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Department of Computer Engineering

Course - System Programming and Compiler Construction (SPCC)

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Class and Batch	TE Computer Engineering - Batch C	
Date	22-04-2024	
Lab #	8	
Aim	Design a two-pass assembler.	
Objective	Implement a two-pass assembler to enhance assembly code translation, achieving robustness, efficiency, and error detection in programming.	
Theory	Introduction of Assembler	
	Assembler is a program for converting instructions written in low-level assembly code into relocatable machine code and generating along information for the loader.	
	It is necessary to convert user written programs into a machinery code. This is called as translation of the high level language to low level that is machinery language. This type of translation is performed with the help of system software. Assembler can be defined as a program that translates an assembly language program into a machine language program. Self assembler is a program that runs on a computer and produces the machine codes for the same computer or same machine. It is also known as resident assembler. A cross assembler is an assembler which runs on a computer and produces the machine codes for other computer.	
	Assembler Machine code	

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It generates instructions by evaluating the mnemonics (symbols) in operation field and find the value of symbol and literals to produce machine code. Now, if assembler do all this work in one scan then it is called single pass assembler, otherwise if it does in multiple scans then called multiple pass assembler. Here assembler divide these tasks in two passes:

• Pass-1:

- 1. Define symbols and literals and remember them in symbol table and literal table respectively.
- 2. Keep track of location counter
- 3. Process pseudo-operations
- 4. Defines program that assigns the memory addresses to the variables and translates the source code into machine code

• Pass-2:

- 1. Generate object code by converting symbolic op-code into respective numeric op-code
- 2. Generate data for literals and look for values of symbols
- 3. Defines program which reads the source code two times
- 4. It reads the source code and translates the code into object code.

Firstly, We will take a small assembly language program to understand the working in their respective passes. Assembly language statement format:

[Label] [Opcode] [operand]

Example: M ADD R1, ='3' where, M - Label; ADD - symbolic opcode; R1 - symbolic register operand; (='3') - Literal

Assembly Program:

```
Label Op-code operand LC value(Location counter)
JOHN START
               200
   MOVER
             R1, ='3'
                     200
   MOVEM
             R1, X
                      201
    MOVER
              R2, = '2' 202
L1
   LTORG
                  203
X
    DS
          1
                 204
```



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END 205

Let's take a look on how this program is working:

- 1. **START:** This instruction starts the execution of program from location 200 and label with START provides name for the program.(JOHN is name for program)
- 2. **MOVER:** It moves the content of literal(='3') into register operand R1.
- 3. **MOVEM:** It moves the content of register into memory operand(X).
- 4. **MOVER:** It again moves the content of literal(='2') into register operand R2 and its label is specified as L1.
- 5. LTORG: It assigns address to literals(current LC value).
- 6. **DS(Data Space):** It assigns a data space of 1 to Symbol X.
- 7. **END:** It finishes the program execution.

Working of Pass-1:

Define Symbol and literal table with their addresses. Note: Literal address is specified by LTORG or END.

Step-1: START 200

(here no symbol or literal is found so both table would be empty)

Step-2: MOVER R1, ='3' 200

(='3' is a literal so literal table is made)

Literal	Address		
='3'			



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Step-3: MOVEM R1, X 201

X is a symbol referred prior to its declaration so it is stored in symbol table with blank address field.

Symbol	Address
X	

Step-4: L1 MOVER R2, ='2' 202

L1 is a label and ='2' is a literal so store them in respective tables

Symbol	Address
X	
L1	202
Literal	Address
='3'	
='2'	

Step-5: LTORG 203

Assign address to first literal specified by LC value, i.e., 203

Literal	Address
='3'	203
='2'	



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Step-6: X DS 1 204

It is a data declaration statement i.e X is assigned data space of 1. But X is a symbol which was referred earlier in step 3 and defined in step 6. This condition is called Forward Reference Problem where variable is referred prior to its declaration and can be solved by back-patching. So now assembler will assign X the address specified by LC value of current step.

Symbol	Address
X	204
L1	202

Step-7: END 205

Program finishes execution and remaining literal will get address specified by LC value of END instruction. Here is the complete symbol and literal table made by pass 1 of assembler.

Symbol	Address		
X	204		
L1	202		
Literal	Address		
='3'	203		

Now tables generated by pass 1 along with their LC value will go to pass-2 of assembler for further processing of pseudo-opcodes and machine op-codes.



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Working of Pass-2:

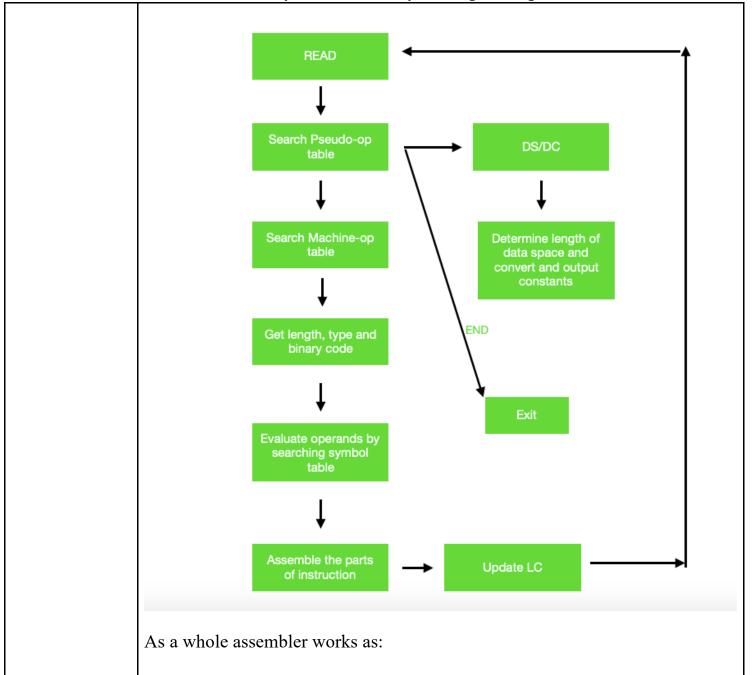
Pass-2 of assembler generates machine code by converting symbolic machine-opcodes into their respective bit configuration(machine understandable form). It stores all machine-opcodes in MOT table (op-code table) with symbolic code, their length and their bit configuration. It will also process pseudo-ops and will store them in POT table(pseudo-op table). Various Data bases required by pass-2:

- 1. MOT table(machine opcode table)
- 2. POT table(pseudo opcode table)
- 3. Base table(storing value of base register)
- 4. LC (location counter)

Take a look at flowchart to understand:



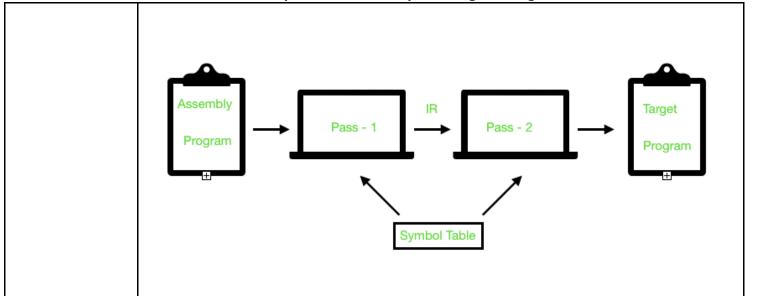
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Implementation/Code

```
import sys
def RemoveSpaces(x):
        return x
def RemoveCommas(x):
    if x[-1] == ",":
        return x[: len(x) - 1]
    else:
        return x
def CheckLiteral(element):
    if element[ : 2] == "='":
        return True
    else:
        return False
def CheckSymbol(Elements):
    global SymbolTable, Opcodes
    if (len(Elements) > 1) and ([Elements[-1], None, None, "Variable"] not in
SymbolTable) and (Elements[-1] != "CLA") and (Elements[-2] not in ["BRP", "BRN",
"BRZ"]) and (Elements[-1][ : 2] != "='") and (Elements[-1][ : 3] != "REG") and
(not Elements[-1].isnumeric()):
        return True
    else:
        return False
def CheckLabel(Elements):
```



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```
global SymbolTable, Opcodes
    if (len(Elements) >= 2) and (Elements[1] in Opcodes):
        if Elements[0] not in SymbolTable:
            return True
    else:
        return False
Opcodes = ["CLA", "LAC", "SAC", "ADD", "SUB", "BRZ", "BRN", "BRP", "INP", "DSP",
"MUL", "DIV", "STP", "DATA", "START"]
AssemblyOpcodes = {"CLA" : "0000", "LAC" : "0001", "SAC" : "0010", "ADD" : "0011",
"SUB" : "0100", "BRZ" : "0101", "BRN" : "0110",
                   "BRP": "0111", "INP": "1000", "DSP": "1001", "MUL": "1010",
"DIV" : "1011", "STP" : "1100"}
SymbolTable = []
LiteralTable = []
Variables = []
Declarations = []
AssemblyCode = []
location_counter = 0
stop_found = False
end found = False
file = open("/Users/pranaysinghvi/Library/CloudStorage/OneDrive-Personal/SPIT
College/3)Class/Semester 6/5)SPCC/1)Experiment/8 /Assembly Code Input.txt", "rt")
# ERROR 1 : Checking for missing START statement
for line in file:
    # Checking for comments
    if line[ : 2] != "//":
        if line.strip() != "START":
            print("STARTError : 'START' statement is missing. " + "( Line " +
str(location_counter) + " )")
            sys.exit(0)
        else:
            file.seek(0, 0)
            break
# First Pass
for line in file:
    if line[ : 2] != "//":
        Elements = line.strip().split(" ")
        Elements = list(filter(RemoveSpaces, Elements))
        Elements = list(map(RemoveCommas, Elements))
```



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```
for i in range(len(Elements)):
            if Elements[i][ : 2] == "//":
                Elements = Elements[ : i]
                break
       # ERROR 2 : Checking for too many operands
       # If the instruction doesn't contain a Label
        if (len(Elements) >= 3) and (Elements[0] in Opcodes):
            print("TooManyOperandsError : Too many operands used for the '" +
Elements[0] + "' assembly opcode. " + "( Line " + str(location_counter) + " )")
            sys.exit(0)
       # If the instruction contains a Label
       elif (len(Elements) >= 4) and (Elements[1] in Opcodes):
            print("TooManyOperandsError : Too many operands used for the '" +
Elements[1] + "' assembly opcode. " + "( Line " + str(location_counter) + " )")
            sys.exit(0)
        if (len(Elements) == 1) and (Elements[0] in ["LAC", "SAC", "ADD", "SUB",
"BRZ", "BRN", "BRP", "INP", "DSP", "MUL", "DIV"]):
            print("LessOperandsError : Less operands used for the '" + Elements[0]
+ "' assembly opcode. " + "( Line " + str(location_counter) + " )")
            sys.exit(0)
       # If the instruction contains a Label
       elif (len(Elements) == 2) and (Elements[1] in ["LAC", "SAC", "ADD", "SUB",
"BRZ", "BRN", "BRP", "INP", "DSP", "MUL", "DIV"]):
            print("LessOperandsError : Less operands used for the '" + Elements[1]
+ "' assembly opcode. " + "( Line " + str(location_counter) + " )")
            sys.exit(0)
       # ERROR 4 : Checking for invalid opcodes
        if stop_found is False:
            if len(Elements) == 3:
                # If the instruction contains a Label
                if Elements[1] not in Opcodes:
                    print("InvalidOpcodeError : '" + Elements[1] + "' is an
invalid opcode. " + "( Line " + str(location_counter) + " )")
                    sys.exit(0)
            if (len(Elements) == 2) and (Elements[1] == "CLA"):
            elif len(Elements) == 2:
                # If the instruction doesn't contain a Label
                if Elements[0] not in Opcodes:
                    print("InvalidOpcodeError : '" + Elements[0] + "' is an
invalid opcode. " + "( Line " + str(location_counter) + " )")
                  sys.exit(0)
```



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```
# Check for STP
        if (len(Elements) == 3) and (Elements[1] == "DATA"):
            stop_found = True
        # Check for END
        if (len(Elements) == 1) and (Elements[0] == "END"):
            end_found = True
            for i in range(len(LiteralTable)):
                if LiteralTable[i][1] == -1:
                    LiteralTable[i][1] = location_counter
                    location_counter += 1
            break
        if not stop_found:
            # Check for Literal
            for x in Elements:
                if CheckLiteral(x):
                    LiteralTable.append([x, -1])
            # Check for Labels
            if CheckLabel(Elements):
                SymbolTable.append([Elements[0], location_counter, None, "Label"])
            # Check for Symbols
            if CheckSymbol(Elements):
                SymbolTable.append([Elements[-1], None, None, "Variable"])
        elif stop_found:
            if (Elements[0] != "STP") and (Elements[0] != "END"):
                # ERROR 5 : Checking for multiple definations
                if Elements[0] not in Variables:
                    Variables.append(Elements[0])
                    Declarations.append((Elements[0], Elements[2]))
                else:
                    print("DefinationError : Variable '" + Elements[0] + "'
defined multiple times. " + "( Line " + str(location_counter) + " )")
                    sys.exit(0)
                # ERROR 6 : Checking for redundant declarations
                if [Elements[0], None, None, "Variable"] not in SymbolTable:
                    print("RedundantDeclarationError : " + Elements[0] + "
declared but not used.")
                    sys.exit(0)
                location = SymbolTable.index([Elements[0], None, None,
"Variable"])
```



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```
SymbolTable[location][1] = location counter
               SymbolTable[location][2] = Elements[2]
       location_counter += 1
# ERROR 7 : Checking for missing END statement
if end found is False:
   print("ENDError : 'END' statement is missing." + "( Line " +
str(location_counter) + " )")
   sys.exit(0)
# ERROR 8 : Checking for undefined variables
for x in SymbolTable:
   if x[1] is None and x[3] == "Variable":
       print("UndefinedVariableError : Variable '" + x[0] + "' not defined.")
       sys.exit(0)
# Printing Tables after First Pass
print(">>> Opcode Table <<<\n")</pre>
print("----")
for key in AssemblyOpcodes:
   print(key.ljust(20) + AssemblyOpcodes[key].ljust(6))
print("-----")
print("\n>>> Literal Table <<<\n")</pre>
print("LITERAL ADDRESS")
print("----")
for i in LiteralTable:
   print(i[0].ljust(12) + str(i[1]).ljust(7))
print("----")
print("\n>>> Symbol Table <<<\n")</pre>
print("SYMBOL ADDRESS VALUE TYPE")
print("-----
for i in SymbolTable:
   print(i[0].ljust(16) + str(i[1]).ljust(12) + str(i[2]).ljust(10) +
i[3].ljust(10))
print("-----
print("\n>>> Data Table <<<\n")</pre>
print("VARIABLES VALUE")
print("----")
```



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```
for i in Declarations:
    print(i[0].ljust(14) + str(i[1]).ljust(10))
print("----\n")
# Second Pass
file.seek(0, 0)
print(">>> MACHINE CODE <<<\n")</pre>
for line in file:
    if line[ : 2] != "//":
        Elements = line.strip().split(" ")
        Elements = list(filter(RemoveSpaces, Elements))
        Elements = list(map(RemoveCommas, Elements))
        s = ""
        for i in range(len(Elements)):
            if Elements[i][ : 2] == "//":
                Elements = Elements[: i]
                break
        # To terminate machine code conversion
        if (len(Elements) == 3) and (Elements[1] == "DATA"):
            break
        if Elements[0] == "STP":
            AssemblyCode.append("00 "+ AssemblyOpcodes["STP"] + " 00 00 00")
            print("00 " + AssemblyOpcodes["STP"] + " 00 00 00")
        # If the CLA opcode has a Label before it
        elif (len(Elements) == 2) and (Elements[1] == "CLA"):
            for i in range(len(SymbolTable)):
                if Elements[0] == SymbolTable[i][0]:
                    AssemblyCode.append(str(SymbolTable[i][1]).rjust(2, "0") + " "
+ AssemblyOpcodes["CLA"] + " 00 00 00")
                    print(str(SymbolTable[i][1]).rjust(2, "0") + " "+
AssemblyOpcodes["CLA"] + " 00 00 00")
        elif Elements[0] != "START":
            if (len(Elements) == 1) and (Elements[0] == "CLA"):
                AssemblyCode.append("00 " + AssemblyOpcodes["CLA"] + " 00 00 00")
                print("00 " + AssemblyOpcodes["CLA"] + " 00 00 00")
            # If there is no Label
            elif (len(Elements) == 2) and (Elements[1] != "CLA"):
                print("00 " + AssemblyOpcodes[Elements[0]], end = " ")
                s = "00 " + AssemblyOpcodes[Elements[0]] + " "
```



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```
# Dealing with Literals
                if CheckLiteral(Elements[1]):
                    for i in range(len(LiteralTable)):
                        if LiteralTable[i][0] == Elements[1]:
                            AssemblyCode.append(s + "00 00 " +
str(LiteralTable[i][1]).rjust(2, "0"))
                            print("00 00 " + str(LiteralTable[i][1]).rjust(2,
"0"))
                # Dealing with Lables (BRP, BRZ, BRN)
                elif Elements[0] in ["BRP", "BRN", "BRZ"]:
                    for i in range(len(SymbolTable)):
                        if SymbolTable[i][0] == Elements[1]:
                            AssemblyCode.append(s +
str(SymbolTable[i][1]).rjust(2, "0") + " 00 00")
                            print(str(SymbolTable[i][1]).rjust(2, "0") + " 00 00")
                # Dealing with Registers
                elif Elements[1][ : 3] == "REG":
                    AssemblyCode.append(s + "00 " + Elements[1][-1].rjust(2, "0")
+ " 00")
                    print("00 " + Elements[1][-1].rjust(2, "0") + " 00")
                # Dealing with Variables
                else:
                    for i in range(len(SymbolTable)):
                        if SymbolTable[i][0] == Elements[1]:
                            AssemblyCode.append(s + "00 00 " +
str(SymbolTable[i][1]).rjust(2, "0"))
                            print("00 00 " + str(SymbolTable[i][1]).rjust(2, "0"))
            # If the instruction conatins a Label
            elif len(Elements) == 3:
                for i in range(len(SymbolTable)):
                        if SymbolTable[i][0] == Elements[0]:
                            print(str(SymbolTable[i][1]).rjust(2, "0") + " " +
AssemblyOpcodes[Elements[1]], end = " ")
                            s = str(SymbolTable[i][1]).rjust(2, "0") + " " +
AssemblyOpcodes[Elements[1]] + " "
                # Dealing with Literals
                if CheckLiteral(Elements[2]):
                    for i in range(len(LiteralTable)):
                        if LiteralTable[i][0] == Elements[2]:
                            AssemblyCode.append(s + "00 00 " +
str(LiteralTable[i][1]).rjust(2, "0"))
                            print("00 00 " + str(LiteralTable[i][1]).rjust(2,
"0"))
                # Dealing with Lables (BRP, BRZ, BRN)
                elif Elements[1] in ["BRP", "BRN", "BRZ"]:
```



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```
for i in range(len(SymbolTable)):
                        if SymbolTable[i][0] == Elements[2]:
                           AssemblyCode.append(s +
str(SymbolTable[i][1]).rjust(2, "0") + " 00 00")
                           print(str(SymbolTable[i][1]).rjust(2, "0") + " 00 00")
                elif Elements[2][ : 3] == "REG":
                   AssemblyCode.append(s + "00 " + Elements[2][-1].rjust(2, "0")
+ " 00")
                   print("00 " + Elements[2][-1].rjust(2, "0") + " 00")
               # Dealing with Variables
               else:
                   for i in range(len(SymbolTable)):
                        if SymbolTable[i][0] == Elements[2]:
                           AssemblyCode.append(s + "00 00 " +
str(SymbolTable[i][1]).rjust(2, "0"))
                           print("00 00 " + str(SymbolTable[i][1]).rjust(2, "0"))
file.close()
file = open("./Machine Code.txt", "x")
file.write("----\nMACHINE CODE\n----\n\n")
for x in AssemblyCode:
    file.write(x + "\n")
file.close()
```



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Output	pranaysinghvi ersonal/SPIT >>> Opcode Ta	College/3)Cla	Book-Air 5 ass/Semest)SPCC % /opt/home er 6/5)SPCC/1)Exp	brew/bin/pyt eriment/8_/E
	ASSEMBLY OPCO	DE OPCODE	Ē		
	CLA LAC SAC ADD SUB BRZ BRN BRP INP DSP MUL DIV STP	0000 0001 0010 0011 0100 0101 0111 1000 1001 1010 1011 1100	-		
	LITERAL A	DDRESS			
	='1' 2 ='35' 2 ='5' 3 ='600' 3	9 0			
	>>> Symbol Ta	ble <<<			
	SYMB0L	ADDRESS	VALUE	TYPE	
	LoopOne X A Loop Subtraction B C D Division E Zero	1 27 22 5 6 23 24 25 12 26 16	None 0 250 None 125 90 88 None 5	Label Variable Variable Label Variable Variable Variable Variable Label Label	
1	Positive	19 	None None	Label Label	



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>>> Data Ta	ble <<<
VARIABLES	VALUE
A	250
В	125
C	90
D	88
E	5
X	0
>>> MACHINE	E CODE <<<
01 0000 00	00 27
00 0001 00	00 22
00 0011 00	00 28
00 0100 00	00 29
05 0111 06	00 00
06 0100 00	00 30
00 0011 00	00 23
00 1010 00	00 24
00 0100 00	00 25
00 1010 00	00 31
00 0101 12	1011 00 00 26
00 0000 00	00 00
00 0001 00	01 00
00 0111 19	00 00
16 0010 00	00 27
00 1001 00	00 27
00 1100 00	00 00
19 0000 00	00 00
00 1001 00	01 00
00 1001 00	02 00
<pre>pranaysingh</pre>	vi@Pranays-MacBoo



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Conclusion	In conclusion, I successfully developed a two-pass assembler, ensuring accurate conversion of assembly code to machine code with improved efficiency and error detection capabilities, enhancing programming reliability and productivity.
References	[1] Geeksforgeeks (25 Sep, 2023) Introduction of Assembler https://www.geeksforgeeks.org/introduction-of-assembler/