

s.size() or s.length()

String Concatenation

```
string s1 = "new";  
string s2 = "delhi";
```

s1. append(s2)

2004 LL 81; // New Delhi
2004 LL 82; // Delhi

```
char s1[] = "new";
char s2[] = "delhi";
cout << s1 + s2;
           error
```

```
string s1 = "new";
string s2 = "delhi";
cout << s1 + s2; //newdelhi
```

String $s3 = s1 + s2$;
 ↳ new delhi"
 cout << s3 ;
 "new"
 cout << s1 ;
 "delhi"
 cout << s2 ;

$S_1 = S_1 + S_2$;
 $\hookrightarrow S_1$ is modified

string s = "hello"
 ↑
 s[0] or s.front()

0 1 2 3 4 5

n	e	e	0	0	10
---	---	---	---	---	----

↑

int n = s.size(); // 5
low + 1 < s[n-1]; // 0

```
cout << s.back(); // 0
```

String Comparison

string s1 = "abc"; s2 > s1
string s2 = "adc";

$$98 - 100 = \underline{\underline{-2}}$$

- * String obj i.e var. of String-type
can be compared using
Relational operators

S1: compare (s2)

```
graph TD
    A["S1: compare (s2)"] --> B["0  
equal"]
    A --> C["+ve  
 $s1 > s2$ "]
    A --> D["-ve  
 $s1 < s2$ "]
```

Searching in a String

⇒ string :: npos

"find" find
 find
 0 1 2 3 4 5 6 7 8 9
 b c d a b c d a b c ;
 find
 8 find

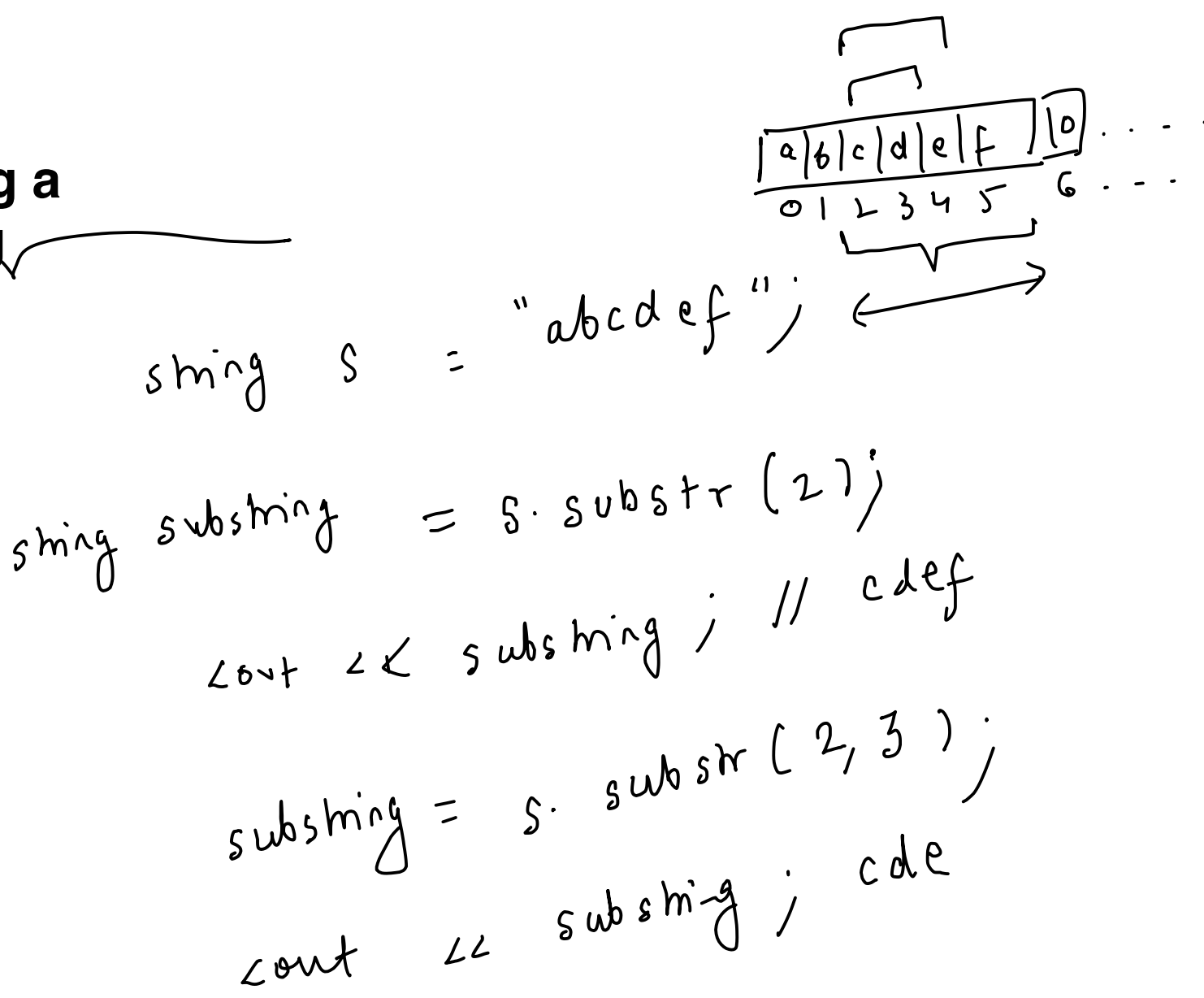
```

string subString = "abc"
if ( s.find ( subString )
    // subString is found
    != string::npos ) {

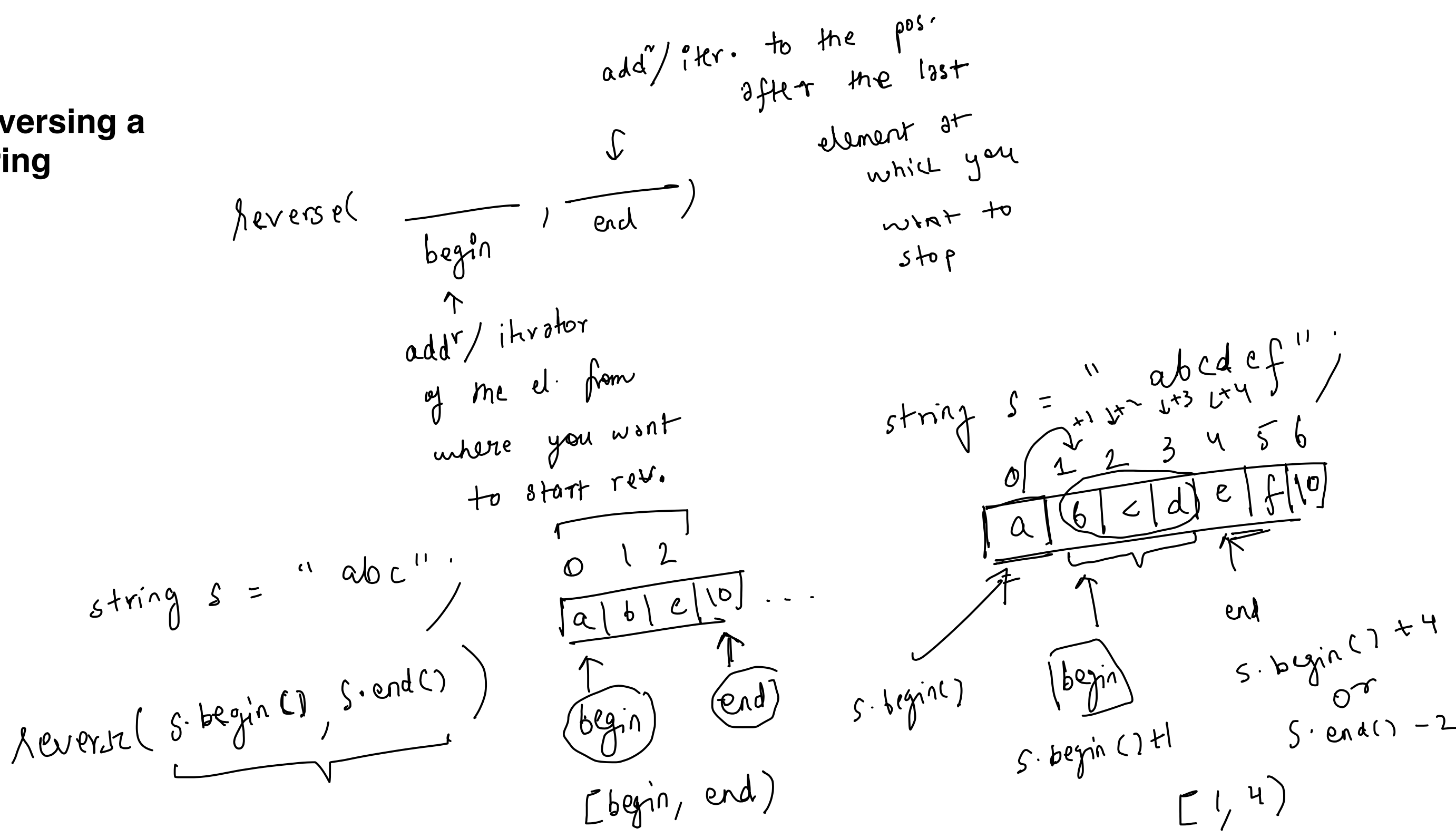
```

\rightarrow unsigned int
 $[0, \frac{\cdot}{\uparrow} \text{npos}]$
size_t
 $\equiv_{(int)}^{-2^{b1} \text{ to } 2^{b1}-1}$

Extracting a Substring



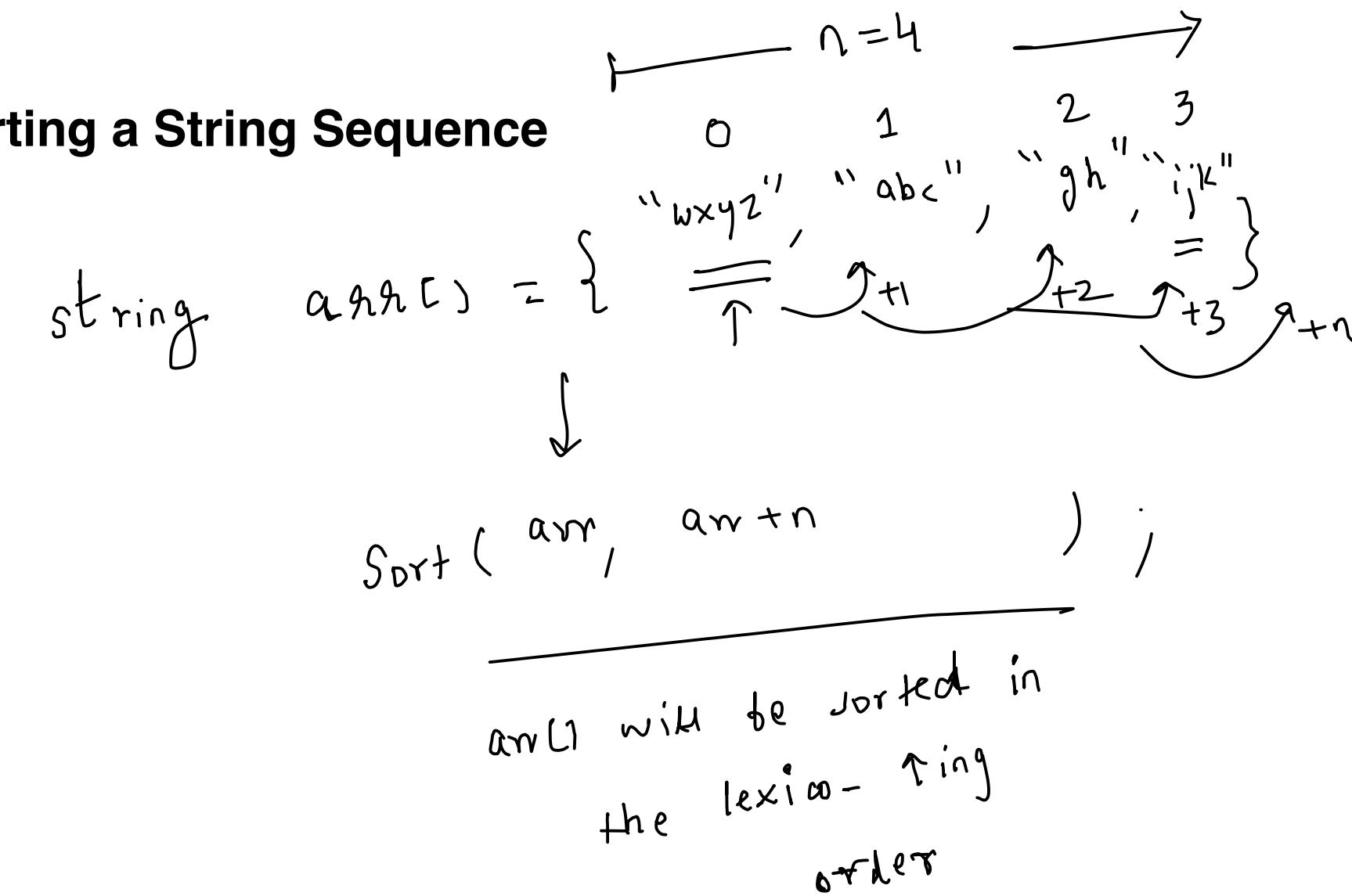
Reversing a String



Sorting a String



Sorting a String Sequence



$$\frac{n(n+1)}{2} \text{ substrings}$$

$$\sim n^2$$

Generate Sub-Strings I

Given a string, design an algorithm to generate all of its **sub-strings**.

We define a **sub-string** of an string as a **contiguous** part of the given string.

Example

Input : "abcde"

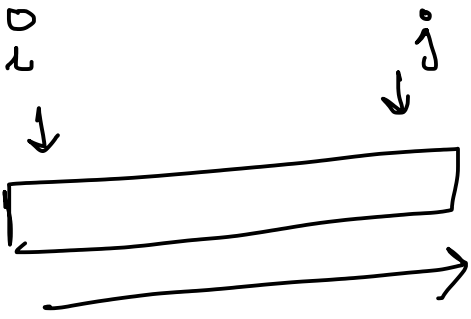
Output : "a", "ab", "abc", "abcd", "abcde"
"b", "bc", "bcd", "bcde"
"c", "cd", "cde"
"d", "de"
"e"

$n = 5$

$0 \leq i \leq n-1$
 $0 \leq j \leq n-1$

$i \quad j$

a b c d e



0th $\left(\begin{array}{l} \boxed{a} \\ a\ b \\ ab\ c \\ abc\ d \\ \boxed{abcde} \end{array} \right)$ 1st $\left(\begin{array}{l} \boxed{b} \\ b\ c \\ bc\ d \\ \boxed{bcde} \end{array} \right)$ 2nd $\left(\begin{array}{l} \boxed{c} \\ c\ d \\ \boxed{cde} \end{array} \right)$

3rd $\left(\begin{array}{l} \boxed{d} \\ d\ e \\ \boxed{de} \end{array} \right)$ 4th $\left(\begin{array}{l} \boxed{e} \end{array} \right)$ len = ?

str.substr(i, len)

$i \quad j$

$j - i + 1$

$i \quad j$

$4 - 2 + 1 = 3$

$2 - 0 + 1 = 3$

$3 - 3 + 1 = 2$

$$\sum_{i=0}^{n-1} (n-i) = \frac{n(n+1)}{2} \sim n^2$$

Generate Sub-Strings II

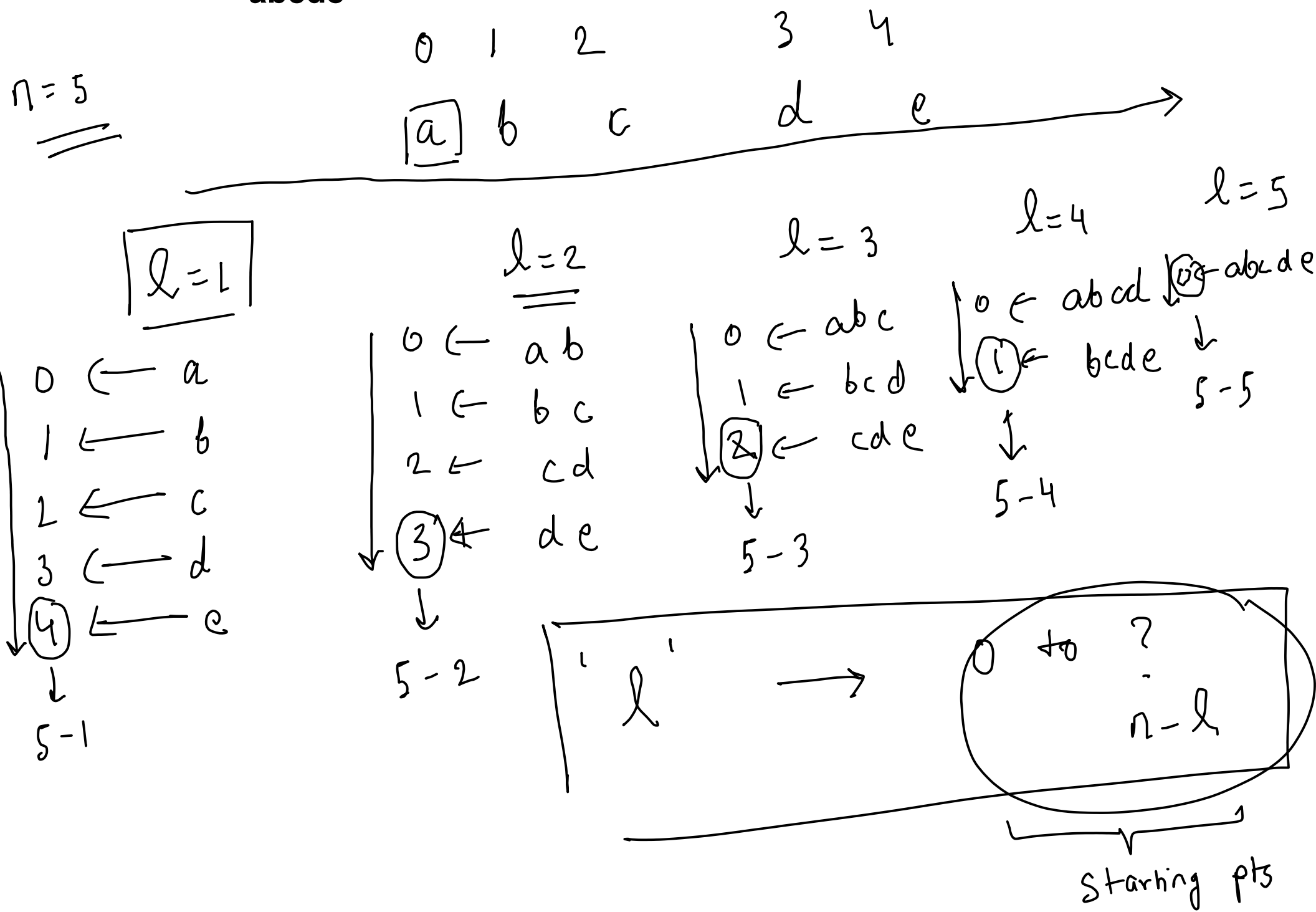
Given a string, design an algorithm to generate all of its **sub-strings** length-wise.

We define a **sub-string** of an string as a **contiguous** part of the given string.

Example

Input : "abcde"

Output : "a", "b", "c", "d", "e"
"ab", "bc", "cd", "de"
"abc", "bcd", "cde"
"abcd", "bcde"
"abcde"



↗ Check Good Strings

Given a string **str**, design an algorithm to check it is a **good string**.

We define a good string as a string only contains **vowels** [a, e, i, o, u].

Example

Input : "uoiea"
Output : True

$a \in \text{Vowel} \rightarrow \text{false}$
 ↓
consonant / not a vowel

C

Range - Based For Loop
or
for - each loop

```
string s = "abc";
```

→ container

```

for ( char ch : s ) {
    cout << ch;
}

```

$$// a, b, c$$

Longest Good Sub-String

Longest Good Sub-String


Given a string **str**, design an algorithm to find the length of its **longest good sub-string**.

We define a good sub-string of **str** as a sub-string that only contains **vowels** [a, e, i, o, u].

Example

Input : "cbāeīcdeiou"

Output : 4

3
Input : "cbāeīcdeiou"
Output : 4
= 



→ maxSoFar = 0 1 2 3 4
cnt = 0 1 2 3 0 1 2 3 4

