Example D.1

记公理系统LPGL的公式构成为

$$egin{aligned} u &:= c \mid x \mid (u \cdot u) \mid (u + u) \mid !u \ P &:= p \mid \neg P \mid P
ightarrow P \mid \Box P \mid u : P \end{aligned}$$

we try to proof $\mathbf{LPGL} \vdash u : \Box P \to P$

Proof.

- $(1)\; u: \Box P \to \Box P \quad Factivity$
- $(2) \neg \Box P \rightarrow \neg u : \Box P \quad 1$
- (3) $\neg u: \Box P \rightarrow \Box (\neg u: \Box P)$ weak negative introspection
- $(4) \square (\neg u : \square P) \rightarrow \square (u : \square P \rightarrow P) \quad TAUT, NEC, K$
- $(5) \neg \Box P \rightarrow \Box (u : \Box P \rightarrow P) \quad 2, 3, 4$
- (6) $\Box P \rightarrow \Box (u:\Box P \rightarrow P) \quad TAUT, NEC, K$
- $(7) \square (u:\square P \rightarrow P) \quad 5,6$
- (8) $u: \Box P \to P$ 7, reflexivity rule

Example D.2

记公理系统 $hybrid-S4_f$ 的公式构成为

$$F ::= i \mid P \mid \hat{\neg} F \mid F \hat{\rightarrow} F \mid \hat{\Box} F \mid @_i F$$

外化公理: $\neg P \rightarrow \hat{\neg} P, (P \rightarrow Q) \hat{\rightarrow} (P \hat{\rightarrow} Q)$

Let $@_i$ be $@_{\mathbf{LPGL}}$, we try to proof $@_i(u:\Box P \to P)$ is a theorem.

Proof.

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(1) @_i(u : \Box P \to \Box P) Factivity on LPGL
(2) @_i((u:\Box P \to \Box P) \to (\neg \Box P \to \neg u:\Box P)) TAUT on LPGL
(4) @_i(u: \Box P \to \Box P) \hat{\to} @_i(\neg \Box P \to \neg u: \Box P) \quad 2, (*), K_{@}
(5) @_i(\neg \Box P \rightarrow \neg u : \Box P) 1, 4, MP
(6) @_i(\neg u: \Box P \to \Box(\neg u: \Box P)) weak negative introspection on LPGL
(7) @_i(\neg u: \Box P \to (u: \Box P \to P)) TAUT on LPGL
(8) @_i(\Box(\neg u:\Box P) \rightarrow \Box(u:\Box P \rightarrow P)) \quad 7, NEC, K \ on \ \mathbf{LPGL}
(9) \ @_i \bigg( (\neg \Box P \to \neg u : \Box P) \to \Big( (\neg u : \Box P \to \Box (\neg u : \Box P)) \to
       \big((\Box(\neg u:\Box P)\to\Box(u:\Box P\to P))\to(\neg\Box P\to\Box(u:\Box P\to P))\big)\Big)\Big)
                                                                                 TAUT on LPGL
(11) @_i(\neg \Box P \to \neg u : \Box P) \hat{\to} \Big( @_i(\neg u : \Box P \to \Box (\neg u : \Box P)) \hat{\to}
       \big(@_i(\Box(\neg u:\Box P)\to\Box(u:\Box P\to P))\hat{\to}@_i(\neg\Box P\to\Box(u:\Box P\to P))\big)\Big)
                                                                          10, (*), @NEC, K_{\odot}
(12) @_i(\neg \Box P \rightarrow \Box (u : \Box P \rightarrow P)) \quad 5, 6, 8, 11, MP
(13) @_i(P \rightarrow (u : \Box P \rightarrow P)) TAUT on LPGL
(14) @_i(\Box P \to \Box (u:\Box P \to P)) \quad 13, NEC, K \ on \ \mathbf{LPGL}
(15) @_i \Big( (\neg \Box P \to \Box (u : \Box P \to P)) \to
       \big((\Box P 	o \Box (u:\Box P 	o P)) 	o \Box (u:\Box P 	o P)\big) \big) TAUT \ on \ \mathbf{LPGL}
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我们没法使用下层逻辑的规则,但是公理可以用,规则被外化到主语言中了,用term上的函数表示

 $ig(@_i(\Box P o\Box(u:\Box P o P))\hat{ o}@_i\Box(u:\Box P o P)ig)\quad 16,(*),K_@$

用term内化了模型中系统的可证的公式的具体的证明过程,这是有趣的。因为term提供了一种证明的核心直觉,非常的简洁,清晰的刻画了证明的关键步骤。

term不仅能够内化外部证明,也能在给出下层系统之后,内化内部证明 成为可以用的推理规则

(19) $@_i(u: \Box P \to P)$ reflexivity rule on **LPGL**

 $(17) \otimes_i (\neg \Box P \rightarrow \Box (u : \Box P \rightarrow P)) \hat{\rightarrow}$

 $(18) @_i \square (u : \square P \to P) \quad 12, 14, 17, MP$

Example D.3

下层逻辑LPGL:

$$\hat{u} ::= c \mid x \mid (u \cdot u) \mid (u + u) \mid !u$$

$$P ::= p \mid \neg P \mid P \rightarrow P \mid \Box P \mid \hat{u} : P$$

函数:

 $f_{nec}(t)$ 对应**LPGL**中的NEC规则;

 $f_r(t)$ 对应**LPGL**中的reflexivity rule规则;

 $f_k(t)$ 对应**LPGL**中的K规则;

$$(1) c : (\hat{u} : \Box P \to \Box P) \quad Factivity$$

$$(2) d: ((\hat{u}: \Box P \to \Box P) \to (\neg \Box P \to \neg \hat{u}: \Box P)) \quad TAUT$$

$$(3) \ dc: (\neg \Box P
ightarrow \neg \hat{u}: \Box P) \quad 2,1,Application$$

(4)
$$h: (\neg \hat{u}: \Box P \to \Box (\neg \hat{u}: \Box P))$$
 weak negative introspection

(5)
$$j: (\neg \hat{u}: \Box P \rightarrow (\hat{u}: \Box P \rightarrow P))$$
 $TAUT$

$$(6) \ f_k f_{nec}(j) : (\Box (\lnot \hat{u} : \Box P)
ightarrow \Box (\hat{u} : \Box P
ightarrow P))$$

$$(7) \ m: \bigg((\neg \Box P \to \neg \hat{u}: \Box P) \to \Big((\neg \hat{u}: \Box P \to \Box (\neg \hat{u}: \Box P)) \to$$

$$\left(\left(\Box(\neg \hat{u}:\Box P)\to\Box(\hat{u}:\Box P\to P)\right)\to \left(\neg\Box P\to\Box(\hat{u}:\Box P\to P)\right)\right)\right)$$

TAUT

$$(8) \ m(dc)h(f_kf_{nec}(j)): (
eg \square P
ightarrow \square(\hat{u}:\square P
ightarrow P)) \quad 7,3,4,6, Application$$

(9)
$$n:(P \to (\hat{u}:\Box P \to P))$$
 $TAUT$

$$(10)\ o: \Big((\neg \Box P \to \Box (\hat{u}: \Box P \to P)) \to$$

$$ig((\Box P
ightarrow \Box (\hat{u}:\Box P
ightarrow P))
ightarrow \Box (\hat{u}:\Box P
ightarrow P)ig) \quad TAUT$$

$$(11) \ oig(m(dc)h(f_kf_{nec}(j))ig)(f_kf_{nec}(n)): \Box(\hat{u}:\Box P o P) \quad 10,8,9,application$$

$$(12) \ f_r\Big(oig(m(dc)h(f_kf_{nec}(j))ig)(f_kf_{nec}(n))\Big): (\hat{u}:\Box P o P) \quad 11, ext{reflexivity rule}$$

Example D.4

上层逻辑hybrid-**LPS** 4_f :

$$t ::= c \mid x \mid t \cdot t \mid t + t \mid !t \mid !_i t \mid ?_i t \mid f_{nec}(t) \mid f_r(t)$$
 $F ::= i \mid P \mid \hat{\neg} F \mid F \hat{\rightarrow} F \mid \hat{\Box} F \mid @_i F \mid t : F$

并使用函数 $q_{nec}(t),q_k(t)$ 简化证明过程:

 $g_k(t)$ 对应hybrid-**LPS4**f中的 $K_@$ 规则;

Proof. Let $@_i$ be $@_{\mathbf{LPGL}}$

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(1) \ c: @_{i}(\hat{u}: \Box P \to \Box P) \quad Factivity \ on \ \mathbf{LPGL}
(2) \ d: @_{i}((\hat{u}: \Box P \to \Box P) \to (\neg \Box P \to \neg \hat{u}: \Box P)) \quad TAUT \ on \ \mathbf{LPGL}
(3) \ g_{k}(d)c: @_{i}(\neg \Box P \to \neg \hat{u}: \Box P) \quad 2, (*), g_{k}, 1, Application
(4) \ h: @_{i}(\neg \hat{u}: \Box P \to \Box (\neg \hat{u}: \Box P)) \quad \text{weak negative introspection } on \ \mathbf{LPGL}
(5) \ j: @_{i}(\neg \hat{u}: \Box P \to (\hat{u}: \Box P \to P)) \quad TAUT \ on \ \mathbf{LPGL}
(6) \ f_{k}f_{nec}(j): @_{i}(\Box (\neg \hat{u}: \Box P) \to \Box (\hat{u}: \Box P \to P))
(7) \ m: @_{i}\left((\neg \Box P \to \neg \hat{u}: \Box P) \to \left((\neg \hat{u}: \Box P \to P) \to \Box (\hat{u}: \Box P \to P))\right) \to \left((\Box (\neg \hat{u}: \Box P) \to \Box (\hat{u}: \Box P \to P)) \to (\neg \Box P \to \Box (\hat{u}: \Box P \to P))\right)\right)
TAUT \ on \ \mathbf{LPGL}
(8) \ g_{k}\left(g_{k}(g_{k}(m)(g_{k}(d)c))h\right)f_{k}f_{nec}(j): @_{i}(\neg \Box P \to \Box (\hat{u}: \Box P \to P))
(9) \ n: @_{i}(P \to (\hat{u}: \Box P \to P)) \quad TAUT \ on \ \mathbf{LPGL}
(10) \ o: @_{i}\left((\neg \Box P \to \Box (\hat{u}: \Box P \to P)) \to \Box (\hat{u}: \Box P \to P)\right)\right)
TAUT \ on \ \mathbf{LPGL}
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$$egin{aligned} ig((oxdot P
ightarrow oxdot (u:oxdot P
ightarrow P))
ightarrow oxdot (u:oxdot P
ightarrow P)) & TAUT \ on \ \mathbf{LPGL} \ ig(11) \ g_k igg(g_k(o) \Big(g_k ig(g_k(m)(g_k(d)c)ig)h\Big) f_k f_{nec}(j)\Big) igg) f_k f_{nec}(n): \ ig(a_i oxdot (a:oxdot P
ightarrow P)) & 10, (*), g_k, 8, g_k, 9, application \end{aligned}$$

(12)
$$f_r \bigg(g_k \Big(g_k(o) \Big(g_k \Big(g_k(m) (g_k(d)c) \Big) h \Big) f_k f_{nec}(j) \Big) \bigg) f_k f_{nec}(n) \bigg) :$$

$$@_i(\hat{u} : \Box P \to P) \qquad \qquad 11, \text{ reflexivity rule } on \ \mathbf{LPGL}$$

这个例子体现出混合算子的必要性。让系统中的定理转化为公理系统中的定理。从而能够公理化方法推理。

四个例子构成平行关系,总结内化定理。

给出的证明项,要求都是从公理出发的,这保证了证明的可靠性,不会出现跳步