

Lambda Cases (lcases)

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Contents

1	Introduction	2
2	Language Description	2
2.1	Basic Expressions	3
2.1.1	Literals and Identifiers	3
2.1.2	Tuples and Lists	4
2.2	Operators	5
2.2.1	Function Application Operators	5
2.2.2	Function Composition Operators	5
2.2.3	Arithmetic Operators	6
2.2.4	Comparison and Boolean Operators	7
2.2.5	Monadic Operators	7
2.2.6	Complete Table, Precedence, Associativity and Grammar	7
2.3	Functions	8
2.3.1	Parenthesis Function Application	8
2.3.2	Function Expressions	9
2.3.3	Special Function Parameters	9
2.3.4	Expressions	9
2.4	Value Definitions	10
2.5	Types	10
2.5.1	Type expressions	10
2.5.2	Tuple Types	12
2.5.3	Or Types	13
2.6	Type Logic	14
2.6.1	Type Predicate	14
2.6.2	Type Proposition	14
2.6.3	Type Theorem	14
2.7	Grammar	14
2.7.1	Tokens	14
2.7.2	Core Grammar	14
3	Notable differences from Haskell	15
4	Parser implimentation	15
4.1	AST Types	15
4.2	Parsers	15
5	Translation to Haskell	15
6	Running examples	15
7	Conclusion	15

1 Introduction

Haskell is a delightful language. Yet, for some reason, it doesn't seem to have its rightful place in terms of popularity in industry. Why is it so? Is it inherently hard to learn and therefore only the brave enough students and corporations dare to use it, or could it be that the syntax is perplexing to the amateur eye? It is my belief that with some syntax changes that give a greater familiarity to the new user, there would be no language more compelling than (the new) Haskell. In an attempt to achieve that familiarity, I present some new syntax, of which some is closer to the imperative/OOP style (to attract more already experienced programmers from these languages), some is closer to mathematics (in which most programmers should be experienced) and some is closer to natural language (in which we are all already experienced).

2 Language Description

An lcases program consists of a set of value, type and predicate definitions along with type theorems. The "main" value determines the program's behaviour. Constants and functions are all considered values and they have no real distinction other than the fact that functions have a function type and constants don't. Functions (just like "values") can be passed to other functions as arguments or can be returned as a result of other functions.

Program example: extended euclidean algorithm

```
// type definitions

tuple_type Coeffs
value (previous, current) : Int x Int

tuple_type GcdAndCoeffs
value (gcd, a, b) : Int x Int x Int

// algorithm

extended_euclidean: (Int, Int) => GcdAndCoeffs
  = (init_a_coeffs, init_b_coeffs) -> ee_recursion

init_a_coeffs, init_b_coeffs: all Coeffs
  = (1, 0), (0, 1)

ee_recursion: (Coeffs, Coeffs, Int, Int) => GcdAndCoeffs
  = (a_coeffs, b_coeffs, x, cases) =>
    0 => (x, a_coeffs.previous, b_coeffs.previous)
    y =>
      ee_recursion(next <- a_coeffs, next <- b_coeffs, y, x -> mod <- y)
      where
        next: Coeffs => Coeffs
          = fields => (current, previous - x / y * current)

// reading, printing and main

read_two_ints : (Int x Int)WithIO
  = print <- "Please give me 2 ints";
  get_line :> split_words o> apply(from_string)to_all o> ints =>
```

```

ints -> length -> cases =>
  2 => ints -> with_io
  ... => io_error <- "You didn't give me 2 ints"

print_gcd_and_coeffs : GcdAndCoeffs => (EmptyValue)WithIO
= fields => print("Gcd: " + gcd + "\nCcoefficients: a = " + a + ", b = " + b)

main : (EmptyValue)WithIO
= read_two_ints :> ints =>
  extended_euclidean(ints.1st, ints.2nd) -> print_gcd_and_coeffs

```

2.1 Basic Expressions

2.1.1 Literals and Identifiers

Literals

- *Examples*

```

1 2 17 42 -100
1.61 2.71 3.14 -1234.567
'a' 'b' 'c' 'x' 'y' 'z' '.' ',' '\n'
"Hello World!" "What's up, doc?" "Alrighty then!"

```

- *Description*

We have literals for the four basic types: Int, Real, Char, String. These are the usual integers, real numbers, characters and strings. The exact specification of literals is the same as in the Haskell report.

- *Grammar*

$\langle literal \rangle ::= \langle literal \rangle$

TODO add the grammar from the haskell report

Identifiers

- *Examples*

```

x y z
a1 a2 a3
funny_identifier
unnecessarily_long_identifier
apply()to_all

```

- *Description*

An identifier is a string used as the name of a value. It is first used in the definition of the value (see "value definition" section 2.4) and later used in the definition of other values that use that defined value. An identifier starts with a lower case letter and is followed by lower case letters or underscores. It is also possible to have a pair of parentheses in the middle of an identifier (see "parenthesis function application" section 2.3.1). Finally, an identifier can be ended with a digit.

- *Grammar*

$\langle identifier \rangle ::= [a-z]([a-z_]|'()'[a-z_])^*[0-9]$

2.1.2 Tuples and Lists

Tuples

- *Examples*

```
(1, "What's up, doc?")
(2, "Alrighty then!", 3.14)
(x, y, z, w)
(1, my_function, (x, y, z) => (x ^ 2 + y ^ 2 + z ^ 2) ^ (1/2))
```

- *Description*

Tuples are used to group many values (of possibly different types) into one. The type of the tuple can be either the product of the types of the subvalues or a defined `tuple_type` which is equivalent to the aforementioned product type i.e. the product type is in the definition of the `tuple_type` (see "tuple_type" section 2.5.2). For example, the type of the second example above could be:

```
Int x String x Real
```

or:

```
MyType
```

assuming "MyType" has been defined in a similar way to the following:

```
tuple_type MyType
value (my_int, my_string, my_real) : Int x String x Real
```

- *Grammar*

$\langle tuple \rangle ::= '(\langle value-expression \rangle (', \sqcup' \langle value-expression \rangle)+ ')$

Lists

- *Examples*

```
[1, 2, 17, 42, -100]
[1.61, 2.71, 3.14, -1234.567]
["Hello World!", "What's up, doc?", "Alrighty then!"]
[x, y, z, w]
```

- *Description*

Lists are used to group many values of the same type into one. The type of the list is `ListOf(A)s` where `A` is the type of every value inside. Therefore, the types of the first three examples are:

```
ListOf(Int)s
ListOf(Real)s
ListOf(String)s
```

And the last list is only legal if `x`, `y`, `z` and `w` all have the same type. Assuming they do and it's the type `T`, the type of the list is:

```
ListOf(T)s
```

- *Grammar*

$\langle list \rangle ::= '[' [\langle value-expression \rangle (', \sqcup' \langle value-expression \rangle)^* ']'$

2.2 Operators

2.2.1 Function Application Operators

The function application operators " \rightarrow " and " \leftarrow " are a different way to apply functions to arguments than the usual parenthesis function application. Each one applies the function from the corresponding direction. The operators are meant to look like arrows that point from the argument to the function. These operators are very useful for chaining many function applications without the clutter of having to open and close parentheses for each one of the functions.

For example, assuming we have the following functions with the behaviour suggested by their names and types:

```
apply()to_all : (A => B, ListOf(A)s) => ListOf(B)s
string_length: String => Int
filter : (A => Bool, ListOf(A)s) => ListOf(A)s
is_odd : Int => Bool
sum_ints : ListOf(Int)s => Int
```

And a list of strings:

```
strings : ListOf(String)s
```

Here is a simple way to get the number of characters in all the strings that have odd length:

```
chars_in_odd_length_strings : Int
  = strings -> apply(string_length)to_all -> filter(is_odd) -> sum_ints
```

Ofcourse this can be done equivalently using the other operator:

```
chars_in_odd_length_strings : Int
  = sum_ints <- filter(is_odd) <- apply(string_length)to_all <- strings
```

These operators can also be used together to put a function between two arguments if that function is commonly used that way in math (or if it looks better for a certain function). For example the "mod" function can be used like so:

```
x -> mod <- y
```

Instead of:

```
mod(x, y)
```

Operator	Type
\rightarrow	$(A, A \Rightarrow B) \Rightarrow B$
\leftarrow	$(A \Rightarrow B, A) \Rightarrow B$

2.2.2 Function Composition Operators

The function composition operators " $\circ>$ " and " $\circ<$ " are used to compose funtions, each one in the corresponding direction. The use of the letter 'o' is meant to be similar to the mathematical function composition symbol ' \circ ' and the symbols ' $>$ ', ' $<$ ' are used so that the operator points from the function which is applied first to the function which is applied second. A neat example using function composition is the following. Assuming we have the following functions with the behaviour suggested by their names and types:

```
split_words : String => ListOf(String)s
apply()to_all : (A => B, ListOf(A)s) => ListOf(B)s
reverse_word: String => String
merge_words : ListOf(String)s => String
```

We can reverse the all the words in a string like so:

```
reverse_words : String => String
  = split_words o> apply(reverse_word)to_all o> merge_words
```

Ofcourse this can be done equivalently using the other operator:

```
reverse_words : String => String
  = merge_words <o apply(reverse_word)to_all <o split_words
```

2.2.3 Arithmetic Operators

The usual arithmetic operators work as they are expected, similarly to mathematics and other programming languages. However, they are generalized. The examples below show their generality:

```
>> 1 + 1
2
>> 1 + 3.14
4.14
>> 'a' + 'b'
"ab"
>> 'w' + "ord"
"word"
>> "Hello " + "World!"
"Hello World!"
>> 5 * 'a'
"aaaaa"
>> 5 * "hi"
"hihihihihi"
>> "1,2,3" - ', '
"123"
```

The generality can also be seen from their types in the table below:

Operator	Type
\wedge	$(A)\text{ToThe}(B)\text{Gives}(C) \dashrightarrow (A, B) \Rightarrow C$
$*$	$(A)\text{And}(B)\text{MultiplyTo}(C) \dashrightarrow (A, B) \Rightarrow C$
$/$	$(A)\text{Divides}(B)\text{To}(C) \dashrightarrow (B, A) \Rightarrow C$
$+$	$(A)\text{And}(B)\text{AddTo}(C) \dashrightarrow (A, B) \Rightarrow C$
$-$	$(A)\text{SubtractsFrom}(B)\text{To}(C) \dashrightarrow (B, A) \Rightarrow C$

Let's analyze further the example of addition. The type can be read as such: the '+' operator has the type $(A, B) \Rightarrow C$, provided that the type proposition $(A)\text{And}(B)\text{AddTo}(C)$ holds. This proposition being true, means that addition has been defined for these three types (see section "type logic" 2.6 for more on type propositions). For example, by the examples above we can see that the following propositions are true (in the order of the examples):

```
(Int)And(Int)AddTo(Int)
(Int)And(Real)AddTo(Real)
(Char)And(Char)AddTo(String)
(Char)And(String)AddTo(String)
(Int)And(Char)MultiplyTo(String)
(Int)And(String)MultiplyTo(String)
(Char)SubtractsFrom(String)To(String)
```

This allows us to use the familiar arithmetic operators in types that are not necessarily numbers but it is somewhat intuitively obvious what they should do in those other types. Furthermore, their behaviour can be defined by the user for new user defined types!

2.2.4 Comparison and Boolean Operators

2.2.5 Monadic Operators

2.2.6 Complete Table, Precedence, Associativity and Grammar

Operator	Type	Description
\rightarrow	$(A, A \Rightarrow B) \Rightarrow B$	Right function application
\leftarrow	$(A \Rightarrow B, A) \Rightarrow B$	Left function application
$\circ>$	$(A \Rightarrow B, B \Rightarrow C) \Rightarrow (A \Rightarrow C)$	Right function composition
$<\circ$	$(B \Rightarrow C, A \Rightarrow B) \Rightarrow (A \Rightarrow C)$	Left function composition
\wedge	$(A)\text{ToThe}(B)\text{Gives}(C) \rightarrow (A, B) \Rightarrow C$	General exponentiation
$*$	$(A)\text{And}(B)\text{MultiplyTo}(C) \rightarrow (A, B) \Rightarrow C$	General multiplication
$/$	$(A)\text{Divides}(B)\text{To}(C) \rightarrow (B, A) \Rightarrow C$	General division
$+$	$(A)\text{And}(B)\text{AddTo}(C) \rightarrow (A, B) \Rightarrow C$	General addition
$-$	$(A)\text{SubtractsFrom}(B)\text{To}(C) \rightarrow (B, A) \Rightarrow C$	General subtraction
$=$ $/=$	$(A)\text{HasEquality} \rightarrow (A, A) \Rightarrow \text{Bool}$	Equality operators
$>$ $<$ $>=$ $<=$	$(A)\text{HasOrder} \rightarrow (A, A) \Rightarrow \text{Bool}$	Order operators
$\&$ $ $	$(\text{Bool}, \text{Bool}) \Rightarrow \text{Bool}$	Boolean operators
$:>$	$(E)\text{IsAnEnvironment} \rightarrow (E(A), A \Rightarrow E(B)) \Rightarrow E(B)$	Monad bind
$;$	$(E)\text{IsAnEnvironment} \rightarrow (E(A), E(B)) \Rightarrow E(B)$	Monad then

Operator	Precedence	Associativity
\rightarrow	10	Left
\leftarrow	9	Right
$\circ>$ $<\circ$	8	Left
\wedge	7	Right
$*$ $/$	6	Left
$+$ $-$	5	Left
$=$ $/=$ $>$ $<$ $>=$ $<=$	4	None
$\&$	3	Left
$ $	2	Left
$:>$ $;$	1	Left

Grammar

$$\begin{aligned}
\langle \text{operator-expression} \rangle &::= \langle \text{operator-argument} \rangle \langle \text{operator} \rangle \langle \text{right-operator-argument} \rangle \\
\langle \text{operator-argument} \rangle &::= \langle \text{literal} \rangle \mid \langle \text{identifier} \rangle \mid \langle \text{parentheses-function-application} \rangle \mid \langle \text{parentheses-expression} \rangle \\
\langle \text{right-operator-argument} \rangle &::= \langle \text{operator-argument} \rangle \mid \langle \text{function-expression} \rangle \\
\langle \text{operator} \rangle &::= \text{'_'} \langle \text{op} \rangle \text{'_'} \\
\langle \text{op} \rangle &::= \text{'->'} \mid \text{'<-'} \mid \text{'o>'} \mid \text{'<o'} \mid \text{'^'} \mid \text{'*'} \mid \text{'/'} \mid \text{'+'} \mid \text{'-'} \mid \text{'='} \mid \text{'/='} \mid \text{'>'} \mid \text{'<'} \mid \text{'>='} \mid \text{'<='} \mid \text{'\&'} \mid \text{'|'} \mid \text{'>:'} \mid \text{';'},
\end{aligned}$$

2.3 Functions

2.3.1 Parenthesis Function Application

Examples

```

f(x)
f(x, y, z)
(x)to_string
apply(f)to_all
apply(f)to_all(1)

```

Description

Function application in lcases can be done in many different ways in an attempt to maximize readability. In this section, we discuss the ways function application can be done with parenthesis. In the first two examples, we have the usual mathematical function application which is also used in most programming languages and should be familiar to the reader. That is, function application is done with the arguments of the function in parenthesis separated by commas and **appended** to the function identifier.

We extend this idea by allowing the arguments to be **prepended** to the function identifier (third example). Finally, it is also possible to have the arguments **inside** the function identifier provided the function has been **defined with parentheses inside the identifier**. For example, below is the definition of "apply()to_all":

```

apply()to_all: (A => B, ListOf(A)s) => ListOf(B)s
= (f, cases) =>
  empty => empty
  non_empty:l => non_empty:(f <- l.head, apply(f)to_all <- l.tail)

```

The actual definition doesn't matter at this point, what matters is that the identifier is "apply()to_all" with the parentheses **included**. This is very useful for defining functions where the argument in the middle makes the function application look and sound more like natural language.

It is possible that many parentheses pairs are present in a single function application (last example). The arguments are always inserted to the function from **left to right**. Therefore, when multiple parentheses pairs are present the arguments of the leftmost parentheses are inserted first then the next ones to the right and so on.

Grammar

$$\begin{aligned}
\langle \text{parenthesis-function-application} \rangle &::= \\
&\quad \langle \text{arguments} \rangle (\langle \text{identifier-with-arguments} \rangle \mid \langle \text{identifier} \rangle) \\
&\mid (\langle \text{identifier-with-arguments} \rangle \mid \langle \text{identifier} \rangle) \langle \text{arguments} \rangle \\
&\mid \langle \text{identifier-with-arguments} \rangle \\
\langle \text{arguments} \rangle &::= \text{'('} \langle \text{value-expression} \rangle (\text{'_',_'} \langle \text{value-expression} \rangle)^* \text{'_',_'}
\end{aligned}$$

$$\langle \text{identifier-with-arguments} \rangle ::= [\text{a-z}] \langle \text{id-char-or-paren-id-char} \rangle^* (\langle \text{arguments} \rangle [\text{a-z_}] \langle \text{id-char-or-paren-id-char} \rangle^*) + [[0-9]]$$

$$\langle \text{id-char-or-paren-id-char} \rangle ::= [\text{a-z_}] \mid '()' [\text{a-z_}]$$

2.3.2 Function Expressions

Examples

```
a => 17 * a + 42
```

```
(x, y, z) => (x ^ 2 + y ^ 2 + z ^ 2) ^ (1 / 2)
```

Description

Function expressions are used to define functions or pass anonymous functions as arguments to other functions. They are comprised by their parameters and their body. The parameters are either only one in which case a single identifier is used, or they are many in which case many identifiers are used in parentheses, separated by a comma. The parameters and the body are separated by an arrow. The body is an operator expression.

Grammar

$$\langle \text{function-expression} \rangle ::= \langle \text{parameters} \rangle \text{'>'} \langle \text{operator-expression} \rangle$$

$$\langle \text{parameters} \rangle ::= \langle \text{identifier} \rangle \mid \text{'('} \langle \text{identifier} \rangle (\text{' ,' } \langle \text{identifier} \rangle) + \text{'('}$$

2.3.3 Special Function Parameters

```
x => body
(x, y, z) => body
cases => body
(x, cases, z) => body
```

2.3.4 Expressions

Examples

```
42
```

```
x
```

```
funny_identifier
```

```
[1, 2, 3]
```

```
"Hello world!"
```

```
x -> mod <- y
```

```
1.61 * 2.71 + 3.14
```

```
a => 17 * a + 42
```

```
(x, y, z) => (x ^ 2 + y ^ 2 + z ^ 2) ^ (1 / 2)
```

```
n -> +1 -> ^2 -> *3 -> print
```

```
f(x, y, z) + g(1, 2, 3)
```

Description

The base of expressions, are literals and identifiers, those can be combined either with operators, or by normal function application with mathematical notation. Finally, on top of that there can be added one or more abstractions (parameters) in the beginning of the expressions with an arrow.

Grammar

2.4 Value Definitions

Examples

```
foo : Int
    = 42
```

```
val1, val2, val3 : Int, Bool, Char
    = 42, true, 'a'
```

```
int1, int2, int3 : all Int
    = 1, 2, 3
```

```
succ : Int => Int
    = +1
```

```
f : (Int, Int, Int) => Int
    = (a, b, c) => a + b * c
```

Description

To define a new value you give it a name, a type and an expression. It is possible to group value definitions by separating the names, the types and the expressions with commas. It is also possible to use the keyword "all" to give the same type to all the values.

Grammar

$\langle \text{value-definitions} \rangle ::= \langle \text{identifiers} \rangle \langle \text{':'} \rangle (\langle \text{types} \rangle \mid \text{'all'} \langle \text{type} \rangle) \langle \text{'='} \rangle \langle \text{value-expressions} \rangle$

$\langle \text{identifiers} \rangle ::= \langle \text{identifier} \rangle (\langle \text{' ,' } \rangle \langle \text{identifier} \rangle)^*$

$\langle \text{types} \rangle ::= \langle \text{type} \rangle (\langle \text{' ,' } \rangle \langle \text{type} \rangle)^*$

$\langle \text{value-expressions} \rangle ::= \langle \text{value-expression} \rangle (\langle \text{' ,' } \rangle \langle \text{value-expression} \rangle)^*$

2.5 Types

2.5.1 Type expressions

Examples

Int

String => String

Int x Int

Int x Int => Real

A => A

(A => B, B => C) => (A => C)

((A, A) => A, A, ListOf(A)s) => A

((B, A) => B, B, ListOf(A)s) => B

(T)HasStringRepresantion --> T => String

Description

Examples	Description
Int Char String	Base types
A => A (A => B, B => C) => (A => C)	Polymorphic types. A, B, C ... are type variables

Differences from Haskell

lcases	haskell	difference description
A => A	a => a	Type variables for polymorphic types are

Grammar

$\langle type \rangle ::= \langle type-application \rangle \mid \langle product-type \rangle \mid \langle function-type \rangle$

$\langle type-application \rangle ::= [\langle types-in-paren \rangle] \langle type-identifier \rangle (\langle types-in-paren \rangle ([A-Za-z])^*)^* [\langle types-in-paren \rangle]$

$\langle types-in-paren \rangle ::= ' (' \langle type \rangle (' , ' \langle type \rangle)^* ') '$

$\langle type-identifier \rangle ::= [A-Z] ([A-Za-z])^*$

$\langle product-type \rangle ::= \langle product-subtype \rangle (' \sqcap \sqcup ' \langle product-subtype \rangle)^+$

$\langle product-subtype \rangle ::= ' (' (\langle function-type \rangle \mid \langle product-type \rangle) ') ' \mid \langle type-application \rangle$

$$\langle \text{function-type} \rangle ::= \langle \text{input-types-expression} \rangle ' \sqsubseteq \Rightarrow \sqsubseteq ' \langle \text{one-type} \rangle$$

$$\langle \text{input-types-expression} \rangle ::= \langle \text{one-type} \rangle \mid \langle \text{two-or-more-types-in-paren} \rangle$$

$$\langle \text{one-type} \rangle ::= \langle \text{type-application} \rangle \mid \langle \text{product-type} \rangle \mid '(\langle \text{function-type} \rangle)'$$

$$\langle \text{two-or-more-types-in-paren} \rangle ::= '(\langle \text{type} \rangle (', ' \langle \text{type} \rangle) +)'$$

2.5.2 Tuple Types

Definition Examples

```
tuple_type Name
value (first_name, last_name) : String x String

tuple_type ClientInfo
value (name, age, nationality) : Name x Int x String

tuple_type Date
value (day, month, year) : Int x Int x Int

tuple_type (A)And(B)
value (a_value, b_value) : A x B

tuple_type (ExprT)WithPosition
value (expr, line, column) : ExprT x Int x Int
```

Usage Examples

```
giorgos_info : ClientInfo
= (("Giorgos", "Papadopoulos"), 42, "Greek")

john_info : ClientInfo
= (("John", "Doe"), 42, "American")

name_to_string : Name => String
= fields => "First Name: " + first_name + "\nLast Name: " + last_name

print_name_and_nationality : ClientInfo => (EmptyValue)WithIO
= fields => print(name -> name_to_string + "\nNationality: " + nationality)

print_error_in_expr : (SomeDefinedExprT)WithPosition => (EmptyValue)WithIO
= ewp =>
  print(
    "Error in the expression:" + es +
    "\nAt the position: (" + ls + ", " + cs + ")"
  )
  where
  es, ls, cs : all String
  = ewp.expr->to_string, ewp.line->to_string, ewp.column->to_string
```

Description

Tuple types group many values into a single value. They are specified by their name, the names of their sub-values and the types of their subvalues. They generate projection functions for all of their subvalues by using a `'.'` before the name of the subvalue. For example the `ClientInfo` type above generates the following functions:

```
.name : ClientInfo => String
.age : ClientInfo => Int
.nationality : ClientInfo => String
```

These functions shall be named "postfix functions" as they can just be appended to their argument.

Definition Grammar

$\langle \text{tuple-type-definition} \rangle ::=$
 $\text{'tuple_type_'} \langle \text{type-application} \rangle \text{'\nvalue_'} \text{'('} \langle \text{identifier} \rangle \text{'(, ' } \langle \text{identifier} \rangle \text{'*) ' } \text{'_:'} \langle \text{product-type} \rangle$

2.5.3 Or Types

Examples

```
or_type Bool
values true | false
```

```
or_type Possibly(A)
values the_value:A | no_value
```

```
or_type ListOf(A)s
values non_empty:HeadAndTailOf(A)s | empty
```

```
tuple_type HeadAndTailOf(A)s
value (head : A, tail : ListOf(A)s)
```

```
is_empty : ListOf(A)s => Bool
= cases =>
  empty => true
  non_empty:anything => false
```

```
get_head : ListOf(A)s => Possibly(A)
= cases =>
  empty => no_value
  non_empty:list => the_value:list.head
```

Description

Values of an `or_type` are one of many cases that possibly have other values inside. The cases which have other values inside are followed by a semicolon and the type of the internal value. The same syntax can be used for matching that particular case in a function using the "cases" syntax, with the difference that after the colon, we write the name given to the value inside. `Or_types` and basic types are the only types on which the "cases" syntax can be used. The cases of an `or_type` which have a value inside create functions. For example, the case "non_empty" of a list creates the function "non_empty:" for which we can say:

```
non_empty: : HeadAndTailOf(A)s => ListOf(A)s
```

Similarly:

```
the_value: : A => Possibly(A)
```

These functions shall be named "prefix functions" as they are prepended to their argument. For example:

```
head_and_tail : HeadAndTailOf(Int)s
  = (1, [2, 3, 4])
```

```
list : ListOf(Int)s
  = non_empty:head_and_tail
```

These functions can be used like any other function as arguments to other functions. For example:

```
heads_and_tails_to_lists : ListOf(HeadAndTailOf(A)s)s => ListOf(ListOf(A)s)s
  = apply(non_empty:)to_each
```

Definition Grammar

$\langle or\text{-}type\text{-}definition \rangle ::=$
 $\text{'or_type_'} \langle type\text{-}application \rangle \text{'\nvalues_'} \langle identifier \rangle [\text{'\text{:}\text{'}} \langle type \rangle] (\text{'_|_'} \langle identifier \rangle [\text{'\text{:}\text{'}} \langle type \rangle])^*$

2.6 Type Logic

2.6.1 Type Predicate

2.6.2 Type Proposition

2.6.3 Type Theorem

2.7 Grammar

2.7.1 Tokens

Keywords

cases use_fields tuple_type or_type

Value names

Type names

2.7.2 Core Grammar

Program

$\langle program \rangle ::= (\langle value\text{-}definitions \rangle | \langle type\text{-}def \rangle)^+$

$\langle value\text{-}definitions \rangle ::= \langle identifiers \rangle \text{'_:_'} (\langle types \rangle | \text{'all'} \langle type \rangle) \text{'\n_=\text{'}} \langle value\text{-}expressions \rangle$

$\langle identifiers \rangle ::= \langle identifier \rangle (\text{'_,'} \langle identifier \rangle)^*$

$\langle types \rangle ::= \langle type \rangle (\text{'_,'} \langle type \rangle)^*$

$\langle value\text{-}expressions \rangle ::= \langle value\text{-}expression \rangle (\text{'_,'} \langle value\text{-}expression \rangle)^*$

Types

Value Expressions

$\langle \text{value-expression} \rangle \quad ::= [\langle \text{input-expr} \rangle] \langle \text{cases-or-where} \rangle \mid \langle \text{op-expr} \rangle$

$\langle \text{cases-or-where} \rangle \quad ::= \langle \text{cases-expr} \rangle \mid \langle \text{where-expr} \rangle$

$\langle \text{where-expr} \rangle \quad ::= \text{'let'} \langle \text{spicy-nl} \rangle (\langle \text{value-definitions} \rangle \langle \text{spicy-nls} \rangle) + \text{'in'} \langle \text{value-expression} \rangle \langle \text{spicy-nl} \rangle$

$\langle \text{cases-expr} \rangle \quad ::= \text{'cases'} (\langle \text{case} \rangle) + [\langle \text{default-case} \rangle]$

3 Notable differences from Haskell

4 Parser implimentation

The parser was implemented using the parsec library.

4.1 AST Types

4.2 Parsers

5 Translation to Haskell

6 Running examples

7 Conclusion

8 To be removed or incorporated

Addition/Subtraction:

```
+ : (A)HasAddition => (A, A) => A
- : (A)HasSubtraction => (A, A) => A
```

Equality and ordering:

```
= : (A)HasEquality => (A, A) => Bool
<= : (A)HasOrder => (A, A) => Bool
>= : (A)HasOrder => (A, A) => Bool
```

$(\text{fmap})\langle \text{inside} \rangle \text{ --- } (W)\text{IsAWrapper} \text{ --> } (A \Rightarrow B, W(A)) \Rightarrow W(B)$ — Apply inside operator
 $\langle \text{<*>} \rangle \langle \text{wrapped_inside} \rangle \text{ --- } (W)\text{IsAWrapper} \text{ --> } (W(A \Rightarrow B), W(A)) \Rightarrow W(B)$ — Order operators

better as postfix functions

Examples in Haskell

```
data ClientInfo =  
    ClientInfoC String Int String  
  
data WithPosition a =  
    WithPositionC a Int Int  
  
data Pair a b =  
    PairC a b
```

Examples in Haskell

```
{-# language LambdaCase #-}  
  
data Bool =  
    Ctrue | Cfalse  
  
data Possibly a =  
    Cwrapper a | Cnothing  
  
data ListOf_s a =  
    Cnon_empty (NonEmptyValueListOf_s a) | Cempty  
  
data NonEmptyValueListOf_s a =  
    CNonEmptyValueListOf_s a (ListOf_s a)  
  
is_empty :: ListOf_s a => Bool  
is_empty = \case  
    Cempty => Ctrue  
    Cnon_empty (CNonEmptyValueListOf_s head tail) => Cfalse  
  
get_head :: ListOf_s a => Possibly a  
get_head = \case  
    Cempty => Cnothing  
    Cnon_empty (CNonEmptyValueListOf_s head tail) => Cwrapper head
```

Examples in Haskell

```
foo :: Int  
foo = 42  
  
val1 :: Int  
val1 = 42  
val2 :: Bool  
val2 = true  
val3 :: Char  
val3 = 'a'  
  
int1 :: Int  
int1 = 1  
int2 :: Int  
int2 = 2  
int3 :: Int
```



```
int3 = 3
```

```
succ :: Int => Int
succ = \x => x + 1
```

```
f :: Int => Int => Int => Int
f = \a b c => a + b * c
```

Or Types the following have automatically generated functions:

```
is_case:
```



Hi

- *Examples*

- *Description*

hi

- *Grammar*

$\langle identifier \rangle ::= [a-z] ([a-z_] | '()' [a-z_])^* [[0-9]]$