

# Type-Level Property Based Testing

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- These can be modelled using dependent types
  - Dependent types are difficult to get right
- How do we increase confidence in our dependent types?

This is not a proof technique

But hopefully, it helps us catch errors faster and provides guarantees that our model behaves as intended

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- Embedded controllers for automatic doors, ATMs, network protocols, etc...



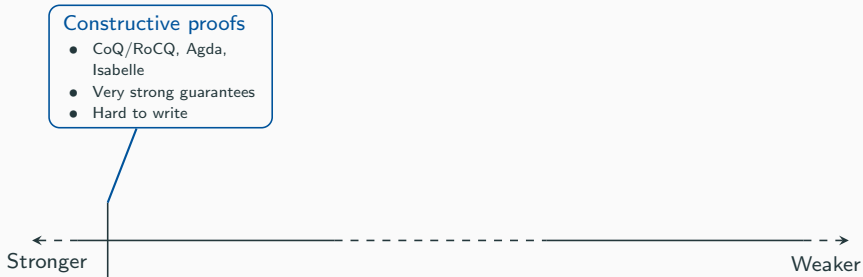
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- These are all stateful
- And we would very much like them to be correct

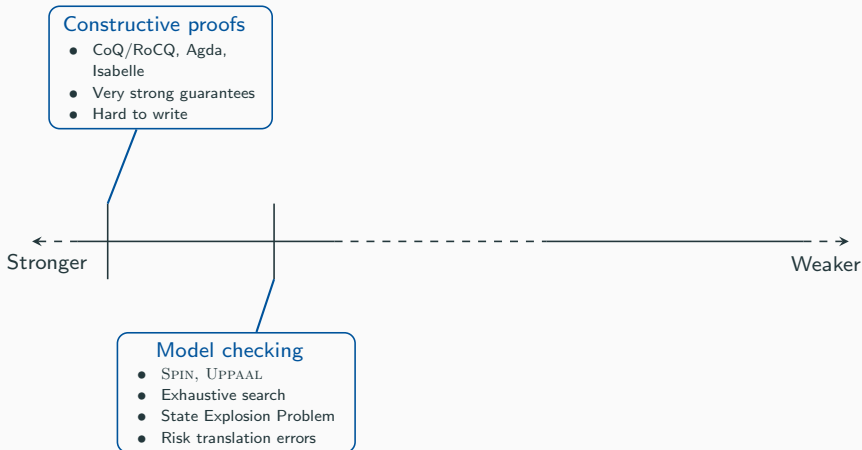
# Spectrum of Verification



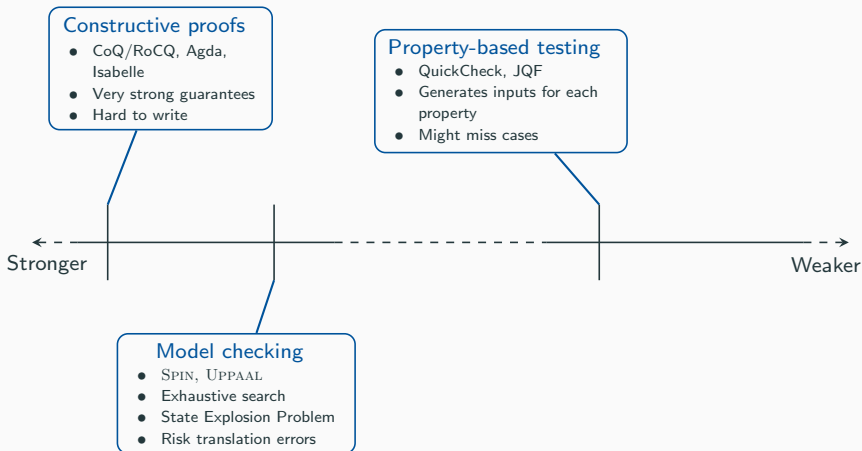
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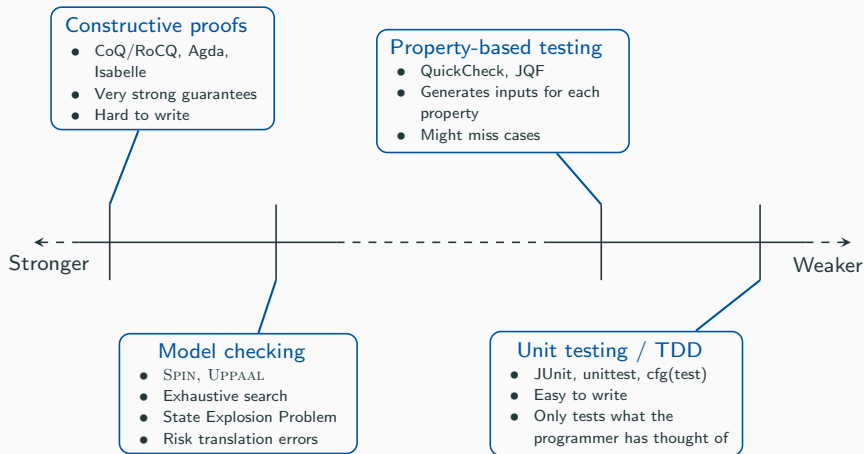
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like Agda and Idris



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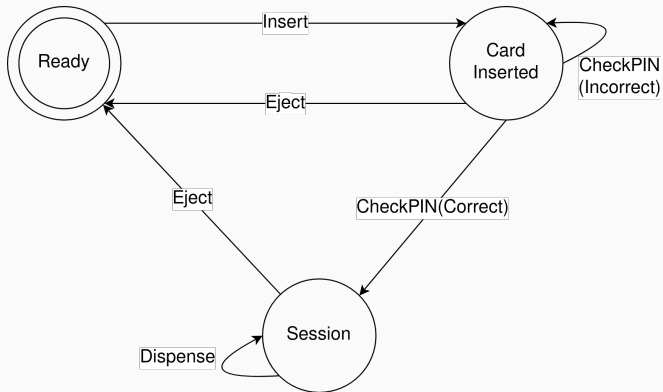
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```
head :: [a] -> a
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```

```
head : Vect (S k) a
      -> a
-- always safe because
-- the length must be
-- at least 1
```

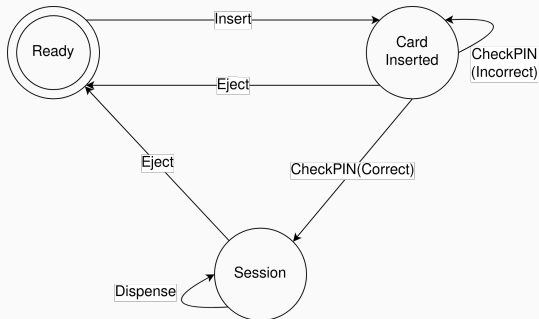
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# The ATM state machine



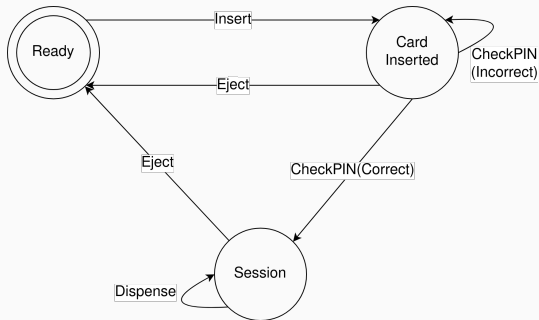
# Datatype for the ATM states

```
data ATMState
  = Ready
  | CardInserted
  | Session
```



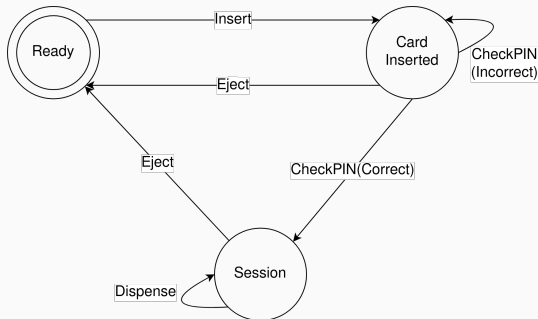
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Everything which does not have a result returns Unit — ()



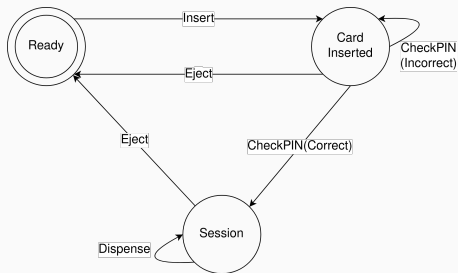
# State Transition Function

ChkPINfn : PINok -> ATMState

ChkPINfn Correct = Session

ChkPINfn Incorrect =

→ CardInserted



# Dependent State Transition

```
data ATM : (t : Type) -> ATMState -> (t -> ATMState)
    -> Type where
  CheckPIN : (pin : Int)
    -> ATM PINok CardInserted ChkPINfn
    :
  (>>=) : ATM a s1 s2f
    -> ((x : a) -> ATM b (s2f x) s3f)
    -> ATM b s1 s3f
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```

# ATM Indexed State Monad

```
data ATM : (t : Type) -> ATMState -> (t -> ATMState) ->
  ↪ Type where
  CheckPIN : (pin : Int)
             -> ATM PINok CardInserted ChkPINfn
  Insert   : ATM () Ready (const CardInserted)
  Dispense : (amt : Nat) -> ATM () Session (const Session)
  Eject    : ATM () st (const Ready)
  Pure     : (x : t) -> ATM t (stFn x) stFn
  (>>=)    : ATM a s1 s2f -> ((x : a) -> ATM b (s2f x) s3f)
  ↪ -> ATM b s1 s3f
```

## Why Is This Neat?

- We declare our intended start and end state in the type  
`prog : ATM () Ready (const Ready)`

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- We declare our intended start and end state in the type  
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- And the type-checker verifies that we don't use operations incorrectly

```
prog = do                                -- We start in Ready
  Insert  ----- Ready to CardInserted
  Correct <- CheckPIN 1234  --- CI to Session
    | Incorrect => <...>  ----- (or stay in CI)
  Dispense 42  ----- Stay in Session
  Eject  ----- Return to Ready
```

# Dependent Types Only Get Some Things Right

Rejected by the type-checker:

```
badProg : ATM ()
          Ready (const Ready)

badProg = do
  Insert
  let pin = 1234
  Correct <- CheckPIN pin
    | Incorrect => InsertCard
  Dispense 42
  -- We never Eject, so we
  -- never come back to
  -- `Ready'
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    | Incorrect => InsertCard
  Dispense 42
  -- We never Eject, so we
  -- never come back to
  -- `Ready'
```

Accepted by the type-checker:

```
loopProg : ATM ()
          Ready (const Ready)

loopProg = do
  InsertCard
  let pin = 4321
  loopIncorrect pin
where
  loopIncorrect : Nat -> ATM ()
                CardInserted
                (const Ready)

  loopIncorrect p = do
    Incorrect <- CheckPIN p
    | Correct => -- <...>
    loopIncorrect p
```



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- We express boolean properties, which are then tested over random inputs (typically 100)

```
reverse (reverse xs) == xs
```

## What About Dependent Types?

QuickCheck works very well for regular types, what about dependent types?

# Generating Dependent Types is Tricky

Consider generating arbitrary vectors:

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Arbitrary t => Arbitrary (Vect n t) where
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How do we solve this?

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- Specifically: dependent pairs

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record DPair a (p : a -> Type) where
  constructor MkDPair
  fst : a
  snd : p fst
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- Specifically: dependent pairs

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record DPair a (p : a -> Type) where
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```

- As long as we know how to generate an ``Arbitrary a``, we can generate an ``Arbitrary (x : a ** p x)``

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Arbitrary (n : Nat ** Vect n a) where
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    arbVect <- nArbitrary length_
    pure (_ ** arbVect)
```

## Arbitrary ATMs?

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With a bit of work, yes!

- Store the operation and its result instance — for example whether `PINok` was successful

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(<resT> ** <stFn> ** MkOpRes <op> <res : resT> ...)
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TraceStep (OpRes <stT> ...) <resSt : stT>
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TraceStep (OpRes <stT> ...) <resSt : stT>
```

- A chain of these make up a *trace*

## ATM: from CardInserted

This is still QuickCheck, we can control the frequency of generated instances:

```
options CardInserted = do
  -- we need a PIN, even though we control the result
  let arbPIN = 0
  let op1 = ( _ ** _ ** MkOpRes (CheckPIN arbPIN) Correct)
  let op2 = ( _ ** _ ** MkOpRes (CheckPIN arbPIN) Incorrect)
  let op3 = ( _ ** _ ** MkOpRes Eject ())

  frequency $ [(1, pure op1), (4, pure op2), (1, pure op3)]
```



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- In Idris2, the `So` type is inhabited iff its argument evaluates to `True`
- In other words, we can run property based testing as part of the type checking process!  
`So (quickCheck <property>)`
- Idris2 is built on Quantitative Type Theory, which has erasure, meaning the tests are removed from the compiled program

## QuickCheck Spots the Error!

```
0 PROP_eventuallyReady : Fn (ATMTrace Ready 10) Bool
PROP_eventuallyReady = MkFn
  (\case (MkATMTrace _ trace) => elem Ready (map (.resSt) trace))
```

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```
-- Error: While processing right hand side of
--      EventuallyReady_OK. When unifying:
--      So True
-- and:
--      So (QuickCheck PROP_eventuallyReady)
-- Mismatch between: True and False.
```

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```
MkQCRes (Just False) <log> ""  
Falsifiable, after 4 tests:  
Starting @ Ready:  
[ (<ATM 'Insert ~ ()'>, CardInserted)  
  , (<ATM 'CheckPIN 0 ~ Incorrect'>, CardInserted)  
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  , (<ATM 'CheckPIN 0 ~ Incorrect'>, CardInserted)  
  , (<ATM 'CheckPIN 0 ~ Incorrect'>, CardInserted)  
]\n""
```



- Now that we know there's an error, we can fix things!

```
ChkPINfn : (retries : Nat) -> PINok -> ATMState
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```
ChkPINfn 0      Correct  = Session
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# Fixing Things

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- Carrying this through to the generators, our property passes: the file reloads successfully and the REPL reports

```
> quickCheck PROP_eventuallyReady
```

```
MkQCRes (Just True) <log> "OK, passed 100 tests"
```

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- With most verification tools, we have to translate between representations
  - Spec, model, and implementation are independent
- This results in the risk of translation mistakes
  - The verification tool might not support the same types as the implementation language
  - Might think we're verifying the same thing, when in actual fact the semantics have changed between representations

In our case, the specification *is* the model; *everywhere*

```
trace : (steps : Nat) -> (st : stT) -> Gen (Vect steps
  ↪ (TraceStep opT))
trace 0 _ = pure []
trace (S j) st = do
  (_ ** stFn ** opR@(MkOpRes op res)) <- arbitrary
  let nextSt = stFn res
  pure $ (MkTS opR nextSt) :: !(trace j nextSt)
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And this works for anything expressed in terms of states and operations with results — ISMs generalise

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- We have tested the dependent types which help guide us when writing the implementation, not the implementation itself since it is being kept in check by the types
- Testing gives us confidence that we are not being misled
- Dependent types are type-level programs, let's test them!

# Thank You

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Paper



[doi:10.1145/3678000.3678206](https://doi.org/10.1145/3678000.3678206)

- Running tests at the type level puts a lot of strain on the compiler, so there may be interesting optimisations to explore there
- Can we do more? ARQ with Sliding Window? Protocols with crash-stop failures?
- What kinds of properties can we test? There are parallels to Model Checking, so how does this compare to LTL or TLA<sup>+</sup>?

# Generic ISM Datatype

```
op : forall st . (t' : Type) -> st -> (t' -> st) -> Type

data Prog : {0 stT : _} -> (opT : (t' : _) -> stT -> (t'
  ↪ -> stT) -> Type) -> (t : Type) -> (from : stT) -> (to
  ↪ : t -> stT) -> Type where
  Pure : (x : t) -> Prog opT t (stFn x) stFn
  Op : {0 opT : (t' : _) -> stT -> (t' -> stT) -> Type} ->
    ↪ opT t st stFn -> Prog opT t st stFn
  (>=) : Prog opT resT1 st1 stFn1 -> ((x : resT1) -> Prog
    ↪ opT resT2 (stFn1 x) stFn2) -> Prog opT resT2 st1
    ↪ stFn2
```

## Operation-Result Pairs

```
record OpRes {0 stT : _} (opT : (t' : _) -> stT -> (t' ->
  ↳ stT) -> Type) (resT : Type) (currSt : stT) (0 nsFn :
  ↳ resT -> stT) where
  constructor MkOpRes
  op : opT resT currSt nsFn
  res : resT
  {auto opShow : Show (opT resT currSt nsFn)}
  {auto rShow : Show resT}
```



# TraceStep

```
record TraceStep (opT : (t' : _) -> stT -> (t' -> stT) ->
  ↳ Type) where
  constructor MkTS
  {0 stepRT : _}
  {0 stepSt : stT}
  {0 stepFn : stepRT -> stT}

opRes : OpRes opT stepRT stepSt stepFn
resSt : stT

{auto showStT : Show stT}
```

```
data Trace : (opT : (t' : _) -> stT -> (t' -> stT) ->
  ↳ Type) -> stT -> Nat -> Type where
  MkTrace : Show stT => (initSt : stT) -> {bound : Nat}
    -> (trace : Vect bound (TraceStep opT))
    -> Trace opT initSt bound
```