Lambda Cases (lcases)

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Contents

1	Introduction	1						
2 Language Description								
	2.1 Values	3						
	2.1.1 Literals	3						
	2.1.2 Identifiers	3						
	2.1.3 Operators	3						
	2.1.4 Expressions	3						
	2.1.5 Definitions	4						
	2.2 Types	5						
	2.2.1 Type expressions	5						
	2.2.2 Tuple Types	6						
	2.2.3 Or Types	7						
	2.3 Type Logic	8						
	2.3.1 Type Predicate	8						
	2.3.2 Type Theorem	8						
	2.4 Grammar	8						
	2.4.1 Tokens	8						
	2.4.2 Core Grammar	9						
3	Parser implimentation	9						
	B.1 AST Types	10						
	3.2 Parsers	10						
4	Translation to Haskell	10						
5	Running examples	10						
6	Conclusion	10						
7	To be removed or incorporated	10						

1 Introduction

Haskell is a delightful language. Yet, for some reason, it doesn't seem to have it's rightful place in terms of popularity in industry. Why is it so? Is it inherently hard to learn and therefore only the brave enough students and corporations dare to use it, or could it be that the syntax is perplexing to the amateur eye? It is my belief that with some syntax changes that give a greater familiarity to the new user, there would be no language more compelling than (the new) Haskell. In an attempt to achieve that familiarity, I present some new syntax, of which some is closer to the imperative/OOP style (to attract more already experienced programmers from these languages), some is closer to mathematics (in which most programmers should be experienced) and some is closer to natural language (in which we are all already experienced).

2 Language Description

An leases program consists of a set of value, type and predicate definitions along with type theorems. The "main" value determines the program's behaviour. Constants and functions are all considered values and they have no real distinction other than the fact that functions have a function type and constants don't. Functions (just like "values") can be passed to other functions as arguments or can be returned as a result of other functions.

Program example: extended euclidean alogirthm

```
// type definitions
tuple_type PrevCoeffs
value (prev_prev, prev : Int, Int)
tuple_type GcdAndCoeffs
value (gcd, a, b : Int, Int, Int)
// algorithm
extended_euclidean: (Int, Int) -> GcdAndCoeffs
  = (init_a_coeffs, init_b_coeffs) ==> ee_recursion
init_a_coeffs, init_b_coeffs: all PrevCoeffs
  = (1, 0), (0, 1)
ee_recursion: (PrevCoeffs, PrevCoeffs, Int, Int) -> GcdAndCoeffs
  = (a_coeffs, b_coeffs, x, cases) ->
    0 -> (x, a_coeffs.prev_prev, b_coeffs.prev_prev)
    y ->
      ee_recursion(next <== a_coeffs, next <== b_coeffs, y, x ==> mod <== y)
      where
      next: PrevCoeffs -> PrevCoeffs
        = fields -> (prev, prev_prev - x ==> div <== y * prev)
// reading, printing and main
read_two_ints : (Int x Int)WithIO
  = print <== "Please give me 2 ints";</pre>
    get_line >>= split_words o> apply(from_string)to_all o> ints ->
    ints ==> length ==> cases ->
      2 -> ints ==> with_env
      ... -> io_error <== "You didn't give me 2 ints"
print_gcd_and_coeffs : GcdAndCoeffs -> (Empty)WithIO
  = fields -> print("Gcd: " + gcd + "\nCoefficients: a = " + a + ", b = " + b)
main : (Empty)WithIO
  = read_two_ints >>= ints ->
    extended_euclidean(ints.1st, ints.2nd) ==> print_gcd_and_coeffs
```

2.1 Values

2.1.1 Literals

Literals are the same as haskell

Examples	Type
1, 2, 17, 42, -100	Int
1.61, 2.71, 3.14, -1234.567	Real
'a', 'b', 'c', 'x', 'y', 'z', '.', ',', '\n'	Char
[1, 2, 3], ['a', 'b', 'c'], [1.61, 2.71, 3.14]	ListOf(Int)s, ListOf(Char)s, ListOf(Real)s
"Hello World!", "What's up, doc?", "Alrighty then!"	String

2.1.2 Identifiers

An identifier is a string of lower case letters or underscore.

Grammar

 $\langle \mathit{identifier} \rangle ::= (\ [a-z_] \)^*$

2.1.3 Operators

Operator	Type	Description	Associativity
==>	(A, A -> B) -> B	Right function application	Left
<==	(A -> B, A) -> B	Left function application	Left
0>	(A -> B, B -> C) -> (A -> C)	Right function composition	Left
<0	(B -> C, A -> B) -> (A -> C)	Left function composition	Right
^	(A) ToThe (B) Gives $(C) \Rightarrow (A, B) \rightarrow C$	General exponentiation	Right
*	(A)And(B)MultiplyTo(C) \Rightarrow (A, B) \rightarrow C	General multiplication	Left
/	(A)Divides(B)To(C) \Rightarrow (A, B) \rightarrow C	General division	Left
+	$(A)And(B)AddTo(C) \Rightarrow (A, B) \rightarrow C$	General addition	Left
-	(A)SubtractsFrom(B)To(C) \Rightarrow (B, A) \rightarrow C	General subtraction	Left
= /=	(A)HasEquality => (A, A) -> Bool	Equality operators	None
> < >= <=	(A)HasOrder \Rightarrow (A, A) \Rightarrow Bool	Order operators	None
&	(Bool, Bool) -> Bool	Boolean operators	Left
>>=	(E)IsAnEnvironment => (E(A), A -> E(B)) -> E(B)	Monad bind	Left
;	(E)IsAnEnvironment \Rightarrow (E(A), E(B)) \rightarrow E(B)	Monad then	Left

2.1.4 Expressions

Examples

42

х

```
funny_identifier
[1, 2, 3]
"Hello world!"

1.61 * 2.71 + 3.14

a -> 17 * a + 42

(x, y, z) -> (x^2 + y^2 + z^2)^(1/2)

n==>(+ 1)==>(^2)==>(* 3)==>print

f(x, y, z) + g(1, 2, 3)
```

Description

The base of expressions, are literals and identifiers, those can be combined either with operators, or by normal function application with mathematical notation. Finally, on top of that there can be added one of more abstractions (parameters) in the beginning of the expressions with an arrow.

Grammar

2.1.5 Definitions

Examples

Description

To define a new value you give it a name, a type and an expression. It is possible to group value definitions by seperating the names, the types and the expressions with commas. It is also possible to use the keyword "all" to give the same type to all the values.

Grammar

$$\langle value\text{-}definitions \rangle ::= \langle identifiers \rangle \text{`\sqcup:$} \text{`$(\langle types \rangle | `all' $\langle type \rangle) ``\n$$} \text{`$(value\text{-}expressions)$}$$

$$\langle identifiers \rangle ::= \langle identifier \rangle \text{ (`$,$$$} \text{`$(identifier) })^*$$

$$\langle types \rangle ::= \langle type \rangle \text{ (`$,$$$$} \text{`$(type) })^*$$

$$\langle value\text{-}expressions \rangle ::= \langle value\text{-}expression \rangle \text{ (`$,$$$} \text{`$(value\text{-}expression) })^*$$

Abstractions

2.2 Types

2.2.1 Type expressions

Examples

Int

String -> String

Int x Int

Int x Int -> Real

A -> A

$$(A \rightarrow B, B \rightarrow C) \rightarrow (A \rightarrow C)$$

$$((A, A) \rightarrow A, A, ListOf(A)s) \rightarrow A$$

$$((B, A) \rightarrow B, B, ListOf(A)s) \rightarrow B$$

(T)HasStringRepresantion => T -> String

Description

Examples	Description
Int	
Char	Base types
String	
A -> A	Polymorphic types. A, B, C are type variables
(A -> B, B -> C) -> (A -> C)	

Differences from Haskell

lcases	haskell	difference description
A -> A	a -> a	Type variables for polymorphic types are

Grammar

```
 \langle type-application \rangle \mid \langle product-type \rangle \mid \langle function-type \rangle \\ \langle type-application \rangle ::= \left[ \langle types-in-paren \rangle \right] \langle type-identifier \rangle (\langle types-in-paren \rangle (\left[ A-Za-z \right])^*)^* \left[ \langle types-in-paren \rangle \right] \\ \langle types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle)^* ')' \\ \langle type-identifier \rangle ::= \left[ A-Z \right] (\left[ A-Za-z \right])^* \\ \langle product-type \rangle ::= \langle product-subtype \rangle (' \sqcup x_{\sqcup}' \langle product-subtype \rangle) + \\ \langle product-subtype \rangle ::= '(' (\langle function-type \rangle \mid \langle product-type \rangle) ')' \mid \langle type-application \rangle \\ \langle function-type \rangle ::= \langle input-types-expression \rangle ::= \langle one-type \rangle \mid \langle two-or-more-types-in-paren \rangle \\ \langle one-type \rangle ::= \langle type-application \rangle \mid \langle product-type \rangle \mid '(' \langle function-type \rangle ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ', ' \langle type \rangle) + ')' \\ \langle two-or-more-types-in-paren \rangle ::= '(' \langle type \rangle (', ', ' \langle ty
```

2.2.2 Tuple Types

Definition Examples

Description

Tuple types group many values into a single value. They are specified by their name, the names of their subvalues and the types of their subvalues. They generate projection functions for all of their subvalues by using a '.' before the name of the subvalue. For example the ClientInfo type above generates the following functions:

```
.name : ClientInfo -> String
.age : ClientInfo -> Int
.nationality : ClientInfo -> String
```

These functions shall be named "postfix functions" as the can just be appended to their argument.

Definition Grammar

```
\langle tuple-type-definition \rangle ::=
\text{`tuple\_type}_{\sqcup}' \ \langle type-application \rangle \ \text{``nvalue}_{\sqcup}' \ \text{``} \ \langle identifier \rangle \ \text{(`,}_{\sqcup}' \ \langle identifier \rangle) * \ \text{`)'} \ \text{`}_{\sqcup}:_{\sqcup}' \ \langle product-type \rangle
```

2.2.3 Or Types

Examples

Description

Values of an or_type are one of many cases that possibly have other values inside. The cases which have other values inside are followed by a semicolon and the type of the internal value. The same syntax can be used for matching that particular case in a fucntion using the "cases" syntax, with the difference that after the colon, we write the name given to the value inside. Or_types and basic types are the only types on which the "cases" syntax can be used. The cases of an or_type which have a value inside create functions. For example, the case "non_empty" of a list creates the function "non_empty:" for which we can say:

```
non_empty: : HeadAndTailOf(A)s -> ListOf(A)s
Similarly:
```

the_value: : A -> Possibly(A)

These functions shall be named "prefix functions" as they are prepended to their argument. For example:

These functions can be used like any other function as arguments to other functions. For example:

```
\label{lem:heads_and_tails_to_lists} $$ = apply(non_empty:)to_each $$
```

Definition Grammar

```
\langle or\text{-}type\text{-}definition\rangle ::= \\ \text{`or\_type}_{\square}\text{`}\langle type\text{-}application\rangle \text{``} \text{`} \text{`} \text{`} \text{`} \text{`} \text{identifier}\rangle \text{ [`:'}\langle type\rangle \text{]} \text{ (`}_{\square}\text{'}_{\square}\text{'} \text{`} \text{`} \text{identifier}\rangle \text{ [`:'}\langle type\rangle \text{]})*}
```

- 2.3 Type Logic
- 2.3.1 Type Predicate
- 2.3.2 Type Theorem
- 2.4 Grammar
- 2.4.1 Tokens

Keywords

```
cases use_fields tuple_type or_type
```

Value names

Type names

2.4.2 Core Grammar

Program

```
 \langle program \rangle \qquad ::= (\langle value\text{-}definitions \rangle \mid \langle type\text{-}def \rangle) + \\ \langle value\text{-}definitions \rangle \qquad ::= \langle identifiers \rangle \; `_{\square} :_{\square} ' \; (\langle types \rangle \mid \text{`all'} \; \langle type \rangle) \; `_{\square} = ' \; \langle value\text{-}expressions \rangle \\ \langle identifiers \rangle \qquad ::= \langle identifier \rangle \; (\; `,_{\square} ' \; \langle identifier \rangle \; )^* \\ \langle types \rangle \qquad ::= \langle type \rangle \; (\; `,_{\square} ' \; \langle type \rangle \; )^* \\ \langle value\text{-}expressions \rangle \qquad ::= \langle value\text{-}expression \rangle \; (\; `,_{\square} ' \; \langle value\text{-}expression \rangle \; )^*
```

Types

Value Expressions

3 Parser implimentation

The parser was implemented using the parsec library.

- 3.1 AST Types
- 3.2 Parsers
- 4 Translation to Haskell
- 5 Running examples
- 6 Conclusion
- 7 To be removed or incorporated

```
Addition/Subtraction:
```

```
+ : (A)HasAddition => (A, A) -> A
- : (A)HasSubtraction => (A, A) -> A

Equality and ordering:

= : (A)HasEquality => (A, A) -> Bool
<= : (A)HasOrder => (A, A) -> Bool
>= : (A)HasOrder => (A, A) -> Bool

(fmap)<inside> — (W)IsAWrapper => (A -> B, W(A)) -> W(B) — Apply inside operator
(<*>)<wrapped_inside> — (W)IsAWrapper => (W(A -> B), W(A)) -> W(B) — Order operators
```

better as postfix functions

Examples in Haskell

```
data ClientInfo =
   ClientInfoC String Int String

data WithPosition a =
   WithPositionC a Int Int

data Pair a b =
   PairC a b

Examples in Haskell
{-# language LambdaCase #-}
```

```
data Bool =
   Ctrue | Cfalse

data Possibly a =
   Cwrapper a | Cnothing

data ListOf_s a =
   Cnon_empty (NonEmptyListOf_s a) | Cempty

data NonEmptyListOf_s a =
```

```
CNonEmptyListOf_s a (ListOf_s a)
is_empty :: ListOf_s a -> Bool
is_empty = \case
  Cempty -> Ctrue
  Cnon_empty (CNonEmptyListOf_s head tail) -> Cfalse
get_head :: ListOf_s a -> Possibly a
get_head = \case
  Cempty -> Cnothing
  Cnon_empty (CNonEmptyListOf_s head tail) -> Cwrapper head
Examples in Haskell
```

```
foo :: Int
foo = 42
val1 :: Int
val1 = 42
val2 :: Bool
val2 = true
val3 :: Char
val3 = 'a'
int1 :: Int
int1 = 1
int2 :: Int
int2 = 2
int3 :: Int
int3 = 3
succ :: Int -> Int
succ = \x -> x + 1
f :: Int -> Int -> Int
f = \langle a b c - \rangle a + b * c
```

Or Types the following have automatically generated functions:

is_case:

