

## I. Unions and Intersections

(a)  $(X \cap Y)$  is a subset of  $X \Rightarrow$  always countable

(b)  $Y$  is a subset of  $(X \cup Y) \Rightarrow$  always  $\times$

(c)  $A_i$  is a subset of  $\bigcup_{i \in X} A_i \Rightarrow$  always  $\times$

(d) Sometimes:

$$\textcircled{1} X = \mathbb{N}. \quad A_i = \mathbb{R} / \{i\} \quad \bigcap_{i \in X} A_i = \mathbb{R} / \mathbb{N} \quad \text{uncountable}$$

$$\textcircled{2} X = \mathbb{N} \quad A_i = \{(i, x) \mid x \in \mathbb{R}\} \quad \bigcap_{i \in X} A_i = \emptyset \quad \text{countable}$$

(e) Sometimes:

$$\textcircled{1} Y = \mathbb{R}. \quad B_i = \mathbb{N}. \quad \bigcup_{i \in Y} B_i = \mathbb{N} \quad \text{countable}$$

$$\textcircled{2} Y = \mathbb{R} \quad B_i = \mathbb{N} \cup \{i\} \quad \bigcup_{i \in Y} B_i = \mathbb{R} \quad \text{uncountable}$$

$$\textcircled{f)} \bigcap_{i \in Y} B_i \text{ is a subset of } B_i. \Rightarrow \text{always } \checkmark$$

## II. Count it

(a) finite

(b)  $8\mathbb{Z}$ , countably infinite.

(c) ?

(d) when given the length, we can calculate the rank in some order

eg. apple > apply ...

So we can link them with 1, 2, ...

(e) ?

(f) ?

### III. Fixed points

(a) Suppose  $\text{Find}(p)$  finds the fixed point of program  $p$ .

(b) Consider program  $F(p)$ :

$$F(p) = \begin{cases} p, & \text{if } p \text{ halts on input } 0 \\ \text{null}, & \text{if } p \text{ doesn't halt on input } 0 \end{cases}$$

if  $\text{Find}(p)$  can calculate the fixed point of  $F$ , then

$\text{Find}(p)$  can find all the programs that halts on input 0  $\times$

(c) if we have  $\text{FindFixed}(F)$ :

$$\text{FindFixed}(F) = \begin{cases} n & \text{if } \exists n, F(n) = n \\ \text{null} & \text{otherwise.} \end{cases} \quad (\text{skipped})$$

#### IV. Unprogrammable Programs

(a).  $\text{def TestHalt}(G, x)$   
     $\text{def } F(x):$   
         $G(x) \quad \Rightarrow \quad X$   
         $\text{return } 0$   
     $\text{return } P(F, x, 0)$

(b)  $\text{def } G(x):$   
     $\text{if } x == 0:$   
         $\text{return } 0.$   
     $\text{else:}$   
         $\text{return } F(x).$

if  $F(x)$  halts on 0, then  $F, G$  halt on the same set of inputs

if  $F(x)$  doesn't halt on 0, then  $F, G$  don't halt on the same ---

$\Rightarrow$ . we can define  $\text{TestHalt}(F) = P(F, G) \quad X$   
to test if  $F$  halts on 0.

#### V. Counting.

(a)  $3 \times 2^{18}$

(b)  $\frac{104!}{2^{52}}$

