RLWE Homomorphic Encryption Scheme

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1 Introduction

This document provides a mathematical explanation of the RLWE (Ring Learning With Errors) homomorphic encryption scheme. The provided Python code demonstrates key generation, encryption, and various operations on encrypted data.

See more:

- "Somewhat Practical Fully Homomorphic",
- "Homomorphic Encryption: a Toy Implementation in Python"

2 RLWE Encryption Parameters

The RLWE encryption scheme operates in a polynomial ring modulo a polynomial modulus $poly_mod$. The ciphertext modulus is denoted as q, and the plaintext modulus is denoted as t. The encryption scheme uses a base T for polynomial switching and a modulus p for modulus switching.

2.1 Parameters

The following parameters are defined:

- n: Polynomial modulus degree
- q: Ciphertext modulus
- t: Plaintext modulus
- T: Base for polynomial switching
- p: Modulus for modulus switching

3 Key Generation

The key generation process involves generating a public key (a, b) and a secret key s. The public key is used for encryption, and the secret key is used for decryption.

3.1 Generating Secret and Public Keys

The secret key s is a binary polynomial. The public key is generated as follows:

- s: Binary secret key polynomial
- a: Uniformly random polynomial
- \bullet e: Error polynomial sampled from a normal distribution
- $b: a \cdot s e \mod q$

4 Encryption

To encrypt a plaintext message m using the public key (a,b), the following steps are performed:

```
\begin{split} m &\equiv m \mod t \\ \mathrm{delta} &= \frac{q}{t} \\ e1 &= \mathrm{Error\ polynomial} \\ \mathrm{scaled\_m} &= \delta \cdot m \mod q \\ \mathrm{new\_ct} &\equiv \mathrm{new\_ct\_0} \\ \mathrm{new\_ct\_0} &= a \cdot u + e1 + \mathrm{scaled\_m} \mod q \end{split}
```

5 Homomorphic Operations

The RLWE encryption scheme supports various homomorphic operations, such as addition and multiplication, on encrypted data.

5.1 Addition of Ciphertext and Plaintext

To add a ciphertext ct and a plaintext pt, the following operation is performed:

$$\begin{split} m &\equiv \text{pt} \mod t\\ \text{delta} &= \frac{q}{t}\\ e1 &= \text{Error polynomial}\\ \text{scaled_m} &= \delta \cdot m \mod q\\ \text{new_ct} &\equiv \text{new_ct_0}\\ \text{new_ct_0} &= ct_0 + \text{scaled_m} \mod q \end{split}$$

5.2 Addition of Ciphertexts

To add two ciphertexts ct1 and ct2, the following operation is performed:

$$\begin{aligned} \text{new_ct} &\equiv \text{new_ct_0} \\ \text{new_ct_0} &= ct1_0 + ct2_0 \mod q \end{aligned}$$

5.3 Multiplication of Ciphertext and Plaintext

To multiply a ciphertext ct by a plaintext pt, the following operation is performed:

$$\begin{split} & \text{scaled_m} = \delta \cdot m \mod q \\ & \text{new_ct} \equiv \text{new_ct_0} \\ & \text{new_ct_0} = ct_0 \cdot \text{scaled_m} \mod q \end{split}$$

5.4 Multiplication of Ciphertexts

The multiplication of ciphertexts involves more complex steps and relinearization keys:

$$\begin{aligned} &\text{new_ct} \equiv \text{new_ct_0} \\ &\text{new_ct_0} = ct1_0 \cdot ct2_0 + ct1_1 \cdot ct2_1 \mod q \\ &\text{new_ct_1} = ct1_0 \cdot ct2_1 + ct1_1 \cdot ct2_0 \mod q \end{aligned}$$

6 Decryption

To decrypt a ciphertext and obtain the original plaintext message, the following steps are performed:

$$scaled_pt = ct_0 \cdot sk + ct_0 \mod q$$

$$decrypted_poly = \left\lfloor \frac{t \cdot scaled_pt}{q} \right\rfloor \mod t$$

$$Decrypted_Plaintext = decrypted_poly$$

The RLWE homomorphic encryption scheme provides a powerful tool for performing computations on encrypted data while preserving data privacy. The scheme involves key generation, encryption, and various homomorphic operations, making it suitable for secure computations in various applications.

7 Python Code

Thank you to the authors of the original code:

- Mădălina Bolboceanu (https://github.com/mbolboceanu)
- Miruna Roșca (https://github.com/MirunaRosca)
- Radu Ţiţiu (https://github.com/rtitiu)

This code is an updated version of the original code, which can be found here:

• https://github.com/bit-ml/he-scheme

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Thank-you-to-the-author-of-the-original-code:
----- Madalina - Bolboceanu - (@mbolboceanu)
This code is an updated version of the original code, which can be found here:
----https://github.com/bit-ml/he-scheme
import numpy as np
from numpy.polynomial import polynomial as poly
\#——Functions for polynomial evaluations mod poly-mod only—
def polymul_wm(x, y, poly_mod):
    """ Multiply - two - polynomials
---Args:
----x, -y: -two-polynomials -to-be-multiplied.
-----poly_mod:-polynomial-modulus.
----Returns:
\operatorname{A-polynomial-in-Z}[X]/(\operatorname{poly_mod}).
   return poly.polydiv(poly.polymul(x, y), poly_mod)[1]
def polyadd_wm(x, y, poly_mod):
   """Add-two-polynomials
----Args:
----x, -y: -two-polynomials - to-be-added.
-----poly_mod:-polynomial-modulus.
-----Returns:
-----A-polynomial-in-Z[X]/(poly_mod).
    return poly.polydiv(poly.polyadd(x, y), poly_mod)[1]
\# ——Functions for polynomial evaluations both mod poly-mod and mod q—
def polymul(x, y, modulus, poly_mod):
   "" "Multiply - two - polynomials
----x, -y: -two-polynomials -to-be-multiplied.
-----modulus: coefficient modulus.
-----poly_mod:-polynomial-modulus.
----Returns:
return np.int64(
       np.round(poly.polydiv(poly.polymul(x, y) %
                             modulus, poly_mod)[1] % modulus)
```

```
)
def polyadd(x, y, modulus, poly_mod):
   """Add-two-polynomials
---- Args:
······x, ·y: ·two·polynoms·to·be·added.
-----modulus: coefficient modulus.
-----poly_mod:-polynomial-modulus.
----Returns:
return np.int64(
       np.round(poly.polydiv(poly.polyadd(x, y) %
                            modulus, poly_mod)[1] % modulus)
# ----Functions for random polynomial generation --
def gen_binary_poly(size):
   """ Generates - a - polynomial - with - coeffecients - in - [0, -1]
-----size: number of coeffcients, size-1 being the degree of the
-----polynomial.
-----array-of-coefficients-with-the-coeff[i]-being
-----the-coeff-of-x-^-i.
   return np.random.randint(0, 2, size, dtype=np.int64)
def gen_uniform_poly(size, modulus):
   "" Generates - a - polynomial - with - coeffecients - being - integers - in - Z_modulus
----polynomial.
----array of coefficients with the coeff[i] being
-----the-coeff-of-x-^-i.
   return np.random.randint(0, modulus, size, dtype=np.int64)
def gen_normal_poly(size, mean, std):
   "" Generates - a - polynomial - with - coeffecients - in - a - normal - distribution
----of-mean-mean-and-a-standard-deviation-std, -then-discretize-it.
```

```
-----size: number of coeffcients, size -1 being the degree of the
-----polynomial.
----Returns:
-----array of coefficients with the coeff[i] being
------the-coeff-of-x-^-i.
    return np.int64(np.random.normal(mean, std, size=size))
# ----- Function for returning n's coefficients in base b ( lsb is on the lef
def int2base(n, b):
    "" Generates - the - base - decomposition - of - an - integer - n.
----Args:
-----n: integer to be decomposed.
-----b:-base.
----Returns:
----array of coefficients from the base decomposition of n
-----with the coeff[i] being the coeff of b-^-i.
    return [n] if n < b else [n \% b] + int2base(n // b, b)
\# — Functions for keygen, encryption and decryption —
def keygen(size, modulus, poly_mod, std1):
    """ Generate - a - public - and - secret - keys
---- Args:
-----size: -size - of -the -polynoms - for -the -public - and -secret - keys.
-----modulus: coefficient modulus.
-----poly_mod:-polynomial-modulus.
-----std1:-standard-deviation-of-the-error.
---- Returns:
-----Public and secret key.
- - - - " " "
    s = gen_binary_poly(size)
    a = gen_uniform_poly(size, modulus)
    e = gen\_normal\_poly(size, 0, std1)
    b = polyadd(polymul(-a, s, modulus, poly_mod), -e, modulus, poly_mod)
    return (b, a), s
def evaluate_keygen_v1(sk, size, modulus, T, poly_mod, std2):
    "" Generate - a - relinearization - key - using - version - 1.
---- Args:
-----sk:-secret-key.
····size: size of the polynomials.
```

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-----modulus: coefficient modulus.
----T:-base.
-----poly_mod:-polynomial-modulus.
-----std2:-standard-deviation-for-the-error-distribution.
-----Returns:
-----rlk0, -rlk1:-relinearization-key.
   n = len(poly_mod) - 1
   l = np.int64 (np.log(modulus) / np.log(T))
   rlk0 = np.zeros((l + 1, n), dtype=np.int64)
   rlk1 = np.zeros((l + 1, n), dtype=np.int64)
   for i in range(l + 1):
       a = gen_uniform_poly(size, modulus)
       e = gen_normal_poly(size, 0, std2)
       secret_part = T ** i * poly.polymul(sk, sk)
       b = np.int64 (polyadd)
       polymul_wm(-a, sk, poly_mod),
       polyadd_wm(-e, secret_part, poly_mod), modulus, poly_mod))
       b = np.int64(np.concatenate((b, [0] * (n - len(b))))) # pad b
       a = np.int64(np.concatenate((a, [0] * (n - len(a))))) # pad a
       rlk0[i] = b
       rlk1[i] = a
   return rlk0, rlk1
def evaluate_keygen_v2(sk, size, modulus, poly_mod, extra_modulus, std2):
   "" Generate - a - relinearization - key - using - version - 2.
-----Args:
-----sk:-secret-key.
----size: size of the polynomials.
   -----modulus: coefficient modulus.
-----poly_mod:-polynomial-modulus.
-----extra_modulus: the "p" modulus for modulus switching.
-----Returns:
-----rlk0, -rlk1: -relinearization -key.
- - - - - - - " " "
   new_modulus = modulus * extra_modulus
   a = gen_uniform_poly(size, new_modulus)
   e = gen_normal_poly(size, 0, std2)
   secret_part = extra_modulus * poly.polymul(sk, sk)
   b = np.int64 (polyadd_wm)
       polymul_wm(-a, sk, poly_mod),
       polyadd_wm(-e, secret_part, poly_mod), poly_mod)) % new_modulus
```

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return b, a
def encrypt (pk, size, q, t, poly_mod, m, std1):
   "" Encrypt an integer.
----Args:
----pk:-public-key.
----size: size of polynomials.
----q:-ciphertext-modulus.
-----t:-plaintext-modulus.
-----poly_mod:-polynomial-modulus.
-----m: plaintext message, as an integer vector (of length <= size) with entr
----Returns:
-----Tuple-representing-a-ciphertext.
   m = np.array(m + [0] * (size - len(m)), dtype=np.int64) % t
   delta = q // t
   scaled_m = delta * m
   e1 = gen\_normal\_poly(size, 0, std1)
   e2 = gen\_normal\_poly(size, 0, std1)
   u = gen_binary_poly(size)
   ct0 = polyadd(
       polyadd (
           polymul(pk[0], u, q, poly_mod),
           e1, q, poly_mod),
       scaled_m, q, poly_mod
   )
   ct1 = polyadd(
       polymul(pk[1], u, q, poly_mod),
       e2, q, poly_mod
   return (ct0, ct1)
def decrypt(sk, size, q, t, poly_mod, ct):
   """Decrypt-a-ciphertext.
----Args:
-----size:-size-of-polynomials.
----q:-ciphertext-modulus.
-----t:-plaintext-modulus.
----ct:-ciphertext.
----Returns:
-----Integer - vector - representing - the - plaintext.
   scaled_pt = polyadd(
```

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polymul(ct[1], sk, q, poly_mod),
         ct[0], q, poly_mod
    )
    decrypted_poly = np.round(t * scaled_pt / q) % t
    decrypted_poly_1 = list(decrypted_poly)
    bound = len(decrypted_poly_1)
    if bound<size:</pre>
         number_of_zeros_to_pad = size-bound
         list_to_append = [0] * number_of_zeros_to_pad
         decrypted_poly_to_return = np.append(decrypted_poly, list_to_append) #pa
         decrypted_poly_to_return = decrypted_poly
    return np.int64(decrypted_poly_to_return)
#
# ----Function for adding and multiplying encrypted values -
def add_plain(ct, pt, q, t, poly_mod):
    "" " Add - a - ciphertext - and - a - plaintext .
---- Args:
----ct:-ciphertext.
-----pt:-integer-to-add.
----q:-ciphertext-modulus.
-----t:-plaintext-modulus.
-----poly_mod:-polynomial-modulus.
----Returns:
-----Tuple-representing-a-ciphertext.
    size = len(poly\_mod) - 1
    # encode the integer into a plaintext polynomial
    m = np.array(pt + [0] * (size - len(pt)), dtype=np.int64) % t
    delta = q // t
    scaled_m = delta * m
    new_ct0 = polyadd(ct[0], scaled_m, q, poly_mod)
    return (new_ct0, ct[1])
\mathbf{def}\ \mathrm{add\_cipher}\left(\,\mathrm{ct1}\;,\;\;\mathrm{ct2}\;,\;\;\mathrm{q}\;,\;\;\mathrm{poly\_mod}\,\right)\colon
    """Add-a-ciphertext-and-a-ciphertext.
---Args:
\cdots \cdots \cot 1, \cot 2: - \operatorname{ciphertexts}.
-----q:-ciphertext-modulus.
-----poly_mod:-polynomial-modulus.
----Returns:
-----Tuple-representing-a-ciphertext.
```

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_ _ _ , ,, ,, ,,
    new_ct0 = polyadd(ct1[0], ct2[0], q, poly_mod)
    new_ct1 = polyadd(ct1[1], ct2[1], q, poly_mod)
    return (new_ct0, new_ct1)
def mul_plain(ct, pt, q, t, poly_mod):
    "" " Multiply - a - ciphertext - and - a - plaintext.
----ct:-ciphertext.
\hbox{\tt -------pt:-integer-polynomial-to-multiply}.
----q:-ciphertext-modulus.
 \hbox{\tt ------} t: \hbox{\tt -plaintext-modulus}\,.
-----poly_mod:-polynomial-modulus.
---- Returns:
\hbox{\tt ------} Tuple \hbox{\tt -representing -a-ciphertext} \; .
_ _ _ _ ;; ;; ;;
    size = len(poly_mod) - 1
    # encode the integer polynomial into a plaintext vector of size=size
    m = np.array(pt + [0] * (size - len(pt)), dtype=np.int64) % t
    new_c0 = polymul(ct[0], m, q, poly_mod)
    new_c1 = polymul(ct[1], m, q, poly_mod)
    return (new_c0, new_c1)
def multiplication_coeffs (ct1, ct2, q, t, poly_mod):
    """ Multiply -two-ciphertexts.
---- Args:
-----ct1: first ciphertext.
----ct2:-second-ciphertext
----q: ciphertext modulus.
-----t:-plaintext-modulus.
 -----poly_mod:-polynomial-modulus.
-----Returns:
----- Triplet \cdot (c0, c1, c2) \cdot encoding \cdot the \cdot multiplied \cdot ciphertexts.
    c_0 = \text{np.int} 64 \text{ (np.round (polymul_wm (ct1 [0], ct2 [0], poly_mod) * t / q)) } \% \text{ q}
    c_1 = p.int64 (polyadd_wm(polymul_wm(ct1[0], ct2[1], poly_mod), polywol_wm(ct1[0], ct2[1], poly_mod),
    c_2 = np.int64(np.round(polymul_wm(ct1[1], ct2[1], poly_mod) * t / q)) % q
    \mathbf{return} \ c_{-}0 \ , \ c_{-}1 \ , \ c_{-}2
def mul_cipher_v1(ct1, ct2, q, t, T, poly_mod, rlk0, rlk1):
    """ Multiply - two - ciphertexts.
--- Args:
\hbox{\tt ------ct1:-first-ciphertext.}
```

```
----ct2:-second-ciphertext
----q:-ciphertext-modulus.
-----t:-plaintext-modulus.
----T:-base
-----poly_mod:-polynomial-modulus.
-----rlk0, -rlk1: -output-of-the-EvaluateKeygen_v1-function.
····Tuple-representing-a-ciphertext.
            n = len(poly_mod) - 1
            1 = \text{np.int} 64 (\text{np.log}(q) / \text{np.log}(T))  # l = log_T T(q)
            c_0, c_1, c_2 = multiplication\_coeffs(ct1, ct2, q, t, poly\_mod)
            c_2 = \text{np.int} 64 \text{ (np.concatenate ( } (c_2, [0] * (n - \text{len}(c_2))) )) \#pad
            \#Next, we decompose c_{-2} in base T:
            \#more\ precisely , each coefficient of c_2 is decomposed in base T such that c
            Reps = np.zeros((n, l + 1), dtype = np.int64)
            for i in range(n):
                         rep = int2base(c_2[i], T)
                         rep2 = rep + [0] * (1 + 1 - len(rep)) \#pad with 0
                         Reps[i] = np.array(rep2, dtype=np.int64)
            \# Each row Reps[i] is the base T representation of the i-th coefficient c_2[
            # The polynomials c_{-2}(j) are given by the columns Reps[:, j].
            c_20 = np.zeros(shape=n)
            c_21 = np. zeros (shape=n)
            # Here we compute the sums: rlk[j][0] * c_2(j) and rlk[j][1] * c_2(j)
            for j in range(l + 1):
                         c\_20 \ = \ polyadd\_wm\left(\,c\_20\;,\;\; polymul\_wm\left(\,rlk\,0\,[\,j\,]\;,\;\; Reps\,[\,:\,,\,j\,]\;,\;\; poly\_mod\,\right)\;,\;\; poly\_m
                         c_-21 \ = \ polyadd\_wm\left(\,c_-21 \ , \ polymul\_wm\left(\,r\,l\,k\,1\,\left[\,j\,\right] \ , \ Reps\left[:\,,\,j\,\right] \ , \ poly\_mod\,\right) \ , \ poly\_mod\,\right) \ , \ poly\_mod\,\right) \ , \ poly\_mod\,\left[:\,,\,j\,\right] \ , \ poly\_mod\,\left[:\,j\,\right] \ ,
            c_{-}20 = np. int 64 (np. round (c_{-}20)) \% q
            c_21 = np.int64(np.round(c_21)) \% q
            new_c0 = np.int64(polyadd_wm(c_0, c_20, poly_mod)) \% q
            \text{new}_{c1} = \text{np.int} 64 (\text{polyadd\_wm} (\text{c\_1}, \text{c\_21}, \text{poly\_mod})) \% \text{ q}
            return (new_c0, new_c1)
\mathbf{def} \ \mathbf{mul\_cipher\_v2} \ (\ \mathbf{ct1} \ , \ \ \mathbf{ct2} \ , \ \ \mathbf{q} \ , \ \ \mathbf{t} \ , \ \ \mathbf{poly\_mod} \ , \ \ \mathbf{rlk0} \ , \ \ \mathbf{rlk1} \ ) :
            """ Multiply two ciphertexts.
---- Args:
----ct1:-first-ciphertext.
----ct2:-second-ciphertext.
----q:-ciphertext-modulus.
```

```
return (new_c0, new_c1)

reduction switched (polyadd_wm(c_1, c_21, poly_mod)) % q

return (new_c0, new_c1)
```