A.K.A. "Access Control Lists"

- General Match Steps:
  - **Examine-1**: Prefix
  - Examine-2: Wildcard Mask
- Description:
  - Most commonly used for filtering, such as packet filtering or route filtering
  - Default "deny any" at the end of an ACL

#### - IP PREFIX-LISTS

- General Match Steps:
  - **Examine-1:** Prefix and prefix-length (**x.x.x.x/yy**)
  - **Examine-2**: Range of prefixes or range of prefix lengths (**ge** and/or **le**)
- Description:
  - Processing is faster than normal ACLs
  - Default "deny any" at the end of a prefix-list
  - Permit Any: ip prefix-list <name> seq 10 permit 0.0.0.0/0 le 32

#### COMMAND FORMAT

- EIGRP Format:
  - distribute-list {ACL | prefix-list | route-map} {in | out} {interface}
- OSPF Format:
  - distribute-list {ACL | prefix-list | route-map} {in | out}
  - area {number} filter-list prefix {name} {in | out}
- BGP Format:
  - neighbor x.x.x.x distribute-list {ACL} {in | out}
  - neighbor x.x.x.x filter-list {AS-Path-ACL} {in | out}
  - neighbor x.x.x.x route-map {name} {in | out}
  - neighbor x.x.x.x prefix-list {name} {in | out}

## - EIGRP ROUTE FILTERING

- Distribution-List Filter Direction:
  - **In**: Prevent incoming updates from entering EIGRP topology table
  - Out: Prevent routes in routing table from being advertised to neighbors
- Extended ACL Details:
  - Match on prefix and neighbor sending the route

#### OSPF ROUTE FILTERING

- Distribution-List Filter Direction:
  - In: Prevent LSA from becoming a route in own local routing table
  - Out: ASBR; Prevents creation of protocols/routes into External LSAs
- Filter-List Filter Direction (ABR-only):
  - **In**: Filter prefixes being created and flooded *into* the configured area
  - Out: Filter prefixes coming out of the configured area
- Extended ACL Details:
  - Match on prefix and advertised router id of the LSA

#### BGP ROUTE FILTERING

- General Filter Direction:
  - In: Prevent updates from entering BGP table, applied against neighbor
  - Out: Prevent best BGP route from being advertised to neighbors
- Filter-List Filter Direction (beyond CCNP level):

- In: Prevents matches of the AS\_PATH attribute, applied against neighbor
- Out: Prevents matches of the AS\_PATH attribute from being advertised

# 3.10 Configure and verify redistribution between any protocols or routing sources

#### GENERAL CONCEPT

- Redistribution takes the routes from the **routing table** to redistribute
- Requires at least one working physical link within each routing domain
- Redistribute from the protocol into the specified routing domain

#### COMMAND FORMAT

EIGRP Format:

OSPF Format:

BGP Format:

## - COMMAND BREAKDOWN

- **protocol**: Source of the routing information
- **metric**: If default is not set, required information by the protocol
- process-id, as-number: Specify the AS/Process number
- match: If redistributing from OSPF, allows match based on type of OSPF route
- tag: Assigns unitless integer value, can later be matched via route-map
- route-map: Apply logic in the referenced route-map
- metric-type: Defines the external metric type (1 or 2) for the routes redistributed
- subnets: Redistributes subnets of classful networks and must be enabled

# 3.11 Configure and verify manual and auto-summarization with any routing protocol

## GENERAL CONCEPT

- Summarization is a technique used to reduce the size of a routing table
- The specific routes are condensed into a larger route

- Manual summarization involves specifying the exact subnet to summarize
- Auto-summarization occurs automatically and only on classful boundaries
- Auto-summary only supports **contiguous** networks
- Contiguous: A single classful network end to end
- **Discontiguous**: Multiple classful networks end to end
- Summarization can also be used for path manipulation (more specifics preferred)

#### COMMAND FORMAT

- EIGRP Format:
  - ip summary-address eigrp {asn} {prefix} {mask} (interface subcommand)
  - auto-summary (router subcommand)
- OSPF Format:
  - area {number} range {prefix} {mask} {cost cost} (ABR only)
  - summary-address {prefix} {mask} (ASBR only)
- BGP Format:
  - aggregate-address {prefix} {prefix-length} {summary-only}
  - Auto-summary

#### EIGRP SUMMARIZATION

- Summarization can be performed on any router
- With summarization, query scope is reduced
- Summary router adds route to routing table with outgoing interface of null0
- This is to prevent loops in case the summary route comes back around
- Summary routes by default are given an AD of 5!
- For auto summarization, network must be local and cannot be redistributed

#### OSPF SUMMARIZATION

- OSPF allows summarization at ABRs and ASBRs
- This is because ALL routers must have the same LSDB
- Summarization is done for LSAs, not routes
- By default on ABR, the cost is the best cost among the subordinate subnets

#### BGP SUMMARIZATION

- With auto-summary and no mask statement, the network statement logic...
  - The router adds a route for that classful network to the BGP table
  - The classful route is added if:
    - The exact classful route is in the routing table
    - Any subset routes of that classful network are in the routing table
  - The first occurs regardless of the **auto-summary** settings
  - The second occurs ONLY IF auto-summary is configured
- With aggregation...
  - The prefix/prefix-length is the summary route
  - The **summary-only** is used to not advertise the subordinate routes
  - You need to apply network/redistribute before being able to aggregate

# 3.12 Configure and verify policy-based routing

#### GENERAL CONCEPT

- Ingress packets typically routed via normal routing process

- PBR overrides the router's natural destination-based forwarding logic
- PBR feature is tied to use of Route-Maps
- Route-Map Purpose:
  - Define match criteria for PBR packets
  - Define forwarding action for these packets
- Route-Map Packet Forwarding:
  - Outgoing interface (should be point-to-point)
  - IP Next-Hop

#### - COMMAND FORMAT

- ip policy route-map {name} (interface subcommand)
- ip local policy route-map {name} (global command)

#### - ROUTE-MAP DETAILS

- Packets can be matched using a route-map via two criteria:
  - match ip address {EXTENDED-ACL | STANDARD-ACL}
  - match length {MIN-BYTES} {MAX-BYTES}
- There are four set command options:
  - set ip next-hop ip-address [... ip-address]
  - set ip default next-hop ip-address [... ip-address]
  - set interface interface [... interface]
  - set default interface interface [... interface]
- With **default** parameter, attempt to use the "default" routing table first, then PBR!
- Route-Map can also be used to set **IP precedence** or **ToS bits**:
  - set ip precedence {value}
  - set ip tos {value}

# 3.13 Identify suboptimal routing

- GENERAL CONCEPT
  - AD is used to determine best route which can result in suboptimal paths
  - This occurs during redistribution, and if not properly designed, can result in loops
  - EIGRP automatically prevents loops by setting AD of external routes higher (170)
  - Routing loops can also be prevented using tags and route-maps
- MODIFY ADMINISTRATIVE DISTANCE
  - **distance** distance ip-adv-router wc-mask [acl-number-or-name]

# 3.14 Explain ROUTE maps

- GENERAL CONCEPT
  - Route-maps provide if/then/else logic like in programming languages
  - A route-map can call an ACL or prefix-list
  - Route-maps also has the concept of using sequence numbers
  - A single route-map contains one or more route-map commands (entries)
  - To match all packets, omit the **match** command
- COMMAND FORMAT
  - route-map {name} {permit | deny} {seq-num}
  - match <match criteria>

# 3.15 Configure and verify loop prevention mechanisms

## 3.15.a Route tagging and filtering

## - GENERAL CONCEPT

- A route tag is a unitless 32-bit integer that most routing protocols can assign
- Assigned by a route-map set reference or the distribute-list/redistribute command
- By tagging redistributed routes, you can then prevent loops easily

#### LOOP PREVENTION EXAMPLE

- Router(config-router)# redistribute eigrp 100 tag 77 subnets
- Router(config)# route-map stop-loops deny 10
- Router(config-route-map)# match tag 77
- Router(config)# route-map stop-loops permit 20
- Router(config-router)# distribute-list route-map stop-loops in

## 3.15.b Split-horizon

#### GENERAL CONCEPT

- Distance vector protocols are susceptible to routing loops
- Prevents a route received on an interface to flood back out the same way
- In NBMA networks with hub and spoke topologies, this will cause issues

#### DISABLING SPLIT HORIZON

- no ip split-horizon {eigrp {AS}} (interface subcommand)

## 3.15.c Route poisoning

#### GENERAL CONCEPT

- Advertises a route with a "bad" metric
- This enables receiving router to remove the route from its routing table

#### - SPECIFIC DETAILS

- EIGRP sends an infinite metric
- OSPF sends an LSA age metric of 3600 seconds
- RIP sends a metric of 16 hops

# 3.16 Configure and verify RIPv2

#### CONFIGURATION STEPS

- Router(config)# router rip
- Router(config-router)# version 2
- Router(config-router)# network x.x.x.x

#### VERIFICATION

- show ip route
- show ip rip database

# 3.17 Describe RIPng

#### - **COMMONALITIES**

- Uses UDP
- Distance vector routing protocol
- AD of 120
- VLSM is supported
- Uses split horizon rule
- Uses poison reverse
- 30-second periodic full update
- Uses triggered updates
- Uses hop count metric
- Metric of 16 as infinity
- Supports route tags
- Multicast update destination
- Does not form neighborships

#### DIFFERENCES

- UDP port of 521 for RIPng, RIPv2 uses 520
- No auto-summarization for IPv6
- Destination address of FF02::9
- Link-Local Next-Hops
- IPv6 uses IPv6 AH/ESP authentication

#### - CONFIGURATION STEPS

- Router(config)# ipv6 unicast-routing
- Router(config)# ipv6 router rip name
- Router(config-if)# ipv6 rip name enable

#### VERIFICATION

- show ipv6 route
- Show ipv6 protocol
- show ipv6 rip next-hops

# 3.18 Describe EIGRP packet types

- EIGRP A.K.A. "Enhanced Interior Gateway Routing Protocol"
  - Uses DUAL A.K.A. "Diffusing Update Algorithm" to ...
    - Determine the best loop-free path
    - Determine the backup loop-free path
    - Provide FAST convergence
  - Uses bandwidth and delay for metric calculation...
    - $BW = (10^7/BW)*256$ , DELAY = (Delay In Microseconds) \* 256
    - Metric = (K1\*BW+((K2\*BW)/(256-load))+K3\*delay)\*(K5/(reliability+K4))
    - Real (Default) Metric = (Slowest\_BW + All\_Link\_Delays)
  - Uses hello packets to determine neighbors...
    - Used to dynamically discover new neighbors
    - Used to maintain the neighbor relationship
  - Uses three types of tables...

- Topology Table: Contains all successor and feasible successors
- Routing Table: Contains only the successors
- Neighbor Table: Contains neighborship information

#### SUCCESSOR AND FEASIBLE SUCCESSOR ROUTES

- **Feasible Distance**: Metric for a route to choose the best route for the prefix
- **Reported Distance**: Metric for a route the neighbor is *reporting* for the prefix
- The route with the smallest feasible distance is elected as primary route
- Successor Route: The primary route
- Feasible Successor Routes: The elected backup routes
- Feasibility Condition: Route's RD < FD of the successor
- The Feasibility Condition exists to prevent potential loops

#### - HELLO PACKETS

- By default, multicast address used is 224.0.0.10
- Unicast will be used for static neighbors, useful with NBMA topologies
- Neighborship Requirements:
  - Interface primary IP address must be in the same subnet
  - AS number must be the same
  - Connected interface must not be passive
  - Authentication must pass
  - K-values used must match

#### - Default Hello And Dead Timers:

- T1 or Slow Links: 60 seconds by default
- LAN or Fast Links: 5 seconds by default
- Modify: ip hello-interval eigrp {as} {interval} (interface subcommand)

#### UPDATE PACKETS

- When neighbors come up, the routers exchange full topology tables
- Once done, there is **NO** periodic reflooding of the information
- If something changes, only a partial update is sent about the affected prefix
- This change can be as small as a metric value change or about link failures
- Uses RTP A.K.A. "Reliable Transport Protocol" to send updates/ACK messages

#### ACK PACKETS

- Used during the topology exchange alongside update packets
- Generally ACK is sent in unicast by the neighbor

#### QUERY/REPLY PACKETS

- When a route has failed (active), a Query message is sent to all neighbors
- Query messages need an **exact** prefix/length for the Reply
- Queries are stopped on the routers that are one-hop from summarization
- If a neighbor receiving the Query message has a loop free route...
  - Reply message is sent in response
  - Otherwise, the Query is forwarded on to its neighbors

## - Once original router has received a Reply from all neighbors...

- If new loop-free routes were learned, install the best as successor
- If no routes are learned, the active prefix is removed from topology table
- Active Timer = 3 minutes

- Once 3 minute timer was up, the router will cut neighbor relationships
- This was an issue with neighbors sending late query replies
- By default, after 90 seconds, routers send SIA-Query to neighbor...
  - SIA A.K.A. "Stuck In Active"
  - Condition: Neighbor is still alive but waiting for a response to OWN query
  - The neighbor in SIA will send an SIA-Reply
  - If neighbor does not respond, active timer continues to decrement
- This timer can be configured using the following command:
  - timers active time {time} (router subcommand)
- You can limit Query Scope:
  - Stub Routers
  - Route Summarization

# 3.19 Configure and verify EIGRP neighbor relationship and authentication

#### NEIGHBOR CONFIGURATION EXAMPLE

- Router(config)# interface Ethernet1/1
- Router(config-if)# ip address 10.0.0.1 255.255.255.0
- Router(config)# router eigrp 100
- Router(config-router)# network 10.0.0.0 0.0.0.255

#### NEIGHBOR VERIFICATION

- show ip eigrp topology
- show ip eigrp neighbors
- show ip protocols

#### AUTHENTICATION

- Must pass for neighbor to establish
- Routers should use the same pre-shared key
- EIGRP only supports MD5 type authentication, not clear text
- These packets can still be intercepted by man-in-the-middle attacks

#### - AUTHENTICATION CONFIGURATION EXAMPLE

- Router(config)# key chain TEST
- Router(config-keychain)# key 1
- Router(config-keychain-key)# key-string CISCO
- Router(config)# interface gi0/1
- Router(config-if)# ip authentication mode eigrp 1 md5
- Router(config-if)# ip authentication key-chain eigrp 1 TEST

#### ADVANCED AUTHENTICATION OPTIONS

- Keychain can also be configured to use time-based logic
- Multiple keys can be configured for different time periods
- If there are multiple active/valid keys at the same time, lowest key number wins
- Clocks should match on both routers, use NTP to update clocks if required
- Clocks can also be set locally using the "clock set" command in exec mode

#### Time Commands:

- send-lifetime hh:mm:ss (in key configuration mode)
- accept-lifetime hh:mm:ss (in key configuration mode)

#### AUTHENTICATION VERIFICATION

- **show key chain** name
- show clock
- debug eigrp packets

# 3.20 Configure and verify EIGRP stubs

#### STUB ROUTER

- Should not forward traffic between two remote EIGRP-learned subnets
- Does not advertise EIGRP-learned routes from one neighbor to another
- Non-stub routers one-hop away from stub **do not** forward queries to stub
- By default, EIGRP stub routers only advertises:
  - Connected
  - Summary
- A stub router can be configured to advertise redistributed, static routes, or mix
- The "receive-only" parameter means the stub will not advertise **any** prefix

#### - CONFIGURATION STEPS

Router(config-router)# eigrp stub receive-only

#### VERIFICATION

- show ip protocols

# 3.21 Configure and verify EIGRP load balancing

# 3.21.a Equal cost

#### GENERAL CONCEPT

- By default, equal cost/metric routes are installed in the routing table
- To change the maximum number of equal cost routes that can be installed:
  - maximum-paths {number} (router subcommand)
- As for whether load balancing is per packet or per session, CEF decides
- Generally, per flow is enabled by default

#### VERIFICATION

- show ip route
- show ip eigrp topology

## 3.21.b Unequal cost

#### - GENERAL CONCEPT

- Variance: Allows EIGRP to consider additional paths for load sharing
- These paths only need to be similar to the successor route
- By default, variance is set to 1 which means only equal paths are load balanced
- The variance multiplier is an integer between 1 and 128
- **ONLY** feasible successor routes are considered in the unequal load balancing!
- This means that if a route did not meet feasibility condition, it is skipped
- Feasible successor considered if...

- (Variance)\*(Successor Route's FD) > Feasible Successors FD
- CONFIGURATION STEPS
  - variance {multiplier} (router subcommand)
- VERIFICATION
  - show ip protocols

# 3.22 Describe and optimize EIGRP metrics

- GENERAL CONCEPT
  - EIGRP metric can be manipulated to prefer one path over the other
  - There are two ways to manipulate the EIGRP metric:
    - Changing the bandwidth and/or delay values
    - Using offset-list
  - Changing bandwidth value is **NOT** recommended
    - Other technologies are affected such as QoS
    - On WAN subinterfaces, bandwidth value should be chosen carefully
    - Bandwidth on these interfaces should be sum of all CIR
  - Recommended to change the delay value instead
  - EIGRP distance calculation can be changed by configuring the k-values
- CONFIGURATION STEPS
  - Router(config-router)# metric weights {tos} k1 k2 k3 k4 k5
  - Router(config-if)# bandwidth {amount}
  - Router(config-if)# delay {amount}
- VERIFICATION
  - show ip protocols
- OFFSET-LIST CONCEPT
  - Offset-list simply allow a value to be added to the EIGRP metric
  - Behind the scenes, offset-list modifies the delay metric for the specified routes
  - The offset value is added to both the FD and the RD
  - Direction of the update message is also specified
- OFFSET-LIST CONFIGURATION STEPS
  - offset-list {ACL} {in | out} <offset value> {interface} (Router Subcommand)

# 3.23 Configure and verify EIGRP for IPv6

- COMMONALITIES
  - Uses Layer 3 header protocol type of 88
  - Uses successor, feasible successor logic
  - Uses DUAL
  - Uses triggered updates
  - Uses composite metric
  - Metric meaning infinity is 2^32 1
- DIFFERENCES
  - Uses neighbor's link local address as next-hop IP
  - Destination address is FF02::A

- Authentication relies on IPv6 built-in authentication and privacy features
- IPv6 has no concept of classful networks, so no auto-summary feature
- Does not require neighbors to be in the same IPv6 subnet to become neighbors

#### CONFIGURATION STEPS

- Router(config)# ipv6 unicast routing
- Router(config)# ipv6 router eigrp 100
- Router(config-router)# eigrp router-id 1.1.1.1
- Router(config-if)# ipv6 eigrp 100

#### VERIFICATION

- show ipv6 route
- show ipv6 protocols
- show ipv6 eigrp neighbors
- show ipv6 eigrp interfaces details
- show ipv6 eigrp topology [all-links]
- debug ipv6 eigrp notifications

# 3.24 Describe OSPF packet types

## OSPF A.K.A. "Open Shortest Path First"

- Uses link state logic, which can be broken into three branches:
  - Neighbor discovery
  - Topology database exchange
  - Route computation

## Uses concept of areas and a hierarchical design

- All areas must connect into the backbone (area 0)
- ABR connect non-backbone areas to the backbone area
- Detailed topology information not exchanged between areas

#### All internal routers must have the same image of the network

- Link State Databases (topology database) must be the same
- Shortest Path First (SPF) is run on the LSDB to find best path

#### Has two neighborship classes:

- 2-Way Neighbors: Not exchanged topology information
- Fully Adjacent Neighbors: Link state databases matches

#### - ROUTER-ID ASSIGNMENT PROCESS

- Same as EIGRP
- Manually configured router-id preferred over...
- Highest IP address of any up/up loopback interface preferred over...
- Highest IP address of any up/up non-loopback interface
- Not preemptive

#### SAME OSPF ROUTER-ID

- Same Router-ID in the same area:
  - The routers will generate a message saying there's a duplicate router-id
- Same Router-ID in a different area:
  - The routers will flush each other's LSAs and declare an OSPF Flood War

#### OSPF PACKETS/MESSAGES

- **Hello**: Used to form and maintain neighborship

- Database Description (DD): Used to exchange header information of LSAs
- Link-State Request (LSR): Used to request for missing LSAs
- Link-State Update (LSU): Used to send missing LSA information to neighbor
- Link-State Acknowledgement (LSAck): Used to ack received updates

#### - HELLO MESSAGES

- Messages are sent using the multicast address 224.0.0.5
- The following parameters must match:
  - Hello Interval
  - Dead Interval
  - Area ID
  - Subnet Mask
  - Stub Area Flag
  - Authentication

#### - Additional parameters that may be present:

- OSPF Router ID
- List of neighbors reachable on interface
- Router Priority
- Designated Router (DR) IP Address
- Backup DR IP Address
- Change Hello/Dead Intervals:
  - ip ospf hello-interface {value} (interface subcommand)
  - ip ospf dead-interval {value (interface subcommand)}
- Verify Intervals:
  - show ip ospf interface {interface}
- Subsecond Interval Modification:
  - Router(config-if)# ip ospf dead-interval minimal hello-multiplier {x}

#### OSPF FLOODING

- Routers advertise their local and learned LSAs to all OSPF neighbors
- Flooding prevents the looping of LSA as a side-effect of the DD exchange
- Re-flooding occurs every 30 minutes based on each LSA's age variable
- Router that creates the LSA sets this age to 0 seconds
- When a router needs to flush LSA from LSDB, Max Age is set to 3600 seconds

# 3.25 Configure and verify OSPF neighbor relationship and authentication

- DATABASE EXCHANGE WITHOUT DR
  - STEP-1 INIT
    - Router sends a hello packet on its OSPF enabled link
  - STEP-2 2-WAY
    - Router receives a Hello packet
    - Sees its own RID as having been seen by the neighbor
    - Also sees all parameters passing
  - STEP-3 EXSTART

- First Database Descriptor (DD) received
- DD gives high level overview of the LSAs
- Also, an election is held for master and slave router
- The master is in charge of the initial sequence number of DD
- The router with the highest RID is elected as Master
- DD packets are sent using unicast

#### STEP-4 EXCHANGE

- Occurs after master/slave election
- Neighbors continue to multicast DD packets to each other
- This continues until they all have the same LSIDs known in the area

#### - STEP-5 LOADING

- The routers have the same view of all LSIDs
- For any missing LSA, router sends a Link State Request (LSR)
- The router listening to LSR sends a Link State Update (LSU)
- Routers sends LSAcks or implicit ack by returning same LSA

#### - STEP-6\_FULL

- All LSAs have been sent, received, and acknowledged
- Database is fully populated
- Routers run SPF to calculate the best paths for each subnet

## Additional neighbor states (8 Total)

- Down: No neighborship with that router
- Attempt: Layer 2 problem with statically defined neighbors

#### - DATABASE EXCHANGE WITH DR

- Only slight differences
- All the other routers exchange their databases with DR/BDR only
- Communication to DR/BR occurs over 224.0.0.6
- Communication from DR/BDR still occurs over 224.0.0.5

#### NEIGHBOR CONFIGURATION DETAILS:

## - To advertise any particular subnet or enable any interface:

- network {subnet} {wildcard-mask} area {area-number} (router subcommand)
- ip ospf {process-id} area {area-number} (interface subcommand)

#### To change the OSPF router-id:

router x.x.x.x (router subcommand)

#### - NEIGHBOR VERIFICATION

- show ip ospf neighbor
- show ip ospf
- show ip ospf database
- show ip protocols

#### OSPF AUTHENTICATION CONCEPT

## - Configuring authentication is a two-step process:

- Authentication and authentication type must be enabled and selected
- Authentication key must be configured per interface

#### - This can be done in one of two ways:

- Enable at the interface
- Enable on all interfaces in an area by changing the area wide settings

- Any authentication on an interface will override area wide settings
- Authentication Key
  - Cannot be configured area wide
  - Must be configured individually on each interface
- Three types of authentication:
  - **Type-0**: No Authentication
  - **Type-1**: Clear Text Authentication
  - Type-2: MD5 Authentication
- Number of keys
  - Multiple keys can be configured on the same interface
  - OSPF does not support the use of a keychain
  - This can only be done if MD5 authentication is enabled
  - Useful for migration of keys
- AUTHENTICATION CONFIGURATION DETAILS:
  - Enable Authentication Area Wide Settings (router subcommand)
    - area {area-number} authentication (type 1)
    - area {area-number} authentication message-digest (type 2)
  - Enable Authentication Interface Settings (interface subcommand)
    - **ip ospf authentication null** (type 0)
    - ip ospf authentication (type 1)
    - ip ospf authentication message-digest (type 2)
  - Configure Key (interface subcommand)
    - ip ospf authentication-key {key-value} (type 1)
    - ip ospf message-digest-key {key-number} md5 {key-value} (type 2)
- AUTHENTICATION VERIFICATION
  - show ip ospf interface {interface}
  - debug ip ospf hello
  - debug ip ospf adj

# 3.26 Configure and verify network types, area types, and router types

# 3.26.a Point-to-point, multipoint, broadcast, nonbroadcast

Network Type	DR/BDR?	Hello Unicast/Multicast?	Hello/Dead Intervals?
Point-to-Point	No	Multicast	10/40
Point-to-Multipoint	No	Multicast	30/120
Point-to-Multipoint Non-Broadcast	No	Unicast	30/120
Broadcast	Yes	Multicast	10/40
Non-Broadcast	Yes	Unicast	30/120

Loopback	No	N/A	N/A

#### - OSPF NETWORK TYPES

- Determines operational characteristics of OSPF on that interface
  - Whether the router will use multicast to discover neighbors
  - If at least two routers can exist in the subnet attached to interface
  - Whether the router should attempt to elect an OSPF DR on that interface
- Determined automatically by Layer-2 encapsulation on the interface

#### NETWORK TYPE BROADCAST

- Broadcast is default on LAN interfaces:
  - Discovers neighbors dynamically
  - Supports use of DR/BDR
  - Intervals: 10 & 40
- Configuration Details:
  - ip ospf network broadcast (interface subcommand)
- NETWORK TYPE NON-BROADCAST
  - Default on Frame Relay / NBMA physical or multipoint interfaces:
    - Does NOT discover neighbors dynamically
    - Intervals: 30 & 120
  - Statically Configure Neighbors:
    - neighbor ip-address {priority priority} (router subcommand)
  - Configuration Details:
    - ip ospf network non-broadcast (interface subcommand)
  - If network is not full mesh, make sure the DR/BDR can reach everyone
- NETWORK TYPE POINT-TO-POINT
  - Default on any point-to-point interface:
    - Does not elect DR/BDR
    - Discovers neighbors dynamically
    - Intervals: 10 & 40
  - Configuration Details:
    - ip ospf network point-to-point (interface subcommand)
- NETWORK TYPE POINT-TO-MULTIPOINT
  - Not default on any kind of interface:
    - Does not elect DR/BDR
    - Discovers neighbors dynamically
    - Intervals: 30 & 120
  - How does this network type help with partial mesh topologies?
    - Regardless of actual mask, router advertises /32 LSAs for connectivity
    - LSAs received on P-2-MP subinterface allowed to be flooded back out
    - Routers without direct PVC connectivity can still achieve L3 connectivity
  - Configuration Details:
    - ip ospf network point-to-multipoint (interface subcommand)
- NETWORK TYPE POINT-TO-MULTIPOINT NON-BROADCAST
  - Not default on any kind of interface:

- Similar to point-to-multipoint
- Neighbors will not be discovered dynamically
- Not described in RFC 2328

# Advantages over other types?

- Cost to reach each neighbor can be changed per neighbor
- By default, changing cost on an interface changes cost for all neighbors

# - Configuration Details:

ip ospf network point-to-multipoint non-broadcast (interface subcommand)

## HOW TO REMEMBER IT ALL

- If the network type starts with "point" it does not use DR/BDR
- If network type has "non-broadcast", it does not do dynamic neighbors
- Only broadcast and point-to-point use faster times

3.26.b LSA types, area type: backbone, normal, transit, stub, NSSA, totally stub

# - LSA GENERAL CONCEPT

- Each router stores data, composed of individual link-state advertisements
- Each router in an area will have the same link-state database information
- Each router, individually, runs the Shortest Path First (SPF) algorithm
- Each router considers itself to be the root of the tree
- SPF process must examine the individual LSAs and see how they fit together

## LSA TYPE OVERVIEW

- **Type-1**: Router LSA
- Type-2: Network LSA
- **Type-3**: Network Summary LSA
- Type-4: ASBR Summary LSA
- **Type-5**: AS External LSA
- Type-7: NSSA External LSA

## - TYPE-1 LSA: ROUTER LSA

- Each router creates a Type-1 LSA for itself and floods it in the area
- The LSA identifies an OSPF router based on its OSPF router-id
- What does it accomplish? Helps to build SPF tree
- OSPF identifies a Type-1 using a 32-bit link-state identifier (LSID)
- Each router uses its own OSPF router-id as the LSID
- ABRs create multiple Type-1 LSAs for themselves, one per area
- This LSA also includes information about its attached links.
  - No neighbors: Subnet advertised; link to 'a stub network'
  - With DR: The IP of the DR; link to 'a transit network'
  - No DR: It lists the neighbor's RID; link to 'another router (point-to-point)'

#### Router LSA Verification:

show ip ospf database router x.x.x.x

### TYPE-2 LSA: NETWORK LSA

- Generated on multi-access networks
- Depends on the existence of a DR

- Required to map all connected routers on multi-access network, like LAN
- Type-2 LSA also lists the RIDs of the OSPF neighbors connected to the interface
- Network type tells the router's interface whether a DR should be elected
- Routers uses the following election rules to elect a DR:
  - Choose a router with the highest priority value (default 1, max 255)
  - If priority is tied, choose the router with highest RID
  - Choose a BDR, based on next-best priority, if tied, next-best highest RID
- Change Priority:
  - ip ospf priority value (interface subcommand)
- When a DR fails, the current BDR becomes the new DR
- Election is held only for the BDR, not the DR
- Network LSA Verification:
  - show ip ospf database network x.x.x.x
- TYPE-3 LSA: NETWORK SUMMARY LSA
  - OSPF uses concept of areas to reduce memory and compute resources
    - Type-1 and 2 LSAs are not advertised across areas
    - This saves CPU and convergence time
  - ABR A.K.A. "Area Border Routers"
    - Connect different OSPF areas together
    - Generate Type-3 LSAs for each subnet in an area
  - Type-3 LSA does not contain all the detailed topology information
    - Consists of each subnet and cost to reach that subnet
    - Summarizes the information from Type-1 and Type-2 LSAs
  - Type-3 LSAs appear to be another subnet connected to the ABR to other routers
  - ABR assigns an LSID of the subnet being advertised
  - ABR also adds its own RID in the LSA as there could be multiple ABRs
  - Summary LSA Verification:
    - show ip ospf database summary x.x.x.x

#### - LIMITING NUMBER OF LSA

- By default, there is no upper limit for the number of LSAs that a router can learn
- As the LSDB grows, memory and convergence becomes an issue
- To limit the number of LSAs, use the following command:
  - max-lsa number (router subcommand)
- Not recommended, as once LSA count goes over limit all relationships die
- TYPE-4 LSA: ASBR SUMMARY LSA
  - ASBR A.K.A. "Autonomous System Border Router"
  - How do the other areas know how to get to the ASBR?
  - ABR generates a Type-4 LSA
  - This tells the area routers that to get to the ASBR, to go through the ABR
  - "I am your ABR and you can reach the ASBR through ME!"
  - A new Type-4 LSA must be generated by the ABR for every new area

# - TYPE-5: AS EXTERNAL LSA

- Generated by ASBRs
- Type-5 LSA floods every area, no other LSA does this
- Type-4 provides reachability to ASBR for the Type-5 external routes

Please see the filtering section 3.9 for further details

## TYPE-7: NSSA EXTERNAL LSA

- Generally, external routes cannot be injected to stub areas
- This means routers within these areas cannot do redistribution
- Not-so-stubby areas (NSSA) solves this problem with Type-7 LSA
- Routers in NSSA areas can redistribute routes as Type-7 LSA
- Type-5 LSAs are still not allowed
- Type-7 converts to Type-5 LSA once it reaches the ABR

## - **GENERAL AREA NOTES**

- Backbone Area
  - All areas must connect into the backbone
  - ABRs connect non-backbone areas to the backbone
  - Virtual-links can get around this limit, though not recommended design

#### Normal Area

- An area that operates independently
- Able to receive Type-3 and Type-5 LSAs from other areas

#### - Transit Area

- An area that connects the backbone area to another area
- Used for virtual-link setups
- This area MUST NOT be stub

# - STUBBY AREA CONCEPT

- Four types of OSPF stubby areas:
  - Stub
  - Totally Stubby
  - Not-So-Stubby Areas (NSSA)
  - Totally NSSA

# - Differences in types of Stubby Areas

- For all types of stubby areas, ABR always filters Type-5 LSAs
- For totally stubby and totally NSSA areas, ABR also filters Type-3 LSAs
- For stubby and NSSA, ABRs do not filter Type- LSAs

# - CONFIGURATION STEPS FOR STUBBY AREAS

- To configure an area as a stub area, use the following command:
  - area area-id stub (router subcommand)
- All routers in area should be configured the same way!
- The metric for the default route injected by ABRs can be changed from default 1:
  - area area-id default-metric metric (router subcommand)
- To configure an area as totally stubby, use the following command:
  - area area-id stub no-summary (router subcommand)
- The no-summary option is only required on the ABR

# CONFIGURATION STEPS FOR NSSA AREAS

- ABRs do not advertise, by default, a default-route into NSSA areas:
  - area area-id nssa default-information-originate (router subcommand)
- To configure an area as NSSA, use the following command:
  - area area-id nssa (router subcommand)
- To configure an area as totally NSSA, use the following command:

- area area-id nssa no-summary (router subcommand)
- The no-summary option is only required on the ABR

#### STUB VERIFICATIONS

- show ip ospf
- show ip ospf database
- show ip ospf database database-summary

# 3.26.c Internal router, backbone router, ABR, ASBR

#### GENERAL NOTES

- Internal Router: All routers within a single area, not including ABR
- Backbone Router: All routers within the backbone, including ABRs
- ABR: Router that connects a normal area to the backbone area
- ASBR: Router that connects an external autonomous system to the domain

# 3.26.d Virtual link

#### - GENERAL CONCEPT

- OSPF requires all non-backbone areas connect to the backbone area
- With some designs, this is not possible!
- This introduces a reachability problem
- Virtual Link allows two ABRs to form a transit area
  - This allows an area to connect through another non-backbone area
  - This makes the disjointed area believe it is connected to area 0 directly
- ABRs connected over virtual links send all OSPF messages as unicast
- ABRs connected over virtual links also mark the Do Not Age (DNA) bit in LSAs
- The transit area **MUST NOT** be a stub area

# - CONFIGURING VIRTUAL LINKS WITHOUT AUTHENTICATION

- To configure a virtual link between ABRs, use the following command:
  - area area-number virtual-link remote-RID (router subcommand)
- The area-number must refer to the transit area number
- The remote-RID can be verified from the local router using the following:
  - show ip ospf border-routers
- To verify OSPF virtual link:
  - show ip ospf virtual-links

# CONFIGURING VIRTUAL LINKS WITH AUTHENTICATION

- Authentication key is not defined on the interface level but otherwise same
- To configure authentication on virtual links, use the following:
  - area area-number virtual-link remote-RID authentication...
    - null OR
    - authentication-key key-value OR
    - message-digest message-digest-key key-number md5 key-value

# 3.27 Configure and verify OSPF path preference

## GENERAL CONCEPT

Multiple ways to do path manipulation:

- By changing metric
- By using summary route (Section 3.11)
- By filtering (Section 3.9)
- By changing metric-type (Section 3.10 can also be done with default)
- Within an OSPF area, options are limited; the only way is by changing metric
- Metric can be changed by...
  - Changing the auto-cost reference bandwidth
  - Changing the bandwidth of an interface
  - Changing OSPF cost directly at the interface level

#### PATH MANIPULATION CONCEPT

- With multiple areas, common network designs have at least two ABRs
- Ideally, both ABRs advertise the same subnets from one area to another
- To prefer one ABR over another, we can use summarization or filtering
- Filtering is not recommended as there is no high availability
- OSPF Route Hierarchy:
  - O > O IA > E1 > E2 > N1 > N2

## - METRIC TUNING CONFIGURATION STEPS

- OSPF calculates the cost for an interface based on the following formula:
  - reference-bandwidth / interface-bandwidth
- Default reference-bandwidth is 100 Mbps
- The reference bandwidth can be changed by using:
  - auto-cost reference-bandwidth bandwidth (router subcommand)
- This setting is local to the router, but it should be changed on all area routers
- Bandwidth can be changed directly with the following command:
  - bandwidth value (interface subcommand)
- Changing bandwidth is not recommended as it will affect other dependencies
- Cost can also be directly changed on the interface:
  - ip ospf cost value (interface subcommand)
- To verify OSPF cost on any interface:
  - show ip ospf interface brief
  - show ip ospf interface

# 3.28 Configure and verify OSPF operations

# CONFIGURATION STEPS

- STEP-1: Start up OSPF process and define router-id
  - Router(config)# router ospf 1
  - Router(config-router)# router-id 1.1.1.1
- **STEP-2**: Place interfaces in appropriate areas
  - Router(config-router)# network 10.0.0.0 0.0.0.255 area 0 OR
  - Router(config-if)# **ip ospf** 1 **area** 0
- **STEP-3**: Configure any miscellaneous parameters
  - Make sure every interface is in right network type
  - Make sure to enable any edge scenarios such as virtual link
  - See the previous sections for any actual details

# - VERIFICATION

- show ip ospf neighbor
- show ip ospf database
- show ip ospf interface brief
- show ip protocols
- show ip route

# 3.29 Configure and verify OSPF for IPv6

# - COMMONALITIES

- Use IP protocol type 89
- Use link state logic
- Supports VLSM
- LSA flooding and aging concepts
- Area structure concept
- Packet types
- 32-bit LSID
- Uses interface cost metric
- Age = 3600 as the infinity metric
- DR/BDR election process
- Periodic reflooding every 30 minutes

#### - DIFFERENCES

- Multicast address of FF02::5, FF02::6
- OSPFv3 uses IPv6 AH/ESP authentication
- Neighbor checks are the same except that "same subnet" is not checked for IPv6
- IPv6 OSPF supports multiple instances basically VRFs (not supported in IOS)

### - LSA TYPES

- The first two LSAs are only responsible for building SPF trees
- Now, Type-9 contains the actual subnet information
- Type-9 makes SPF recalculation not so disruptive anymore
- New Type-8 link LSA goes to next-hop neighbors to provide link local info
- The DR can send that info to connected neighbors and the area with Type-9

LSA Name	LS Type Code	Flooding Scope	LSA Function Code
Router LSA	0x2001	Area Scope	1
Network LSA	0x2002	Area Scope	2
Inter-Area-Prefix-LSA	0x2003	Area Scope	3
Inter-Area-Router-LSA	0x2004	Area Scope	4
AS-External-LSA	0x4005	AS Scope	5
Group-membership-LSA	0x2006	Area Scope	6
Type-7-LSA	0x2007	Area Scope	7

Link-LSA	0x0008	Link-local scope	8
Intra-Area-Prefix-LSA	0x2009	Area scope	9

# - CONFIGURATION STEPS

- General Configuration:
  - Router(config)# ipv6 unicast-routing
  - Router(config)# ipv6 router ospf cprocess-id>
  - Router(config-if)# ipv6 ospf cprocess-id> area <#>
- The concept and commands related to OSPF stub areas are identical
- Summarize on ABR and ASBRs using the following:
  - area <#> range prefix/length (router subcommand)
- Change Network Types:
  - ipv6 ospf network type (interface subcommand)
- VERIFICATION
  - show ipv6 route
  - show ipv6 ospf
  - show ipv6 protocols
  - show ipv6 ospf neighbors
  - show ipv6 ospf interfaces brief
  - show ipv6 ospf database

# 3.30 Describe, configure, and verify BGP peer relationships and authentication

# - COMPARISON BETWEEN IGP AND BGP

- BGP needs to form neighborship like IGPs
- BGP needs to advertise prefixes, just like IGPs
- BGP also advertises Next Hops for those prefixes
- Neighbor IP address may not be a common subnet for BGP
- BGP uses TCP while IGP uses multicast
- IGPs emphasize fast convergence for efficient routes, BGP scalability
- BGP uses path vector logic which is similar to distance vector
- NLRI A.K.A. "Network Layer Reachability Information"
  - Advertised prefix and length
  - Known as NLRI as it can transport beyond just IPv4 addressing
  - BGP advertises different path attributes for path making decisions

# GENERAL BGP CONCEPT

- Two types of neighbors in BGP: internal and external
  - Neighborship behavior is different between the two types
  - As everything is unicast, you can have neighbors not directly connected
- There are two kinds of AS numbers: public and private
  - Public AS numbers can be advertised over the Internet
  - Private AS numbers should not be advertised over the Internet
- eBGP NEIGHBORSHIP OVERVIEW
  - BGP must complete three steps to get best routes:

- Form neighborship
- Exchange topology information
- Run a best-path algorithm
- BGP forms neighborship using TCP port 179 (unicast)
- eBGP neighbors are assumed to be directly connected

# eBGP Neighborship Requirements

- Following requirements must be met:
  - The ASN number must match (remote-as) for the local/remote router
  - Peer must be reachable via an IGP route
  - BGP router-ids must not be the same for the two routers
  - If configured, MD5 authentication must pass
  - Each router must complete a TCP connection with the BGP peer
  - The source IP address used to reach that peer must match

# The BGP router-id is elected as follows:

- Manually configure the router-id
- Choose highest IP address of any up/up loopback interface
- Choose the highest IP address of any up/up non-loopback interface

# BGP UPDATE-SOURCE

- The local router finds the outgoing interface to be used to reach IP address
- This default behavior can be changed by explicitly specifying the source address
- Useful with two links or redundant Layer 3 paths between the two routers

# - iBGP Neighborship Requirements

- Neighborship requirements are the same as eBGP except for the asn value
- The asn value should be the same AS number as the local AS number
- The exception here is the TTL value for iBGP neighbors is 255 by default
- iBGP also defaults to multihop
- Full mesh is preferred for iBGP despite being able to communicate multihop
  - The routers in between will not have routes going back
  - When the destination and source want to talk, they will be blackholed!

# - **BGP MULTIPATH**

- BGP supports the maximum-paths number-of-paths router subcommand
- The logic is significantly different to IGP multipaths
- BGP uses the best path algorithm to pick only one end route as the best path
- These factors must be considered before load balancing can occur:
  - If the attributes tie from the 'N WLLA OMNI' formula
  - If routes are iBGP learned, next-hop IP address should be different
  - If routes are eBGP, should be learned from the same neighboring AS

# 3.30.a Peer group

# PEER GROUP CONCEPT

- A BGP router may have dozens of BGP neighbors
- Many neighbors may require the same BGP policies
- Configuring per-neighbor policies are burdensome

- Takes more time to configure
- Requires more of CPU when sending/receiving BGP updates

# Peer Groups are groups of peers which the same outbound policies apply

- Only inbound policies can be overridden on a per-neighbor basis
- All configuration options for neighbors can be applied to peer groups
- Updates generated once per group

# - INTERNAL PEER GROUP CONFIGURATION EXAMPLE

- Router(config)# router bgp 300
- Router(config-router)# neighbor my-company peer-group
- Router(config-router)# neighbor my-company remote-as 300
- Router(config-router)# neighbor my-company route-map Attribute out
- Router(config-router)# neighbor my-company filter-list 2 out
- Router(config-router)# neighbor 2.2.2.2 peer-group my-company
- Router(config-router)# neighbor 3.3.3.3 peer-group my-company
- Router(config-router)# neighbor 3.3.3.3 filter-list 4 in

#### EXTERNAL PEER GROUP CONFIGURATION EXAMPLE

- Router(config)# router bgp 300
- Router(config-router)# neighbor Externals peer-group
- Router(config-router)# neighbor Externals route-map set-metric out
- Router(config-router)# neighbor 7.7.7.7 remote-as 150
- Router(config-router)# neighbor 7.7.7.7 peer-group Externals
- Router(config-router)# neighbor 9.9.9.9 remote-as 250
- Router(config-router)# neighbor 9.9.9.9 peer-group Externals

# 3.30.b Active, passive

#### - GENERAL CONCEPT

- Neighbors with the lowest BGP router-id will establish the connection
- The port used will be destination TCP port 179 and the source port random
- This behavior can be modified with a few commands.

# - CONFIGURATION EXAMPLE

- Router(config)# router bgp 500
- Router(config-router)# neighbor 10.0.0.1 transport connection-mode passive

#### - BEFORE AND AFTER SNAPSHOT

#### BEFORE:

```
    Router# show ip bgp neighbors | in host
Local host: 10.0.0.2, Local port: 57717
    Foreign host: 10.0.0.1, Foreign port: 179
```

#### - AFTER:

 Router# show ip bgp neighbors | in host Local host: 10.0.0.2, Local port: 179
 Foreign host: 10.0.0.1, Foreign port: 34121

# 3.30.c States and timers

# - OVERALL PROCESS

- A router tries to establish a TCP connection using destination port 179
- When the TCP connection completes, the router sends BGP Open message
- The BGP Open message operates similarly to Hello messages

- Once all the parameters must the routers become neighbors

# - BGP NEIGHBOR STATES

- Idle
  - The BGP process is administratively down or...
  - The BGP process is awaiting next retry attempt
- Connect
  - The BGP process is waiting for the TCP connection to be completed
- Active
  - TCP connection failed, Connect-Retry timer running
  - Listening for incoming TCP connection
- Opensent
  - TCP connection exists, BGP Open message has been sent to the peer
  - The matching Open message has to yet been received from the router
- Openconfirm
  - An Open message has been sent and received by both sides
- Established
  - All neighbor parameters match and the neighbor relationship works
  - The peers can now exchange Update messages

## - BGP MESSAGE TYPES

- Open
  - Used in neighbor establishment
  - BGP values and capabilities are exchanged
- Update
  - Informs neighbors about withdrawn, exchanged, and new routes
  - Used to exchange PAs and associated NLRI that use those attributes
- Notification
  - Used to signal a BGP error
  - Typically results in a reset to the neighbor relationship
- Keepalive
  - Sent on a periodic basis to maintain the neighbor relationship
  - The lack of receipt within the Hold timer causes down neighbor

# 3.31 Configure and verify eBGP (IPv4 and IPv6 address families)

# 3.31.a eBGP

#### eBGP BASIC CONFIGURATION STEPS

- Router(config)# router bgp asn-number
- Router(config-router)# bgp router-id rid
- Router(config-router)# neighbor ip-address remote-as remote-asn
- Router(config-router)# neighbor ip-address password key
- Please note that the MD5 key is in the TCP header of the packet!

# BGP UPDATE-SOURCE CONFIGURATION STEPS

- Router(config)# interface loopback0
- Router(config-if) ip address loopback-address

- Router(config)# router bgp asn-number
- Router(config-router)# neighbor neighbor-loopback remote-as
- Router(config-router)# neighbor update-source loopback-address
- Router(config-router)# neighbor ebgp-multihop hops

#### iBGP BASIC CONFIGURATION STEPS

- Router(config)# router bgp asn
- Router(config-router)# neighbor neighbor-ip remote-as asn
- Router(config-router)# neighbor neighbor-ip update-source loopback-address
- Router(config-router)# neighbor neighbor-ip password key

#### VERIFICATION

- **show ip bgp prefix** [subnet-mask]
- **show ip bgp neighbors** ip-address **received-routes**
- **show ip bgp neighbors** ip-address **routes**
- show ip bgp neighbors ip-address advertised-routes
- show ip bgp summary

# 3.31.b 4-byte AS number

# - GENERAL CONCEPT

- 4-byte ASNs provide 4,294,967,296 autonomous system numbers
- The original 2-byte ASNs only provide 65536
- Mappable ASNs are legacy ASN values that are mapped in 4-byte
- **ASPLAIN** is a simple decimal representation of the ASN
- **ASDOT**+ breaks into ASN into 16-bit values separated by a dot

#### INTEROPERABILITY

- Legacy peer responds it does not support 4-byte ASN capability or...
- Legacy peer does not respond at all with the capability
- New BGP peer can then use ASN 23456 which is a transit value (AS\_TRANS)
- Interoperable peering is achieved as new peer adapts to the legacy peer

# 3.31.c Private AS

#### - GENERAL CONCEPT

- The range for public and private AS numbers are:
  - Public AS numbers 1 64495
  - Private As numbers 65512 65534
- The following numbers and ranges are reserved:
  - 0, 64496 65511, 65535

# 3.32 Explain BGP attributes and best-path selection

### OVERVIEW OF BGP BEST PATH ALGORITHM

- The best path algorithm takes the following initial steps to determine best path:
  - N: Next-Hop, it should be reachable
  - W: Weight, bigger value is preferred
  - L: Local Preference, bigger is preferred
  - L: Locally Originated Routes

- **A**: AS\_PATH length, smaller length is preferred
- **O**: Origin, Prefer i> e > ?
- M: MED, smaller is preferred
- N: Neighbor Type, prefer eBGP over iBGP
- I: IGP cost to next hop, smaller value is preferred
- When BGP is unable to find best path, the following three are used as tiebreaker:
  - Oldest (longest known) eBGP route
  - Lowest neighbor BGP rid
  - Lowest neighbor IP address
- BGP Path Attribute Classifications
  - Mandatory (Next-Hop, AS\_PATH, Origin Code)
  - Optional Transitive Can leave local AS)
  - Optional Non-Transitive Not allowed to leave local AS (Local Pref)

## NEXT-HOP ATTRIBUTE

- With eBGP, the next-hop IP address is changed by the advertising router
- With iBGP, the next-hop IP address is not changed, but this is configurable
- iBGP peers must have IGP reachability to next-hops. If not:
  - iBGP-learned routes will not be installed in IP Routing Table
  - iBGP-learned routes will not be advertised to any other BGP peers
- The most common ways to resolve this issue are:
  - Advertise links to eBGP peers to the internal network
  - Use the **neighbor** x.x.x.x **next-hop-self** command
- Generally, to prevent potential black holes, iBGP should be full mesh
- A route reflector system can also be used for scalability
- **Synchronization** can also resolve this issue, but not recommended
  - Only BGP routes with IGP routes are preferred
  - This inflates the IGP routing table tremendously
  - Enable with synchronization (router subcommand)

# - WEIGHT ATTRIBUTE

- Not generally considered a path attribute; it is a Cisco proprietary feature
- It is locally significant to a single router and is not advertised to any neighbors
- It can be used to influence the choice of outbound routes
- It is set on inbound routes
- A bigger value is preferred for it
- The default value of weight is 0 for learned routes and 32768 for locally injected
- These default values cannot be modified
- Two Configuration Options:
  - neighbor x.x.x.x route-map name (use set weight)
  - neighbor x.x.x.x weight value (affects all prefixes)

## LOCAL PREFERENCE ATTRIBUTE

- Purpose is to identify the best exit point from the AS to reach a given prefix
- This attribute is advertised to all iBGP neighbors but not to eBGP
- Higher value is preferred
- Default value is 100

- Changing Default Value:
  - default local-preference <0-4294967295> (router subcommand)
- Changes to LOCAL PREF are applied to an eBGP neighbor using:
  - neighbor x.x.x.x route-map name in (router subcommand)
- Within the route-map, the **set local-preference** command is used

## LOCALLY ORIGINATED ROUTES ATTRIBUTE

- A route that was injected locally by the router in question
- If you see a next-hop of 0.0.0.0, that means a locally originated route
- Inject Route Locally:
  - network x.x.x.x mask y.y.y.y (router subcommand)

## - AS PATH ATTRIBUTE

- By default, if not BGP PA have been explicitly set, AS\_PATH is used
- AS\_SEQ is a component of the AS\_PATH attribute
- Lists the ASNs that would be part of an end-to-end path
- BGP uses the AS\_PATH to perform two functions:
  - Choose best route for a prefix on the shortest AS\_PATH
  - Prevent potential routing loops
- When a route is advertised by BGP, the AS number is prepended in AS\_PATH
- When a router receives an update that lists its own ASN, router ignores that route

# - AS\_PATH PREPENDING

- With AS\_PATH prepending, the length of AS\_PATH PA can be made longer
- As mentioned previously, longer AS\_PATH is less preferred
- Ideally, local AS number is prepended, otherwise you risk other routers dropping
- Configured using **set as-path prepend** under a route-map entry

#### - ORIGIN ATTRIBUTE

- When BGP creates a route, one of the mandatory attributes is the Origin Code
- Routes originated by:
  - network command = igp
  - redistributed command = unknown (?)
- Order of preference:
  - igp > egp > ?
- The symbols will be located in show ip bgp command to the far right

# - MULTI-EXIT DISCRIMINATOR (MED)

- Optional non-transitive attribute
- Can be modified to influence the inbound traffic coming from the neighboring AS
- It is not advertised beyond the neighboring AS so it only goes one AS deep!
- A smaller value is preferred
- Default value is 0
- The MED value can be changed using the command set metric under route-map
- The direction will always be outbound
- If the same route came from two different external AS, MED cannot be compared
- You can get the router to compare the MED values if you want:
  - **bgp always-compare-med** (router subcommand)

# 4.0 VPN Technologies

# 4.1 Configure and verify GRE

- GRE A.K.A. "Generic Route Encapsulation"
  - Tunneling mechanism, can contain many different protocols in header
  - Default encapsulation for Cisco IOS Tunnel Interfaces
- CONFIGURATION STEPS
  - Router(config)# interface tunnel 0
  - Router(config-if)# ip address 10.1.3.1 255.255.255.0
  - Router(config-if)# tunnel source Loopback 0
  - Router(config-if)# tunnel destination 3.3.3.2
- VERIFICATION
  - show ip interface brief

# 4.2 Describe DMVPN (single hub)

- DMVPN A.K.A. "Dynamic Multipoint Virtual Private Network"
  - Point-to-multipoint Layer 3 Overlay VPN
  - Logical Hub and Spoke topology
  - Direct spoke to spoke traffic supported
- DMVPN REQUIREMENTS
  - Hub Router, reachable via static, public IP address
  - Spoke Routers, reachable via static, or dynamic, public IP address
  - Multipoint GRE Tunnels
  - IPSec, optional, but typically used
- MULTIPOINT GRE CONFIGURATION (mGRE):
  - Router(config)# interface tunnel 0
  - Router(config-if)# ip address 10.1.3.1 255.255.255.0
  - Router(config-if)# tunnel source Loopback0
  - Router(config-if)# tunnel key 1
  - Router(config-if)# tunnel mode gre multipoint
- NHRP A.K.A. "Next Hop Resolution Protocol"
  - Used for a tunnel to answer the following questions:
    - What is the private tunnel IP address of my peer?
    - What is the public IP address of my peer?
  - Maps a peer's tunnel IP address to peer's public IP address
  - Can populate NHRP cache via static or dynamic entries
- NHS A.K.A. "Next Hop Server" (Hub)
  - Statically defined in spoke
  - Example: NHRP NHS = 10.1.3.2, NHRP Cache: 10.1.3.2 -> 3.3.3.2 (static)
- NHRP MESSAGE TYPES
  - NHRP Registration Request
    - Spokes register their NBMA and VPN IP to NHS
    - Required to build Spoke-to-Hub tunnels

- NHRP Resolution Request
  - Spoke queries for NBMA-to-VPN mappings of other spokes
  - Required to build Spoke-to-Spoke tunnels
- NHRP Redirect
  - NHS answer to Spoke-to-Spoke data-plane packet through it

#### DMVPN PHASES OVERVIEW

- Phase-1: Hub and Spoke only
  - If spoke wanted to send data to each other, they could not
  - Hub can become the bottleneck in CPU and Bandwidth
  - Summarization is possible
  - Now considered obsolete
- Phase-2: Adds Spoke-to-Spoke Capability
  - All spoke routers must know all IP routes of all other spoke routers
  - NO summarization allowed
  - NO default route origination allowed from hub
  - Removes burden from hub in terms of CPU and Bandwidth
  - Spokes now have a mGRE interface like the hub
- Phase-3: Modification of Routing Behavior
  - Hub allowed to summarize all routes from Spokes
  - Sets the next-hop of summarized routes to itself
  - Hub can send NHRP Redirect Messages to Spokes
  - Hub preemptively sends information to Spokes about other Spokes
  - Spokes no longer needs to ask via NHRP Resolution Request
  - Next-hop preservation no longer required
  - Default route from hub is now possible

# DMVPN NEGOTIATION OVERVIEW

- Creates on-demand tunnels between nodes
  - Initial tunnel-mesh is hub-and-spoke (always on)
  - Traffic patterns trigger spoke-to-spoke tunnels
  - Solves management scalability problem
- Maintains tunnels based on traffic patterns
  - Spoke-to-Spoke tunnel is on-demand
  - Spoke-to-Spoke tunnel lifetime is based on traffic
- Requires two IGPs: Underlying and Overlay
  - IPv4/IPv6 supported for both passenger and transport
  - One routing protocol with service provider to exchange public addresses
  - Another routing protocol between corporate routers

# DMVPN PHASE 1 AND 2 CONFIGURATION

- Hub(config)# interface tunnel 0
- Hub(config-if)# tunnel source {ipv4-address | interface}
- Hub(config-if)# tunnel key {#}
- Hub(config-if)# ip nhrp map multicast dynamic
- Hub(config-if)# ip nhrp network-id {#}
- Hub(config-if)# ip mtu 1400
- Hub(config-if)# ip tcp adjust-mss 1360

- Spoke(config)# interface tunnel 0
- Spoke(config-if)# tunnel destination {NBMA IP of Hub}
- This is the phase differentiator remove for Phase 2
- Spoke(config-if)# tunnel key {#}
- Spoke(config-if)# ip nhrp map multicast {NBMA IP of HUB}
- Spoke(config-if)# ip nhrp map {private-ip} {NBMA IP}
- Spoke(config-if)# ip nhrp network-id {#}
- Spoke(config-if)# ip mtu 1400
- Spoke(config-if)# ip tcp adjust-mss 1360

# **DMVPN PHASE 2 NHRP FLOW**

- NHRP Resolution Request initiated by spoke:
- Spoke looks up next-hop address (tunnel IP) for destination network
- Spoke looks into local NHRP database for mapping
  - If it does, it sends the traffic
  - If not, sends the NHRP Resolution Request to its NHS
- The hub forwards to owner of network we are trying to reach
- Owner responds directly to requesting spoke
- Requestor adds NHRP mapping to local database
- Spoke-to-spoke communication has established

#### DMVPN PHASE 3 NHRP FLOW

- Spoke goes to send traffic to another spoke
- Spokes can and should be using a default route!
- Traffic routed to the hub
- Hub sends NHRP Redirect to spoke (traffic continues to flow!)
- Spoke receives redirect and generates NHRP Resolution Request
- Hub receives NHRP Resolution Request and forwards to destination spoke
- Destination spoke sends NHRP Resolution Reply to originating spoke
- Viola! Spoke-to-spoke tunnel as initiated by the hub

## DMVPN PHASE 3 CONFIGURATION

- Hub(config-if)# ip nhrp redirect
- Spoke(config-if)# ip nhrp shortcut

# 4.3 Describe Easy Virtual Networking (EVN)

- EVN A.K.A. "Easy Virtual Networking"
  - An easy method to provide complete isolation between network segments
    - No VRF / VRF-lite
    - No BGP
    - No Vlan Routing / ACLs
  - **Reality**: Automation tool for vrf-lite, relies on 802.1Q encapsulation

#### VRF-LITE CONFIGURATION STEPS

- Router(config)# interface gi1/1
- Router(config-if)# ip address 10.5.5.1 255.255.255.0
- Router(config-if)# ip pim sparse-mode
- Router(config)# interface gi1/1.51
- Router(config-if)# description Subinterface for Group1
- Router(config-if)# encapsulation dot1Q 51
- Router(config-if)# ip vrf forwarding Group1

- Router(config-if)# ip address 10.5.5.1 255.255.255.0
- Router(config-if)# ip pim sparse-mode
- Router(config)# interface gi1/1.52
- Router(config-if)# description Subinterface for Group2
- Router(config-if)# encapsulation dot1Q 52
- Router(config-if)# ip vrf forwarding Group2
- Router(config-if)# ip address 10.5.5.1 255.255.255.0
- Router(config-if)# ip pim sparse-mode

#### **EVN CONFIGURATION STEPS**

- Router(config)# vrf definition group1
- Router(config-vrf)# vnet tag 51
- Router(config)# vrf definition group2
- Router(config-vrf)# vnet tag 52
- Router(config)# interface gi1/1
- Router(config-if)# vnet trunk
- Router(config-if)# ip address 10.5.5.1 255.255.255.0
- Router(config-if)# pim sparse-mode
- Router(config)# interface gi1/2
- Router(config-if)# vrf forwarding group1
- Router(config-if)# **ip address** 10.75.1.2 255.255.255.0

# EVN SUMMARY AND REALITY

- EVN generates subinterfaces with vnet trunk
- EVN adds a vrf forwarding instance to each subinterface

#### VERIFICATION

- show vrf
- ping vrf <VRF> <IP>

# 5.0 Infrastructure Security

# 5.1 Describe IOS AAA using local database

- AAA A.K.A. "Authentication, Authorization, and Accounting"
  - **Authentication**: Verify the identity of the user.
  - **Authorization**: What is the user allowed to do?
  - Accounting: Used for billing and auditing

# GENERAL CONCEPT

- 802.1X is the port-based mechanism that will block or unblock the interface
- EAPoL A.K.A. "Extensible Authentication Protocol over LAN"
  - No traffic allowed beside EAPoL
  - EAPoL is the protocol that verifies the user
- Once authenticated, network access is granted and traffic resumes

## AUTHENTICATION SERVER TYPES

- RADIUS A.K.A. "Remote Authentication Dial-In User Service"
- TACACS+ A.K.A. "Terminal Access Controller Access-Control System"

#### PRIVILEGE LEVELS

- Level-0: Only a few commands are available, most used command is 'enable'
- Level-1: Default exec user level, can use some show commands
- **Level-15**: Also known as 'enable mode' or 'privileged mode'; full access
- CONFIGURATION OPTIONS

- Assign some privilege level 15 commands to level 1
- Move some commands from level 1 to higher level
- Create new privilege level and assign some level 15 commands to it

#### CONFIGURATION STEPS

- Set Command Privilege
- Router(config)# privilege exec level 1 show running-config
- Router(config)# privilege exec level 15 show ip arp
- Enable AAA Locally
- Router(config)# username JUNIOR privilege 8 password CISCO
- Router(config)# aaa new-model
- Router(config)# aaa authentication login default local
- New User Privilege
- Router(config)# privilege exec level 8 configure terminal
- Router(config)# privilege exec level 8 debug ip routing
- Router(config)# privilege exec level 8 undebug all
- Router(config)# privilege exec level 8 show running-config
- Router(config)# privilege configure level 8 interface
- Router(config)# privilege interface level 8 shutdown
- Router(config)# privilege interface level 8 no shutdown

# 5.2 Describe device security using IOS AAA with TACACS+ and Radius

# 5.2.a AAA with TACACS+ and RADIUS

#### - RADIUS CONFIGURATION STEPS

- Router(config)# aaa new-model
- Router(config)# radius-server host 192.168.1.200 auth-port 1812 acct-port 1813
- Router(config)# radius-server key MY KEY
- Router(config)# aaa authentication login default group radius

#### - TACACS+ CONFIGURATION STEPS

- Router(config)# aaa new-model
- Router(config)# tacacs-server host 192.168.1.200 port 49 key MY\_KEY
- Router(config)# aaa authentication login default group tacacs+

#### PORT INFORMATION

- RADIUS uses UDP, reserved ports are 1812 and 1813
- RADIUS legacy unofficial ports are 1645 and 1646
- TACACS+ uses TCP, reserved port is 49

# - AUTHORIZATION CONFIGURATION STEPS

- Router(config)# aaa authorization exec group tacacs+

#### ACCOUNTING CONFIGURATION STEPS

- Router(config)# aaa accounting network default stop-only group tacacs+

#### VERIFICATION

- test aaa group [radius | tacacs+] <username> <password> new-code

# 5.2.b Local privilege authorization fallback

- GENERAL CONCEPT

- If TACACS+ or RADIUS becomes unavailable, need fallback option
  - Router(config)# aaa authentication login default group tacacs+ local
- The **local** keyword allows local database usage!

# 5.3 Configure and verify device access control

# 5.3.a Lines (VTY, AUX, console)

- GENERAL CONCEPT
  - The lines all provided access to the CLI, in one form or another.
  - Line VTY
    - Provided SSH or Telnet access to the device
  - Line AUX
    - Provided CLI access via aux cable
    - Only showed output during complete power up
  - Line CON
    - Provided CLI access through console cable
    - Showed output during power up process as well

# 5.3.b Management plane protection

- MPP A.K.A. "Management Plane Protection"
  - Allows certain protocols on the management interface
    - beep: Beep Protocol
    - **ftp**: File Transfer Protocol
    - http: HTTP Protocol
    - https: HTTPS Protocol
    - **snmp**: Simple Network Management Protocol
    - ssh: Secure Shell Protocol
    - telnet: Telnet Protocol
    - tftp: Trivial File Transfer Protocol
    - tl1: Transaction Language Session Protocol

## CONFIGURATION STEPS

- Router(config)# control-plane host
- Router(config-cp-host)# management-interface FastEthernet0/0 allow ssh
- Router(config-cp-host)# management-interface FastEthernet0/0 allow snmp

# 5.3.c Password encryption

- GENERAL CONCEPT
  - There are two primary ways to encrypt passwords on Cisco devices
    - Router(config)# enable secret PASSWORD
    - Router(config)# service password-encryption
  - Secret enables encryption on the enable password
  - The service command does it globally to all passwords

# 5.4 Configure and verify router security features

# 5.4.a IPv4 access control lists (standard, extended, time-based)

- ACL A.K.A. "Access Control List"
  - Two flavours, standard and extended
  - PERMIT or DENY addresses and/or protocols
  - Explicit DENY at the end of every ACL

## - STANDARD ACL CONFIGURATION STEPS

- Router(config)# access-list 1 permit 192.168.12.0 0.0.0.255
- Router(config)# interface fastEthernet 0/0
- Router(config-if)# ip access-group 1 in

# EXTENDED ACL CONFIGURATION STEPS

- Router(config)# access-list 100 permit tcp 1.1.1.0 0.0.0.255 host 2.2.2.2 eq 80

# - ACCESS-LIST EDITOR CONFIGURATION STEPS

- Router(config)# ip access-list extended 100
- Router(config-ext-nacl)# no 20

# - NAMED ACL CONFIGURATION STEPS

- Router(config)# ip access-list extended DROPICMP
- Router(config-ext-nacl)# deny icmp host 192.168.12.2 1.1.1.0 0.0.0.255

#### VERIFICATION

- show access-lists

#### TIME-BASED ACL CONFIGURATION STEPS

- Router(config)# time-range WH
- Router(config-time-range)# periodic weekdays 9:00 to 17:00
- Router(config)# ip access-list extended NO FACEBOOK
- Router(config-ext-nacl)# deny tcp any host 192.168.23.3 eq 80 time-range WH
- Router(config-ext-nacl)# permit ip any any
- When verifying time based ACLs, check for active keyword

# 5.4.b IPv6 traffic filter

#### GENERAL CONCEPT

- IPv6 standard ACL offers the following functions:
  - Filter traffic based on source and destination address
  - Filter traffic inbound or outbound on a specific interface
  - Add priority to the ACL
  - Implicit deny all at the end of the ACL
- No concept of numbered IPv6 ACLs, just named ACLs
- No concept of standard or extended

# IPv6 ACL CONFIGURATION STEPS

- Router(config)# ipv6 access-list TEST
- Router(config-ipv6-acl)# permit tcp any 2001:aaaa::/64 eq telnet
- Router(config)# interface Ethernet0
- Router(config-if)# ipv6 traffic-filter TEST out

#### VERIFICATION

- show ipv6 access-list
- debug ipv6 packet

# 5.4.c Unicast reverse path forwarding

- GENERAL CONCEPT
  - Normal Scenario: Router looks at destination address of packet before routing
  - uRPF A.K.A. "Unicast Reverse Path Forwarding"
    - Checks source address as well
    - Three modes: Loose, Strict, VRF Mode
- LOOSE MODE
  - Checks to see if the source address route exists on the routing table
- STRICT MODE
  - Same as loose, except also checks if source interface is the same on router
  - This means asynchronous routing will be blocked
- CONFIGURATION STEPS
  - Router(config-if)# ip verify unicast source reachable-via [any | rx]

# 6.0 Infrastructure Services

# 6.1 Configure and verify device management

# 6.1.a Console and VTY

- GENERAL CONSOLE CONFIGURATION STEPS
  - Router(config)# line con 0
  - Router(config-line)# login
  - Router(config-line)# password THE\_PASS

#### GENERAL VTY CONFIGURATION STEPS

- Router(config)# line vty 0 4
- Router(config-line)# login local
- Router(config-line)# access-class 1 in
- Router(config-line)# transport input ssh telnet
- Router(config)# username cindy password cisco
- VERIFICATION
  - show line
  - Note: You can also configure based on the line number as displayed in the above

# 6.1.b Telnet, HTTP, HTTPS, SSH, SCP

- TELNET
  - TCP Port 23
  - Clear text authentication protocol
  - Configuration Example:
    - Router(config)# username cisco password 0 cisco
    - Router(config)# line vty 0 4
    - Router(config-line)# transport input telnet

#### - SSH

- TCP Port 22
- Encrypted and secure
- Configuration Example:
  - Router(config)# hostname Router
  - Router(config)# aaa new-model
  - Router(config)# ip domain-name cisco.com
  - Router(config)# username cisco password 0 cisco
  - Router(config)# crypto key generate rsa
  - Router(config)# line vty 0 4
  - Router(config-line)# transport input ssh

#### - HTTP

- TCP Port 80
- Clear text www protocol
- Used for GUI access to devices
- Configuration Example:
  - Router(config)# ip http server
  - Router(config)# ip http authentication aaa

## - HTTPS

- TCP Port 443
- Encrypted www protocol
- Used for GUI access to devices
- Configuration Example:
  - Router(config)# ip http secure-server
  - Router(config)# crypto ca trustpoint CA-trust-local
  - Router(config-ca)# enrollment url http://host1:80
  - Router(config)# crypto ca authenticate CA-trust-local
  - Router(config)# crypto ca enrollment CA-trust-local
  - Router(config)# ip http secure-trustpoint CA-trust-local

## - SCP

- TCP Port 22
- Secure file transfer, uses same system as SSH
- Triple A and SSH must be configured in order for SCP to work
- Configuration Example:
  - Router(config)# aaa new-model
  - Router(config)# aaa authentication login default local
  - Router(config)# aaa authorization exec default local
  - Router(config)# username user1 privilege 15 password 0 lab
  - Router(config)# ip scp server enable

# 6.1.c (T)FTP

- GENERAL CONCEPT
  - Another means of copying from one location to another
  - Port Information:
    - **TFTP**: UDP 69
    - FTP: TCP 21 for initial authentication, TCP 20 for data transfer
- TFTP CONFIGURATION STEPS

- Router1(config)# tftp-server flash:/c2500-js-1.122-10b
- Router2# copy tftp flash

#### FTP CLIENT CONFIGURATION STEPS

- Router(config)# ip ftp username cisco
- Router(config)# ip ftp password cisco123
- Router# copy running-config ftp:

# - FTP SERVER CONFIGURATION STEPS (DEPRECATED)

Router(config)# ftp-server enable

# 6.2 Configure and verify SNMP

# 6.2.a v2

- SNMP VERSIONS
  - Version-1: Open standard and basic
  - Version-2: Requires community string to gain read/write access
  - **Version-3**: Adds encryption and real authentication
- OID A.K.A. "Object Identifier"
  - Numeric string that can be assigned various properties of a router
  - Example: Interface Counters
- MIB A.K.A. "Management Information Base"
  - Collection of OIDs
  - This can be imported to monitoring tool for it to understand router
- SNMP GENERAL CONCEPT
  - NMS A.K.A. "Network Management System"
    - External server to store logging information
  - SNMP Agents
    - Runs on network devices to monitor it
- COMMUNITY CONFIGURATION STEPS
  - Router(config)# snmp-server community TSHOOT ro
- OPTIONAL CONFIGURATION STEPS
  - Router(config)# snmp-server location Amsterdam
  - Router(config)# snmp-server contact info@networklessons.com
  - Displays information about where the device is located and contact to NMS
- FULL CONFIGURATION STEPS
  - Router(config)# snmp-server host 192.168.12.2 version 2c TSHOOT
  - Router(config)# snmp-server enable traps

# 6.2.b v3

- THREE SECURITY LEVELS
  - noAuthNoPriv = no authentication and no encryption
  - **AuthNoPriv** = authentication but no encryption
  - AuthPriv = authentication AND encryption
  - With noAuthNoPriv, the username will replace the community-string
- THREE NEW ELEMENTS

- **SNMP View:** Can restrict "view" of device, so only specific OIDs are visible
- **SNMP Group**: Can associate view with group, and can specify read/write access
- **SNMP User**: Exact user credentials required to access view based on group

#### GROUP CONFIGURATION STEPS

- Router(config)# snmp-server group MYGROUP v3 priv <access | read | write>
- If you do not specify, it will assume read-only

# - USER CONFIGURATION STEPS

- Router(config)# snmp-server user USER GROUP v3 auth md5 PASS priv aes 128 KEY

#### VERIFICATION

- show snmp user
- show snmp group

#### VIEW CONFIGURATION STEPS

- Router(config)# snmp-server view ALL-ACCESS 1.3.6.1.2.1.2.2.11 included
- Router(config)# snmp-server group MYGROUP v3 priv read ALL-ACCESS

# 6.3 Configure and verify logging

# 6.3.a Local logging, syslog, debugs, conditional debugs

#### LOGGING CONSOLE CONFIGURATION STEPS

- Router(config)# logging console
- This is enabled by default on IOS
- Router# terminal monitor
- Enables logging on screen to SSH/Telnet for that session

# LOCAL HISTORY CONCEPT

- Cisco IOS keeps history of syslog messages
- Can be displayed via show logging
- Messages stored in RAM by default

## INCREASE BUFFER SIZE (RAM) CONFIGURATION STEPS

- Router(config)# logging buffered 16384

#### SYSLOG SERVER CONFIGURATION STEPS

- Store logging messages remotely in a server as RAM data can be lost
- Router(config)# logging 192.168.1.2

# - DEBUG CONFIGURATION STEPS

- See debug information via console
  - Router(config)# logging console debugging
- See debug information via VTY
  - Router(config)# logging monitor debugging
- Filtering can be done based on severity levels, debug being severity 7

# - SEVERITY LEVELS

- Every Alley Cat Eats Wet Noodles In Doors
- Severity-0: Emergencies
- **Severity-1**: Alerts
- Severity-2: Critical
- **Severity-3**: Errors
- Severity-4: Warnings

- Severity-5: Notifications
- Severity-6: Informational
- **Severity-7**: Debugging

#### CONDITIONAL DEBUGGING CONCEPT

- Filter based on specific interfaces and some other items
- You need to turn on some debug item first, then you can turn on the filter

# - CONDITIONAL DEBUGGING CONFIGURATION STEPS

- Router# debug ip rip
- Router# debug condition interface serial 0
- VERIFICATION
  - show debug condition

# - REMOVE CONDITIONAL DEBUGGING STEPS

- Router# undebug condition interface serial 0

# 6.3.b Timestamps

#### CONFIGURATION STEPS

- Router(config)# service timestamps log
- Router(config)# service timestamps log datetime msec
- Router(config)# service sequence-numbers

# 6.4 Configure and verify Network Time Protocol (NTP)

# 6.4.a NTP master, client, version 3, version 4

#### TIME TRACKING PURPOSE

- Logging output
- Debugging output
- User "show" commands
- Network Management/Reporting Tools

#### CLOCK SOURCE

- Internal System Clock
- Most are battery-driven and maintains the time/date across reloads
- This information can be distributed via NTP

#### SYSTEM CLOCK INFORMATION SOURCE

- Manual Configuration
- NTP A.K.A. "Network Time Protocol"
- SNTP A.K.A. "Simple Network Time Protocol"
- VINES A.K.A. "Virtual Integrated Network Service"

#### VERSIONS

- Version 3: Only works with IPv4
- Version 4: Introduced concept of NTP authentication
- Uses UDP port 123 for source and destination
- NTP CLIENTS AND SERVERS

- NTP Client
  - Device that periodically polls a server for time/calendar information
- NTP Server
  - Provides time information to client
  - Considered "authoritative source" based on stratum:
    - Stratum-1: Device directly connected to atomic clock source
    - Stratum-X: Time Server that is X hops away from Stratum-1

#### MODES OF OPERATION

- NTP Server
  - Also called NTP Master
- NTP Client
  - Receives NTP data from Server/Master
- NTP Peers
  - Also called "Symmetric Mode"
  - Two-or-more NTP Servers exchange information
- NTP Broadcast/Multicast
  - Unidirectional "push" from NTP server

#### NTP MASTER CONFIGURATION STEPS

- Router# clock set 07:01:30 28 Feb 2015
- Router(config)# ntp master [stratum]
- Router(config)# ntp peer x.x.x.x (optional)
- Router(config-if)# ntp broadcast (optional)

#### NTP CLIENT CONFIGURATION STEPS

- Router(config)# ntp server x.x.x.x version 4 OR
- Router(config-if)# ntp broadcast client

#### VERIFICATION

- show ntp status
- show ntp associations

# 6.4.b NTP authentication

#### - GENERAL CONCEPT

- Two Protection Methods:
  - Access-Lists
  - NTP Authentication
- Uses an MD5 Hash of configured authentication keys
- Multiple keys can be configured for use against different peers

#### - MASTER CONFIGURATION STEPS

- Router(config)# ntp authentication-key <key-number> md5 <key-string>
- Router(config)# ntp trusted-key <key-number>
- Router(config)# ntp authenticate

# - CLIENT CONFIGURATION STEPS

- Router(config)# ntp authentication-key <key-number> md5 <key-string>
- Router(config)# ntp trusted-key <key-number>
- Router(config)# ntp authenticate
- Router(config)# **ntp server** x.x.x.x **key** <key-number>

# 6.5 Configure and verify IPv4 and IPv6 DHCP

# 6.5.a DHCP client, IOS DHCP server, DHCP relay

- DHCP A.K.A. "Dynamic Host Configuration Protocol"
- DHCPv4 PROCESS (DORA)
  - DHCPv4 Discover
    - Client sends a broadcast requesting DHCP services
    - May include last known IPv4 address or 0.0.0.0
  - DHCPv4 Offer
    - Server responds with an offer, which includes:
      - IPv4 Address
      - Subnet Mask
      - Lease Duration
      - Client's MAC address
    - Offer may include other options:
      - Default Gateway
      - DNS Server Addresses
  - DHCPv4 Request
    - Client responds with request to accept the offered address
    - The request is broadcasted to inform all DHCP servers on LAN
  - DHCPv4 Acknowledgement
    - Server sends ACK to client and closes out communication
- DHCPv4 CONFIGURATION STEPS
  - Router(config)# ip dhcp pool MY\_POOL
  - Router(dhcp-config)# network 192.168.12.0 255.255.255.0
- DHCPv4 EXCLUDING CONFIGURATION STEPS
  - Router(config)# ip dhcp excluded-address 192.168.12.100
- DHCPv6 THREE METHODS
  - **Method-1:** Stateless Address Autoconfiguration (SLAAC)
  - Method-2: SLAAC + Stateless DHCPv6 Server
  - Method-3: Stateful DHCPv6 Server
- COMMUNICATION OCCURS VIA ICMPv6
  - RS A.K.A. "Router Solicitation"
    - Clients send RS to obtain IPv6 addresses dynamically
    - Source IPv6 address for RS is either link local address or unspecified
  - RA A.K.A. "Router Advertisement"
    - Automatically sent out interfaces if IPv6 unicast-routing is enabled
    - Configuring Router as IPv6 Router (ipv6 unicast-routing):
      - Can forward IPv6 packets transiting the router
      - Can configure with dynamic IPv6 routing protocol
      - IPv6 interfaces joins the all-IPv6 multicast group (FF02::2)
    - Used to advertise presence and link specific parameters:
      - Link Prefix
      - Prefix-length

- Default gateway
- Link MTU
- Suggested IPv6 address method
- Sends advertisements:
  - Every 200 seconds automatically
  - In response to RS message
- A, O, AND M FLAGS
  - Address Autoconfiguration Flag (A Flag)
    - Suggests SLAAC
    - Allows host to generate own GUA A.K.A. "Global Unicast Address"
    - Two options for creating Interface ID:
      - EUI-64 Process
      - Random 64-bit Value
  - Other Configuration Flag (O Flag)
    - Suggests Stateless DHCPv6 Server
    - May include DNS server address and domain name
  - Managed Address Configuration Flag (M Flag)
    - Suggests Stateful DHCPv6 Server
    - Similar to DHCP for IPv4
    - Only difference is, host uses RA packet source as default gateway
- FOUR PRIMARY DHCPv6 MESSAGE TYPES
  - SOLICIT (1)
    - Client message to locate servers
  - ADVERTISE (2)
    - Server message to clients to indicate available DHCPv6 service
  - REQUEST (3)
    - Client message to request configuration parameters including IPv6 addresses, from specific DHCPv6 server
  - REPLY (7)
    - Server message containing assigned addresses and configuration parameters
- SLAAC A.K.A. "Stateless Address Autoconfiguration"
  - Host generates own GUA without any services based on provided prefix in RA
  - Flags: A = 1, O = 0, M = 0 (default behavior)
  - With Privacy Extension, host will generate random interface, otherwise EUI-64:
    - MAC address split into two sections, and FF:FE inserted in between
    - The seventh bit of the first byte is inverted
  - Hosts generate Temporary Addresses with Privacy Extension:
    - Generated using random 64-bit value
    - Used as source IPv6 address when device originates connection
  - Hosts will generate a Public Address:
    - Used as a permanent destination address by other hosts
    - Hosts may have multiple Temporary Addresses, but only one Public
  - Client Configuration:

Router(config-if)# ipv6 address autoconfig

#### STATELESS DHCPv6 WITH SLAAC

- Host generates own GUA with services provided by DHCP server
- **Flags**: A = 1, O = 1, M = 0
- Server Configuration:
  - Router(config)# ipv6 dhcp pool <pool-name>
  - Router(config-dhcpv6)# dns-server <ipv6-address>
  - Router(config-dhcpv6)# domain-name <domain>
  - Router(config-if)# ipv6 nd other-config-flag
  - Router(config-if)# ipv6 dhcp server <pool-name> [rapid-commit]
- The rapid-commit option shortens messages from four to two:
  - SOLICIT
  - REPLY
- Relay Agent Configuration:
  - Router(config-if)# ipv6 dhcp relay destination <ipv6-address>

#### STATEFUL DHCPv6

- Host receives prefix and options from DHCP server
- Flags: A = 0, O = 0, M = 1
- Server Configuration:
  - Router(config)# ipv6 dhcp pool <pool-name>
  - Router(config-dhcpv6)# address prefix prefix/length>
  - Router(config-dhcpv6)# dns-server <ipv6-address>
  - Router(config-dhcpv6)# domain-name <domain>
  - Router(config-if)# ipv6 nd managed-config-flag
  - Router(config-if)# ipv6 nd prefix prefix/length> no-autoconfig
  - Router(config-if)# ipv6 dhcp server <pool-name>

# 6.5.b DHCP options (describe)

- GENERAL CONCEPT
  - DHCP includes other configuration options other than just address allocation
  - DHCPv4 can hand out useful data such as default gateway, DNS server, etc.
  - DHCPv6 will not hand out default gateway as the RA does that already

# 6.6 Configure and verify IPv4 Network Address Translation (NAT)

# 6.6.a Static NAT, dynamic NAT, PAT

- NAT A.K.A. "Network Address Translation"
  - Function by which IP addresses within a packet are replaced with another
  - NAT44 refers to IPv4 <> IPv4
  - NAT444 describes architecture rather than technology
    - Two levels of NAT44
    - Enables IPv4 private address space to be used over and over
    - Service providers use NAT444 to extend dwindling IPv4 space
  - Originally introduced to slow down the depletion of available IP address space
  - No guaranteed private address will not leak, requires use of filters

# - FOUR TYPES OF ADDRESSES

- Inside Local (IL)
  - Addresses assigned to inside devices
  - Not advertised to the outside
- Inside Global (IG)
  - Addresses by which inside devices are known to the outside
- Outside Global (OG)
  - Addresses assigned to outside devices
  - These addresses are not advertised
- Outside Local (OL)
  - Addresses by which outside devices are known to the inside

#### - VERIFICATION

show ip nat translations

#### INSIDE/OUTSIDE CONFIGURATION STEPS

- Router(config)# interface FastEthernet1/0
- Router(config-if)# ip nat inside
- Router(config)# interface Serial 2/0
- Router(config-if)# ip nat outside

#### STATIC NAT

- One-to-One Mappings
- Configuration Example:
  - Router(config)# ip nat inside source static 10.1.1.3 204.15.87.1
  - Router(config)# ip nat inside source static 10.1.2.2 204.15.87.2

# - DYNAMIC NAT

- One-to-Many Mappings
- Three Steps:
  - Define a pool of outside global addresses
  - Define an access list with the inside local addresses
  - Link the pool and access list together with a source list

# Configuration Example:

- Router(config)# ip nat pool PoolOne 204.15.86.1 204.15.86.254 prefix 24
- Router(config)# access-list 1 permit 10.1.1.0 0.0.0.255
- Router(config)# **ip nat inside source list** 1 **pool** PoolOne

# Alternate Pool Configurations:

- Using netmask instead of prefix
- **ip nat pool** PoolOne 204.15.86.1 204.15.86.254 **netmask** 255.255.255.0
- Matching host address (10.1.1.23 -> 204.15.86.23)
- ip nat pool PoolOne 204.15.86.1 204.15.86.254 prefix 24 type match-host

# - PAT A.K.A. "Port Address Translation"

- Many-to-Many Mappings
- Allows you to map on a per port basis rather than per IP
- Configuration Example:
  - Using the overload keyword in the source list command enables PAT
  - Router(config)# ip nat inside source list 1 interface Serial0 overload

# 6.7 Describe IPv6 NAT

# 6.7.a NAT64

# - GENERAL CONCEPT

- Commonly used to translate between private and public IPv4 address space
- Replaced NAT-PT as it had tight coupling with DNS in order to work
- AFT A.K.A. "Address Family Translation"
  - Provides communication between IPv6-only and IPv4-only networks
  - Translation is not a long-term strategy, ultimate goal is native IPv6
- Advantages over tunneling:
  - Translation provides a means of gradual and seamless migration to IPv6
  - Content providers can provide services transparently to IPv6 users

#### - STATELESS NAT64

- No bindings or session state maintained
- Not recommended and is generally considered useless
- IPv4 address of destination embedded in IPv6 address
- Specific IPv6 range represents IPv4 systems within IPv6 world
- Allows bidirectional reachability of v6 to / from v4 hosts

#### STATEFUL NAT64

# Three Components:

- NAT64 Prefix
  - Any /32 ~ /96 prefix used with a converted IPv4 address
  - NAT64 prefix can be network-specific (NSP) or well-known (WSP)
  - NSP is assigned by the organization and manually configured
  - WKP for NAT64 is 64:ff9b::/96
  - WKP is prepended to the converted IPv4 if NSP is not configured

# - DNS64 Server

- Functions as a normal DNS server for IPv6 AAAA records
- Attempts to locate an IPv4 record if AAAA record is not available
- If an A record is located, DNS64 converts the IPv4 record to IPv6
- This is done using the NAT64 prefix
- Only required from 6 -> 4

### - NAT64 Router

- NAT64 device advertises the NAT64 prefix into IPv6-only network
- Perform the translation between IPv6-only and IPv4-only networks

#### CONFIGURATION STEPS

- Router(config)# nat64 v6v4 static 2001:db8:feed:1::e 172.16.1.10
- Router(config-if)# nat64 enable

# 6.7.b NPTv6

- NPTv6 A.K.A. "Network Prefix Translation"
  - Stateless technology alongside the deprecated NAT66
  - Concept of translation from ULA to GUA is the subject of ongoing debate

- NAT does not provide security and with IPv6 address depletion is not an issue
- Useful mechanism for implementing address independence in an IPv6 network
- The translation occurs as a 1:1 relationship between the IPv6 addresses
  - Occurs between the inside and outside addresses
  - Only the prefix is replaced
- Described in more depth in RFC 6296

#### CONFIGURATION STEPS

- Only works on ASRs and IOS XE
- Router(config)# interface GigabitEthernet1/1
- Router(config-if)# nat66 inside
- Router(config)# interface GigabitEthernet1/2
- Router(config-if)# nat66 outside
- Router(config)# nat66 prefix inside 2002:AB01::/64 outside 2002:AB02::/64

# 6.8 Describe SLA architecture

- SLA A.K.A. "Service-Level Agreements"
  - Measures network's current performance
  - If performance does not meet threshold, PBR reacts appropriately
  - Measurement can be as simple as using ping to determine response
  - Or it can be sophisticated as measuring jitter (delay variation) of VoIP

#### - SLA OPERATION SETUP

- When defining IP SLA operation, a receiver can be a router or a host
- Receiver may need to be configured as IP "SLA Responder"

## SUMMARY OF AVAILABLE OPERATION TYPES

- ICMP (echo, jitter)
- RTP (VoIP)
- TCP Connection (establishes TCP connections)
- UDP (echo, jitter)
- DNS
- DHCP
- HTTP
- FTP

# 6.9 Configure and verify IP SLA

# 6.9.a ICMP

#### CONFIGURATION STEPS

- **ip sla** sla-ops-number (global command)
- icmp-echo {destination-ip-address | destination-hostname} ... (SLA subcommand)
  - source-ip {ip-address | hostname} OR # optional
  - source-interface interface # optional
- frequency seconds (SLA subcommand)
- **ip sla schedule** sla-ops-number ... (global subcommand)
  - life {forever | seconds} OR

- start-time {hh:mm[:ss] month day | pending | now | after hh:mm:ss}
- ageout seconds # optional
- recurring # optional

#### VERIFICATION

- show ip sla configuration
- show ip sla statistics

# 6.10 Configure and verify tracking objects

# 6.10.a Tracking objects

- CONFIGURATION STEPS
  - To configure track object, use the following command in global:
    - track <object-name> ip sla <sla-ops-number> [state | reachability]
  - To configure delay reaction, use the delay command:
    - delay {down seconds | up seconds} (track subcommand)

#### VERIFICATION

- show track
- show ip sla statistics

# 6.10.b Tracking different entities (for example, interfaces, IP SLA results)

- CONFIGURATION STEPS
  - To configure a static route with a track object, use the following command:
    - ip route prefix mask {interface | next-hop} track object-number (global)
  - To configure PBR to use object tracking, use the following set command:
    - set ip next-hop verify-availability next-hop seg track object-number
    - When tracking object is down, PBR acts as if set command does not exist

# 6.11 Configure and verify Cisco NetFlow

- GENERAL CONCEPT
  - Accounting feature on top of an existing switching path such as CEF
  - Provides statistics on packets flowing through router:
    - INGRESS: IP, IP to MPLS, Frame Relay, ATM
    - **EGRESS**: IP, MPLS to IP. Traffic sourced from the router not captured
  - Single flow is a combination of seven key-fields
- FLOW PARAMETERS
  - Source IP Address
  - Destination IP Address
  - Source Port Number
  - Destination Port Number
  - Layer 3 Protocol Type
  - Type of Service (ToS)
  - Logical Interface (ifIndex)
- LIFE CYCLE

- Router notices network traffic and creates flows in NetFlow cache locally
- Active flows expire
- Router exports flow-records to collector
- Expiration of Flow:
  - Inactivity Timer (15 Seconds Default)
  - Active Timer Expired (30 Minutes Default)
  - NetFlow Cache Full (FIFO)
  - TCP RST or FIN Flag Seen

#### FLOW PROCESSING

- Pre Processing
  - Packet Sampling
    - Statistical sampling of network traffic
    - Set up for traffic engineering or capacity planning
    - E.G. 1 in 100 packets
  - Filtering
    - Sets up specific subset of network traffic for class-based analysis
- Post Processing
  - Aggregation Schemes
    - Sets up extra aggregation caches
    - Different combinations of fields determine which flows are grouped
  - Export to multiple destinations
    - Sets up identical streams of NetFlow data to be sent

# GOOD TO KNOW TERMINOLOGY

- Flow: Actual traffic that occurred
- Flow-Record: Info about that traffic
- **Export-Packet**: Transport to the collector
- Packet Header: Has version, sequence number, and more
- **Schema**: Cognitive framework that helps organize and interpret info
- Template Record (Schema): Defines fields and structure in Flow-Record
- Options Template Record (Schema): Defines fields and structures in Options
- Options-Record (Data): Configuration of the Netflow, such as sampling

# 6.11.a Netflow v5, v9

- VERSION 5
  - Fixed format, cannot be extended or added to
  - IPv4 only
  - Added BGP AS information and sequence numbers
  - Export data from main cache only
  - No real concept of 'ingress' and 'egress' flows
  - Collector engine reverses the info behind the scenes without additional config
- VERSION 9
  - Add additional information to flows and template based
  - Added support for MPLS, BGP next-hop, and IPv6 headers

Exports data from main and aggregation cache

# 6.11.b Local retrieval

## GENERAL CONFIGURATION

- Router(config)# snmp-server community public ro
- Router(config)# ip flow-capture icmp
- Router(config)# ip flow-capture {capture-type}
- Router(config)# ip multicast netflow rpf-failure
- Router(config)# ip flow-cache entries 1024
- Router(config)# ip flow-cache timeout active 15
- Router(config)# ip flow-cache timeout inactive 30
- Router(config)# interface fa0/0
- Router(config-if)# ip flow ingress
- Router(config-if)# ip flow egress

#### GENERAL VERIFICATION

- show ip flow interface
- show ip cache flow
- show ip cache verbose flow

#### - SAMPLING CONFIGURATION

- Router(config)# flow-sampler-map {sampler-name}
- Router(config-sampler)# mode random one-out-of 10
- Router(config-if)# flow-sampler {sampler-name} egress

#### TOP TALKERS CONFIGURATION

- Router(config)# ip flow-top-talkers
- Router(config-flow-top-talkers)# top 10
- Router(config-flow-top-talkers)# sort-by packets
- Router(config-flow-top-talkers)# match protocol 1

# TOP TALKERS VERIFICATION

- show ip flow top-talkers

# 6.11.c Export (configuration only)

## GENERAL CONFIGURATION

- Router(config)# ip flow-export destination {ip-address} {port}
- Router(config)# ip flow-export version {5 | 9}

#### AGGREGATION CONFIGURATION

- Router(config)#ip flow-aggregation cache destination-prefix
- Router(config-flow-cache)# export destination {ip-address} {port}
- Router(config-flow-cache)# export version 9
- Router(config-flow-cache)# cache entries 1024
- Router(config-flow-cache)# mask destination minimum 26
- Router(config-flow-cache)# enabled

#### VERIFICATION

- show ip flow export
- show ip flow export template
- show ip cache flow aggregation destination-prefix