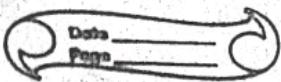


Push-down Automata



$Q:$

$q_0:$

$S:$

$\delta:$

$F:$

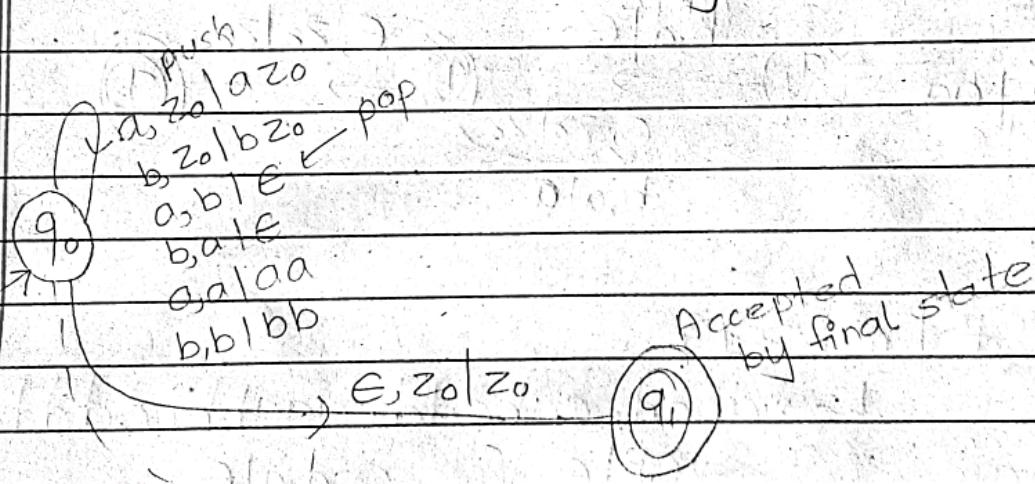
$\Gamma:$ set of state stack alphabets

$z_0:$ start stack symbol

$\delta: Q \times \Sigma \times \Gamma^* \rightarrow Q \times \Gamma^*$

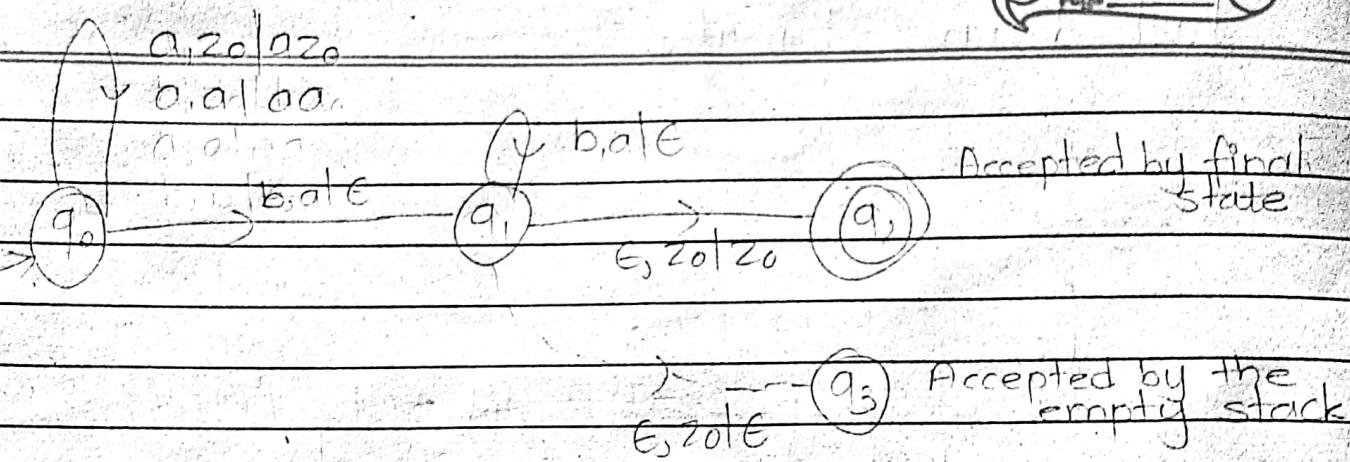
$\delta: Q \times (\Sigma \cup \{e\}) \times \Gamma^* \rightarrow Q \times \Gamma^*$

* Construct a PDA containing equal no. of (a/b) ^{set of all string}



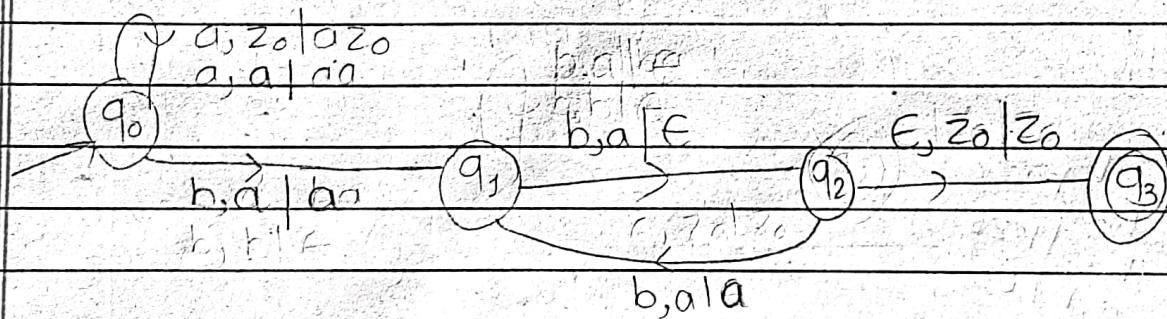
$$L = \{a^n b^n \mid n \geq 1\}$$

$$L = \{ab, aabb, aaabbh, aaaabbhh, \dots\}$$



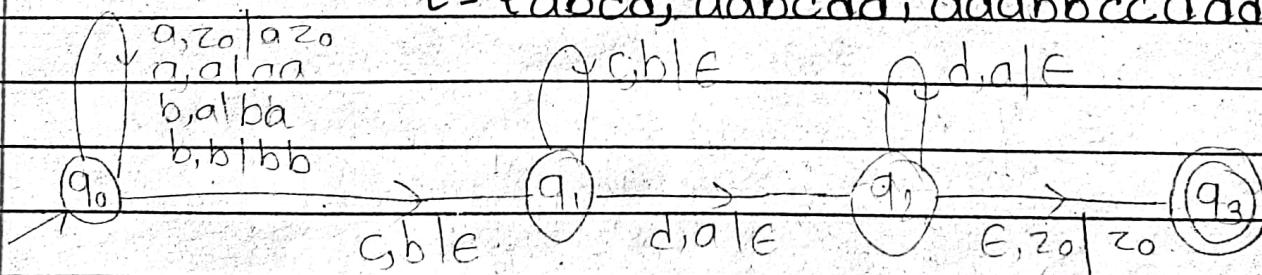
* $L = \{a^n b^{2n} \mid n \geq 1\}$

$L = \{abb, aabb, aabb, \dots\}$



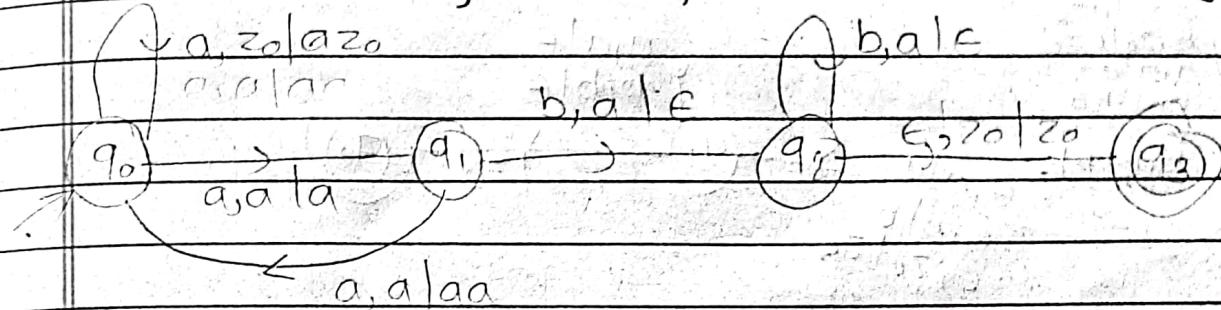
* $L = \{a^m b^n c^m d^n \mid m, n \geq 1\}$

$L = \{abcd, aabcdd, aaabbcccd, \dots\}$



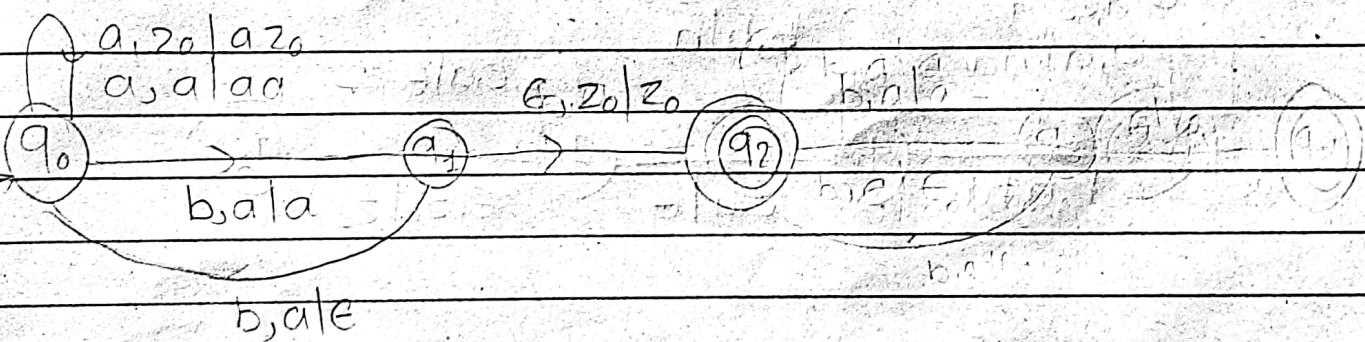
④ $L = \{a^{2n}b^n \mid n \geq 1\}$

$L = \{aab, aaaabb, aaaaaabb, \dots\}$



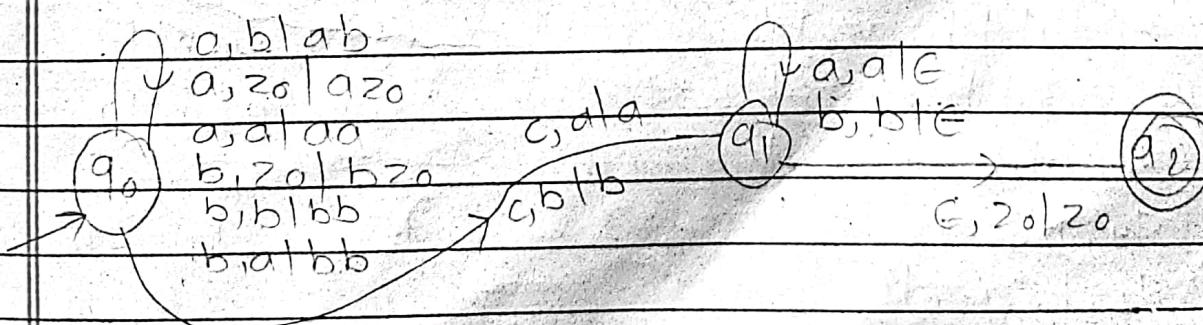
④ $L = \{a^n b^{2n+1} \mid n \geq 1\}$

$L = \{abbb, aabbhh, aaahhhhhh\}$



④ $L = \{w c w^R \mid w \in (a,b)^*\}$

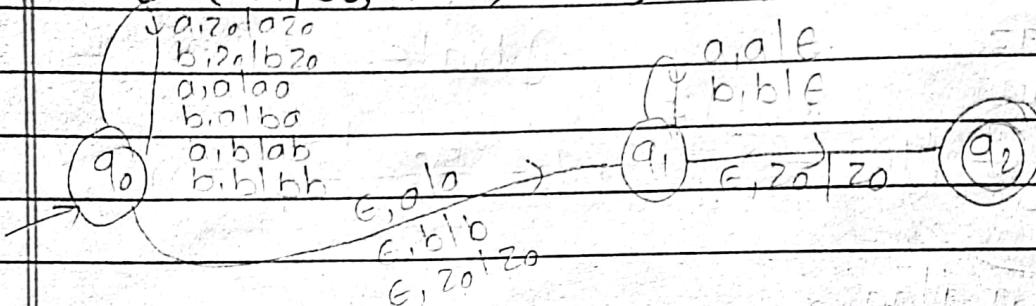
$L = \{aca, bcb, aacaa, bbccb, bbacabb, \dots\}$



Instantaneous description of PDA

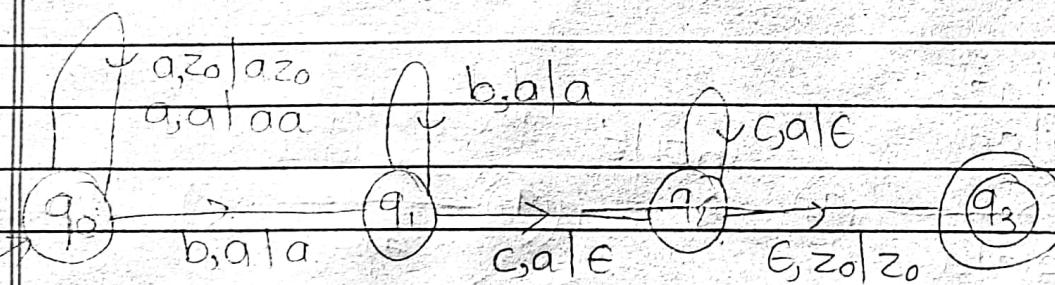
① $L = \{ww^R \mid w \in \{a,b\}^*\}$

$L = \{aa, bb, aaaa, baab, bbbb, \dots\}$



② $L = \{a^n b^m c^n \mid m, n \geq 1\}$

$L = \{abc, aabbcc, abbc, aabbcc \dots\}$



Instantaneous description of PDA

$$(S, \Sigma, \Gamma)$$

$$\delta(q_0, aaahhh, z_0)$$

$$\vdash (q_0, aahhh, a z_0)$$

$$\vdash (q_0, ahbb, aa z_0)$$

$$\vdash (q_0, bbb, aaa z_0)$$

$$\vdash (q_1, bb, aa z_0)$$

$$\vdash (q_1, b, a z_0)$$

$$\vdash (q_1, \epsilon, z_0)$$

$$\vdash (q_2, z_0) \text{ or } (q_2, \epsilon)$$

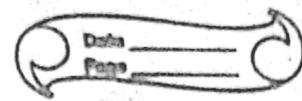
accepted by state

$$\boxed{\delta(q_0, w, r) \vdash^* (q_f, \Gamma)}$$

$$\delta(q_0, w, r) \vdash^* (P, \epsilon)$$

accepted by empty state

string



Equivalence of CFG and PDA

$\text{CFG} \equiv \text{PDA}$

$$G = (V, T, P, S)$$

$$\stackrel{\cong}{=} \text{PDA} = (\{q_0\}, T, V \cup T, S, q_0, \delta, \phi)$$

δ can be defined as

$$\textcircled{1} \quad A \rightarrow \alpha$$

$$\delta(q_0, \epsilon, A) \ni (q_0, \alpha)$$

$$\textcircled{2} \quad a \in T$$

$$\delta(q_0, a, a) = (q_0, \epsilon)$$

example: $E \rightarrow E + T \mid T$

$$T \rightarrow T^* F \mid F$$

$$F \rightarrow (E) \mid a$$

$$\text{PDA} = (\{q_0\}, q_0, \{+, *\}, \{(\), \), \}, a), \{E, T, F, +, *, (\), \), a\}, E, \phi)$$

set of initial state initial input alphabet stack alphabet initial stack symbol
 final state

δ can be defined as:

$$\delta(q_0, \epsilon, E) \ni \{(q_0, E + T), (q_0, T)\}$$

$$\delta(q_0, \epsilon, T) \ni \{(q_0, T^* F), (q_0, F)\}$$

$$\delta(q_0, \epsilon, F) \ni \{(q_0, (E)), (q_0, a)\}$$

Instantaneous description of given string:

$$\delta(q_0, +, +) \ni (q_0, \epsilon)$$

$$\delta(q_0, *, *) \ni (q_0, \epsilon)$$

$$\delta(q_0, (,)) \ni (q_0, \epsilon)$$

$$\delta(q_0, ., .) \ni (q_0, \epsilon)$$

$$\delta(q_0, a, a) \ni (q_0, \epsilon)$$

$$\delta(q_0, a + a * a, E) \ni \dots$$

$$w = a + a * a$$

$$\vdash (q_0, a + a * a, E + T) \ni \dots$$

$$\vdash (q_0, a + a * a, T + T) \ni \dots$$

$$\vdash (q_0, a + a * a, F + T) \ni \dots$$

$$\vdash (q_0, a + a * a, a + T) \ni \dots$$

$$\vdash (q_0, a + a * a, + T) \ni \dots$$

$$\vdash (q_0, a * a, T) \ni \dots$$

$$\vdash (q_0, a * a, T^* F) \ni \dots$$

$$\vdash (q_0, a * a, F * F) \ni \dots$$

$$\vdash (q_0, a, F) \ni \dots$$

$$\vdash (q_0, a, a) \ni \dots$$

$$\vdash (q_0, \epsilon, \epsilon) \ni \dots$$

Hence, accepted.

(*) $S \rightarrow aSa \mid bSb \mid \epsilon$

$$\text{PDA} = (\{q_0\}, q_0, \{a, b\}, \{S, a, b\}, S, \emptyset)$$

$$\delta(q_0, \epsilon, S) \Rightarrow (q_0, aSa)$$

$$\delta(q_0, \epsilon, S) \Rightarrow (q_0, bSb)$$

$$\delta(q_0, \epsilon, S) \Rightarrow (q_0, \epsilon)$$

$$\delta(q_0, a, a) \Rightarrow (q_0, \epsilon)$$

$$\delta(q_0, b, b) \Rightarrow (q_0, \epsilon)$$

$$w = abba$$

$$\vdash \delta(q_0, abba, S)$$

$$\vdash \delta(q_0, abba, aSa)$$

$$\vdash \delta(q_0, bba, Sa)$$

$$\vdash \delta(q_0, bba, bSa)$$

$$\vdash \delta(q_0, ba, bSa)$$

$$\vdash \delta(q_0, ba, \epsilon ba) \rightarrow$$

$$\vdash \delta(q_0, ba, ba)$$

$$\vdash \delta(q_0, a, a)$$

$$\vdash \delta(q_0, \epsilon, \epsilon)$$

(*) PDA \Rightarrow CFG

The set of variable in the grammar:

1) S is a start variable

2) All the symbols of the form $[p, x, q]$,

$p, q \in Q$ and $x \in \Gamma$

i.e. $V = \{S\} \cup \{[p, x, q]\}$

Terminal $T = \Sigma$

The production of G_1 as follows:

for all states $q \in Q$, $S \rightarrow [q_0, z_0, q]$ $i = a$ production of G

for any states $q, r \in Q$, $X \in \Gamma$ and $a \in \Sigma \cup \{\epsilon\}$

if $\delta(q, a, X) = (r, \epsilon)$ then $[q, X, p] \xrightarrow{qa} [r, \epsilon, p]$

for any states $q, r \in Q$, $X \in \Gamma$ and $a \in \Sigma \cup \{\epsilon\}$

if $\delta(q, a, X) \ni (r, Y_1, Y_2, \dots, Y_k)$ where

$Y_1, Y_2, \dots, Y_k \in \Gamma$ and $k \geq 0$

$[q, a, q_k] \Rightarrow [r, Y_1, q_1] [q_1, Y_2, q_2] \dots [q_{k-1}, Y_k, q_k]$

$$L = \{a^n b^n \mid n \geq 1\}$$

$$1) \delta(q_0, a, z_0) \ni (q_0, a z_0)$$

$$2) \delta(q_0, a, a) \ni (q_0, aa)$$

$$3) \delta(q_0, b, a) \ni (q_1, \epsilon)$$

$$4) \delta(q_1, b, a) \ni (q_1, \epsilon)$$

$$5) \delta(q_1, \epsilon, z_0) \ni (q_1, \epsilon)$$

$$S \rightarrow [q_0, z_0, q_0] \mid [q_0, z_0, q_1]$$

$$[q_0, z_0, q_0] \xrightarrow{a} [q_0, a, q_0] \quad [q_0, z_0, q_0]$$

$$[q_0, z_0, q_0] \xrightarrow{a} [q_1, a, q_1] \quad [q_1, z_0, q_0]$$

$$[q_0, z_0, q_1] \xrightarrow{a} [q_0, a, q_0] \quad [q_0, z_0, q_1]$$

$$[q_0, z_0, q_1] \xrightarrow{a} [q_0, a, q_1] \quad [q_1, z_0, q_1]$$

$$[q_0, a, q_0] \xrightarrow{a} [q_0, a, q_0] \quad [q_0, a, q_0]$$

$$[q_0, a, q_0] \xrightarrow{a} [q_0, a, q_1] \quad [q_1, a, q_0]$$

$$[q_0, a, q_1] \xrightarrow{a} [q_0, a, q_0] \quad [q_0, a, q_1]$$

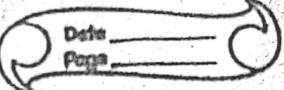
$$[q_0, a, q_1] \xrightarrow{a} [q_0, a, q_1] \quad [q_1, a, q_1]$$

$$[q_0, a, q_1] \xrightarrow{b}$$

$$[q_1, a, q_1] \xrightarrow{b}$$

$$[q_1, z_0, q_1] \xrightarrow{c}$$

$$l = 2a^n b^n \quad (n \geq 1)$$



$$\textcircled{5} \quad \delta(q_1, e, z_0) \Rightarrow (q_2, z_0)$$

तर्फः

$$S \rightarrow [q_0, z_0, q_0] \quad [q_0, z_0, q_1] \quad [q_0, z_0, q_2]$$

$$[q_0, z_0, q_0] \rightarrow a \quad [q_0, a, q_0] \quad [q_0, z_0, q_0]$$

$$[q_0, z_0, q_0] \rightarrow a \quad [q_0, a, q_1] \quad [q_1, z_0, q_0]$$

$$[q_0, z_0, q_0] \rightarrow a \quad [q_0, a, q_2] \quad [q_2, z_0, q_0]$$

$$[q_0, z_0, q_1] \rightarrow a \quad [q_0, a, q_0] \quad [q_0, z_0, q_1]$$

$$[q_0, z_0, q_1] \rightarrow a \quad [q_0, a, q_1] \quad [q_1, z_0, q_1]$$

$$[q_0, z_0, q_1] \rightarrow a \quad [q_0, a, q_2] \quad [q_2, z_0, q_1]$$

$$[q_0, z_0, q_2] \rightarrow a \quad [q_0, a, q_0] \quad [q_0, z_0, q_2]$$

$$[q_0, z_0, q_2] \rightarrow a \quad [q_0, a, q_1] \quad [q_1, z_0, q_2]$$

$$[q_0, z_0, q_2] \rightarrow a \quad [q_0, a, q_2] \quad [q_2, z_0, q_2]$$

$$[q_0, a, q_0] \rightarrow a \quad [q_0, a, q_0] \quad [q_0, a, q_0]$$

$$[q_0, a, q_0] \rightarrow a \quad [q_0, a, q_1] \quad [q_1, a, q_0]$$

$$[q_0, a, q_0] \rightarrow a \quad [q_0, a, q_2] \quad [q_2, a, q_0]$$

$$[q_0, a, q_1] \rightarrow a \quad [q_0, a, q_0] \quad [q_0, a, q_1]$$

$$[q_0, a, q_1] \rightarrow a \quad [q_0, a, q_1] \quad [q_1, a, q_1]$$

$$[q_0, a, q_1] \rightarrow a \quad [q_0, a, q_2] \quad [q_2, a, q_1]$$

$$[q_0, a, q_2] \rightarrow a \quad [q_0, a, q_0] \quad [q_0, a, q_2]$$

$$[q_0, a, q_2] \rightarrow a \quad [q_0, a, q_1] \quad [q_1, a, q_2]$$

$$[q_0, a, q_2] \rightarrow a \quad [q_0, a, q_2] \quad [q_2, a, q_2]$$

$$[q_0, a, q_1] \rightarrow b$$

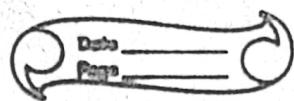
$$[q_1, a, q_1] \rightarrow b$$

$$[q_1, z_0, q_0] \rightarrow e \quad [q_2, z_0, q_0] \quad [q_2, z_0, q_1]$$

$$[q_1, z_0, q_1] \rightarrow f \quad [q_2, z_0, q_1] \quad [q_2, z_0, q_2]$$

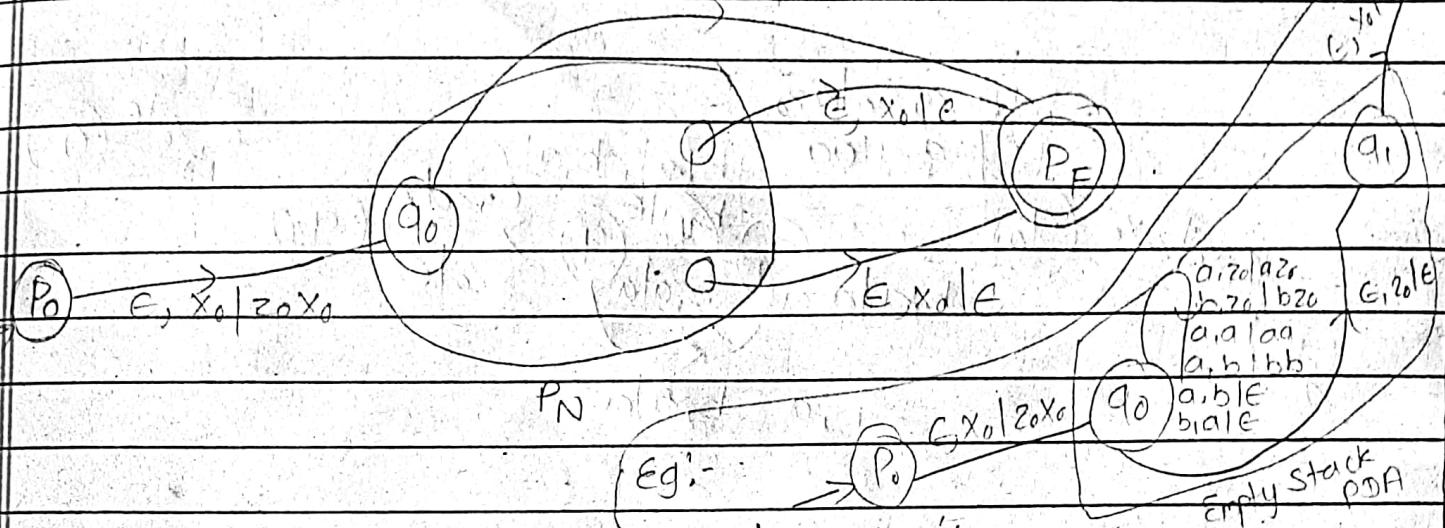
$$[q_2, z_0, q_2] \rightarrow e \quad [q_2, q_2, q_2] \quad [q_2, z_0, q_2]$$

Conversion



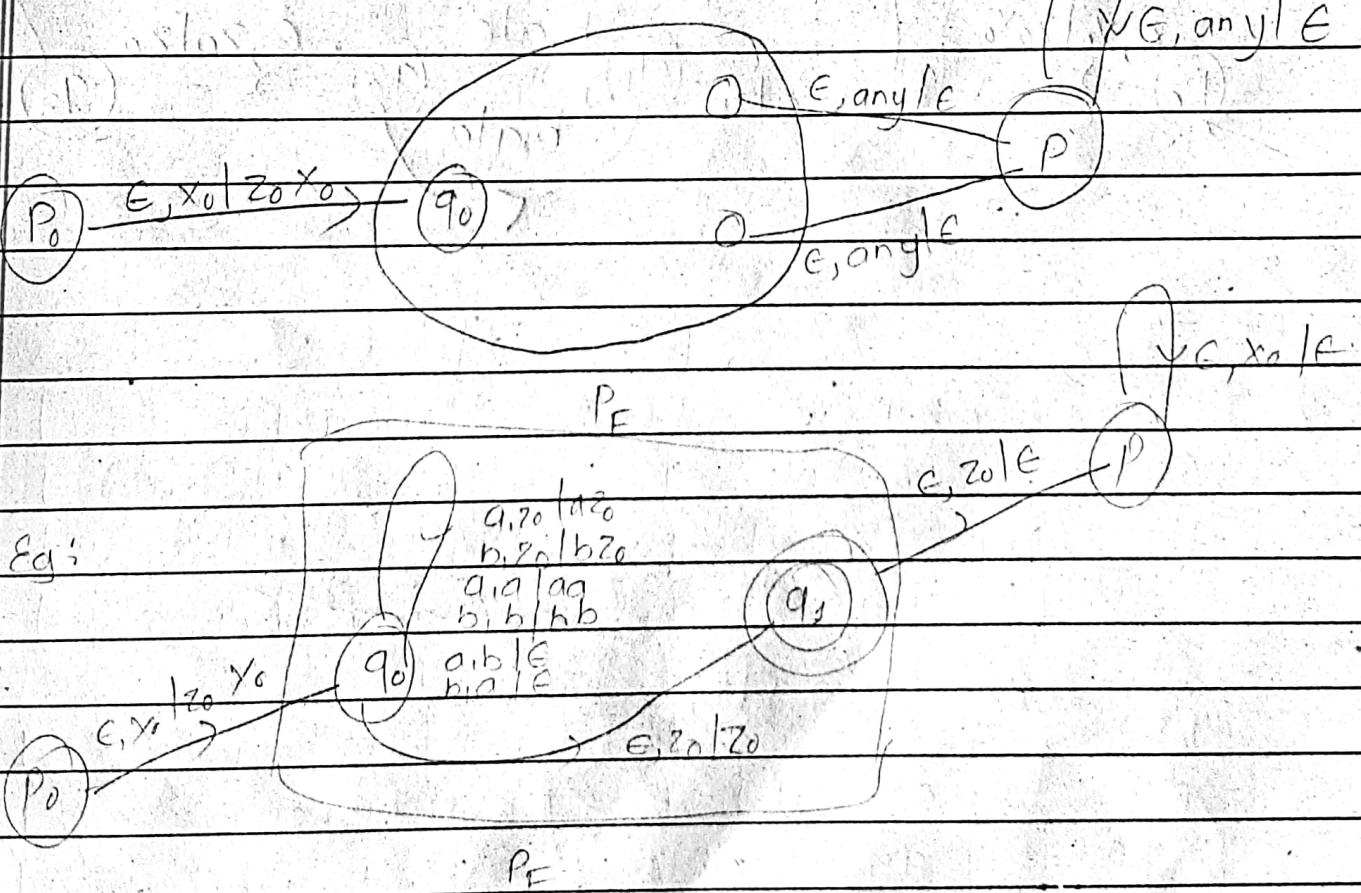
From empty stack PDA to final state PDA

$\epsilon, x_0 | F$

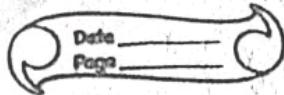


From final state PDA to empty stack PDA

$\epsilon, \text{any} | E$



a b b



$$L = \{a^n b^{2n} \mid n \geq 1\}$$

