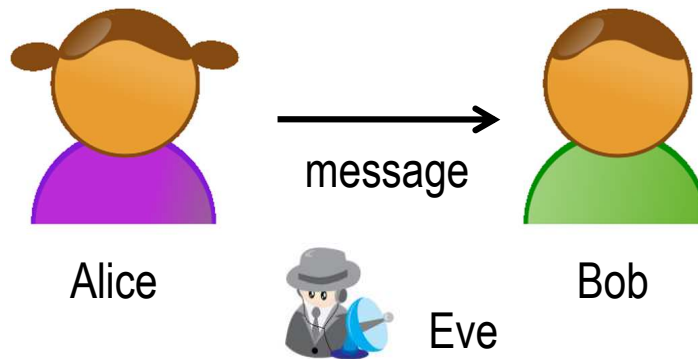


# **Classical Cryptosystems**

Applied Cryptography  
Andrei Bulatov

## Notation



Plaintext (明文) – Alice produced  $\xrightarrow{\text{encode}}$  ciphertext

Ciphertext (密文) – Bob received  $\xrightarrow{\text{decode}}$  plaintext (same as Alice)

Key

Protocol: (K, E, D)

K – key generation algorithm

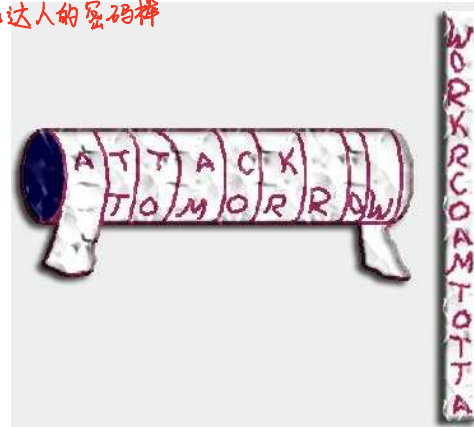
E – encryption algorithm

D – decryption algorithm  $\swarrow$  usually doing same in reverse order

Typically, if we know encryption algorithms, then it's easy to know decryption algorithm.

## Three Types of Cryptosystems

- Steganography** 信息隐藏, 隐写术, 速记式加密
  - eg. A king write msg on a slave's hand, and send the slave to another king. nobody knows that.
  - eg. Send a jpeg picture which is slightly modified and contains a msg.
  - 'Security by obscurity' 隐晦
    - send the msg in such a way that nobody knows.
    - (nobody know the msg is sent)
- Transposition cryptosystems:** 转置密码
  - E permutes (transposes) the letters of plaintext 置换/转置明文系统
  - D applies the converse transposition 应用相反的换位.
  - Example: Spartans Scytale 斯巴达人的密码棒



## Three Types of Cryptosystems (cntd)

- Substitution cryptosystems 替换式密码

E substitutes each letter of the plaintext with another letter or symbol

D applies the converse substitution

Example: Caesar cipher 凯撒密码

He made messages secret by shifting each letter three letters forward.

Thus we can replace letters by integers from 0 to 25.

Then E adds 3 modulo 26 to every letter.

To decrypt a message, D subtracts 3 from each letter



# Caesar Cipher

● Encrypt 'SEND MORE MEN AND AMUNITION'

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

S E N D M O R E M E N A N D A M U N I T I O N

18 4 13 3 12 14 17 4 12 4 13 0 13 3 0 12 20 13 8 19 8 14 13

21 7 16 6 15 17 20 7 15 7 16 3 16 6 3 15 23 16 11 22 11 17 16

V H Q G P R U H P H Q D Q G D P X Q L W L R Q

## Drawbacks of Classical Cryptosystems

- Too few keys 26!  
If the type of the cryptosystem is known it can be bruteforced
- Kerchoff's Principle: 柯克霍夫原则  
System should be secure even if algorithms are known,  
as long as key is secret
- Problem: How to increase the number of keys?

篱笆密码法

重新防御

## Transposition: Railfence and Redefence Ciphers

- Railfence cipher:

`SEND MORE MEN AND AMUNITION`

S				M				M				N				U				I		
	E		D		O		E		E		A		D		M		N		T		O	
		N				R				N				A				I				N

`SMMNUIEDOEEADMNTONRNAIN`

- Redefence Cipher

2	S				M				M				N				U				I		
1		E		D		O		E		E		A		D		M		N		T		O	
3			N				R				N				A				I				N

`EDOEEADMNTOSMMNUINRNAIN`

## Substitution: Linear Cipher

- Similar to Caesar cipher, but instead of adding 3, computes a linear function on letters. Say,

$$E: X \rightarrow 4X + 21 \pmod{26}$$



## Substitution: Playfair (人名)

### ● Keysquare:

logrithm + 未出现字母

L	O	G	A	R
I	T	H	M	B
C	D	E	F	K
N	P	Q	S	U
V	W	X	Y	Z

### ● Encryption

'SEND MORE MEN AND AMUNITION'

SE ND MO RE ME NA ND AM UN IT IO NA



QF PC TA GK HF SL PC MF NP TH TL SL

'QFPCTAGKHFSLPCMFNPHTLTL'

# possible keys ↑

## Substitution: Checkerboard

	W	H	I	T	E
B	E	N	C	R	Y
L	P	T	IJ	O	A
A	B	D	F	G	H
C	K	L	M	Q	S
K	U	V	W	X	Z

Plaintext: THIS IS A BETTER CIPHER

Ciphertext: LHAELICE LICE LE AW EW LH LH BW BT BI LI LW AE BW BT

先横向

再纵向

## Drawbacks of Classical Cryptosystems

### ● Frequencies analysis

Different letters have different probabilities to appear in a text

### ● Example

Ciphertext:

VXEVLWXWLRQ

FLSKHUV FDQ

RIWHQ EH EURNHQ

EB IUHTXHQLHV

DQDOBVLV

*freq in the left  
ciphertext*

Frequencies (in %%):

*avg freq in English*

A	0	6.9	J	0	0.8	S	2	6.8
B	4	0.9	K	2	0.9	T	2	9
C	0	4	L	10	3.9	U	6	2.8
D	6	4.2	M	0	3	V	12	2.1
E	8	13.1	N	2	8	W	8	2.1
F	6	2.7	O	2	8	X	6	1
G	0	2	P	0	2.2	Y	0	2.5
H	14	3	Q	12	1	Z	0	0.8
I	4	7.9	R	6	8.2			

# Frequencies Analysis

s u b s t i t u t i o n      c i p h e r s      c a n  
 V X E V W L W X W L R Q      F L S K H U V      F D Q  
    t? on → tion

o f t e n      b e      b r o k e n      b y  
 R I W H Q      E H      E U R N H Q      E B  
 very likely to be  
 "often"

f r e q u e n c i e s      a n a l y s i s  
 I U H T X H Q F L H V      D Q D O B V L V  
    analysis

E: H Q V

T, R, N, O: H @ V E L W "Q" may be "n"

b is guessed from "E H" to "be"

y is guessed from "E B" to "by"

# Smoothing Frequencies: Grandpre (人名) 格朗普雷加密法

	1	2	3	4	5	6	7	8
1	A	B	A	S	H	I	N	G
2	Y	O	K	O	H	A	M	A
3	C	O	E	X	I	S	T	S
4	D	E	A	T	H	F	U	L
5	J	A	C	K	P	O	T	S
6	Q	U	I	V	E	R	E	D
7	W	I	T	C	H	I	N	G
8	Z	O	D	I	A	C	A	L

Plaintext: Y O C A N O T B R E A K M E

Ciphertext: 21 22 47 31 11 17 77 24 37 12 66 33 13 23 27 42

for letters occurs more frequently, use more numbers to "smooth" the frequency.

# Smoothing Frequencies: <sup>(人名)</sup>Vegenere Cipher <sup>维吉尼尔加密法</sup>

Plaintext: SEND MORE MEN AND MUNITION

Key: KEY *need to count the length of the key.  
then use frequency analysis.*

Equivalent to shifts by 10 4 24 letters

S	E	N	D	M	O	R	E	M	E	N	A	N	D	M	U	N	I	T	I	O	N	
<i>key</i>	<i>k</i>	<i>e</i>	<i>y</i>	<i>k</i>	<i>e</i>	<i>y</i>	<i>k</i>	<i>e</i>	<i>y</i>	<i>k</i>	<i>e</i>	<i>y</i>	<i>k</i>	<i>e</i>	<i>y</i>	<i>k</i>	<i>e</i>	<i>y</i>	<i>k</i>	<i>e</i>	<i>y</i>	<i>k</i>
10	4	24																				

C I L N Q M B I K I L K R B W Y L S X G Y R

$$C_i \equiv P_i + K_{(i \bmod 3)} \pmod{26}$$

## Smoothing Frequencies: Vegenere Cipher (cntd)

- Idea: The longer key the better
- Codebooks 密码本 (缺点: 要保存好, 被偷了就完蛋)  
 ↓  
 soln: just use any books
- Autokey 自动密钥  
 ↳ use ciphertext or part of the ciphertext to be the key  
 (plaintext) (plaintext)
- Enigma "谜"
- One-time pad 一次性密码本



→ Enigma