An Introduction to Network Centrality

with Applications in R

David Schoch

2018-12-07

Contents

Preface

Chapter 1

Introduction

Chapter 2

Centrality Indices

The purpose of network centrality is to identify important actors or entities in a network. Structural importance is determined by so called *measures of centrality*, commonly defined in terms of indices $c : V \to \mathbb{R}$ interpreted as

$$c(u) > c(v) \iff u \text{ is more central than } v.$$

Since the meaning of structural importance is by no means unambiguous, a vast amount of different indices have been introduced over time. In addition, any mapping $c : V \to \mathbb{R}$ induces a ranking of nodes, but not every such ranking might represent a plausible concept of structural importance.

2.1 Degree

Degree centrality is the most simple form of a centrality index. It is defined as

$$c_d(u) = \{v : \{u, v\} \in E\} = |N(u)|$$

Degree centrality is a purely *local* measure since it only depends on the direct neighborhood of a node. A simple application example is popularity in friend- ship networks, i.e. "who has the most friends?". The definition of degree centrality can easily be adapted for directed and weighted networks. For directed networks,

$$c_d^+(u) = \{v : (u, v) \in E\} = |N^+(u)| \quad \text{and} \quad c_d^-(u) = \{v : (v, u) \in E\} = |N^-(u)|$$

are the *out-degree* and *in-degree*, respectively. *Weighted degree*, sometimes also refered to as *strength*, is defined as

$$c_{wd}(u) = \sum_{v \in N(u)} w_{uv}.$$

2.2 Betweenness (and variants)

Betweenness was independently developed by Anthonisse (1971) and Freeman (1977) and is an extension of *stress centrality*, introduced by Shimbel (1953). Shimbel assumes that the number of shortest paths containing a node u is an estimate for the amount of "stress" the node has to sustain in a network. The more shortest paths run through a node the more central it is. Formally, stress centrality is defined as

$$c_{stress}(u) = \sum_{s,t \in V} \sigma(s, t|u),$$

where $\sigma(s,t|u)$ is the number of shortest paths from s to t passing through u. Instead of the absolute number of shortest paths, betweenness centrality quantifies their relative number. The relative count is given by

$$\delta(s,t|u) = \frac{\sigma(s,t|u)}{\sigma(s,t)},$$

where $\sigma(s,t)$ is the total number of shortest paths connecting s and t. The expression $\delta(s,t|u)$ can be interpreted as the extent to which u controls the communication between s and t. Betweenness is then defined as

$$c_b(u) = \sum_{s,t \in V} \delta(s,t|u)$$

The interpretation of betweenness is not only restricted to communication. More generally, betweenness quantifies the influence of vertices on the trans- fer of items or information through the network with the assumption that it follows shortest paths. Many different variants of shortest path betweenness have been proposed to incorporate additional assumptions, e.g. the specific lo- cation of a vertex u on a shortest s,t-path or its length. Some of these variants are given in the following table.

proximal source	$c_{bs}(u) = \sum_{s,t \in V} \delta(s,t|u) \cdot A_{su}$
proximal target	$c_{bt}(u) = \sum_{s,t \in V} \delta(s,t|u) \cdot A_{ut}$
k-bounded distance	$c_{bk}(u) = \sum_{dist(s,t) \leq k} \delta(s,t|u)$
length-scaled	$c_{bd}(u) = \sum_{s,t \in V} \frac{\delta(s,t|u)}{dist(s,t)}$
linearly-scaled	$c_{bl}(u) = \sum_{s,t \in V} \delta(s,t|u) \frac{dist(s,u)}{dist(s,t)}$

Details on these variants can be found in Brandes (2008). Other variants of the betweenness concept rely on different assumptions of transfer in networks besides shortest paths.

A measure based on network flow was defined by Freeman et al. (1991). The authors assume information as a flow process and assign to each edge a non-negative value representing the maximum amount of information that can be passed between the endpoints. The aim is then to measure the extent to which the maximum flow between two vertices s and t depends on a vertex u. Denote by $f(s,t)$ the maximum (s,t)-flow w.r.t. constraints imposed by edge capacities and the amount of flow which must go through u by $f(s,t|u)$. Similar to shortest path betweenness, *flow betweenness* is then defined as

$$c_f(u) = \sum_{s,t \in V} \frac{f(s,t|u)}{f(s,t)}.$$

The index was introduced as a betweenness variant for weighted networks but it can easily applied to unweighted ones. In the case of simple undirected and unweighted networks, the maximum (s,t)-flow is equivalent to the number of edge disjoint (s,t)-paths and $f(s,t|u)$ is the number of paths u lies on.

A variant based on walks instead of paths was proposed by Newman (2005). His *random walk betweenness* calculates the expected number of times a random (s,t)-walk passes through a vertex u, averaged over all s and t. Newman shows, that his variant of betweenness can also be calculated with a current-flow analogy by viewing a graph as an electrical network. Random walk betweenness is then equiva- lent to the amount of current that flows through u. The index is thus also known as *current flow betweenness*. Details and formal definitions of his versions can be found in the literature (Newman, 2005; Brandes and Fleischer, 2005).

Two variants based on the *randomzed shortest path* (RSP) framework (Yen et al., 2008; Saerens et al., 2009; Kivimäki et al., 2014) were proposed by Kivimäki et al. (2016). The variants are refered to as *simple RSP betweenness* and *RSP net betweenness*. The derivation of both indices is rather involved and goes beyond the scope of this tutorial. The interested reader should consult the original work for details. One aspect worth mentioning, though, is that both variants include a tuning parameter β. Both variants converge to the traditional version of betweennes for $\beta \to \infty$. For $\beta \to 0$, simple RSP betweenness converges to the expected number of visits to a node over all absorbing walks with respect to the unbiased random walk probabilities. This means for undirected networks, that the index converges to degree. RSP net betweenness, on the other hand, converges to Newmann's random walk betweenness.

All variants of betweenness can be described in a more general form considering a flow of information analogy. Depending on the assumption of how information is 'flowing' between s and t, the set $P(s,t)$ contains all possible information channels to transmit the piece of information. This set might contain all shortest (s,t)-paths if the information has to be trans- mitted as fast as possible or all random (s,t)-walks when the delivery time does not matter. In principle, any kind of trajectory on a graph can be thought of as an information channel. The set $P(s,t|u)$ contains all information channels where the vertex u is in a position to control the information flow. For shortest path betweenness, u is in a controlling position if it is part of an information channel and for proximal target betweenness if it presents the information to the target t. In the former case $P(s,t|u)$ comprises all elements of $P(s,t)$ that contain u as an intermediary and in the latter all elements that contain the edge (u,t) . Again, the position of control could be defined as any location on a trajectory. A measure of relative betweenness is then defined with aggregation rules over the two specified sets, commonly the fraction of their cardinalities. This fraction can also be weighted according to specified rules, e.g. as in length scaled betweenness. Aggregating over all possible sources and targets, we can thus define a *generic betweenness* index as

$$c_{bg}(u) = \sum_{s,t \in V} \frac{|P(s,t|u)|}{|P(s,t)|} \cdot w(s,t).$$

Thus, many other variants are possible, for instance also k-betweenness mentioned by Borgatti and Everett (2006), where $P(s,t)$ is the set of all (s,t)-paths of length at most k.

There exist many more betweenness-like indices that where not covered here, but are listed below in no particular order:

- *communicability betweenness* based on the matrix exponential (Estrada et al., 2009)
- α and β betweenness, which are closely related to current flow betweenness (Avrachenkov et al., 2013, 2015),
- *ranking betweenness*, which combines betweenness with the idea of PageRank (Agryzkov et al., 2014)
- range-limited forms of betweenness (Ercsey-Ravasz et al., 2012)
- *bridgeness* (Jensen et al., 2016)
- *super mediator* (Saito et al., 2016)
- *BridgeRank* (Salavati et al., 2018)

2.3 Closeness (and variants)

2.4 Eigenvector (and variants)

2.5 Others

Bibliography

Agryzkov, T., Oliver, J. L., Tortosa, L., and Vicent, J. (2014). A new betweenness centrality measure based on an algorithm for ranking the nodes of a network. *Applied Mathematics and Computation*, 244(Supplement C):467–478.

Anthonisse, J. (1971). The rush in a directed graph. *Stichting Mathematisch Centrum. Mathematische Besliskunde*, BN 9/71.

Avrachenkov, K., Litvak, N., Medyanikov, V., and Sokol, M. (2013). Alpha current flow betweenness centrality. In *International Workshop on Algorithms and Models for the Web-Graph*, pages 106–117. Springer.

Avrachenkov, K. E., Mazalov, V. V., and Tsynguev, B. T. (2015). Beta Current Flow Centrality for Weighted Networks. In Thai, M. T., Nguyen, N. P., and Shen, H., editors, *Computational Social Networks*, Lecture Notes in Computer Science, pages 216–227. Springer International Publishing.

Borgatti, S. P. and Everett, M. G. (2006). A Graph-theoretic perspective on centrality. *Social Networks*, 28(4):466–484.

Brandes, U. (2008). On variants of shortest-path betweenness centrality and their generic computation. *Social Networks*, 30(2):136–145.

Brandes, U. and Fleischer, D. (2005). Centrality Measures Based on Current Flow. In *STACS*, volume 3404, pages 533–544. Springer.

Ercsey-Ravasz, M., Lichtenwalter, R. N., Chawla, N. V., and Toroczkai, Z. (2012). Range-limited centrality measures in complex networks. *Physical Review E*, 85(6):066103.

Estrada, E., Higham, D. J., and Hatano, N. (2009). Communicability betweenness in complex networks. *Physica A: Statistical Mechanics and its Applications*, 388(5):764–774.

Freeman, L. C. (1977). A Set of Measures of Centrality Based on Betweenness. *Sociometry*, 40(1):35–41.

Freeman, L. C., Borgatti, S. P., and White, D. R. (1991). Centrality in valued graphs: A measure of betweenness based on network flow. *Social Networks*, 13(2):141–154.

Jensen, P., Morini, M., Karsai, M., Venturini, T., Vespignani, A., Jacomy, M., Cointet, J.-P., Mercklé, P., and Fleury, E. (2016). Detecting global bridges in networks. *Journal of Complex Networks*, 4(3):319–329.

Kivimäki, I., Lebichot, B., Saramäki, J., and Saerens, M. (2016). Two betweenness centrality measures based on Randomized Shortest Paths. *Scientific Reports*, 6.

Kivimäki, I., Shimbo, M., and Saerens, M. (2014). Developments in the theory of randomized shortest paths with a comparison of graph node distances. *Physica A: Statistical Mechanics and its Applications*, 393(Supplement C):600–616.

Newman, M. E. J. (2005). A measure of betweenness centrality based on random walks. *Social Networks*, 27(1):39–54.

Saerens, M., Achbany, Y., Fouss, F., and Yen, L. (2009). Randomized Shortest-Path Problems: Two Related Models. *Neural Computation*, 21(8):2363–2404.

Saito, K., Kimura, M., Ohara, K., and Motoda, H. (2016). Super mediator – A new centrality measure of node importance for information diffusion over social network. *Information Sciences*, 329(Supplement C):985–1000.

Salavati, C., Abdollahpouri, A., and Manbari, Z. (2018). BridgeRank: A novel fast centrality measure based on local structure of the network. *Physica A: Statistical Mechanics and its Applications*, 496:635–653.

Shimbel, A. (1953). Structural parameters of communication networks. *The bulletin of mathematical biophysics*, 15(4):501–507.

Yen, L., Saerens, M., Mantrach, A., and Shimbo, M. (2008). A Family of Dissimilarity Measures Between Nodes Generalizing Both the Shortest-path and the Commute-time Distances. In *Proceedings of the 14th ACM SIGKDD International Conference on Knowledge Discovery and Data Mining*, KDD '08, pages 785–793, New York, NY, USA. ACM.