

Note: Do not write on the backs; only front pages are digitized and graded.

1. (20 points) Answer with True or False.

- $3n^2 + n^3 \in \Omega(n^3)$  True
- $3n^2 + n^3 \in O(n^4)$  True

2. (10 points) Use the definition of  $O$  to show that  $5n^2 + 10n \in O(n^2)$ .

$$5n^2 + 10 \leq c \cdot n^2 \quad / \quad 5n^2 + 10 \leq (5n^2 + 10n^2) = 15n^2 \quad c=15 \text{ while } n \geq 1$$

3. (20 points) Use the definition of  $\Theta$  to show that  $3n^3 + 5n - 5 \in \Theta(n^3)$ .

$$3n^3 + 5n - 5 \leq c \cdot n^3 \quad / \quad 3n^3 + 5n - 5 \leq 3n^3 + 5n^3 - 5n^3 = 3n^3 \quad c=3 \text{ while } n \geq 1$$

$$3n^3 + 5n - 5 \geq c \cdot n^3 \quad / \quad 3n^3 + 5n - 5 \geq 1 \cdot n^3 \quad c=1 \text{ while } n \geq 1$$

since it is both in  $O(n^3)$  and  $\Omega(n^3)$ , it is also in  $\Theta(n^3)$

4. (10 points) Use  $\Theta$  to specify the running time of the following procedure in terms of  $n$ .

```
sum = 0
for i in range(n):
    for j in range(5):
        sum = sum + i + j
```

$$T(n) = n \cdot n \in \Theta(n^2)$$

5. (10 points) Use  $\Theta$  to specify the running time of the following procedure in terms of  $n$ .

```
sum = 0
for i in range(n):
    j = n
    while j > 1:
        sum = sum + i * j
        j = j / 2
```

$$T(n) = n \cdot 2 \log n \in \Theta(n \log n)$$

6. (10 points) Is the running time of the procedure in Question 4 is in  $\Omega(n)$ ? Explain briefly.

If the running time is specified in terms of  $\Theta(n^2)$ , it is by extension also in  $O(n^2)$  and  $\Omega(n^2)$ . Thus it is not in  $\Omega(n)$ .

7. (10 points) What is the space complexity of the procedure in Question 4?

$$S(n) = 1 \quad \Theta(1)$$

8. (10 points) Explain what the following function does. The input is a list of numbers.

```
def foo(L):
    if len(L) == 0:
        return 1
    return L[0] * foo(L[1:])
```

$foo(L)$  returns 1 if  $len(L) == 0$ , but recursively returns  $L[0] \cdot L[1:]$ .

$foo(L)$  returns the product of all values in  $L$  if  $len(L) > 0$ , else returns 1

- Smallest input - "solve" fork  
- assume (k-1)

9. (10 points) Use mathematical induction to explain that the following function correctly adds up all numbers that are divisible by 5 in the input list.

```
def bar(L):
    if len(L) == 0:
        return 0
    if L[0] % 5 == 0:
        return L[0] + bar(L[1:])
    return bar(L[1:])
```

If L is empty [len(L) == 0], the list returns 0 since no numbers can be added up

If  $\text{len}(L) = (k-1)$ ,  $L[0] \% 5 == 0$  (is evenly divisible by 5), bar(L) returns

$L[0]$  and recursively checks and adds  $L[0]$  and any other L member that fulfills the condition of " $L[0] \% 5 == 0$ "

else, bar(L) recursively checks  $L[1:]$  for c1

Assuming it works for (k-1), if  $\text{len}(L) = k$ , it will add all values that are divisible by 5, as explained above.

10. (10 points) The input of the following function is a binary tree T. Assume that such a tree T has 3 attributes: T.left, T.right (both of which are also binary trees), and T.color (which is a string of either "red", "green", or "blue"). If T is empty, the function *is\_empty*(T) returns True (if not, it returns False). Complete defining the following Python function so that it correctly counts the number of red nodes in the input binary tree T.

def count\_red\_nodes(T): recursion  
# fill in your codes below to count the number of red nodes in the binary tree T

```
totRed = 0
if is_empty(T) == True:
    return 0
if is_empty(T.left) == False:
    if T.left.color == "red":
        totRed = totRed + 1
    return totRed + count_red_nodes(T.left)
else if is_empty(T.right) == False:
    if T.right.color == "red":
        totRed = totRed + 1
    return totRed + count_red_nodes(T.right)
return totRed
```