

# Introduction to Quantum Computing

## A Comprehensive Guide for Computer Scientists

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# Outline

**Mermin-Peres magic square**

**Quantum principle**

**Quantum computation, AI and more**

# Outline

**Mermin-Peres magic square**



Quantum principle

Quantum computation, AI and more

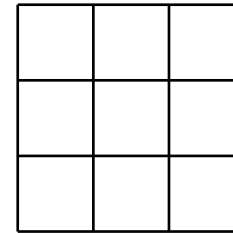
# The magic square game

This is a **collaborative** game where two players work together to pass a test.

## Players

-  Rowena
-  Colin

## Board



Fill 1 or  $-1$

- Game starts, the teacher tells Rowena and Colin a row/column number  $i/j$ .
  - Rowena gets  $i (= 1, 2, 3)$
  - Colin gets  $j$
- Rowena and Colin fill the  $i$ -th row/ $j$ -th column with numbers (-1 or 1).
- The teacher checks if the row/column product is 1/-1 and the numbers at the intersection are the same. If yes, Rowena and Colin pass the exam; otherwise, they fail.

-1		
1	-1	-1
1		

win!

-1		
1	-1	1
1		

loss! inconsistency

1		
1	-1	-1
1		

loss! column product is 1

# Winning strategy?

Rowena and Colin can agree on any kind of strategy **prior** to the game, but are not permitted to communicate during the game.

## Question

Is there a 100% winning strategy for Rowena and Colin?

(5 min free discussion)

# The optimal strategy for those who haven't studied quantum

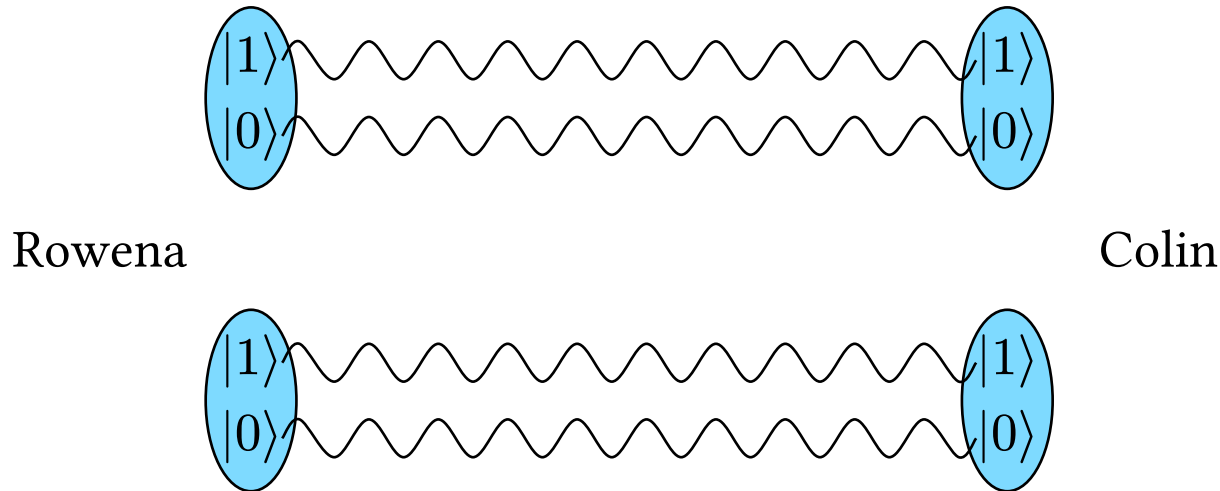
Sharing a predetermined table.

1	1	1
-1	1	-1
1	-1	?

success rate =  $8/9 \approx 0.889$

# Quantum solution to the Magic square game

**superposition** - both 0 and 1



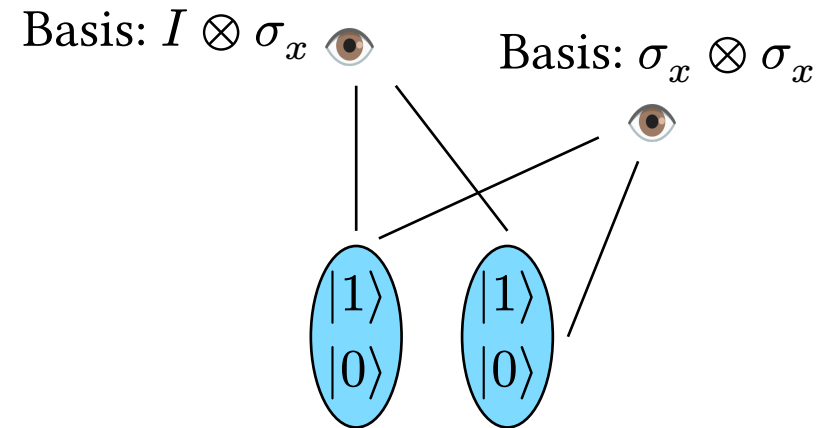
- Step 1: Share two **entangled** pairs of **qubits** (or **Bell pairs**) (Mermin, 1990). Each qubit can be in any **superposition** of  $|0\rangle$  and  $|1\rangle$  (this is the **Dirac notation**).
- Step 2: **Measure** the qubits, and send the outcome to the teacher.



# Quantum solution to the Magic square game

Measure on different **bases**!

$$\begin{pmatrix} I \otimes \sigma_z & \sigma_z \otimes I & \sigma_z \otimes \sigma_z \\ \sigma_x \otimes I & I \otimes \sigma_x & \sigma_x \otimes \sigma_x \\ -\sigma_x \otimes \sigma_z & -\sigma_z \otimes \sigma_x & \sigma_y \otimes \sigma_y \end{pmatrix}$$



Finally verified by experiments (Xu et al., 2022).

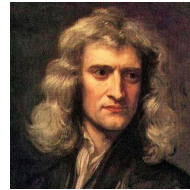
# Outline

Mermin-Peres magic square

**Quantum principle**

Quantum computation, AI and more

# Classical view of the world: probability is due to lack of knowledge



Newton's law:  $f = ma$

- Our world is a deterministic machine, and the future is predictable.

# How to understand probability?



Boltzmann distribution:  $p(E) = e^{-\frac{E}{kT}}$

- Probability is due to a lack of knowledge about underlying **hidden variables**.

# Probabilistic view of state

Probabilistic **state** is a real vector

$$\vec{p} = \begin{pmatrix} p_1 \\ p_2 \\ \vdots \\ p_n \end{pmatrix}, \text{ s.t. } \sum_i p_i = 1$$

where  $n$  is the number of possible outcomes.

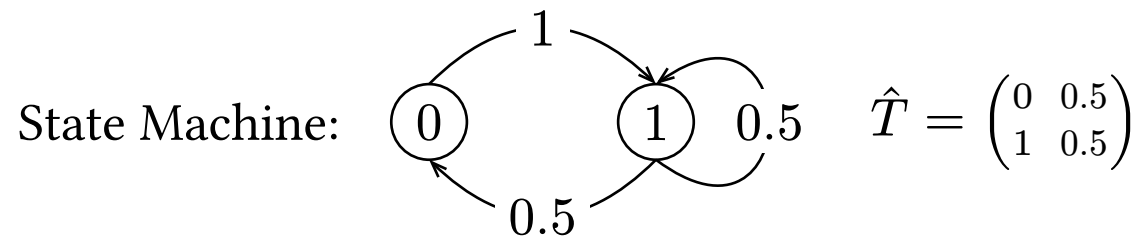
- bit:  $x = 0$  or  $1$
- p-bit (Chowdhury et al., 2023):  $p(x) = \begin{pmatrix} p \\ 1-p \end{pmatrix}$

# Probabilistic view of operations

Probabilistic **operation** is a transition matrix  $\hat{T} = (t_{ij})$

$$\vec{p}' = \hat{T}\vec{p} \text{ s.t. } t_{ij} \geq 0 \text{ and } \sum_i t_{ij} = 1 \text{ by conservation of probability}$$

Example:  $t_{10} = 1, t_{01} = 0.5, t_{11} = 0.5$



# Example: Probabilistic NOT gate

- NOT gate (classical):

$$y = \neg x = \begin{cases} 1 & \text{if } x = 0 \\ 0 & \text{if } x = 1 \end{cases}$$

- NOT gate on a probabilistic bit:

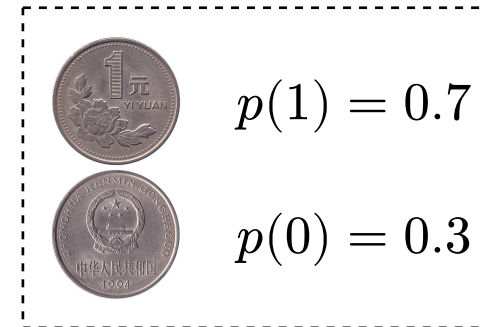
$$y = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} p \\ 1-p \end{pmatrix} = \begin{pmatrix} 1-p \\ p \end{pmatrix}$$

- Probabilistic NOT gate on a bit:

$$y = \begin{pmatrix} 1-p & p \\ p & 1-p \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1-p \\ p \end{pmatrix}$$



↓ (flip with probability 0.3)

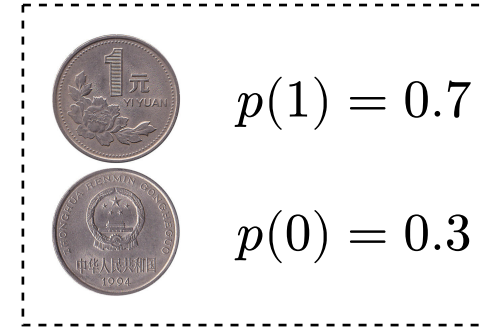


Probabilistic gates are a **superset** of deterministic gates.

# Probabilistic view of measurement

Probabilistic **measurement** is a sampling  
 $x \sim \vec{p}(x)$ , after measurement,  $\vec{p} \rightarrow \delta(x)$

$$\delta(x) = \begin{pmatrix} \vdots \\ 0 \\ 1 \\ 0 \\ \vdots \end{pmatrix}$$



$(p = 0.7) \swarrow$    $\searrow (p = 0.3)$



State = 1



State = 0

Note: The state is changed by measurement!



# Probabilistic state space



Q: How many numbers are needed to represent an  $n$ -bit probability distribution:

$$p(x_1, x_2, \dots, x_n)$$

# Is PTM more powerful than Turing machine?

Answer:  $2^n$

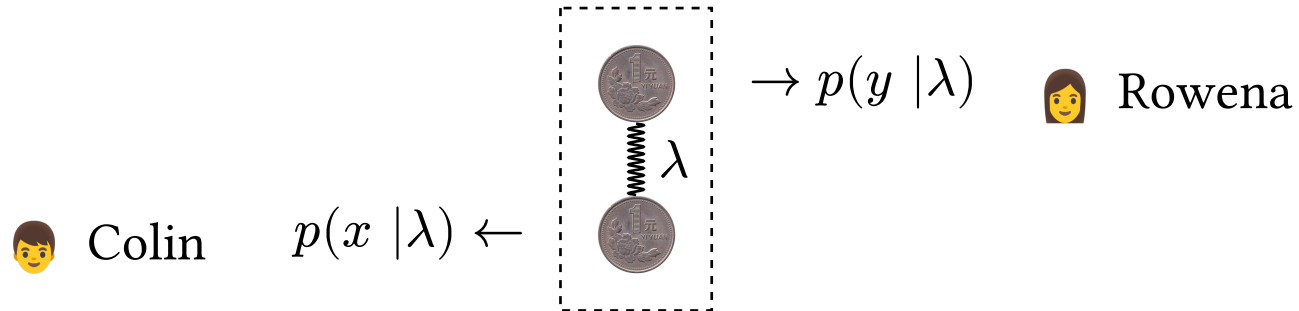
Probabilistic Turing machine (PTM) = Turing machine + Coin  $\stackrel{?}{=}$  Turing machine

PTM is probably not more powerful than a Turing machine, e.g., linear congruential generators can simulate random numbers.

$$X_{n+1} = (aX_n + c) \bmod m$$

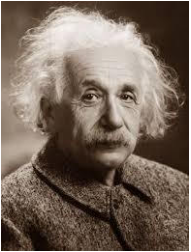
# Hidden variable models

Let  $x$  and  $y$  be the bits that Colin and Rowena share.



Q: If Colin and Rowena share coins with probability distribution  $p(x, y | \lambda)$ , can they have a higher success rate than  $8/9$ ?

# Hidden variable models cannot explain our world



God does not play dice.

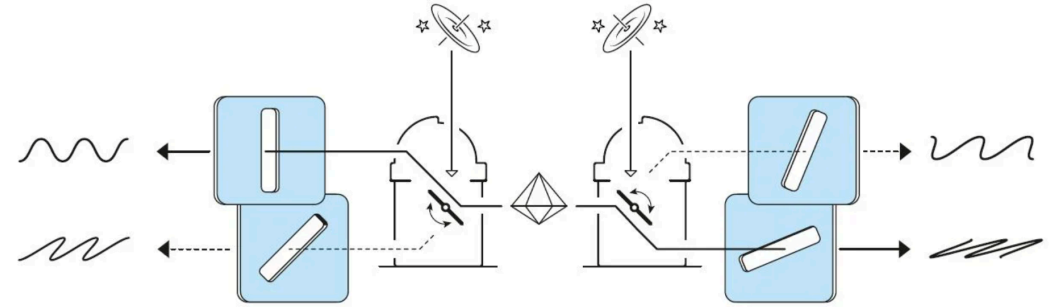


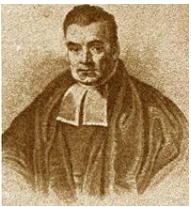
Figure 1: Bell's inequality experiment, Nobel prize in physics, 2022.

Probabilistic model with hidden variable  $\lambda$ :

$$\text{left} \sim p(A) = p(A|\lambda)$$

$$\text{right} \sim p(B) = p(B|\lambda)$$

Even with an infinite number of hidden variables, it is still impossible to explain the outcome!



Classical statistics cannot help. 🙄

# Quantum view of state

A quantum state is the square root of a probability distribution:

$$|\psi\rangle = \sqrt{\vec{p}} \text{ (hand-waving)}$$

Quantum **state** is a vector of complex numbers

$$\text{Ket: } |\psi\rangle = \begin{pmatrix} \psi_1 \\ \psi_2 \\ \vdots \\ \psi_n \end{pmatrix} = \begin{pmatrix} \sqrt{p_1} e^{i\varphi_1} \\ \sqrt{p_2} e^{i\varphi_2} \\ \vdots \\ \sqrt{p_n} e^{i\varphi_n} \end{pmatrix} \text{ s.t. } \langle\psi|\psi\rangle = 1$$

Hermitian conjugate (bra):  $\langle\psi| = (\psi_1^* \ \psi_2^* \ \dots \ \psi_n^*)$

# Normalization of probability

$$\langle \psi | \psi \rangle = (\psi_1^* \ \psi_2^* \ \dots \ \psi_n^*) \begin{pmatrix} \psi_1 \\ \psi_2 \\ \vdots \\ \psi_n \end{pmatrix} = \sum_i \psi_i^* \psi_i = \sum_i p_i = 1$$

For complex numbers,  $\psi_i^* = \sqrt{p_i} e^{-i\varphi_i}$  and  $\psi_i^* \psi_i = p_i$ .

To access probability, we use basis vectors:

$$|i\rangle = \begin{pmatrix} \vdots \\ 0 \\ 1 \\ 0 \\ \vdots \end{pmatrix}, p_i = |\langle i | \psi \rangle|^2$$

# Quantum view of operations

Probabilistic operation:  $\hat{T}\vec{p} = \vec{p}'$ , s.t.  $\sum_i t_{ij} = 1$

Quantum **operation** is a unitary operator.

$\hat{U}|\psi\rangle = |\psi'\rangle$  s.t.  $\hat{U}\hat{U}^\dagger = \hat{I}$  conservation of probability

The normalization condition requires  $(\langle\psi|U^\dagger)(U|\psi\rangle) = 1$  for any state  $|\psi\rangle$ . Hence, a quantum operation must be a **unitary operator**, i.e.  $\hat{U}\hat{U}^\dagger = \hat{I}$ .

# Quantum computation, from classical gates to quantum gates

- Probabilistic NOT gate:

$$y = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} p \\ 1-p \end{pmatrix} = \begin{pmatrix} 1-p \\ p \end{pmatrix}$$


- Quantum NOT (or X) gate:

$$y = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \sqrt{p} \\ \sqrt{1-p}e^{i\varphi} \end{pmatrix} = \begin{pmatrix} \sqrt{1-p}e^{i\varphi} \\ \sqrt{p} \end{pmatrix}$$

$$|0\rangle \text{ --- } [X] \text{ --- } |1\rangle$$

- Hadamard gate:

$$y = \begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} \end{pmatrix}$$

$$|0\rangle \text{ --- } [H] \text{ --- } \text{SCHRODINGER'S CAT IS ALIVE} \text{ --- } \text{Schrödinger's Cat Box}$$




# Unitary matrix vs. transition matrix

- *Remark:* Unitary matrices are reversible.
- *Remark:* Bennett showed (Bennett, 1982) that  $\text{RTM} = \text{TM}$ .
- *Remark:* Landauer's principle showed (Reeb & Wolf, 2014) that erasing information requires thermal dissipation.
- *Remark:* Two qubit quantum gates are universal (Dawson & Nielsen, 2005). Stronger: almost any single arbitrary two qubit gate is universal (Deutsch et al., 1995). Spectral gap theory shows any gate can be approximated exponentially fast (Bourgain & Gamburd, 2011).
- *Remark:*  $\text{QTM} \geq \text{PTM}$ . Unitary circuits cannot be simulated classically in polynomial time

# Quantum Measurement

Probabilistic measurement:

$$x \sim \vec{p}(x) = \delta(x)^T \vec{p}, \text{ after measurement, } \vec{p} \rightarrow \delta(x)$$

Quantum **measurement** follows the Born's rule:

$$x \sim |\langle x | \psi \rangle|^2, \text{ after measurement, } |\psi\rangle \rightarrow |x\rangle$$

## Example: Quantum measurement

$$|\psi\rangle = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{pmatrix}$$

$$\langle 0|\psi\rangle = \frac{1}{\sqrt{2}}, p(0) = |\langle 0|\psi\rangle|^2 = \frac{1}{2}$$

$$\langle 1|\psi\rangle = \frac{1}{\sqrt{2}}, p(1) = |\langle 1|\psi\rangle|^2 = \frac{1}{2}$$

With 50% probability, we get 0, and the state collapses to  $|0\rangle$ ;

With 50% probability, we get 1, and the state collapses to  $|1\rangle$ .

## Example: Quantum measurement

$$|\psi\rangle = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} \end{pmatrix}$$

$$\langle 0|\psi\rangle = \frac{1}{\sqrt{2}}, p(0) = |\langle 0|\psi\rangle|^2 = \frac{1}{2}$$

$$\langle 1|\psi\rangle = -\frac{1}{\sqrt{2}}, p(1) = |\langle 1|\psi\rangle|^2 = \frac{1}{2}$$

With 50% probability, we get 0, and the state collapses to  $|0\rangle$ ;

With 50% probability, we get 1, and the state collapses to  $|1\rangle$ .

# The difference!

Unlike classical probability,  $|x\rangle$  can be any state! (not just the computational basis  $|i\rangle$ )

Basis is a **key** to access the information in the quantum state. Any set of **normalized, orthogonal states** can be a basis.

	$z$ -basis	$x$ -basis
Outcome: $\lambda = +1$	$ 0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$	$ x; +\rangle = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{pmatrix}$
Outcome: $\lambda = -1$	$ 1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$	$ x; -\rangle = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} \end{pmatrix}$
Orthogonality	$\langle 0 1\rangle = 0$	$\langle x; + x; -\rangle = 0$
Normalization	$\langle 0 0\rangle = 1$	$\langle x; + x; +\rangle = 1$

# Phase matters

$$\text{State: } |\psi\rangle = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}}e^{i\varphi} \end{pmatrix}$$

$$\begin{aligned} \text{Expectation value: } \mathbb{E}(\sigma_z) &= \sum_i p_i \lambda_i \\ &= 1|\langle 0|\psi\rangle|^2 - 1|\langle 1|\psi\rangle|^2 \\ &= \left| (1 \ 0) \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}}e^{i\varphi} \end{pmatrix} \right|^2 - \left| (0 \ 1) \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}}e^{i\varphi} \end{pmatrix} \right|^2 \\ &= \frac{1}{2} - \frac{1}{2} = 0 \end{aligned}$$

# Phase matters

Expectation value ( $\sigma_x$ ):  $\mathbb{E}(\sigma_x) = 1|\langle x; +|\psi\rangle|^2 - 1|\langle x; -|\psi\rangle|^2$

$$= \left| \begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{pmatrix} \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} e^{i\varphi} \end{pmatrix} \right|^2 - \left| \begin{pmatrix} \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix} \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} e^{-i\varphi} \end{pmatrix} \right|^2$$
$$= \cos(\varphi)$$

# That's all about quantum mechanics



Congratulations! You have learned all about quantum mechanics!

1. Quantum states are complex vectors.
2. Quantum operations are unitary operators.
3. Quantum measurement is described by the Born's rule.



# Quantum observables

Quantum **observables** can be represented by Hermitian operators.

$$\langle \psi | \hat{A} | \psi \rangle \text{ s.t. } \hat{A} = \hat{A}^\dagger$$

# Measurement bases correspond to unitary operators

$$\begin{array}{cc} \begin{array}{c} z\text{-basis} \\ (|0\rangle, |1\rangle) = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \\ = I \end{array} & \begin{array}{c} x\text{-basis} \\ (|x; +\rangle, |x; -\rangle) = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \\ = H \end{array} \end{array}$$

Question: These bases are unitary—is this an accident?

# Observable represented by Hermitian operator

Let  $A^\dagger = A$  be a Hermitian operator. Then we have

$$U \Lambda_A U^\dagger = A,$$

where  $\Lambda_A$  is a **real** diagonal matrix of eigenvalues and  $U$  is a unitary matrix.

- $U$  is a unitary operator representing the measurement basis.
- $\Lambda_A$  contains the real values for labeling the outcomes of the measurement.

# Why bother? Convenient for calculating expectation values

$$\langle \psi | A | \psi \rangle = \langle \psi | U \Lambda_A U^\dagger | \psi \rangle = \sum_i \lambda_i |\langle \psi | i \rangle|^2,$$

where

- $\lambda_i$  is the value for the  $i$ th eigenvalue of  $A$ , representing the measurement outcome.
- $|i\rangle$  is the  $i$ th eigenvector of  $A$ , representing the measurement basis.

# Popular observables: Pauli matrices

Pauli matrices are both Hermitian and unitary.

$$\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

The  $x$ ,  $y$ , and  $z$  bases are on “equal footing”

$$\langle x; \pm | y; \pm \rangle = \langle x; \pm | z; \pm \rangle = \langle y; \pm | z; \pm \rangle = \frac{1}{2}$$

Q: Is there a fourth basis satisfying the same properties?

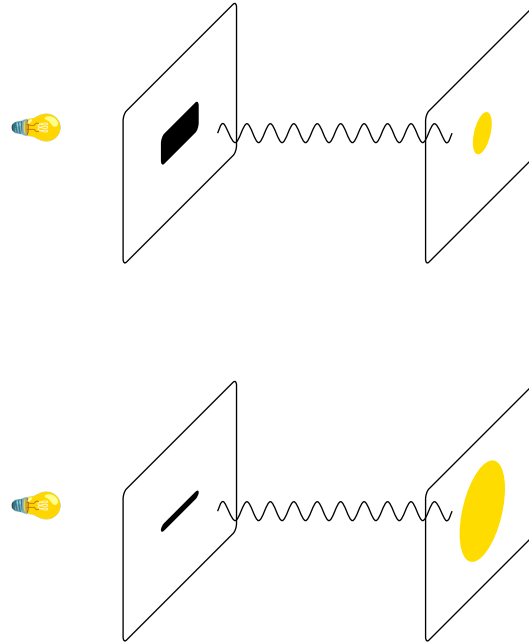
# Key facts about quantum mechanics

- **Ubiquitous**: Quantum mechanics explains all phenomena in the microscopic world, e.g., superconductivity. The only missing piece in this course is the quantum dynamics.
- **Computational advantage**: Quantum many-body systems are intractable for classical computers(Arute et al., 2019), (Pan et al., 2022).
- **Delicate**: Quantum probabilities always reduce to classical probabilities at the macroscopic limit(Zhang et al., 2024).

# Uncertainty principle

Information contained in different bases (e.g.,  $\sigma_x$  and  $\sigma_z$ ) cannot be simultaneously known precisely. (**Uncertainty principle**)

# Uncertainty principle - position and momentum



The less we know about the position,  
the more we know about the momentum.



# Returning to the magic square game

Qubits are measured in different **bases**.

$$\begin{pmatrix} I \otimes \sigma_z & \sigma_z \otimes I & \sigma_z \otimes \sigma_z \\ \sigma_x \otimes I & I \otimes \sigma_x & \sigma_x \otimes \sigma_x \\ -\sigma_x \otimes \sigma_z & -\sigma_z \otimes \sigma_x & \sigma_y \otimes \sigma_y \end{pmatrix}$$

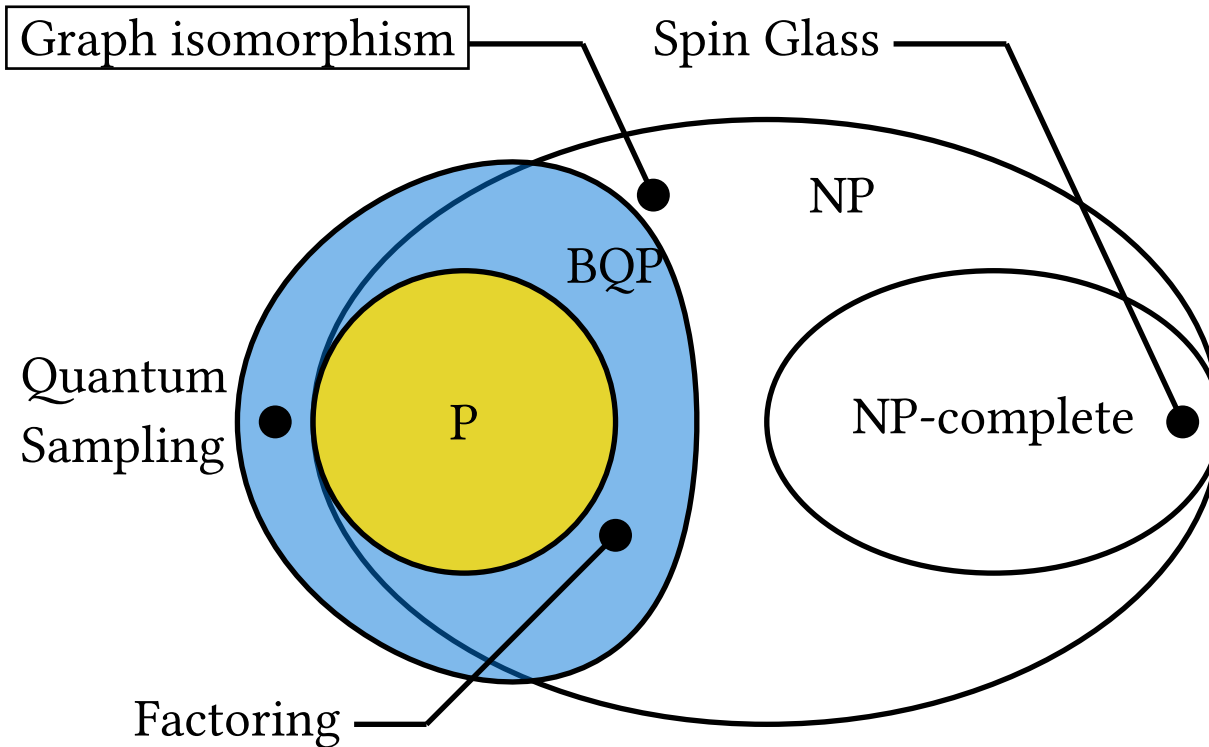
- Operators in the same row/column are **commuting** (can be measured simultaneously with certainty).
- The operators in the same row/column multiply to  $+1/-1$ , e.g.

$$\langle \psi | (\sigma_x \otimes I) (I \otimes \sigma_x) (\sigma_x \otimes \sigma_x) | \psi \rangle = 1$$

$$\langle \psi | (\sigma_z \otimes I) (I \otimes \sigma_x) (-\sigma_z \otimes \sigma_x) | \psi \rangle = -1$$

- The bell states, when measured, always give the same result.

# Quantum complexity classes



**P:** Polynomial time solvable

**NP:** Polynomial time verifiable

**NP-complete:** The hardest problems in NP

**BQP:** Polynomial time solvable on a quantum computer

# Outline

Mermin-Peres magic square

Quantum principle

**Quantum computation, AI and more**

# Modeling probabilistic distributions with quantum circuits

Born machine (Liu & Wang, 2018)

$N = 20000$   
 $\chi = 90.0\%$

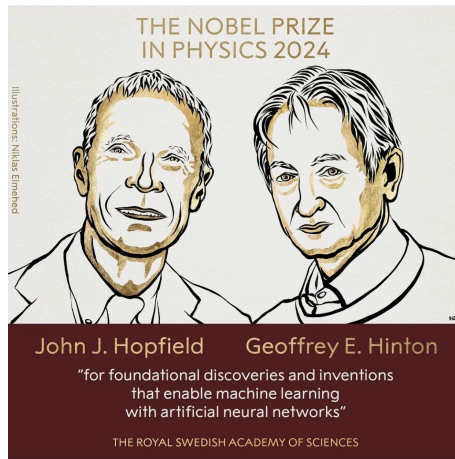


Contextual recurrent neural network(Anschuetz et al., 2023)

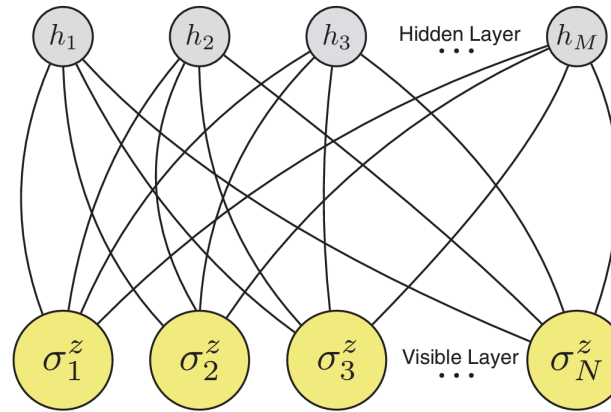
Input	“Debemos limpiar la cocina.”
Truth	“We must clean up the kitchen.”
CRNN	“We must clean the kitchen.”
GRU	“We have to turn the right address.”

# What can AI do for quantum physics?

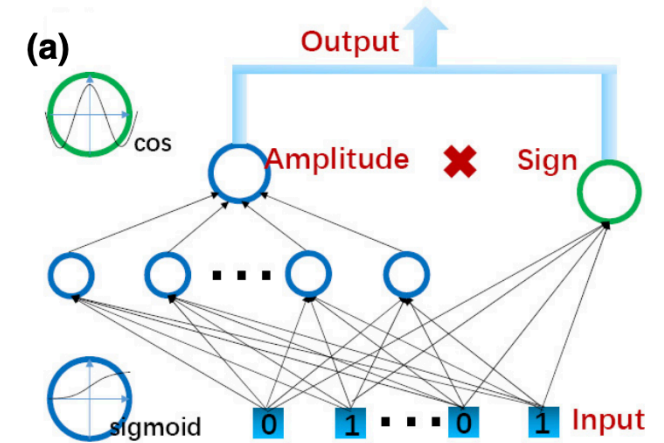
Restricted Boltzmann Machine (RBM) for quantum states.



(Hinton & Salakhutdinov,  
2006)



(Carleo & Troyer, 2017)



(Cai & Liu, 2018)

# Hardware

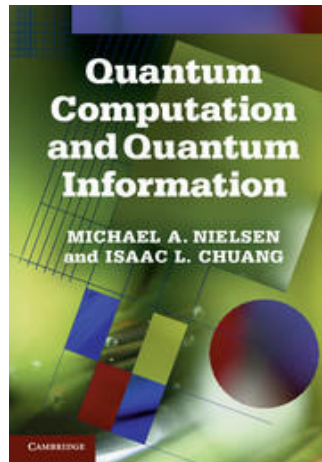
Nuclear magnetic resonance (NMR) qubits (3 qubits)

**Location:** W3-6F, Modern Matter Lab

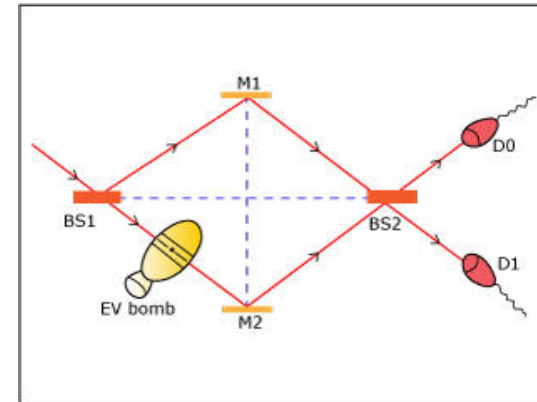


# Learning resources

Book: Nielsen & Chuang: Quantum Computation and Quantum Information



MIT Open Course: Quantum Physics I, II, III



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