B461 Database Concepts Assignment 3 Fall 2023

This assignment tests the following concepts:

- Pure SQL
- Relational Algebra (RA)
- Joins and semi-joins
- Pure SQL to RA SQL and RA Expressions

with particular focus on the last two lectures.

To turn in your assignment, you will need to upload to Canvas a single file with name assignment3.sql which contains the necessary SQL statements that solve the problems in this assignment. The assignment3.sql file must be so that the AI's can run it in their PostgreSQL environment. You should use the script file to construct the assignment3.sql file. (note that the data to be used for this assignment is included in this file.) In addition, you will need to upload a separate assignment3.txt file that contains the results of running your queries. Finally, you need to upload a file assignment3.pdf that contains the solutions to the problems that require it.

- Include all problems with the blue bullet •in assignment3.sql
- Include all problems with the red bullet •in assignment3.pdf

Database schema and instances

For the problems in this assignment we will use the following database schema:¹

Person(pid, pname, city)
Company(cname, headquarter)
Skill(skill)
worksFor(pid, cname, salary)
companyLocation(cname, city)
personSkill(pid, skill)
hasManager(eid, mid)
Knows(pid1, pid2)

In this database we maintain a set of persons (Person), a set of companies (Company), and a set of (job) skills (Skill). The pname attribute in Person is the name of the person. The city attribute in Person specifies the city in which the person lives. The cname attribute in Company is the name of the company. The headquarter attribute in Company is the name of the city wherein the company has its headquarter. The skill attribute in Skill is the name of a (job) skill.

A person can work for at most one company. This information is maintained in the worksFor relation. (We permit that a person does not work for any company.) The salary attribute in worksFor specifies the salary made by the person.

The city attribute in companyLocation indicates a city in which the company is located. (Companies may be located in multiple cities.)

A person can have multiple job skills. This information is maintained in the personSkill relation. A job skill can be the job skill of multiple persons. (A person may not have any job skills, and a job skill may have no persons with that skill.)

A pair (e, m) in hasManager indicates that person e has person m as one of his or her managers. We permit that an employee has multiple managers and that a manager may manage multiple employees. (It is possible that an employee has no manager and that an employee is

¹The primary key, which may consist of one or more attributes, of each of these relations is underlined.

not a manager.) We further require that an employee and his or her managers must work for the same company.

The relation Knows maintains a set of pairs (p_1, p_2) where p_1 and p_2 are pids of persons. The pair (p_1, p_2) indicates that the person with pid p_1 knows the person with pid p_2 . We do not assume that the relation Knows is symmetric: it is possible that (p_1, p_2) is in the relation but that (p_2, p_1) is not.

The domain for the attributes pid, pid1, pid2, salary, eid, and mid is integer. The domain for all other attributes is text.

We assume the following foreign key constraints:

- pid is a foreign key in worksFor referencing the primary key pid in Person;
- cname is a foreign key in worksFor referencing the primary key cname in Company;
- cname is a foreign key in companyLocation referencing the primary key cname in Company;
- pid is a foreign key in personSkill referencing the primary key pid in Person;
- skill is a foreign key in personSkill referencing the primary key skill in Skill;
- eid is a foreign key in hasManager referencing the primary key pid in Person;
- mid is a foreign key in hasManager referencing the primary key pid in Person;
- pid1 is a foreign key in Knows referencing the primary key pid in Person; and
- pid2 is a foreign key in Knows referencing the primary key pid in Person

Pure SQL and RA SQL

In this assignment, we distinguish between Pure SQL and RA SQL.

Pure SQL Rules:

- SELECT, FROM, WHERE
- Set operations: UNION, INTERSECT, EXCEPT
- Predicates: EXISTS, NOT EXISTS
- Predicates: IN, NOT IN
- Predicates: [NOT] (ALL, SOME)
- Window functions PARTITION, RANK etc. and aggregate functions SUM, AVG etc. are not allowed
- VIEWs, WITH clauses that obey the rules above

RA SQL Rules:

- SELECT, FROM, WHERE
- The WHERE clause can **only** be used with constants.
- Set operations: UNION, INTERSECT, EXCEPT
- JOINS, CROSS JOINS, NATURAL JOINS
- Commas are **not** allowed
- Window functions PARTITION, RANK etc. and aggregate functions SUM, AVG etc. are **not** allowed
- VIEWs, WITH clauses that obey the rules above

In particular, any other elements of SQL that are not mentioned in the lecture slides (like LIMIT) are **not** allowed for Pure SQL **or** RA SQL

1 Theoretical problems related to query translation and optimization

1. In the translation algorithm from Pure SQL to RA we assumed that the argument of each set predicate was a (possibly parameterized) Pure SQL query that did not use a union, intersect, nor an except operation.

In this problem, you are asked to extend the translation algorithm from Pure SQL to RA such that the argument of set predicate is a (parameterized) pure SQL query containing UNION, INTERSECT, EXCEPT operations.

More specifically, consider the following types of queries using the [not] exists set predicate.

Observe that there are six cases to consider:

```
(a) exists (... union ...)
(b) exists (... intersect ...)
(c) exists (... except ...)
(d) not exists (... union ...)
(e) not exists (... intersect ...)
(f) not exists (... except ...)
```

• Show how such SQL queries can be translated to equivalent RA expressions in standard notation. In the translation, you should take into account that projections do not in general distribute over intersections and over set differences. [15 pts]

To get practice, first consider the following special case where n = 1, m = 1, and k = 1. I.e., the following case: ²

²Once you can handle this case, the general case is a similar.

2 Translating Pure SQL to RA SQL

In this section, you are asked to *translate* Pure SQL queries into RA SQL queries as well as standard RA expressions using the *translation* algorithm.

You are required to show the intermediate steps that you took during the translation.

You can use the following notation to denote relation names in RA expressions:

P, P_1, P_2, \cdots	Person
C, C_1, C_2, \cdots	Company
S, S_1, S_2, \cdots	Skill
W, W_1, W_2, \cdots	worksFor
cL, cL_1, cL_2, \cdots	companyLocation
pS, pS_1, pS_2, \cdots	personSkill
hM, hM_1, hM_2, \cdots	hasManager
K, K_1, K_2, \cdots	Knows

Note: Please make note of the following example, and use it as a template to construct your answers for this section. You should write all the RA expressions in Latex or a word editor. Images of handwritten notes will NOT be accepted.

Example 1 Consider the query "Find each (p, c) pair where p is the pid of a person who works for a company c located in Bloomington and whose salary is the lowest among the salaries of persons who work for that company.

A possible formulation of this query in Pure SQL is

Translation of 'and' in the 'where' clause.

```
select q.pid, q.cname
from
       (select w.*
       from
             worksfor w
        where w.cname in (select cl.cname
                          from
                                 companyLocation cl
                          where cl.city = 'Bloomington')
        intersect
       select w.*
       from worksfor w
       where w.salary <= ALL (select w1.salary
                               from
                                     worksfor w1
                               where w1.cname = w.cname)) q;
   Translation of 'in' and '<= ALL'.
select q.pid, q.cname
from
       (select w.*
             worksfor w, companyLocation cl
       where w.cname = cl.cname and cl.city = 'Bloomington'
        intersect
        (select w.*
        from worksfor w
        except
        select w.*
               worksfor w, worksfor w1
        from
        where w.salary > w1.salary and w1.cname = w.cname)) q;
   Move 'constant' condition.
select q.pid, q.cname
     (select w.*
from
       from worksfor w,
        (select cl.* from companyLocation cl where cl.city = 'Bloomington') cl
        where w.cname = cl.cname
        intersect
        (select w.*
        from worksfor w
        except
        select w.*
        from worksfor w, worksfor w1
        where w.salary > w1.salary and w1.cname = w.cname)) q;
```

Introduction of natural join and join.

This RA SQL query can be formulated as an RA expression in standard notation as follows:

$$\pi_{W.pid,W.cname}(\mathbf{E} \cap (W - \mathbf{F}))$$

where

$$\mathbf{E} = \pi_{W.*}(W \bowtie \sigma_{city = \mathbf{Bloomington}}(cL))$$

and

$$\mathbf{F} = \pi_{W.*}(W \bowtie_{W.salary>W_1.salary \land W_1.cname=W.cname} W_1).$$

2. "Find the pid of each person that knows at least 2 people, such that at least 1 of them works at Apple or Netflix ."

A possible way to write this query in Pure SQL is

- (a) Translate the Pure SQL query to RA SQL using the translation algorithm step-by-step. [12 pts]
- (b) Write the RA Expression of the translated RA SQL query in standard notation. [5 pts]
- 3. "Return the pair (p, c) where p is the pid of a person, and c is the cname of the company where p works, such that (1) p is managed by someone who has at-least 2 skills and (2) p does not know anyone that lives in Seattle."

A possible way to write this query in Pure SQL is

```
select p.pid, c.cname
from person p , company c, worksfor w
where p.pid = w.pid and c.cname = w.cname
and exists (
    select 1 from hasManager m, personSkill ps1, personSkill ps2
    where m.eid = p.pid and ps1.pid = m.mid and ps2.pid = m.mid
    and ps1.skill <> ps2.skill
)
and p.pid not in (
    select k1.pid1 from knows k1
    where k1.pid2 in (select p1.pid from person p1 where p1.city='Seattle')
);
```

- (a) Translate the Pure SQL query to RA SQL using the translation algorithm step-by-step. [12 pts]
- (b) Write the RA Expression of the translated RA SQL query in standard notation. [5 pts]

4. "Return each skill that is the skill of at least 2 persons, such that at least 1 of them lives in Bloomington"

A possible way to write this query in Pure SQL is

```
select s.skill from skill s
where exists (
    select 1 from personskill ps1, personskill ps2
    where ps1.pid <> ps2.pid
    and ps1.skill = ps2.skill and s.skill = ps1.skill
    and exists (
        select 1 from person p where p.city = 'Bloomington'
        and (ps1.pid = p.pid or ps2.pid = p.pid)
    );
```

- (a) Translate the Pure SQL query to RA SQL using the translation algorithm step-by-step. [12 pts]
- (b) Write the RA Expression of the translated RA SQL query in standard notation. [5 pts]
- 5. "Return the pair (p, s) where p is the pid of a person that works at a company headquartered in MountainView and s is the minimum salary among all people that know p."

A possible way to write this query in Pure SQL is

```
with mv
    as (select P.pid
        from person P, worksfor W, company C
        where P.pid = W.pid and W.cname = C.cname
               and C.headquarter = 'MountainView'),
    all_that_know_p
    as (select mv.pid, W.salary
        from mv MV, knows K, worksfor W
        where K.pid2 = mv.pid and W.pid = K.pid1
        order by 1),
    min_salary
    as (select distinct ATP.pid, ATP.salary
        from all_that_know_p ATP
        where not exists (select 1 from all_that_know_p ATP1 where ATP.salary >
        ATP1.salary and ATP1.pid = ATP.pid))
select * from min_salary;
```

(a) • Translate the Pure SQL query to RA SQL using the translation algorithm step-by-step. [12 pts]

- (b) Write the RA Expression of the translated RA SQL query in standard notation. [5 pts]
- 6. "Return each cname such that
 - (1) at least 1 person working there has the OperatingSystems skill
 - (2) at least 2 persons working there live in different cities"

A possible way to write this query in Pure SQL is

- (a) Translate the Pure SQL query to RA SQL using the translation algorithm step-by-step. [12 pts]
- (b) Write the RA Expression of the translated RA SQL query in standard notation. [5 pts]