

Lambda Cases(lcases)

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1 Introduction

Haskell is a delightful language. Yet, for some reason, it doesn't seem to have it's rightful place in terms of popularity in industry. Why is it so? Is it inherently hard to learn and therefore only the brave enough students and corporations dare to use it, or could it be that the syntax is perplexing to the amateur eye? It is my belief that with some syntax changes that give a greater familiarity to the new user, there would be no language more compelling than (the new) Haskell. In an attempt to achieve that familiarity, I present some new syntax, of which some is closer to the imperative/OOP style (to attract more already experienced programmers from these languages), some is closer to mathematics (in which most programmers should be experienced) and some is closer to natural language (in which we are all already experienced).

2 Language Description

Program example: extended euclidean algorithm

```
// type definitions

tuple_type PrevCoeffs
value (prev_prev, prev : Int, Int)

tuple_type GcdAndCoeffs
value (gcd, a, b : Int, Int, Int)

// algorithm functions

extended_euclidean: (Int, Int) -> GcdAndCoeffs
= (init_a_coeffs, init_b_coeffs) ==> ee_recursion

init_a_coeffs, init_b_coeffs: all PrevCoeffs
= (1, 0), (0, 1)

ee_recursion: (PrevCoeffs, PrevCoeffs, Int, Int) -> GcdAndCoeffs
= (a_coeffs, b_coeffs, x, cases) ->
  0 -> (x, a_coeffs.prev_prev, b_coeffs.prev_prev)
  y ->
    ee_recursion(next <== a_coeffs, next <== b_coeffs, y, x ==> mod <== y)
    where
      next: PrevCoeffs -> PrevCoeffs
      = fields -> (prev, prev_prev - x ==> div <== y * prev)

// reading, printing and main

read_two_ints : (Int x Int)WithIO
= print <== "Please give me 2 ints";
  get_line >>= split_words o> apply(from_string)to_all o> ints ->
  ints ==> length ==> cases ->
    2 -> ints ==> with_env
    ... -> io_error <== "You didn't give me 2 ints"

print_gcd_and_coeffs : GcdAndCoeffs -> (Empty)WithIO
= fields ->
  print("Gcd: " + gcd + "\nCoefficients: a = " + a + ", b = " + b)

main : (Empty)WithIO
= read_two_ints >>= ints ->
  extended_euclidean(ints.1st, ints.2nd) ==> print_gcd_and_coeffs
```

2.1 Values

As with Haskell, an `lcases` program is a set of value/type (and other) definitions, with the "main" value determining the program's behaviour. Values (constants) and functions have no real distinction other than the fact functions have a function type. Functions (just like values) can be passed to other functions as arguments or can be returned as a result of other functions.

2.1.1 Literals

Literals are the same as haskell

Examples	Type
1, 2, 17, 42, -100	Int
1.61, 2.71, 3.14, -1234.567	Real
'a', 'b', 'c', 'x', 'y', 'z', '.', ',', '\n'	Char
[1, 2, 3], ['a', 'b', 'c'], [1.61, 2.71, 3.14]	ListOf(Int)s, ListOf(Char)s, ListOf(Real)s
"Hello World!", "What's up, doc?", "Alrighty then!"	String

2.1.2 Identifiers

An identifier is a string of lower case letters or underscore.

Grammar

$\langle identifier \rangle ::= (\langle lower-case-letter \rangle | _)^*$

2.1.3 Operators

Operator	Type	Description
<code>==></code>	$(A, A \rightarrow B) \rightarrow B$	Right function application
<code><==</code>	$(A \rightarrow B, A) \rightarrow B$	Left function application
<code>o></code>	$(A \rightarrow B, B \rightarrow C) \rightarrow (A \rightarrow C)$	Right function composition
<code><o</code>	$(B \rightarrow C, A \rightarrow B) \rightarrow (A \rightarrow C)$	Left function composition
<code>^</code>	$(A) \text{ToThe}(B) \text{Gives}(C) \Rightarrow (A, B) \rightarrow C$	General exponentiation
<code>*</code>	$(A) \text{And}(B) \text{MultiplyTo}(C) \Rightarrow (A, B) \rightarrow C$	General multiplication
<code>/</code>	$(A) \text{Divides}(B) \text{To}(C) \Rightarrow (A, B) \rightarrow C$	General division
<code>+</code>	$(A) \text{And}(B) \text{AddTo}(C) \Rightarrow (A, B) \rightarrow C$	General addition
<code>-</code>	$(A) \text{SubtractsFrom}(B) \text{To}(C) \Rightarrow (B, A) \rightarrow C$	General subtraction
<code>= / =</code>	$(A) \text{HasEquality} \Rightarrow (A, A) \rightarrow \text{Bool}$	Equality operators
<code>> < >= <=</code>	$(A) \text{HasOrder} \Rightarrow (A, A) \rightarrow \text{Bool}$	Order operators
<code>& </code>	$(\text{Bool}, \text{Bool}) \rightarrow \text{Bool}$	Boolean operators
<code>>>=</code>	$(E) \text{IsAnEnvironment} \Rightarrow (E(A), A \rightarrow E(B)) \rightarrow E(B)$	(Monad) right bind
<code>=<<</code>	$(E) \text{IsAnEnvironment} \Rightarrow (A \rightarrow E(B), E(A)) \rightarrow E(B)$	(Monad) left bind

2.1.4 Expressions

Examples

42

x

funny_identifier

[1, 2, 3]

"Hello world!"

1.61 * 2.71 + 3.14

a -> 17 * a + 42

(x, y, z) -> (x² + y² + z²)^(1/2)

n==>(+ 1)==>(^2)==>(* 3)==>print

f(x, y, z) + g(1, 2, 3)

Description

The base of expressions, are literals and identifiers, those can be combined either with operators, or by normal function application with mathematical notation. Finally, on top of that there can be added one or more abstractions (parameters) in the beginning of the expressions with an arrow.

Grammar

2.1.5 Definitions

Examples

```
foo : Int
    = 42
```

```
val1, val2, val3 : Int, Bool, Char
    = 42, true, 'a'
```

```
int1, int2, int3 : all Int
    = 1, 2, 3
```

```
succ : Int -> Int
    = x -> x + 1
```

```
f : (Int, Int, Int) -> Int
    = (a, b, c) -> a + b * c
```

Description

To define a new value you give it a name, a type and an expression. It is possible to group value definitions by separating the names, the types and the expressions with commas. It is also possible to use the keyword "all" to give the same type to all the values.

Grammar

$\langle \text{value-definitions} \rangle ::= \langle \text{identifiers} \rangle \text{'}_:_ \text{' } (\langle \text{types} \rangle \mid \text{'all' } \langle \text{type} \rangle) \text{'}\backslash \mathbf{n}_ = \text{' } \langle \text{value-expressions} \rangle$

$\langle \text{identifiers} \rangle ::= \langle \text{identifier} \rangle (\text{'},_ \text{' } \langle \text{identifier} \rangle)^*$

$\langle \text{types} \rangle ::= \langle \text{type} \rangle (\text{'},_ \text{' } \langle \text{type} \rangle)^*$

$\langle \text{value-expressions} \rangle ::= \langle \text{value-expression} \rangle (\text{'},_ \text{' } \langle \text{value-expression} \rangle)^*$

Abstractions

x -> body
(x, y, z) -> body
cases -> body
(x, cases, z) -> body

2.2 Types

2.2.1 Type expressions

Examples

Int

String -> String

Int x Int

Int x Int -> Real

A -> A

(A, A) -> A

((A, A) -> A, A, ListOf(A)s) -> A

((B, A) -> B, B, ListOf(A)s) -> B

(R)IsReducable => ((B, A) -> B, B, R(A)) -> B

(T)HasStringRepresantion => T -> String

Description

Differences from Haskell

lcases	haskell	difference description
A -> A	a -> a	Type variables for polymorphic types are

Grammar

$\langle type \rangle ::= \langle func-type \rangle \mid \langle prod-type \rangle \mid \langle type-app \rangle$
 $\langle func-type \rangle ::= \langle input-types-expr \rangle \text{ ' } \sqcup \text{ } \rightarrow \text{ ' } \sqcup \text{ ' } \langle output-type \rangle$
 $\langle prod-type \rangle ::= \langle prod-sub-type \rangle (\text{ ' } \sqcup \text{ ' } \langle prod-sub-type \rangle) +$
 $\langle type-app \rangle ::= [\langle t-inputs \rangle \text{ ' } == \text{ ' }] \langle type-name \rangle [\text{ ' } < == \text{ ' } \langle t-inputs \rangle]$
 $\langle input-types-expr \rangle ::= \langle many-ts-in-paren \rangle \mid \langle one-type \rangle$
 $\langle output-type \rangle ::= \langle prod-type \rangle \mid \langle type-app \rangle$
 $\langle prod-sub-type \rangle ::= \text{ ' } (\text{ ' } \langle func-type \rangle \mid \langle prod-type \rangle \text{ ' }) \text{ ' } \mid \langle type-app \rangle$
 $\langle one-type \rangle ::= \text{ ' } (\text{ ' } \langle func-type \rangle \text{ ' }) \text{ ' } \mid \langle prod-type \rangle \mid \langle type-app \rangle$
 $\langle t-inputs \rangle ::= \langle many-ts-in-paren \rangle \mid \text{ ' } (\text{ ' } \langle type \rangle \text{ ' }) \text{ ' } \mid \langle type-name \rangle$
 $\langle many-ts-in-paren \rangle ::= \text{ ' } (\text{ ' } \langle type \rangle (\text{ ' } , \text{ ' } \langle type \rangle) + \text{ ' }) \text{ ' }$

2.2.2 Tuple Types

Examples

```
tuple_type ClientInfo
value (name, age, nationality : String, Int, String)

john_info: ClientInfo
  = ("John Doe", 42, "American")

giorgos_info: ClientInfo
  = ("Giorgos Papadopoulos", 42, "Greek")

print_name_and_nationality : ClientInfo -> (Empty)WithIO
  = ci -> print("Name: " + ci.name + "\nNationality: " + ci.nationality)

tuple_type (ExprT)WithPosition
value (expr, line, column : ExprT, Int, Int)

print_error_in_expr : (SomeDefinedExprT)WithPosition -> (Empty)WithIO
  = ewp ->
    print(
      "Error in the expression:" + es +
      "\nAt the position: (" + ls + ", " + cs + ")"
    )
  where
    es, ls, cs : all String
      = ewp.expr==>to_string, ewp.line==>to_string, ewp.column==>to_string
```

Description

Tuple types group many values into a single value. They are specified by their name, the names of their sub-values and the types of their subvalues. They generate projection functions for all of their subvalues by using a `'.'` before the name of the subvalue. For example the `ClientInfo` type above generates the following functions:

```
.name : ClientInfo -> String
.age : ClientInfo -> Int
.nationality : ClientInfo -> String
```

These functions shall be named "postfix functions" as they can just be appended to their argument.

2.2.3 Or Types

Examples

```
or_type Bool
values true | false
```

```
or_type Possibly(A)
values the_value:A | nothing
```

```
or_type ListOf(A)s
values non_empty:HeadAndTailOf(A)s | empty
```

```
tuple_type HeadAndTailOf(A)s
value (head : A, tail : ListOf(A)s)
```

```
is_empty : ListOf(A)s -> Bool
= cases ->
  empty -> true
  non_empty:anything -> false
```

```
get_head : ListOf(A)s -> Possibly(A)
= cases ->
  empty -> nothing
  non_empty:list -> the_value:list.head
```

Description

Values of an `or_type` are one of many cases that possibly have other values inside. The cases which have other values inside are followed by a semicolon and the type of the internal value. The same syntax can be used for matching that particular case in a function using the "cases" syntax, with the difference that after the colon, we write the name given to the value inside. `Or_types` and basic types are the only types on which the "cases" syntax can be used. The cases of an `or_type` which have a value inside create functions. For example, the case "non_empty" of a list creates the function "non_empty:" for which we can say:

```
non_empty: : HeadAndTailOf(A)s -> ListOf(A)s
```

Similarly:

```
the_value: : A -> Possibly(A)
```

These functions shall be named "prefix functions" as they are prepended to their argument. For example:

```
head_and_tail : HeadAndTailOf(Int)s
= (1, [2, 3, 4])
```

```
list : ListOf(Int)s
= non_empty:head_and_tail
```

These functions can be used like any other function as arguments to other functions. For example:

```
heads_and_tails_to_lists : ListOf(HeadAndTailOf(A)s)s -> ListOf(ListOf(A)s)s
= apply(non_empty:)to_each
```

2.3 Type Logic

Type Predicate

Type Theorem

2.4 Grammar

2.4.1 Tokens

Keywords

cases use_fields tuple_type or_type

Value names

Type names

$\langle type-name \rangle ::= \langle upper-case-letter \rangle (\langle upper-case-letter \rangle | \langle lower-case-letter \rangle)^*$

2.4.2 Core Grammar

Program

$\langle program \rangle ::= (\langle value-definitions \rangle | \langle type-def \rangle)^+$

$\langle value-definitions \rangle ::= \langle identifiers \rangle ' _ : _ ' (\langle types \rangle | 'all' \langle type \rangle) ' _ = ' \langle value-expressions \rangle$

$\langle identifiers \rangle ::= \langle identifier \rangle (' , ' \langle identifier \rangle)^*$

$\langle types \rangle ::= \langle type \rangle (' , ' \langle type \rangle)^*$

$\langle value-expressions \rangle ::= \langle value-expression \rangle (' , ' \langle value-expression \rangle)^*$

Types

Value Expressions

$\langle value-expression \rangle ::= [\langle input-expr \rangle] \langle cases-or-where \rangle | \langle op-expr \rangle$

$\langle cases-or-where \rangle ::= \langle cases-expr \rangle | \langle where-expr \rangle$

$\langle where\text{-}expr \rangle ::= 'let' \langle spicy\text{-}nl \rangle (\langle value\text{-}definitions \rangle \langle spicy\text{-}nls \rangle) + 'in' \langle value\text{-}expression \rangle \langle spicy\text{-}nl \rangle$

$\langle cases\text{-}expr \rangle ::= 'cases' (\langle case \rangle) + [\langle default\text{-}case \rangle]$

3 Parser implimentation

The parser was implemented using the parsec library.

3.1 AST Types

3.2 Parsers

4 Translation to Haskell

5 Running examples

6 Conclusion

7 To be removed or incorporated

Addition/Subtraction:

```
+ : (A)HasAddition => (A, A) -> A
- : (A)HasSubtraction => (A, A) -> A
```

Equality and ordering:

```
= : (A)HasEquality => (A, A) -> Bool
<= : (A)HasOrder => (A, A) -> Bool
>= : (A)HasOrder => (A, A) -> Bool
```

$(fmap)\langle inside \rangle \text{ --- } (W)IsAWrapper \Rightarrow (A \rightarrow B, W(A)) \rightarrow W(B)$ — Apply inside operator
 $(\langle *> \rangle)\langle wrapped_inside \rangle \text{ --- } (W)IsAWrapper \Rightarrow (W(A \rightarrow B), W(A)) \rightarrow W(B)$ — Order operators

better as postfix functions

Examples in Haskell

```
data ClientInfo =
  ClientInfoC String Int String
```

```
data WithPosition a =
  WithPositionC a Int Int
```

```
data Pair a b =
  PairC a b
```

Examples in Haskell

```
{-# language LambdaCase #-}
```

```
data Bool =  
    Ctrue | Cfalse
```

```
data Possibly a =  
    Cwrapper a | Cnothing
```

```
data ListOf_s a =  
    Cnon_empty (NonEmptyListOf_s a) | Cempty
```

```
data NonEmptyListOf_s a =  
    CNonEmptyListOf_s a (ListOf_s a)
```

```
is_empty :: ListOf_s a -> Bool  
is_empty = \case  
    Cempty -> Ctrue  
    Cnon_empty (CNonEmptyListOf_s head tail) -> Cfalse
```

```
get_head :: ListOf_s a -> Possibly a  
get_head = \case  
    Cempty -> Cnothing  
    Cnon_empty (CNonEmptyListOf_s head tail) -> Cwrapper head
```

Examples in Haskell

```
foo :: Int  
foo = 42
```

```
val1 :: Int  
val1 = 42  
val2 :: Bool  
val2 = true  
val3 :: Char  
val3 = 'a'
```

```
int1 :: Int  
int1 = 1  
int2 :: Int  
int2 = 2  
int3 :: Int  
int3 = 3
```

```
succ :: Int -> Int  
succ = \x -> x + 1
```

```
f :: Int -> Int -> Int -> Int  
f = \a b c -> a + b * c
```

Or Types the following have automatically generated functions:

```
is_case:
```