Data Compression Tutorial

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(1) Huffman

"data" $2^3 = 8$

Map & Code creation

$P(d) = 0.25$

$P(a) = 0.5 \rightarrow$ highest

$P(t) = 0.25$

d

a

t

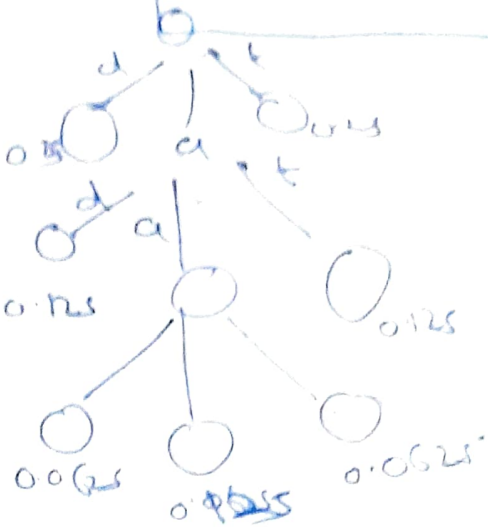
d

ad

a a

a t

t



Code:

d - 000, ad - 001, at - 01

aad - 110, aat - 110

aat - 111

(2) $n = 25, m = 7$

$q = \left\lfloor \frac{25}{7} \right\rfloor = 3 \Rightarrow r = 25 - 21 = 4$

$n = qm + r$

$\Rightarrow 25 = 3 \times 7 + 4$

So we have: 1110 (q | r)

$n = 25, m = 7$

$\lceil \log_2 7 \rceil = 3$

$\lfloor \log_2 7 \rfloor = 2$

$2^3 - 7 = 1 \rightarrow 2$ bit

$4 + 2^3 - 7 = 5 \rightarrow 3$ bit

Code	n
1110000	0
1110001	1
1110010	2
1110111	3
1110100	4

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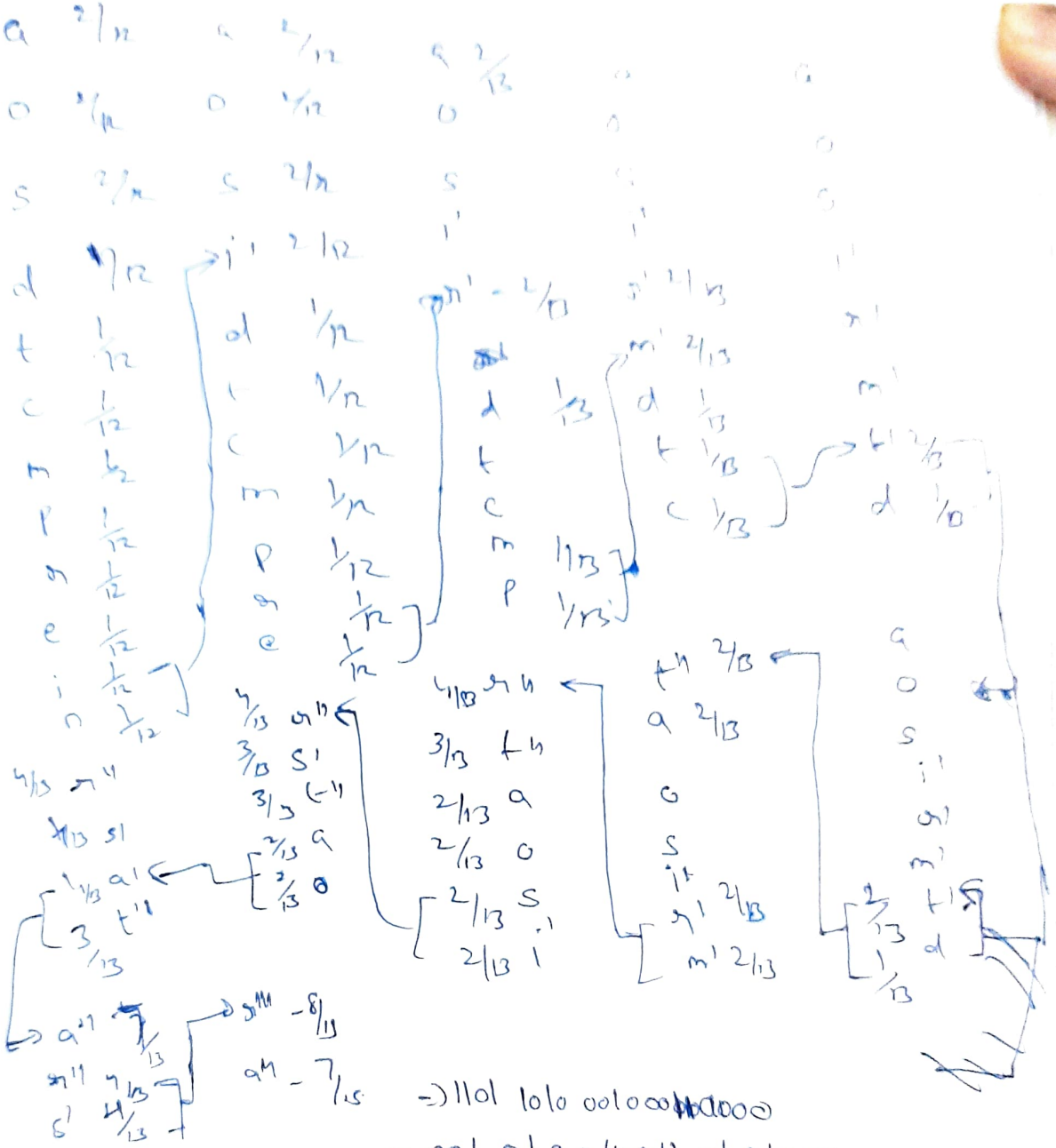
③ Huffman coding



=> 1101 1010 0010 0000 0000

0 001 0100 0110 1010 111

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(4) Adaptive Huffman Coding

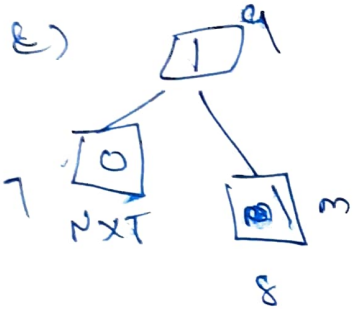
Str: "Malayalam"

$m = 2^e + r$ (let) , $0 \le r < 2^e$

$26 = 2^4 + 10$

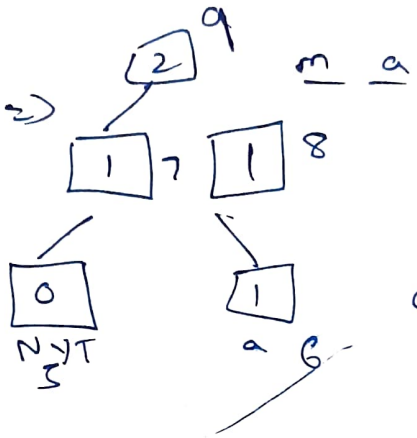
$e = 4, \quad r = 10$

Code : 01100



m

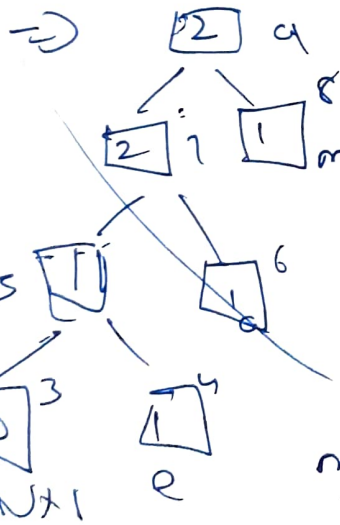
k=13



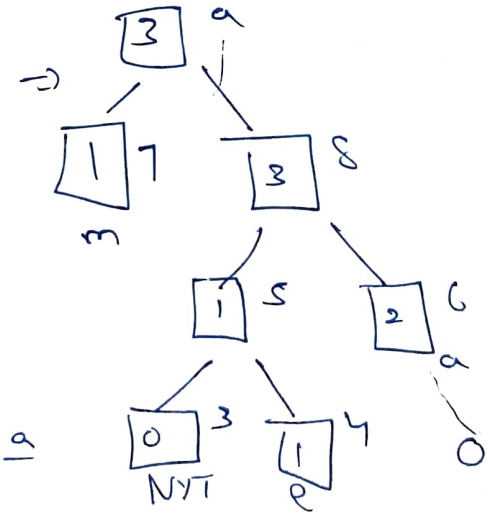
m a

0 11 000 00000

m NY T

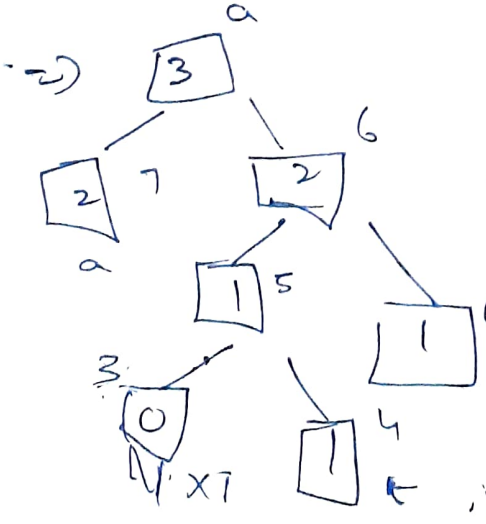


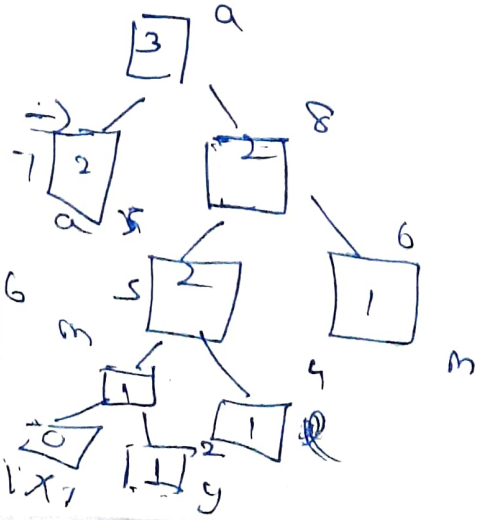
m a l a



01100 0 00000 0

m a e a





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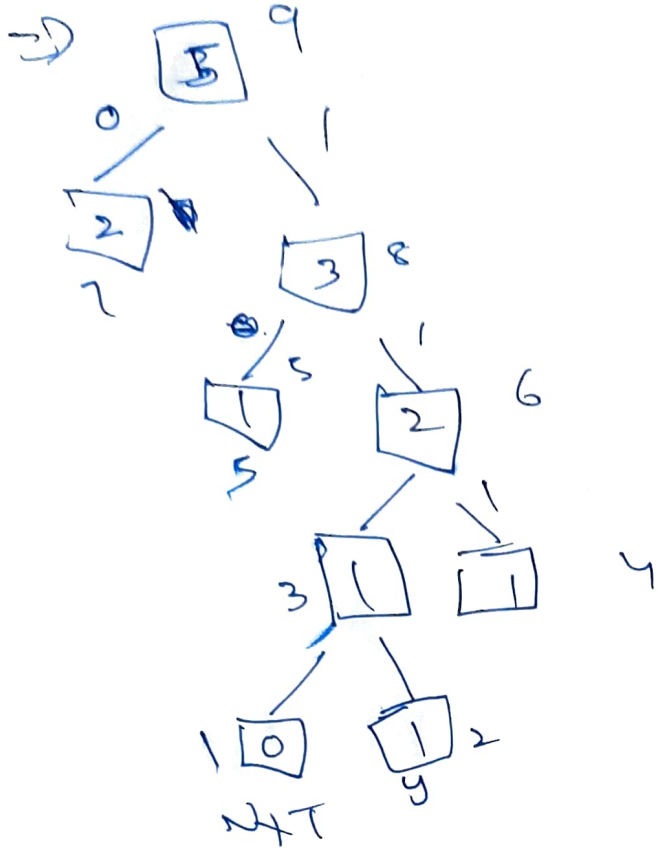
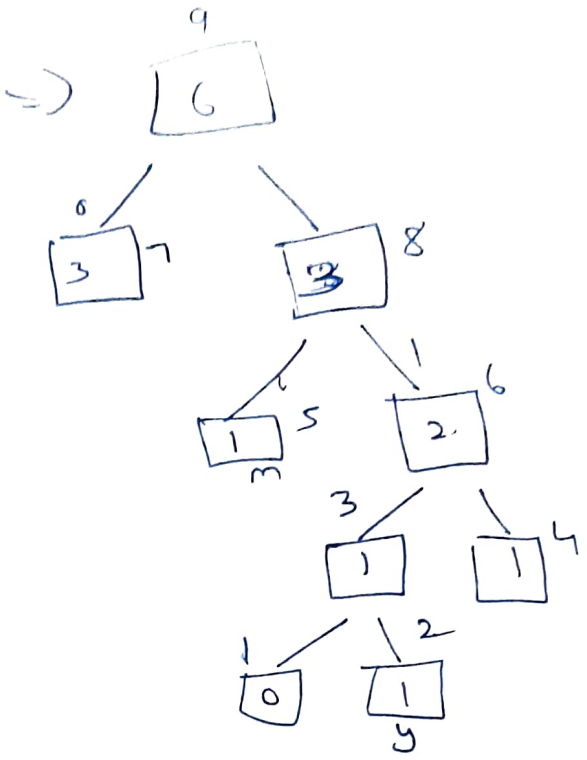
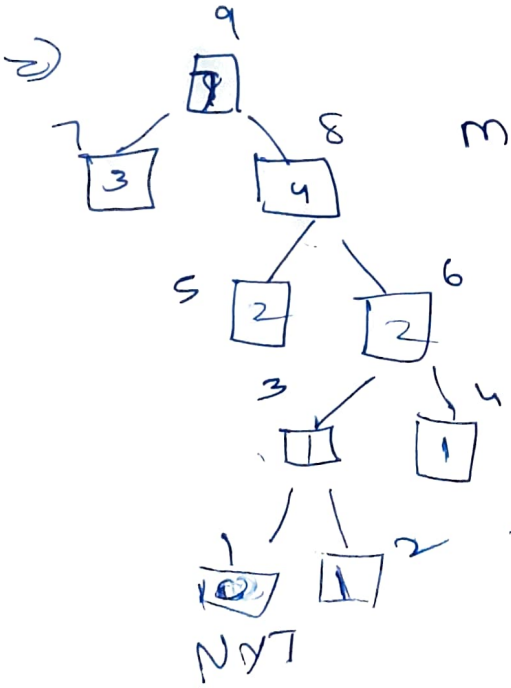
2) malay → 01100 0 00000 0 0101c 0 110 1110





malayr

01100 0 00000 0 0101c 0 110 111000



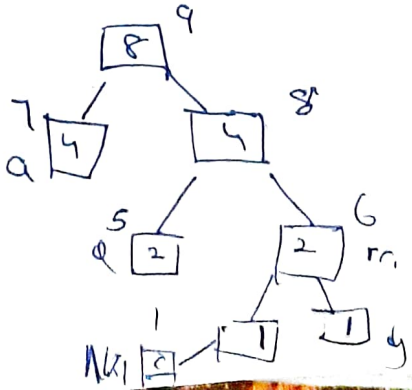
malayala

→ 01100 0 00000 0 01010 0 11. c

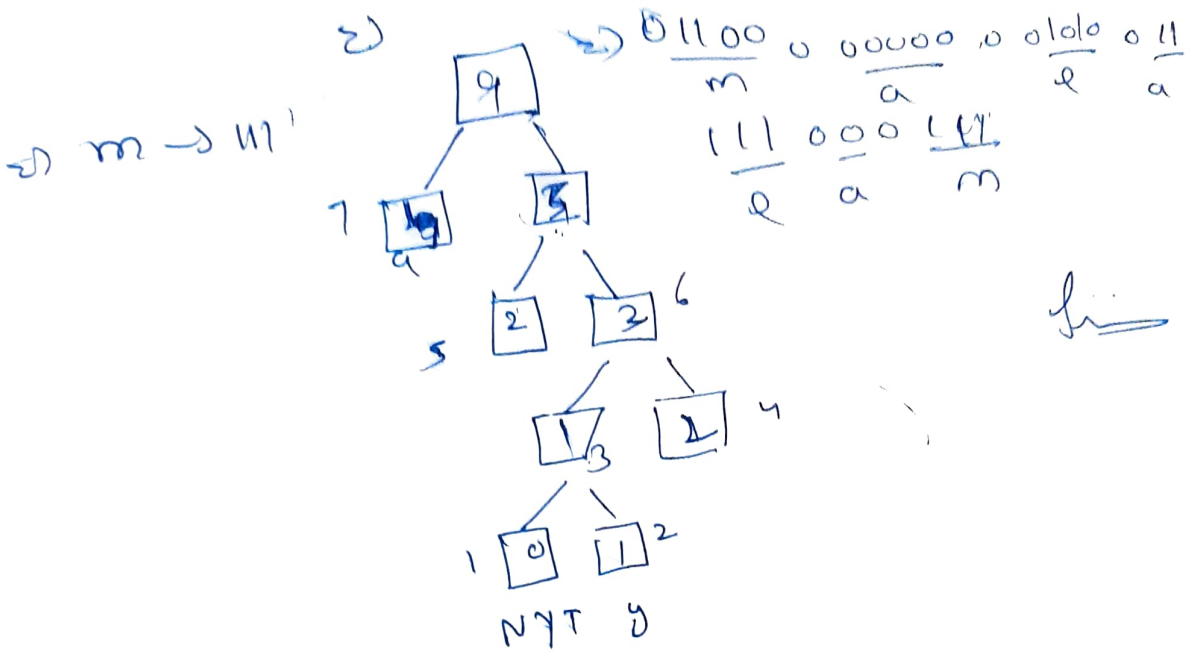
1110 0 0 0111.

→ ~~Prata~~

→ malayala a → 0



malayalam.


m → 111
9
7
5
4
2
3
6
1
1
4
3
1
0
1
2
NYT
m 01100
a 00000
l 01010
y 11011
111 000 111
l a m