

27-1 Escape problem

An $n \times n$ **grid** is an undirected graph consisting of n rows and n columns of vertices, as shown in Figure 27.12. We denote the vertex in the i th row and the j th column by (i, j) . All vertices in a grid have exactly four neighbors, except for the boundary vertices, which are the points (i, j) for which $i = 1$, $i = n$, $j = 1$, or $j = n$.

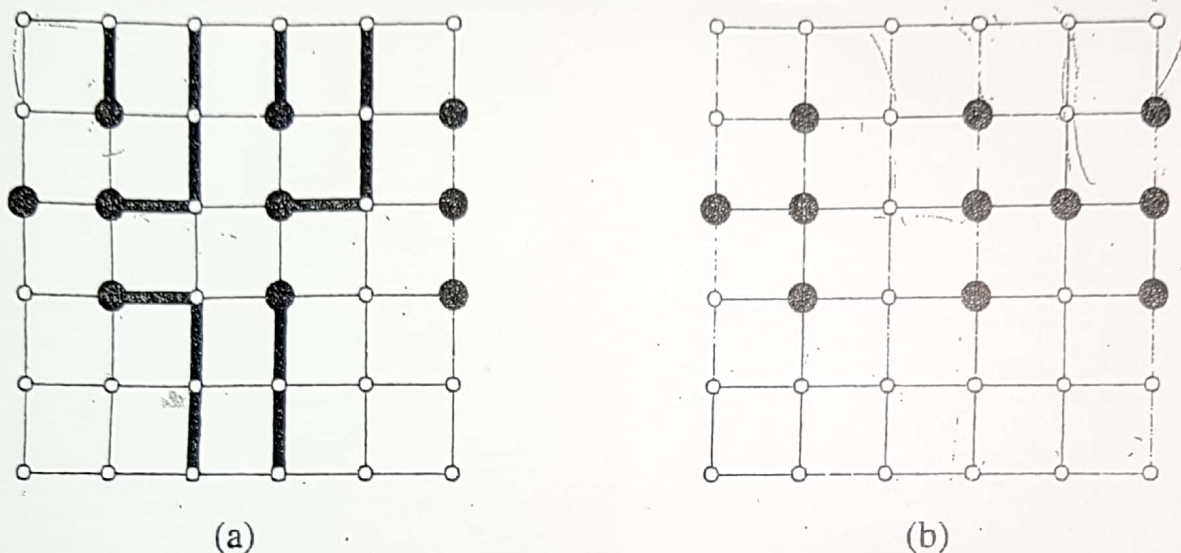


Figure 27.12 Grids for the escape problem. Starting points are black, and other grid vertices are white. (a) A grid with an escape, shown by shaded paths. (b) A grid with no escape.

Given $m \leq n^2$ starting points $(x_1, y_1), (x_2, y_2), \dots, (x_m, y_m)$ in the grid, the **escape problem** is to determine whether or not there are m vertex-disjoint paths from the starting points to any m different points on the boundary. For example, the grid in Figure 27.12(a) has an escape, but the grid in Figure 27.12(b) does not.

- Consider a flow network in which vertices, as well as edges, have capacities. That is, the positive net flow entering any given vertex is subject to a capacity constraint. Show that determining the maximum flow in a network with edge and vertex capacities can be reduced to an ordinary maximum-flow problem on a flow network of comparable size.
- Describe an efficient algorithm to solve the escape problem, and analyze its running time.