

Lab Final – Verilog

Developing code for a **D-flip flop** with an **asynchronous active low reset input**:
The design source is as follows:

```
module d_flipflop(d, clk, reset_n, q);

input d;
input clk;
input reset_n;
output reg q;

always@(posedge clk or negedge reset_n)

begin

    if(~reset_n) // If the active low signal is asserted of the reset;

        q <= 1'b0;
    else
        q <= d;

end

endmodule
```

The equivalent testbench code is the following:

```
`timescale 1ns / 1ps

module d_flipflop_tb;

// Testbench signals (reg for inputs, wire for outputs)
reg d;
reg clk;
```

```

reg reset_n;
wire q;

// Instantiate the DUT (Device Under Test)
d_flipflop uut (
    .d(d),
    .clk(clk),
    .reset_n(reset_n),
    .q(q)
);

// Clock generation: 10 ns period
always #5 clk = ~clk;

initial begin
    // Initialize signals
    clk = 0;
    d = 0;
    reset_n = 1;

    // Apply asynchronous reset
    #3 reset_n = 0; // Assert reset (active low)
    #7 reset_n = 1; // Deassert reset

    // Apply test vectors
    #10 d = 1;
    #10 d = 0;
    #10 d = 1;

    // Apply reset again while clock is running
    #7 reset_n = 0;
    #5 reset_n = 1;

    #20 $finish;
end

endmodule

```

Differentiation between a Moore and Mealy FSM

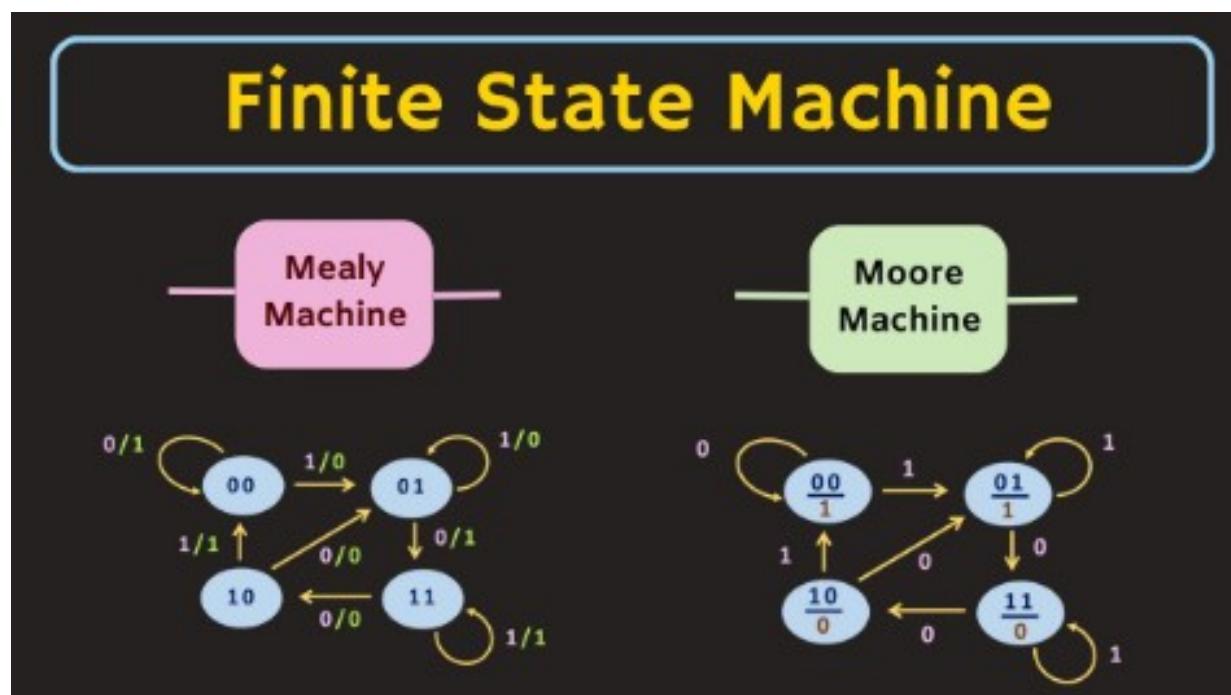
Table 1: Sequence Detector - Mealy FSM

In a Mealy type FSM, the **number of states** in the detector is the **same** as the **number of bits in the sequence**.

Table 2: Sequence Detector - Moore FSM

In a Moore type FSM, the **number of states** in the detector is **one more (+1)** than the **number of bits in the sequence**.

| Moore Machine | Mealy Machine |
|---|---|
| Output depends only upon the present state. | Output depends on the present state as well as present input. |
| <u>Moore machine</u> also places its output in the state. | <u>Mealy Machine</u> places its output on the transition. |



Task: Design a Moore Finite State Machine (FSM) for Sequence Detection

Problem Statement:

Design a **non-overlapping** sequence detector that detects the binary sequence 1011. The detector should have:

- **Input:** din (1-bit serial input, sampled on positive clock edge)
- **Output:** detect (1-bit, goes high for one clock cycle when the sequence is detected)
- **Reset:** Asynchronous active-low reset (reset_n)

Specifications:

1. Use **Moore machine** style (output depends only on current state).
2. When the **exact sequence** 1011 is received, detect should be **1 in the same clock cycle** when the last bit (1) is received.
3. After detection, the FSM should restart looking for a new sequence from the next bit (non-overlapping detection).
4. Assume sequence bits arrive one per clock cycle.

Writing the **design source code**:

```
module moore_fsm(
    input clk;
    input reset_n;
    input x; // Input bit
    output y // there is no comma at the end here
);

reg[2:0] state_reg, state_next;
localparam A = 3'b000;
localparam B = 3'b001;
localparam C = 3'b010;
localparam D = 3'b011;
localparam E = 3'b100;

// state register
always @ (posedge clk, negedge reset_n)
```

```

begin
    if(~reset_n)
        state_reg <= A;
    else
        state_reg <= state_next;
end

// Next state logic
always @ (*) // start the always block by defining your begin and end statements
in the beginning
begin
    case(state_reg) // make sure to include the endcase (IMPORTANT)
        A: if(x)
            state_next = B
        else
            state_next = A
        B: if(x)
            state_next = B
        else
            state_next = C
        C: if(x)
            state_next = D
        else
            state_next = A
        D: if(x)
            state_next = E
        else
            state_next = A
        E: if(x)
            state_next = A
        else
            state_next = A
        default: state_next = state_reg;
    endcase
end

// Output logic
assign y = (state_reg == E);

endmodule

```

Developing the **test bench code** for this design source:

```
module moore_fsm_tb;

// Outputs: wire
// Inputs: reg

reg clk;
reg reset_n;
reg x;
wire y;

// Instantiate the Unit Under Test
moore_sequence_detector dut(
    .clk(clk),
    .reset_n(reset_n),
    .x(x),
    .y(y)
);

// Generate a clock signal
always #5 clk = ~clk;

reg [7:0] simple_test = 8'b10111011;
// Should detect ONLY once (non-overlapping)
integer i; // iterative variable to loop through all the input bit stream and assign to the
input var;

initial
begin
    clk = 0;
    reset_n = 0; // Assert reset (active low)
    x = 0;
    #20 reset_n = 1; // Release reset after 20ns

    for (i = 7; i >= 0; i = i-1) // Count DOWN from 7 to 0
    begin
        @(posedge clk);
```

```
x = simple_test[i];
end

repeat (5) @(posedge clk);
$stop;
end
```

For an active low asynchronous reset; the following code block is crucial in the testbench:

```
reset_n = 0; // Assert reset (active low)

x = 0; // hardcode the default input bit

#20 reset_n = 1; // Release reset after 20ns
```

◆ FSM PRACTICE QUESTION (LAB FINAL LEVEL)

Problem Title:

A university has installed a **smart entry gate** that controls access using an **ID card and a security PIN**. The system operates synchronously on a clock and has an **active-low asynchronous reset**.

System Inputs

- **card** :
 - 1 → Valid ID card detected
 - 0 → No card / invalid card
- **pin** :
 - 1 → Correct PIN entered
 - 0 → Incorrect PIN or no PIN
- **clk** : System clock
- **reset_n** : Active-low asynchronous reset

System Outputs

- **gate_open :**
1 → Gate opens
0 → Gate remains closed
 - **alarm :**
1 → Alarm triggered
0 → Alarm off
-

◆ Functional Requirements

1. Idle State

- System starts in the **IDLE** state.
- Gate is closed and alarm is off.

2. Card Verification

- If a **valid card** is detected (**card** = 1), the system moves to a **PIN_WAIT** state.
- If **card** = 0, the system remains in IDLE.

3. PIN Verification

- In the **PIN_WAIT** state:
 - If the correct PIN is entered (**pin** = 1), the gate opens for **one clock cycle**.
 - After opening, the system returns to IDLE.
 - If an incorrect PIN is entered (**pin** = 0), the system moves to an **ALARM** state.

4. Alarm State

- In the **ALARM** state:
 - The alarm is activated.
 - The system remains in this state until reset is asserted.
 - The gate must remain closed.

5. Reset Behavior

- When `reset_n = 0`, the system must **immediately return to IDLE**, regardless of the clock.
-

◆ Tasks (Exactly Exam Style)

You are required to:

1. **Identify and list all states** of the FSM.
 2. **Draw a complete state transition diagram** (clearly labeled with inputs and outputs).
 3. **Specify whether the FSM is Moore or Mealy**, with justification.
 4. **Write synthesizable Verilog code** to implement the FSM in Vivado.
 - Use **behavioral modeling**
 - Use **non-blocking assignments**
 - Use **asynchronous active-low reset**
 5. Assume **one-hot or binary encoding** (your choice, but be consistent).
-

◆ Important Exam Hints (Read Carefully)

- Outputs must be **deterministic for each state**
- Avoid latches
- Reset must override all logic
- Follow the **three-block FSM style** if possible:
 - State register
 - Next-state logic
 - Output logic

Any time an output is asserted for “one clock cycle”, Moore FSMs are usually preferred.

- If an input has *no meaning* in a state, **ignore it**.
- Any output that must last exactly one clock cycle is best implemented as a separate state.
- If outputs describe system behavior over a clocked interval, use Moore; if outputs must react instantly to inputs, use Mealy.
- If a state exists *only* to assert an output for one cycle, its exit must be unconditional.

That means that, the input to transition into another state should simply be xx (don’t cares).

- In **both Moore and Mealy FSMs**, the **output logic is written inside an always @(*) block**.
What changes is **what the output depends on, not how it is coded structurally**.
- If an output is assigned inside an **always** block, it must be declared as **reg**.

The **design source code** is as follows:

```
module fsm_practice(
    input clk,
    input reset_n,
    input [1:0] i,
    output reg [1:0] O
);
    reg [1:0] state_reg, state_next;
    localparam I = 2'b00;
    localparam PW = 2'b01;
    localparam GO = 2'b10;
    localparam A = 2'b11;

    // 1 State register
    always @(posedge clk or negedge reset_n) begin
        if (!reset_n)
```

```

    state_reg <= I;
else
    state_reg <= state_next;
end

// [2] Next-state logic
always @(*) begin
    state_next = state_reg; // default

    case (state_reg)
        I: begin
            if (i[1] == 1'b1)
                state_next = PW;
            else
                state_next = I;
        end

        PW: begin
            if (i[0] == 1'b1)
                state_next = GO;
            else
                state_next = A;
        end

        GO: begin
            state_next = I; // single-cycle pulse
        end

        A: begin
            state_next = A; // absorb state
        end
    endcase
end

// [3] Output logic (Moore)
always @(*) begin
    case (state_reg)
        I: O = 2'b00;
        PW: O = 2'b00;
        GO: O = 2'b10;
        A: O = 2'b01;
    end

```

```

    default: O = 2'b00;
endcase
end

endmodule

```

The **test bench code** for the above task is as follows:

```

module fsm_practice_tb;

reg clk;
reg reset_n;
reg [1:0] i;
wire [1:0] O;

// Instantiate DUT
fsm_practice dut (
    .clk(clk),
    .reset_n(reset_n),
    .i(i),
    .O(O)
);

// Clock generation (10ns period)
always #5 clk = ~clk;

initial begin
    // Initial conditions
    clk = 0;
    reset_n = 0;
    i = 2'b00;

    // Hold reset
    #20;
    reset_n = 1;

    // ---- FSM stimulus ----

    // IDLE -> PIN_WAIT (card inserted)
    @(posedge clk);

```

```
i = 2'b10;

// PIN_WAIT -> GO (pin correct)
@(posedge clk);
i = 2'b11;

// GO -> IDLE (automatic)
@(posedge clk);
i = 2'b00;

// IDLE -> PIN_WAIT
@(posedge clk);
i = 2'b10;

// PIN_WAIT -> ALARM (wrong pin)
@(posedge clk);
i = 2'b10;

// Stay in ALARM
repeat (2) @(posedge clk);

// Reset from ALARM
reset_n = 0;
#10;
reset_n = 1;

repeat (3) @(posedge clk);
$stop;
end

endmodule
```

In Verilog behavioral modelling, there are three (3) core modules that we studied and implemented:

1. ALU
2. MUX
3. Decoder

ALU

Design source:

```
module alu( // Four (4) inputs and one (1) output
    input [7:0] operand_1,
    input [7:0] operand_2,
    input [2:0] command,
    input enable,
    output reg [15:0] out // declared as a reg since it will be used inside the
always block
);
```

```
// Declare parameters for the different operations
parameter ADD = 3'b000;
parameter INC = 3'b001;
parameter SUB = 3'b010;
parameter MUL = 3'b011;
parameter DIV = 3'b100;
parameter SHL = 3'b101;
parameter SHR = 3'b110;
parameter AND = 3'b111;
```

**// A procedural block using behavioral modeling for a combinational circuit
(ALU):**

```
always @ (operand_1, operand_2, command)
begin
    case(command)
        ADD : out = operand_1 + operand_2;
        INC : out = operand_1 + 1;
        SUB : out = operand_1 - operand_2;
        MUL : out = operand_1 * operand_2;
```

```

        DIV : out = operand_1 / operand_2;
        SHL : out = operand_1 << 1'b1;
        SHR : out = operand_1 >> 1'b1;
        AND : out = operand_1 & operand_2;
        default: result = 8'b0; // Default case: Prevents latch formation
    endcase
end

endmodule

```

Testbench:

```

`timescale 1ns / 1ps

module alu_tb;

// Inputs to DUT → reg
reg [7:0] operand_1;
reg [7:0] operand_2;
reg [2:0] command;
reg enable;

// Output from DUT → wire
wire [15:0] out;

// Instantiate the DUT
alu dut (
    .operand_1(operand_1),
    .operand_2(operand_2),
    .command(command),
    .enable(enable),
    .out(out)
);
initial begin
    // Initial values
    operand_1 = 8'd10;
    operand_2 = 8'd5;
    command  = 3'b000;
    enable   = 0;

```

```
#10 enable = 1; // Enable output

// ADD
#10 command = 3'b000; // ADD → 10 + 5 = 15

// INC
#10 command = 3'b001; // INC → 10 + 1 = 11

// SUB
#10 command = 3'b010; // SUB → 10 - 5 = 5

// MUL
#10 command = 3'b011; // MUL → 10 * 5 = 50

// DIV
#10 command = 3'b100; // DIV → 10 / 5 = 2

// SHL
#10 command = 3'b101; // SHL → 10 << 1 = 20

// SHR
#10 command = 3'b110; // SHR → 10 >> 1 = 5

// AND
#10 command = 3'b111; // AND → 10 & 5 = 0

// Disable output (tri-state)
#10 enable = 0;

#10 $stop;
end

endmodule
```

MUX

Design source:

```
module mux_8x1(
    input [7:0]i,
    input [2:0]s,
    output reg y
);

    // Procedural block: Behavioral modelling of a combinational circuit (MUX)
    always @ (*)
    begin
        case(s)
            // y is getting assigned 1-bit inputs
            3'b000 : y = i[0];
            3'b001 : y = i[1];
            3'b010 : y = i[2];
            3'b011 : y = i[3];
            3'b100 : y = i[4];
            3'b101 : y = i[5];
            3'b110 : y = i[6];
            3'b111 : y = i[7];
            default : y = 1'b0;
        endcase

    end
endmodule
```

Testbench:

```
`timescale 1ns / 1ps

module mux_8x1_tb;
```

// Inputs → reg

```
reg [7:0] i;
reg [2:0] s;
```

```
// Output → wire
wire y;

// Instantiate the DUT
mux_8x1 dut (
    .y(y),
    .i(i),
    .s(s)
);

initial begin
    // Apply input pattern
    i = 8'b1010_1101; // Example data inputs

    // Select each input line one by one
    s = 3'b000; #10; // y = i[0]
    s = 3'b001; #10; // y = i[1]
    s = 3'b010; #10; // y = i[2]
    s = 3'b011; #10; // y = i[3]
    s = 3'b100; #10; // y = i[4]
    s = 3'b101; #10; // y = i[5]
    s = 3'b110; #10; // y = i[6]
    s = 3'b111; #10; // y = i[7]

    $stop;
end

endmodule
```

Decoder

Design source:

```
module decoder_3x8(
    input [2:0] in,      // 3-bit input
    output reg [7:0] out // 8-bit output (active high)
);
    always @(*) begin
        case(in)
            3'b000: out = 8'b00000001;
            3'b001: out = 8'b00000010;
            3'b010: out = 8'b00000100;
            3'b011: out = 8'b00001000;
            3'b100: out = 8'b00010000;
            3'b101: out = 8'b00100000;
            3'b110: out = 8'b01000000;
            3'b111: out = 8'b10000000;
            default: out = 8'b00000000; // Safety default
        endcase
    end
endmodule
```

Testbench:

```
module decoder_tb;
    // 1. Make test inputs
    reg [2:0] in;

    // 2. Get output from decoder
    wire [7:0] out;

    // 3. Connect to our decoder
    decoder_3x8 my_decoder (
```

```
.in(in),  
.out(out)  
);  
  
// 4. Start testing  
initial begin  
    // Test input 000  

```

```
endmodule
```

Shift Registers

SISO (Serial In-Serial Out)

```
module SISO(
    input clk,
    input reset_n,
    input din,
    output dout
);

reg [3:0] shift;

always @ (posedge clk or posedge reset_n)
begin
    if(reset_n)
        q <= 4'b0000;
    else
        q[0] <= q[1];
        q[1] <= q[2];
        q[2] <= q[3];
        q[3] <= din;
end

assign dout = q[3];

endmodule
```

SIPO (Serial In-Parallel Out)

```
module SIPO(
    input clk,
    input reset_n,
    input din,
    output reg [3:0] dout // Defining dout as reg, since it will be used inside the
always block.
);

// Since its a sequential circuit element, it will be defined in an always block
always @ (posedge clk or posedge reset_n)
begin
    if (reset_n) // if reset is asserted
        dout <= 4'b0000;
    else
        dout[3] <= dout[2];
        dout[2] <= dout[1];
        dout[1] <= dout[0];
        dout[0] <= din;
end

endmodule
```

The **difference** between the design source codes of SISO and SIPO are the following three (3):

- A separate reg is defined in SISO, to act as the internal 4-bit shift register.
- The movement of the bits in SISO and SIPO are different.
 - The output in the case of SISO is assigned as q[3].
 - The output in the case of SIPO is assigned as q[0].
- There is no assign statement in the SIPO design source code.
There is an assign statement in the SISO design source code.

