

3) Difference ($A - B$)

* Unsorted

A	3	5	10	4	6	m
	0	1	2	3	4	

B	12	4	7	2	5	n
	0	1	2	3	4	

C	3	10	6									
	0	1	2	3	4	5	6	7	8	9		

$$m * n$$

$$= n * n$$

$$= n^2$$

$$\text{Time} : \Theta(n^2)$$

* Sorted

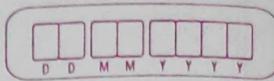
A	3	4	5	6	10	m
	0	1	2	3	4	

B	2	4	5	7	12	n
	0	1	2	3	4	

C	3	6	10								
	0	1	2	3	4	5	6	7	8	9	

$$\Theta(m+n)$$

$$\text{Time} : \Theta(n)$$



12). Finding Single missing Element in a Sorted Array

A	1	2	3	4	5	6	8	9	10	11	12
	0	1	2	3	4	5	6	7	8	9	10

$$\frac{n(n+1)}{2} = \frac{12(13)}{2} = 78$$

* $\text{Sum} = 0;$
 $\text{for } (i=0; i < 11; i++) \{$ // Using sum of first
 $\quad \text{Sum} = \text{Sum} + A[i];$ N Natural no.
 y

$$S = \frac{n * (n+1)}{2};$$

$S - \text{Sum};$
 $\text{printf} (" \text{Missing no is } \%d ", S - \text{Sum});$

A	6	7	8	9	10	11	13	14	15	16	17
	0	1	2	3	4	5	6	7	8	9	10

$$6-0 = 6$$

$$6+6=12$$

$$l = 6$$

$$7-1 = 6$$

$$h = 17$$

$$8-2 = 6$$

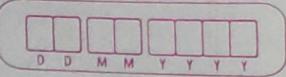
$$n = 11$$

$$9-3 = 6$$

:

$$11-5 = 6$$

$$13-6 = 7$$



```

* diff = l - 0;
for(i=0; i<n; i++) {
    if (A[i] - i != diff) {
        printf("missing el %d\n", i+diff);
        break;
}

```

y

Time - O(n)

122. Finding multiple missing Elements in an array

A	6	7	8	9	11	12	15	16	17	18	19
0	1	2	3	4	5	6	7	8	9	10	
6	6	6	6	7	7	9	9	9	9	9	

$$6+4=10$$

$$7+6=13$$

$$8+6=14$$

$$l=6$$

$$h=19$$

$$n=11$$

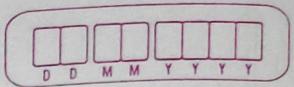
```

diff = l - 0;
for(i=0; i<n; i++) {
    if (A[i] - i != diff) {
        while (diff < A[i] - i)
            printf("%d\n", i+diff);
        diff++;
}

```

y

Time - O(n)



123. Finding missing Element in an Unsorted Array Method 2 using hashing

A	3	7	4	9	12	6	1	11	2	10	
	0	1	2	3	4	5	6	7	8	9	

$i=1$

$h=12$

$n=10$

$O(n^2)$

Time - $O(n)$

H	0	1	0	1	0	1	0	0	1	0	1	2	1	0
	0	1	2	3	4	5	6	7	8	9	10	11	12	

Hashtable / Bitset

```
for (i=0; i<n; i++) {
    H[AC[i]]++;
}
```

```
for (i=1; i<=h; i++) {
    if (H[i]==0)
        printf ("%d\n", i);
```

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

- n

$2n$
Time $O(n)$

124. Finding Duplicates in a Sorted array

A	3	6	8	8	10	12	15	15	15	20	
	0	1	2	3	4	5	6	7	8	9	

$n=10$

(a) \rightarrow wait

* Last Duplicate = 0 || for Printing

```

for (i=0; i<n-1; i++) {
    if (A[i] == A[i+1] && A[i] != LastDuplicate)
    {
        printf("%d\n", A[i]);
        LastDuplicate = A[i];
    }
}

```

y
3 Time - O(n)

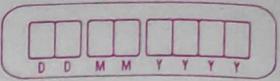
* for (i=0; i<n-1; i++) || for Counting

```

if (A[i] == A[i+1]) {
    j = i+1;
    while (A[j] == A[i])
        j++;
    printf("%d is appearing %d times",
           A[i], j-i);
}

```

y
3 i
y
Time - O(n)



125. Finding Duplicates in Sorted Array using Hashing

A	3	6	8	8	10	12	15	15	15	20	n
	0	1	2	3	4	5	6	7	8	9	+RN
Time - $O(n)$											$2n$

H	0	0	0	3	0	0	0	1	0	2	0	0	1	0	0	0	3	0	0	0	0
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	

0	1	2
---	---	---

```

for(i=0; i<n; i++) {
    H[A[i]]++;
}
for(i=0; i<=max; i++) {
    if(H[i] > 1)
        printf("%d %d", i, H[i]);
}
  
```

126. Finding Duplicates in an unsorted Array

A	8	3	6	4	6	5	6	8	1	2	7	$n = 10$
	0	1	2	3	4	5	6	7	8	9		

$$\frac{n(n-1)}{2} = n-1 + n-2 - \dots + 3+2+1$$

$$\frac{n^2-n}{2} = \text{Time} - O(n^2)$$

D	D	M	M

```

n - for (i=0; i<n-1; i++) {
    count = 1;
    if (A[i] != -1) {
        for (j = i+1; j < n; j++) {
            if (A[i] == A[j]) {
                count++;
                A[j] = -1;
            }
        }
        if (count > 1)
            printf("%d %d", A[i], count);
    }
}

```

Time - $O(n^2)$

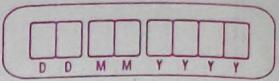
* Using Hashing

A	8	3	6	4	6	5	6	8	2	7	n = 10
0	1	2	3	4	5	6	7	8	9		

H	0	0	0	0	0	0	0	0	0	n
0	1	2	3	4	5	6	7	8		

Time - $O(n)$

Linear



127. Finding a Pair of Elements with sum k $(a+b = k)$

$$A = [6|3|R|10|16|7|5|2|9|14] \quad a+b=10$$

$j^0 \quad j \quad 1 \quad 2 \quad 3 \quad 4 \quad 5 \quad 6 \quad 7 \quad 8 \quad 9$
 $n-1 = 10 - 1 = 9$
 $n-1 + n-2 + \dots + 3 + 2 + 1 = \frac{n(n-1)}{2} = \frac{10(10-1)}{2} = 45$

```

    for (i=0; i<n-1; i++) {
        for (j=i+1; j<n; j++) {
            if (A[i]+A[j] == b)
                printf ("%d + %d = %d\n",
                        A[i], A[i], b);
    }
}

```

3 Time - $O(n^2)$ < Quadratic \Rightarrow wait

* Using Hashing

$$A = \begin{vmatrix} 6 & 3 & 8 & 10 & 16 & 7 & 5 & 2 & 9 & 14 \\ 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \end{vmatrix} \quad ad-bc = 0$$

H	0	0	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16		

```

for (i=0; i<n; i++) {
    if (H[K-A[i]] == 0) {
        printf("%d %d %d", A[i], K-A[i], k);
    }
}

```

$H[ACi]_2^{++}$

y

Time - O(n)

D	D	M	M

128 Finding a Pair of Elements with sum k in Sorted Array

A	1	3	4	5	6	8	9	10	12	14
	0	1	2	3	4	5	6	7	8	9
i										j

$a+b = 10$
 $n = 10$

$i=0, j=n-1;$

```
while (i < j) { // for (i=0, j=n-1; i<j; ) {
    if (A[i] + A[j] == k) {
```

```
        printf ("%d + %d = %d", A[i], A[j], k);
```

i++;

j--;

y

```
    else if (A[i] + A[j] < k)
```

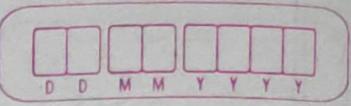
i++;

else

j--;

3

Time - O(n)



129. Finding Max and min in a single scan

A	5	8	3	9	6	2	10	7	-1	4
0	1	2	3	4	5	6	7	8	9	

$$\min = 5 \text{ } \cancel{8} \cancel{2} -1$$

$$\max = \cancel{5} \cancel{8} 9 \text{ } 10$$

$$\min = A[0];$$

$$\max = A[0];$$

for ($i = 1; i < n; i++$) {

 if ($A[i] < \min$)

$$\min = A[i];$$

 else (~~if~~ if ($A[i] > \max$)

$$\max = A[i];$$

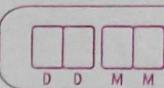
y

Compare $\Rightarrow n-1$ $10, \cancel{9}, \cancel{8}, \cancel{7}, \cancel{2}, 1$

Compare $\Rightarrow 2(n-1)$ $1, 2, \cancel{3}, \cancel{5}, \cancel{8}, \cancel{9}, 10$

Best $n-1$ $O(n)$

Worst $2(n-1)$ $O(n)$



130 Solutions for Quiz 3

1) $A[1 \dots n]$

$\text{reverse}(1, k)$

$\text{reverse}(k+1, n)$

$\text{reverse}(1, n)$

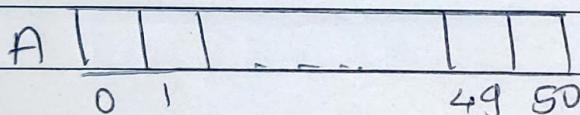
A	2	3	4	5	6	7	8
	4	3	2	5	6	7	8
	4	3	2	8	7	6	5
	5	6	7	8	2	3	4
	1	2	3	4	5	6	7
		*					
		K				n	

→ rotates left by k position

2) 500 elements

Range 0-100

frequency of elements > 50



Array of 50 numbers

D	D	M	M	Y

3) $a = [2, 4, 8, 9, 13, 16, 19, 23, 25]$

0	1	2	3	4	5	6	7	8
j								

$s = 10$

if ($a[j] - a[i] < s$) $j++;$

4) $A = [8, 7, 5, 12, 9, 3, 4]$

0	1	2	3	4	5	6

$O(1)$

5) $A = [7, 12, 15, 6, 4, 13, 9, 5]$

0	1	2	3	4	5	6	7

$n + n = 2n$

$O(n)$