

Chap. 12) I/O Systems

경희대학교 컴퓨터공학과

조 진 성

Revisited: Computer System Operation

I/O operation

- ✓ I/O request via I/O instruction
 - Direct I/O vs. Memory-mapped I/O
 - Communicates with registers in I/O controller
 - Typically, IR (Instruction Register) & DR (Data Register)
- ✓ I/O method
 - Polling vs. Interrupt
 - Programmed I/O vs. DMA (Direct Memory Access)

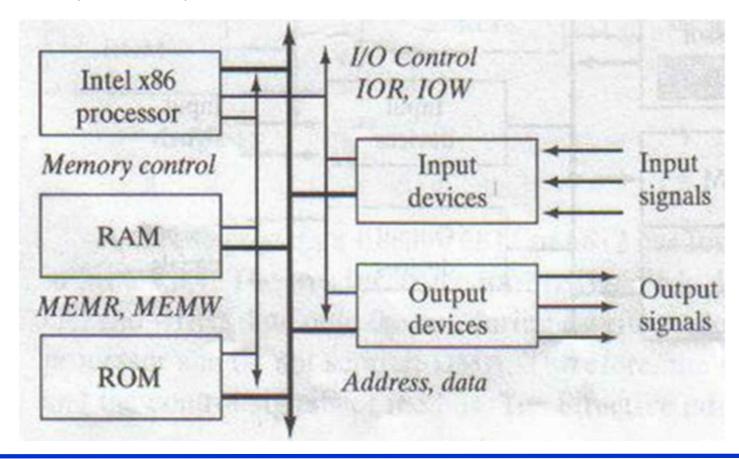


Direct I/O

Separate addresses: memory vs. I/O

Separate instructions

- ✓ MEMR / MEMW (Load / Store)
- ✓ IOR / IOW (IN / OUT)

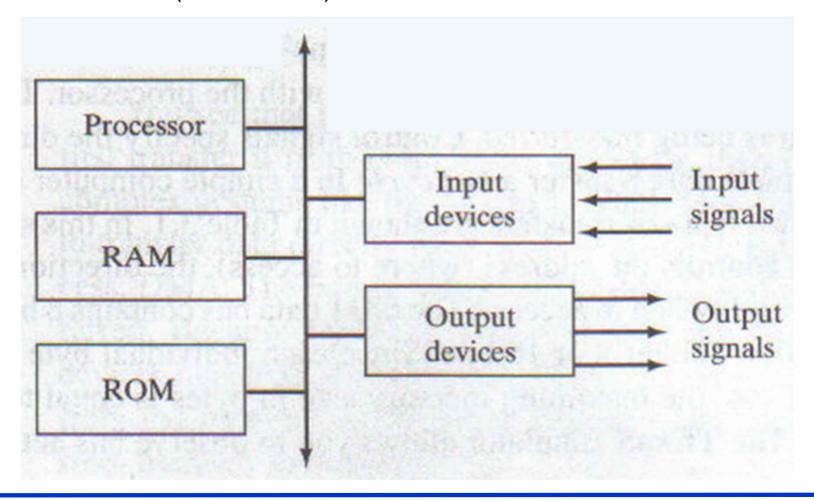




Memory-Mapped I/O

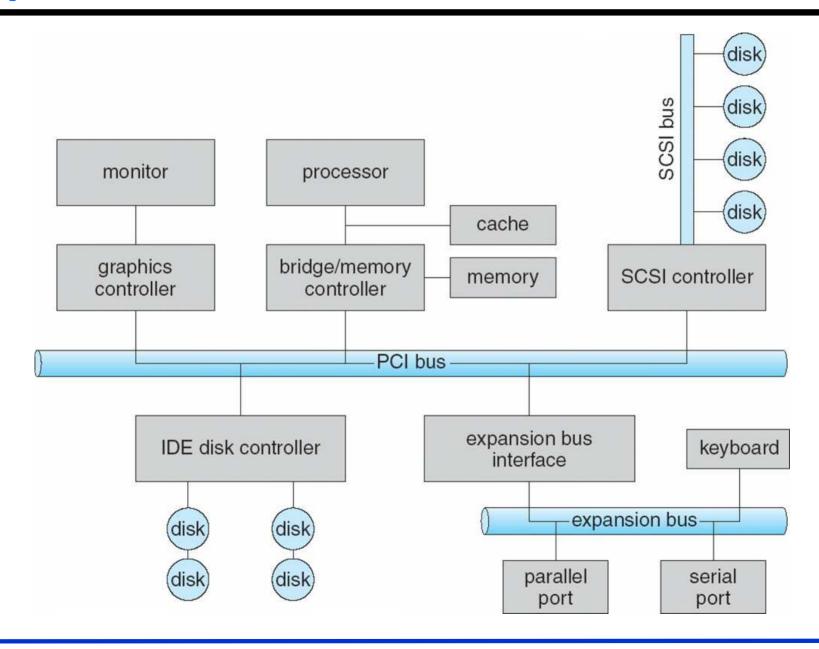
Integrated addresses: memory & I/O Integrated instructions

✓ MEMR / MEMW (Load / Store)





A Typical PC Bus Structure





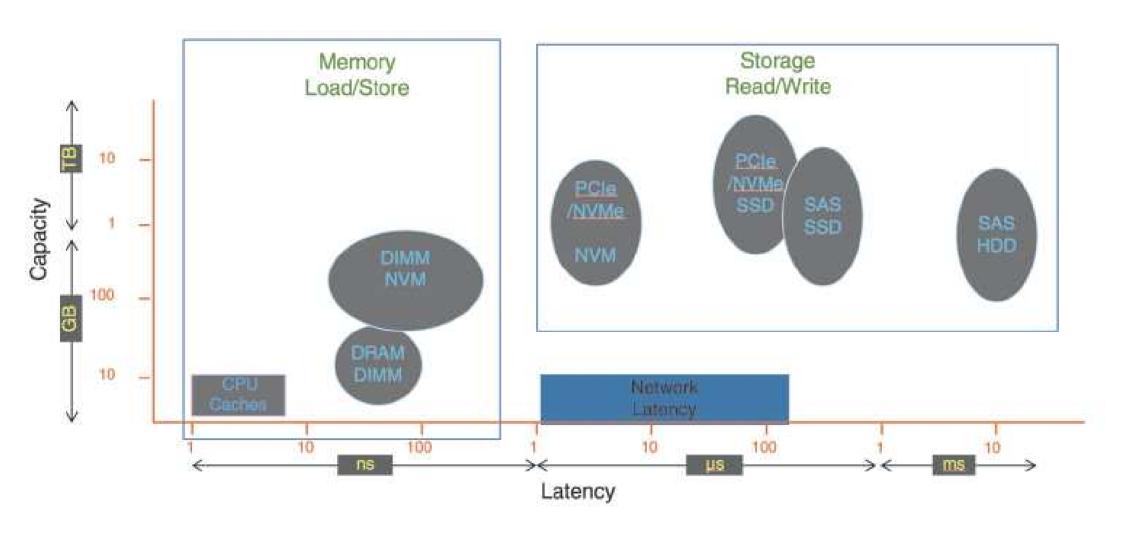
I/O Device Controller

Device controller (or host adapter)

- ✓ I/O devices have components:
 - Mechanical component
 - Electronic component
- ✓ The electronic component is the device controller
 - May be able to handle multiple devices
- ✓ Controller's tasks
 - Convert serial bit stream to block of bytes
 - Perform error correction as necessary
 - Make available to main memory

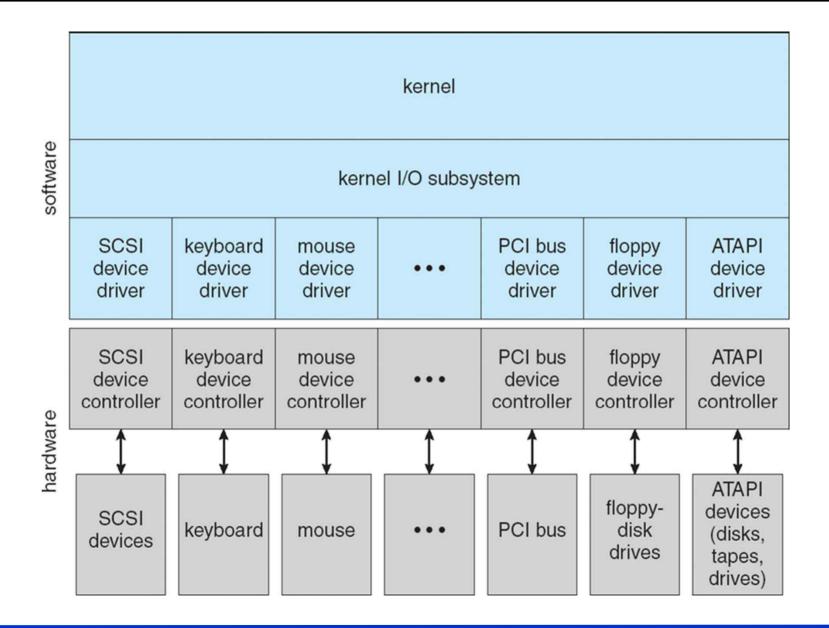


I/O Performance of Storage and Network Latency



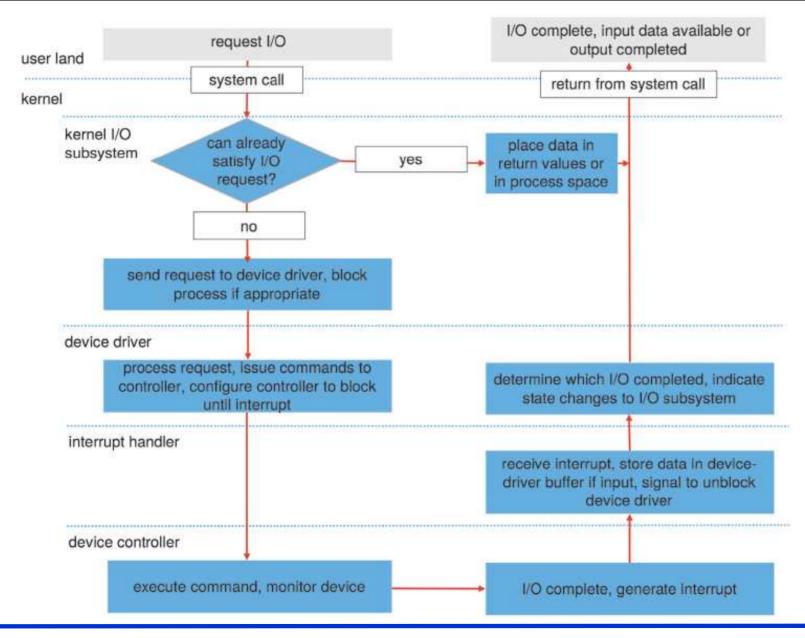


Kernel I/O Structure



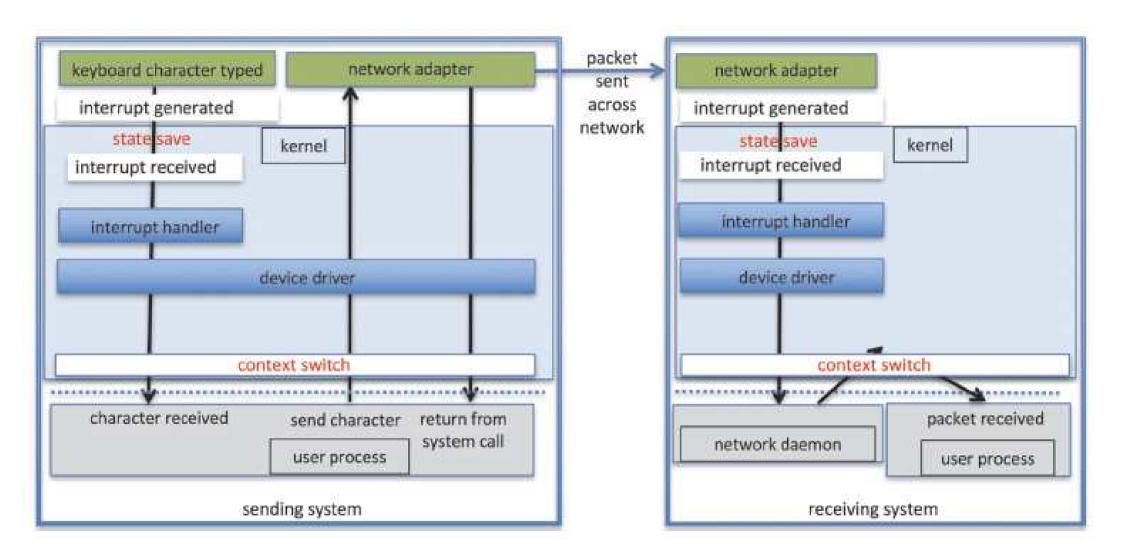


Life Cycle of An I/O Request



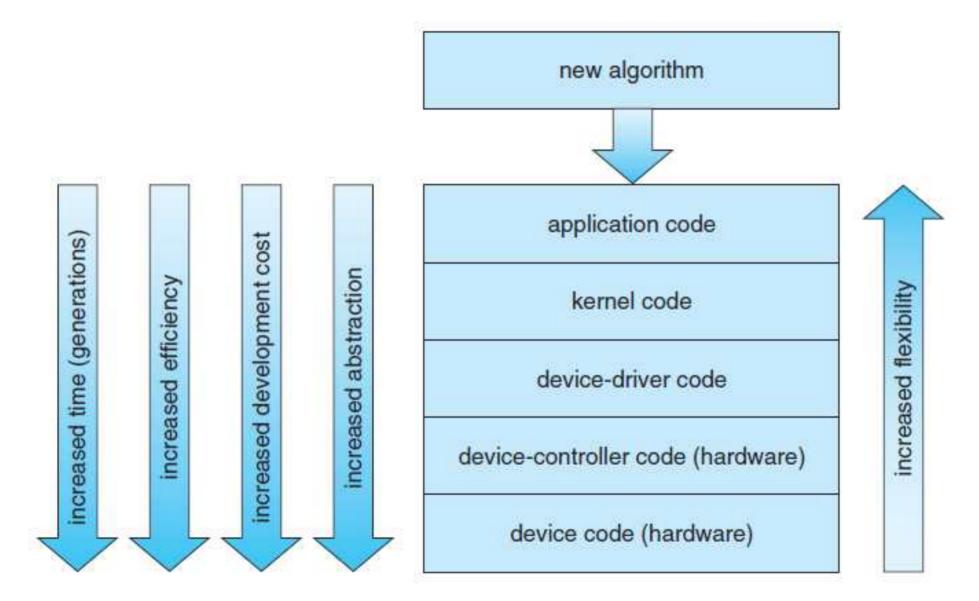


Intercomputer Communications





Device-Functionality Progression





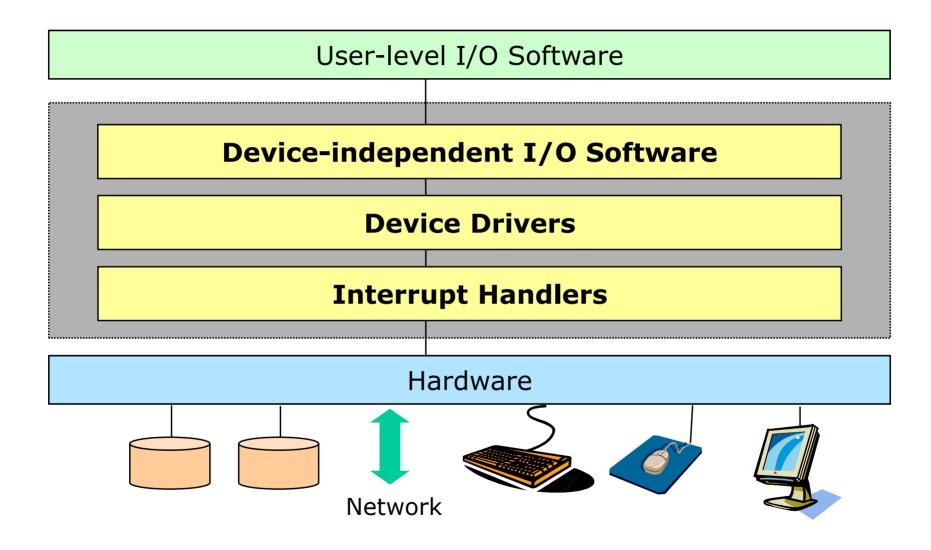
Goals of I/O Software

Goals

- ✓ Device independence
 - Programs can access any I/O device without specifying device in advance
- ✓ Uniform naming
 - Name of a file or device should simply be a string or an integer
- ✓ Error handling
 - Handle as close to the hardware as possible
- ✓ Synchronous vs. asynchronous
 - blocked transfers vs. interrupt-driven
- ✓ Buffering
 - Data coming off a device cannot be stored in final destination
- ✓ Sharable vs. dedicated devices
 - Disks vs. tape drives
 - Unsharable devices introduce problems such as deadlocks



I/O Software Layers





Interrupt Handlers

Handling interrupts

Critical actions

: Acknowledge an interrupt to the PIC.

: Reprogram the PIC or the device controller.

: Update data structures accessed by both the device and the processor.

Reenable interrupts

Noncritical actions

: Update data structures that are accessed only by the processor.

(e.g., reading the scan code from the keyboard)

Return from interrupts

Noncritical deferred actions

: Actions may be delayed.

: Copy buffer contents into the address space of some process (e.g., sending the keyboard line buffer to the terminal handler process).

Bottom half (Linux)



Device Drivers

Device drivers

- ✓ Device-specific code to control each I/O device interacting with deviceindependent I/O software and interrupt handlers
- ✓ Requires to define a well-defined model and a standard interface of how they interact with the rest of the OS
- ✓ Implementing device drivers:
 - Statically linked with the kernel
 - Selectively loaded into the system during boot time
 - Dynamically loaded into the system during execution (especially for hot pluggable devices)



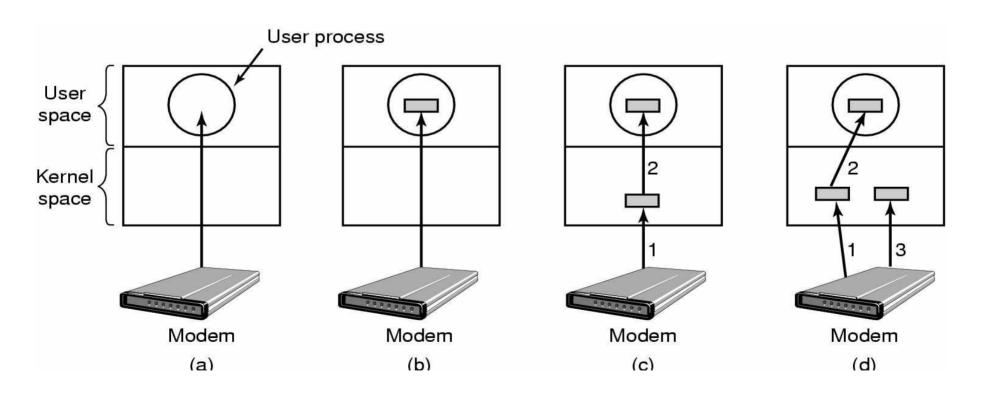
Uniform interfacing for device drivers

- ✓ In Unix, devices are modeled as special files
 - They are accessed through the use of system calls such as open(), read(), write(), close(), ioctl(), etc.
 - A file name is associated with each device
- ✓ Major device number locates the appropriate driver
 - Minor device number (stored in i-node) is passed as a parameter to the driver in order to specify the unit to be read or written
- ✓ The usual protection rules for files also apply to I/O devices



Buffering

- √ (a) Unbuffered
- √ (b) Buffered in user space
- √ (c) Buffered in the kernel space
- √ (d) Double buffering in the kernel





Error reporting

- ✓ Many errors are device-specific and must be handled by the appropriate driver, but the framework for error handling is device independent
- ✓ Programming errors vs. actual I/O errors
- ✓ Handling errors
 - Returning the system call with an error code
 - Retrying a certain number of times
 - Ignoring the error
 - Killing the calling process
 - Terminating the system



Allocating and releasing dedicated devices

- ✓ Some devices cannot be shared
- (1) Require processes to perform open()'s on the special files for devices directly
 - The process retries if open() fails
- (2) Have special mechanisms for requesting and releasing dedicated devices
 - An attempt to acquire a device that is not available blocks the caller

Device-independent block size

- ✓ Treat several sectors as a single logical block
- ✓ The higher layers only deal with abstract devices that all use the same block size



User-Space I/O Software

Provided as a library

- ✓ Standard I/O library in C
 - fopen() vs. open()

Spooling

- ✓ A way of dealing with dedicated I/O devices in a multiprogramming system.
- ✓ Implemented by a daemon and a spooling directory
- ✓ Printers, network file transfers, USENET news, mails, etc.



I/O Systems Layers

