

Cody Kesler

CS 312 – 2

Lab 3

Code:

```
from CS312Graph import *
import time
import sys

class NetworkRoutingSolver:

    def __init__( self, display ):
        pass

    def initializeNetwork( self, network ):
        assert( type(network) == CS312Graph )
        self.network = network

// O(n) – worst case. This loops through one path of nodes(could visit each node)
    def getShortestPath( self, dest_index ):
        self.dest = dest_index
        path_edges = []
        total_length = 0

        self.cur_node = self.queue.nodes[self.dest]
        while(self.cur_node.node_id != self.source):
            cur_loc = self.cur_node.loc
            prev = self.queue.get_prev_node(self.cur_node)

            if(prev != None):
                prev_loc = prev.loc

                else:
                    break

            length = self.queue.get_prev_length(self.cur_node)
            path_edges.append((cur_loc, prev_loc, '{:.0f}'.format(length)))
            self.cur_node = self.queue.get_prev_node(self.cur_node)
            total_length += length

        return {'cost':total_length, 'path':path_edges}

// Thus total is O((n + e)*log(n))
    def computeShortestPaths( self, srcIndex, use_heap=False ):
        self.source = srcIndex
```

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```
t1 = time.time()

if(use_heap):
    //O(n)

    self.queue = Heap(self.network, srcIndex)
else:
    //O(n)

    self.queue = Array(self.network, srcIndex)

//Repeats edges times
while(not self.queue.queue_empty()):
    lowest_node = self.queue.delete_min()
    if(lowest_node == -1):
        break

    if(lowest_node.node_id != 0):
        // O(n) Ends up being n since reaches total for all the nodes
        self.queue.push_neighbors_on_queue(lowest_node)

    neighbors = lowest_node.neighbors

    // O(E) since it visits all the edges for each nodes
    for i in range(0, len(neighbors)):
        if(self.queue.get_distance_to(neighbors[i].dest) >
            self.queue.get_distance_to(lowest_node) + neighbors[i].length):
            self.queue.set_distance_to(neighbors[i].dest,
                self.queue.get_distance_to(lowest_node) + neighbors[i].length)
            self.queue.set_prev_node(neighbors[i].dest, lowest_node)
            self.queue.decrease_key()
```

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```
t2 = time.time()

    return (t2-t1)
class Heap:
//O(n) This only calls insert and makes lists the size of the nodes in the graph
    def __init__(self, graph, src):
        def __init__(self, graph, src):
            self.nodes = graph.nodes
            self.queue = list()
            self.queue_size = 0
            self.node_dist = [sys.maxsize-1]*len(self.nodes)
            self.prev_node = [None]*len(self.nodes)
            self.popped_nodes = list()

            self.insert(CS312GraphEdge(self.nodes[src], self.nodes[src], 0))
            self.node_dist[src] = 0
            self.prev_node[src] = self.nodes[src]
            self.push_neighbors_on_queue(self.nodes[src])

        pass

//O(log(n)) This function gets called on just the 3 neighbors of the function and calls sift up for each of them
    def push_neighbors_on_queue(self, node):
        for i in range(len(node.neighbors)):
            if (node.neighbors[i].dest.node_id not in self.popped_nodes):
                self.insert(node.neighbors[i])
                self.sift_up(self.queue_size-1)
//O(1)
    def insert(self, edge):
        self.queue.append(edge)
        self.queue_size = self.queue_size + 1

//O(Log(n) It repeats the height of the tree which is log(n)
    def sift_up(self, i):
        parent = (i-1) // 2
        while i != 0 and self.node_dist[self.queue[i].dest.node_id] <
```

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```
        self.node_dist[self.queue[parent].dest.node_id]:

    temp = self.queue[parent]
    self.queue[parent] = self.queue[i]
    self.queue[i] = temp
    i = parent
    parent = (i - 1) // 2

//O(Log(n) It uses the sift down function which is Log(n)

def delete_min(self):
    if (self.queue_empty()):
        return -1

    return_node = self.queue[0].dest
    self.queue[0] = self.queue[-1]
    self.queue.pop()
    self.popped_nodes.append(return_node.node_id)
    self.queue_size = self.queue_size - 1
    self.sift_down(0)

    return return_node

// O(log(n) This is Log(n) as it starts as the vertex of the tree and only visits its children until it hits the bottom which is
Log(n) levels thus O(log(n)

def sift_down(self, i):
    if(not self.queue_empty()):
        while (i * 2) < len(self.queue):
            min_child = self.min_child(i)
            if(min_child == -1):
                break
            if (self.node_dist[self.queue[i].dest.node_id] >
                self.node_dist[self.queue[min_child].dest.node_id]):
                temp = self.queue[i]
                self.queue[i] = self.queue[min_child]
                self.queue[min_child] = temp
```

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```
i = min_child

//This function just looks at the two children and returns its lowest value child. O(1)
def min_child(self, i):
    left_child = i * 2 + 1
    right_child = i * 2 + 2
    if left_child >= len(self.queue):
        return -1
    elif right_child >= len(self.queue):
        return left_child
    else:
        if self.node_dist[self.queue[left_child].dest.node_id] <
            self.node_dist[self.queue[right_child].dest.node_id]:
            return left_child
        else:
            return right_child

//O(1)
def queue_empty(self):
    if (len(self.queue) > 0):
        return False
    else:
        return True

//O(1)
def get_distance_to(self, node):
    return self.node_dist[node.node_id]

//O(1)
def get_prev_length(self, node):
    if (self.prev_node[node.node_id] != None):
        return self.node_dist[node.node_id] -
            self.node_dist[self.prev_node[node.node_id]]
    else:
        return 0

//O(1)
def get_prev_node(self, node):
    if (self.prev_node[node.node_id] != None):
```

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```
        return self.nodes[self.prev_node[node.node_id]]
//O(1)
    def set_distance_to(self, node, distance):
        self.node_dist[node.node_id] = distance
//O(1)
    def set_prev_node(self, node, prev_node):
        self.prev_node[node.node_id] = prev_node.node_id
//O(1)
    def decrease_key(self):
        return True

//O(1)
class Array:
//O(n) This only calls insert and makes lists the size of the nodes in the graph
    def __init__(self, graph, src):
        self.graph = graph
        self.nodes = graph.getNodes()
        self.node_dist = list()
        self.prev_node = list()
        self.popped_nodes = list()
        self.queue = list()
        nodes = graph.getNodes()

// This for loop is O(n) as it repeats for the number of nodes
        for i in range(len(nodes)):
            if(nodes[i].node_id == src):
                self.queue.append(CS312GraphEdge(nodes[i], nodes[i], 0))
                self.push_neighbors_on_queue(nodes[i])
                self.node_dist.append(0)
                self.prev_node.append(src)
            else:
                self.node_dist.append(sys.maxsize-1)
                self.prev_node.append(None)

// This is O(n) total
    def push_neighbors_on_queue(self, node):
```

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```
        for i in range(0, len(node.neighbors)):
            if (node.neighbors[i].dest.node_id not in self.popped_nodes):
                self.queue.append(node.neighbors[i])

// O(1)
def queue_empty(self):
    if (len(self.queue) > 0):
        return False
    else:
        return True

// O(
def delete_min(self):
    temp_highest = sys.maxsize
    node_index = -1
    queue_index = -1
    for i in range(len(self.queue)):
        if (self.node_dist[self.queue[i].dest.node_id] < temp_highest
            and self.queue[i].dest.node_id not in self.popped_nodes):
            temp_highest = self.node_dist[self.queue[i].dest.node_id]
            node_index = self.queue[i].dest.node_id
            queue_index = i

    if (node_index == -1):
        return -1
    else:
        return_node = self.queue[queue_index].dest
        self.queue.pop(queue_index)
        self.popped_nodes.append(return_node.node_id)
        return return_node

def get_distance_to(self, node):
    return self.node_dist[node.node_id]

def set_distance_to(self, node, distance):
    self.node_dist[node.node_id] = distance
```

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```
def get_prev_node(self, node):
    if (self.prev_node[node.node_id] != None):
        return self.nodes[self.prev_node[node.node_id]]

def set_prev_node(self, node, prev_node):
    self.prev_node[node.node_id] = prev_node.node_id

def get_prev_length(self, node):
    if (self.prev_node[node.node_id] != None):
        return self.node_dist[node.node_id] - self.node_dist[self.prev_node[node.node_id]]
    else:
        return 0

def decrease_key(self):
    return True
```

Complexity Analysis:

Array: Making the queue is (n) . Inserting is $O(1)$. Deleting the Min is $O(n)$. Decreasing the Key doesn't do much

Thus the total time complexity is $n + n + n(n) + e(1) = O(n^2 + e)$

Heap: Making the queue is n . Inserting comes to $\log(n)$. Deleting in is $\log(n)$. Decreasing key runs at $\log(n)$

Thus the total time complexity is $n + \log(n) + n*\log(n) + e\log(n) = O((n + e)*\log(n))$

Drawing the shortest path yields $O(n)$ added to each of the implementations.

So it is either

$O((n + e)*\log(n) + n)$ Or $O(n^2 + e + n)$

Space Complexity:

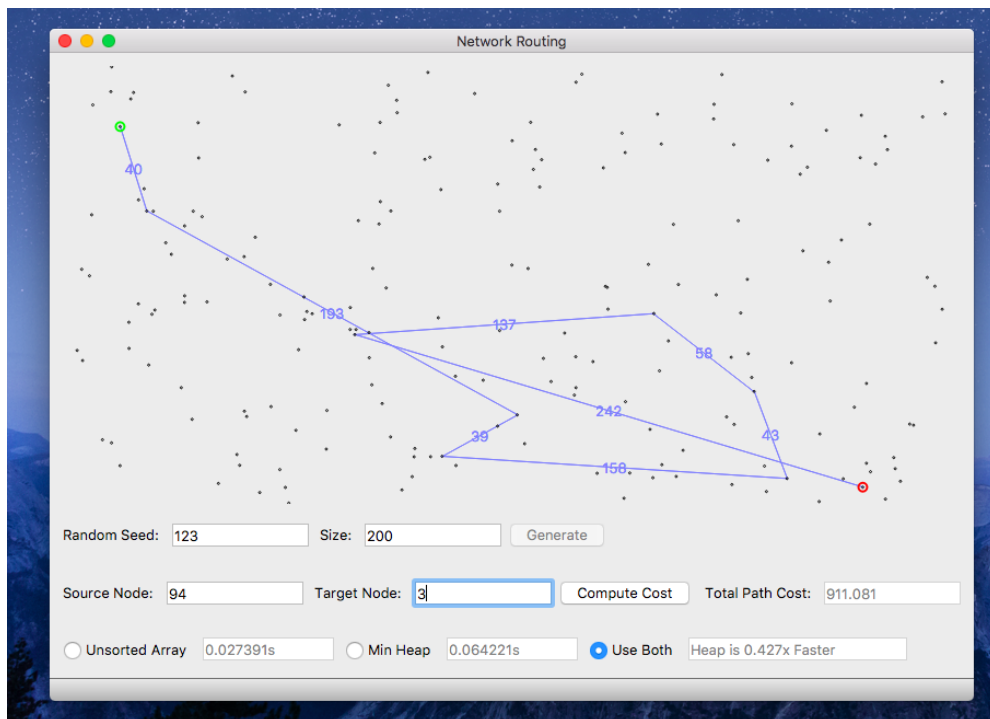
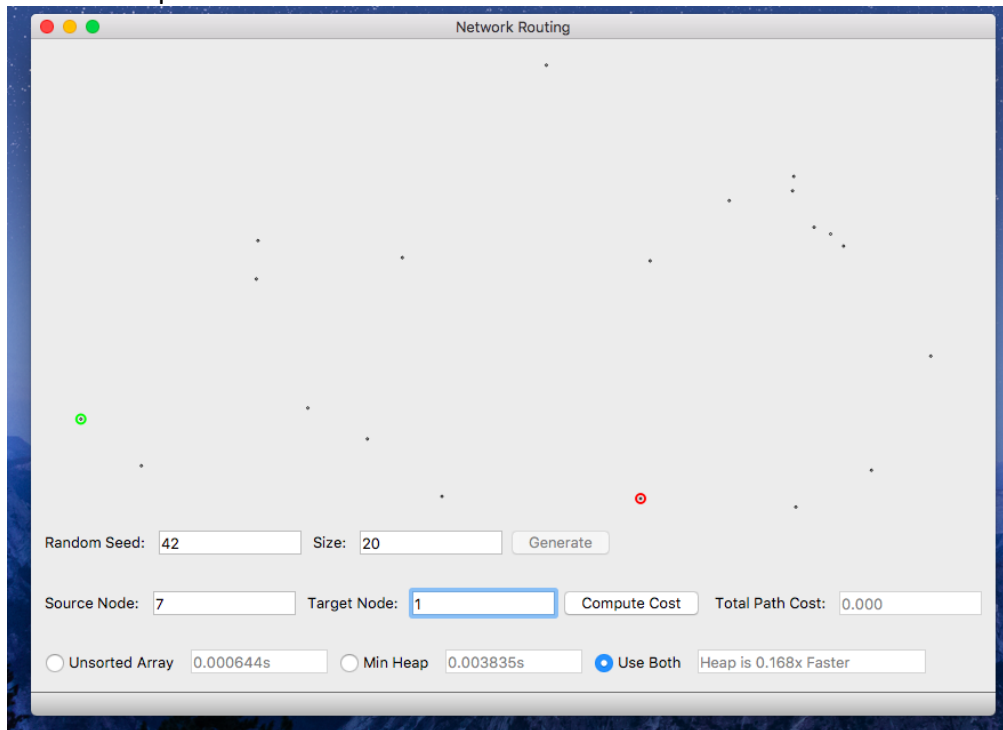
The algorithm keeps 2 arrays of size n for total distances and previous nodes thus $O(2n)$

Array: The array implementation stores another 3 arrays of size n for the graph, nodes, and popped nodes. Thus total space is $O(5n)$

Heap: The heap implementation also stores another 3 arrays. Thus is the same space complexity $O(5n)$.

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Screenshots:

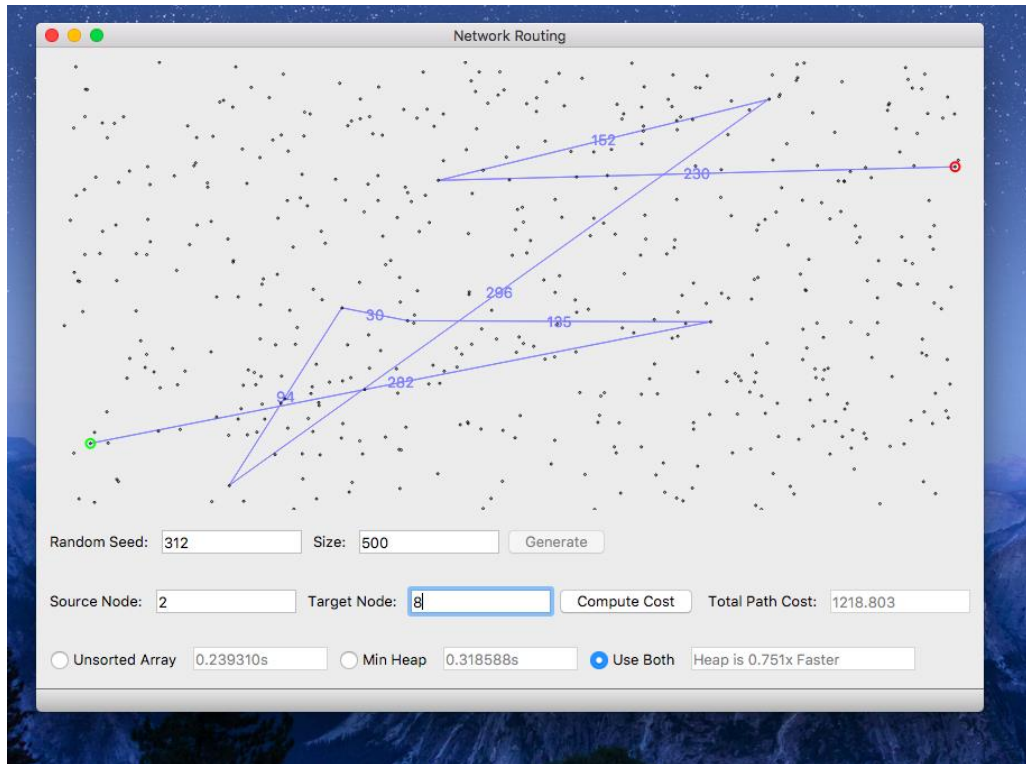
There is no path.



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Graph and Chart:

Array: Seed: 312		Trial1	Trial 2	Trial 3	Trial 4	Trial5	Avg
	100	0.009675	0.0078	0.00789	0.00743	0.00833	0.008225
	1000	1.658	1.903	1.734	1.674	1.926	1.779
	10000	1798.54	1711.89	1726.45	1753.27	1705.66	1739.162
	100000						
	1000000						

Heap: Seed: 312		Trial1	Trial 2	Trial 3	Trial 4	Trial5	Avg
	100	0.01977	0.1006	0.0207	0.0375	0.0239	0.040494
	1000	0.856	0.8695	0.8496	0.8471	1.006	0.88564
	10000	86.68	86.45	86.77	86.34	86.1	86.468
	100000						
	1000000						

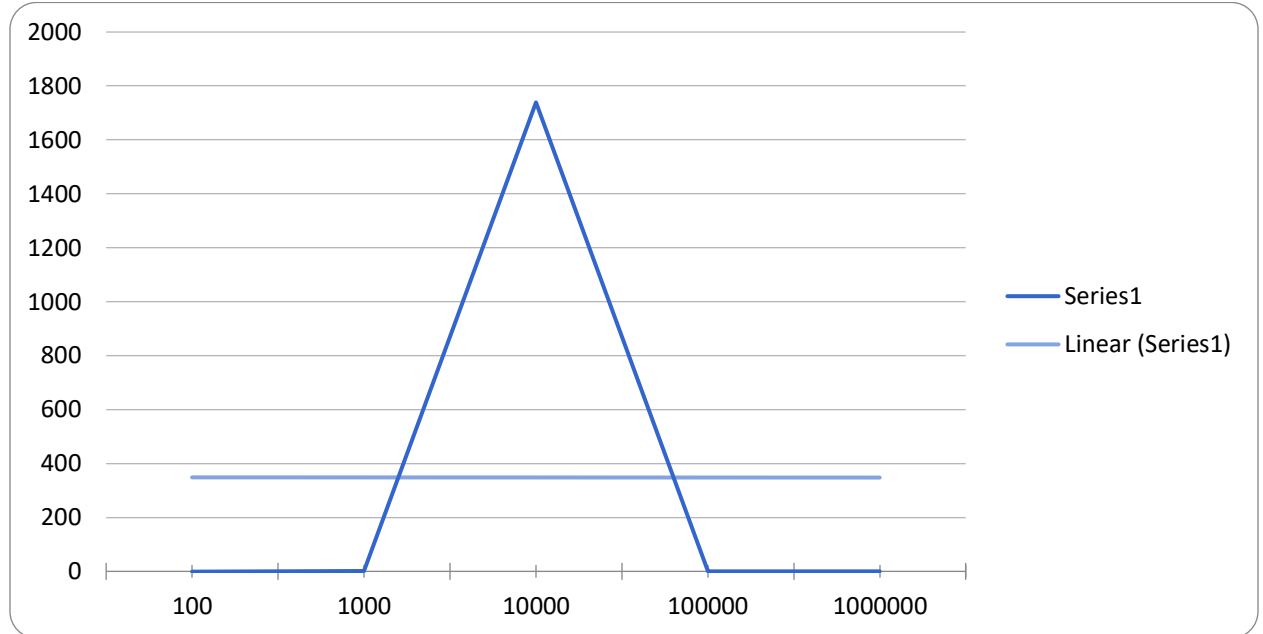
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For some reason the program wouldn't run in an efficient manner on my laptop when running the higher node counts so I couldn't complete the time trials. But these are the graphs projecting the higher node count run times:

Array:



Heap:

