

02 - Languages, Abstract Machines, Runtime systems

Programming Languages

- **Syntax**: form of the program
- **Semantics**: meaning of well-formed program
- **Pragmatics**: conventions, guidelines, use cases

Paradigms

Classification of programming languages, based on features.

Some paradigms (where some are included in others):

- Imperative;
- Object-oriented;
- Concurrent;
- Functional;
- Logic.

Abstract Machine

An Abstract Machine is a collection of **data structures** and **algorithms** which can perform the **storage** and **execution** of programs written in a specific programming language. **Vice versa**, each Abstract Machine **defines** a **language** including all programs which can be executed by the interpreter of M.

Composition

- **Memory**: contains programs and data;
- **Interpreter**:
 - **Memory management**: how and where to store data;
 - **Primitive Data Processing**: executes easy primitive operations;
 - **Sequence control**: handles conditional jumps, returns...
 - **Data transfer control**: handles parameter passing and values returns.

Implementation

Abstract Machine M can be either implemented:

- In **hardware** or in firmware: it is difficult because they have different abstraction standpoints (e.g. calling a function in AM is easy, on hardware is really complex);
- Over an already implemented **host machine** M_0 . The **interpreters** of the two machines can **coincide** (so M is an extension of M_0), or **differ** (so M is interpreted over M_0).

To implement M over M_0 we have two different approaches:

- **Pure interpretation**: real time translation, not efficient (fetch-decode phases, no code optimizations);
- **Pure compilation**: beforehand translation of entire code.

Generally, **compilation** leads to **better performance**, while **interpretation** facilitates interactive **debugging** and **testing**. Modern implementations is to go for **both approaches**: compilation into intermediate program, interpreted

through **Virtual Machine**. Several advantages:

- **Portability**: once a program is compiled to the intermediate program, it can be executed on any platform equipped with the VM;
- **Interoperability**: can create a new language which is compiled for an existing VM;

This kind of implementation can be **iterated**, leading to a hierarchy.

Runtime Systems

The **execution model** of a programming language defines the **sequence of steps** that are taken to **process** the code and **produce** the desired output (fetch-decode-execute).

Runtime systems implement **part** of the **execution model**, providing **runtime support**. Needed both by interpreted and by compiled programs.

Generally made of:

- **Compiler/Interpreter/VM**;
- **Code generated** by the compiler;
- **Code** running in other **threads** during program execution (e.g. garbage collector);
- Language **libraries**;
- **OS** functionalities;

Needed for:

- **Memory management**: stack (push/pop activation records), heap (allocation, garbage collection);
- Interaction with **runtime environment**, which are data needed by the program in execution but not part of the program itself (e.g. environment variables, I/O, threads communication, dynamic type checking/binding...);
- **Debugging**, monitoring, ecc...