

# Correctness of a Simple Compiler

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Based upon the example discussed by Professor Graham Hutton on Computerphile on November 27: [https://youtu.be/T\\_IINWzQhow](https://youtu.be/T_IINWzQhow)

Proof based upon Associate Professor Adam Chlipala's proof: <http://adam.chlipala.net/cpdt/html/StackMachine.html>

## 1 Implementation

Inductive **Expr** :=

| Val ( $n : \mathbf{Z}$ )  
| Add ( $e1\ e2 : \mathbf{Expr}$ ).

Inductive **Op** :=

| PUSH ( $n : \mathbf{Z}$ )  
| ADD.

Fixpoint eval ( $e : \mathbf{Expr}$ ) :=

match  $e$  with  
| Val  $n \Rightarrow n$   
| Add  $x\ y \Rightarrow \text{eval } x + \text{eval } y$   
end.

Fixpoint comp ( $e : \mathbf{Expr}$ ) :=

match  $e$  with  
| Val  $n \Rightarrow [\text{PUSH } n]$   
| Add  $x\ y \Rightarrow \text{comp } x ++ \text{comp } y ++ [\text{ADD}]$   
end.

Fixpoint exec ( $ops : \text{list Op}$ ) ( $stack : \text{list Z}$ ) : list Z :=

match  $ops, stack$  with  
| PUSH  $n :: c, s \Rightarrow \text{exec } c\ (n :: s)$   
| ADD ::  $c, (m :: n :: s) \Rightarrow \text{exec } c\ (n + m :: s)$   
| \_,  $s \Rightarrow s$   
end.

## 2 Examples

Example eg1 : exec (comp (Add (Val 42) (Val 42))) [] = [84]. reflexivity. Qed.

Example eg2 : eval (Add (Val 42) (Val 42)) = 84. reflexivity. Qed.

Example eg3 : exec ([PUSH 1; PUSH 2; PUSH 3]) [] = [3;2;1]. reflexivity. Qed.

Example eg4 : exec ([ADD; PUSH 1; PUSH 2; PUSH 3]) [] = []. reflexivity. Qed.

## 3 Correctness Proof

First we need to strengthening the induction hypothesis as described in: Adam Chlipala's similar proof at <http://adam.chlipala.net/cpdt/html/StackMachine.html>

Lemma correct\_helper :

$\forall e \text{ ops } s, \text{exec (comp } e \text{ ++ ops) } s = \text{exec ops (eval } e \text{ :: } s).$

Proof.

induction e.

- simpl. reflexivity.

- intros. simpl.

rewrite app\_assoc\_reverse.

rewrite IHe1.

rewrite app\_assoc\_reverse.

rewrite IHe2.

simpl. reflexivity.

Qed.

Now the proof follows from the lemma.

Theorem correct :  $\forall e, \text{exec (comp } e) [] = [\text{eval } e].$

Proof.

intros.

pose (correct\_helper e [] []) as H.

rewrite app\_nil\_r in H.

assumption.

Qed.

QED! - We now have the highest degree of certainty (proof) that our implementation meets the specification:  $\forall e, \text{exec (comp } e) [] = [\text{eval } e]$

The Coq source file of this proof is available at: [https://coda-coda.github.io/blog/program\\_correctness\\_based\\_on\\_computerphile/coq\\_simple\\_compilation\\_correctness.v](https://coda-coda.github.io/blog/program_correctness_based_on_computerphile/coq_simple_compilation_correctness.v)