# Strategy for developing efficient programs:

- 1. Design the program well
- 2. Implement the program well\*\*
- 3. Test the program well
- 4. Only after you're sure it's working, measure performance
- 5. If (and only if) performance is inadequate, find the "hot spots"
- 6. Tune the code to fix these
- 7. Repeat measure-analyse-tune cycle until performance ok

(\*\* see "Programming Pearls", "Practice of Programming", etc. etc.) Rapid development of a prototype may be the best way to discover/assess performance issues.

Hence Fred Brooks maxim - "Plan To Throw One Away".

# simple C program which allocate+initialializes an array

```
/home/cs2041/public html/code/performance/cachegrind example.c
#include <stdio.h>
#include <stdlib.h>
#include <assert.h>
void test0(int x, int y, int a[x][y]) {
    int i, j;
    fprintf(stderr, "writing to array i-j order\n");
    for (i = 0: i < x: i++)
        for (j = 0; j < y; j++)
            a[i][j] = i+j;
void test1(int x, int y, int a[x][y]) {
    fprintf(stderr, "writing to array j-i order\n");
    int i, j;
    for (j = 0; j < y; j++)
        for (i = 0; i < x; i++)
            a[i][j] = i+j;
int main(int argc, char*argv[]) {
    int x = atoi(argv[2]):
    int y = atoi(argy[3]);
    fprintf(stderr, "allocating a %dx%d array = %11d bytes\n", x, y, ((long long)x)*y*sizeof (int));
    void *m = malloc(x*y*sizeof (int));
    assert(m):
    switch (atoi(argv[1])) {
    case 0: test0(x, y, m); break;
    case 1: test1(x, y, m); break;
    7
    return 0:
```

Although the loops are almost identical, the first loop runs 20x faster on a large array!

```
$ time ./cachegrind_example 0 32000 32000
allocating a 32000x32000 \text{ array} = 4096000000 \text{ bytes}
writing to array i-j order
real 0m0.893s
user 0m0.364s
sys 0m0.524s
$ time ./cachegrind example 1 32000 32000
allocating a 32000x32000 \text{ array} = 4096000000 \text{ bytes}
writing to array j-i order
real 0m15.189s
user 0m14.633s
sys 0m0.528s
```

The tool valgrind used to detect accesses to uninitialized variables at runtime also can give memory caching infomation.

The memory subsystem is beyond the scope of this course but you can see valgrind explain the performance difference between these loops.

For the first loop D1 miss rate = 24.8%

For the second loop D1 miss rate = 99.9%

Due to the C array memory layout the first loop produces much better caching performance.

Tuning caching performance is important for some application and valgrind makes this much easier.

```
$ valgrind '--tool=cachegrind' ./cachegrind example 0 10000 10000
allocating a 10000x10000 array = 400000000 bytes
writing to array i-j order
==7025==
==7025== I refs: 225,642,966
==7025== I1 misses:
                           882
==7025== LLi misses:
                         875
==7025== I1 miss rate:
                          0.00%
==7025== LLi miss rate:
                          0.00%
==7025==
==7025== D refs:
                    25,156,289 (93,484 rd + 25,062,805 wr)
==7025== D1  misses: 6,262,957 (2,406 rd + 6,260,551 wr)
==7025== LLd misses: 6,252,482 (1,982 rd + 6,250,500 wr)
==7025== D1 miss rate:
                          24.8% ( 2.5%
                                                  24.9%
==7025== LLd miss rate:
                          24.8% ( 2.1%
                                                  24.9% )
==7025==
==7025== LL refs: 6,263,839 (3,288 rd + 6,260,551 wr)
==7025== LL misses: 6,253,357 (2,857 rd + 6,250,500 wr)
==7025== LL miss rate:
                           2.4% ( 0.0%
                                                  24.9% )
```

```
$ valgrind '--tool=cachegrind' ./cachegrind example 1 10000 10000
allocating a 10000x10000 array = 400000000 bytes
writing to array j-i order
==7006==
==7006== I refs: 600,262,960
==7006== I1 misses:
                           876
==7006== LLi misses:
                          869
==7006== I1 miss rate:
                          0.00%
==7006== LLi miss rate:
                          0.00%
==7006==
==7006== D refs: 100,056,288 (43,483 rd + 100,012,805 wr)
==7006== D1 misses:
                    100,002,957 (2,405 rd + 100,000,552 wr)
==7006== LLd misses: 6,262,481 (1,982 rd
                                           + 6,260,499 wr)
==7006== D1 miss rate:
                          99.9% ( 5.5%
                                                   99.9%)
                                           +
==7006== LLd miss rate: 6.2% ( 4.5%
                                                    6.2%
==7006==
==7006== LL refs: 100,003,833 (3,281 rd
                                           + 100,000,552 wr)
==7006== LL misses: 6,263,350 (2,851 rd
                                              6,260,499 wr)
==7006== LL miss rate:
                           0.8% (
                                                    6.2\%
                                   0.0%
                                           +
```

#### Where is execution time being spent?

Typically programs spend most of their execution time in a small part of their code.

This is often quoted as the 90/10 rule (or 80/20 rule or ...):

"90% of the execution time is spent in 10% of the code"

#### This means that

- most of the code has little impact on overall performance
- small parts of the code account for most execution time

We should clearly concentrate efforts at improving execution spped in the 10% of code which accounts for most of the execution time.

#### gcc -p/gprof

Given the -p flag gcc instruments a C program to collect profile information

When the program executes this data is left in the file gmon.out. The program gprof analyzes this data and produces:

- number of times each function was called
- % of total execution time spent in the function
- average execution time per call to that function
- execution time for this function and its children

Arranged in order from most expensive function down. It also gives a *call graph*, a list for each function:

- which functions called this function
- which functions were called by this function

# Program for producing sorted counts of words



Program is slow on large inputs e.g.

We can instrument the program to collect profiling information and examine it with gcc

```
$ gcc -p -g word_frequency0.c -o word_frequency0_profile
$ head -10000 WarAndPeace.txt|word_frequency0_profile >/dev/null
$ gprof word_frequency0_profile
```

% с	umulative	self		self	total	
time	seconds	seconds	calls	ms/call	ms/call	name
88.90	0.79	0.79	88335	0.01	0.01	get
7.88	0.86	0.07	7531	0.01	0.01	put
2.25	0.88	0.02	80805	0.00	0.00	get_word
1.13	0.89	0.01	1	10.02	823.90	read_words
0.00	0.89	0.00	2	0.00	0.00	size
0.00	0.89	0.00	1	0.00	0.00	create_map
0.00	0.89	0.00	1	0.00	0.00	keys
0.00	0.89	0.00	1	0.00	0.00	sort_words

Its clear that only the functions *get* and to a much lesser extent *put* are relevant to performance improvement.

Examine *get* and we find it traverses a linked list. So replace it with a binary tree and the program runs 200x faster on War and Peace.

Was C the best choice for our count words program.

#### Shell - Word Count

#### For comparison Shell:

/home/cs2041/public\_html/code/performance/word\_frequency.sh

```
#!/bin/sh
tr -c a-zA-Z ' '|
tr ' ' '\n'|
tr A-Z a-z|
egrep -v '^$'|
sort|
uniq -c
```

#### Perl - Word Count

```
/home/cs2041/public html/code/performance/word frequency.pl
#!/usr/bin/perl -w
while ($line = <>) {
    =  tr/A-Z/a-z/;
    foreach \$word (\$line =~ /[a-z]+/g) {
        $count{$word}++;
@words = keys %count;
@sorted words = sort {$count{$a} <=> $count{$b}} @words;
foreach $word (@sorted_words) {
    printf "%8d %s\n", $count{$word}, $word;
}
```

## Python - Word Count

```
/home/cs2041/public html/code/performance/word frequency.py
#!/usr/bin/python3
import fileinput, re, collections
count = collections.defaultdict(int)
for line in fileinput.input():
    for word in re.findall(r'\w+', line.lower()):
        count[word] += 1
words = count.keys()
sorted words = sorted(words, key=lambda w: count[w])
for word in sorted words:
    print("%8d %s" % (count[word], word))
```

```
Shell. Perl and Python are slower - but a lot less code.
So faster to write, less bugs to find, easier to maintain/modify
$ time word_frequency1 <WarAndPeace.txt >/dev/null
       0m0.277s
real
user 0m0.268s
sys 0m0.008s
$ time word frequency.sh <WarAndPeace.txt >/dev/null
real 0m0.564s
    0m0.584s
user
sys
      0m0.036s
$ time word frequency.pl <WarAndPeace.txt >/dev/null
real 0m0.643s
user 0m0.632s
sys 0m0.012s
$ time word frequency.py <WarAndPeace.txt >/dev/null
real 0m1.046s
user 0m0.836s
       0m0.024s
sys
$ wc word frequency1.c word frequency.sh word frequency.pl
286 759 5912 word frequency1.c
     19 82 word frequency.sh
      38 325 word_frequency.py
  11
  14
      43
          301 word frequency.pl
```

```
Here is a cp implementation in C using low-level calls to read/write
/home/cs2641/public html/code/serformance/cs0.c
// Written by andrewt@cse.unsw.edu.au
// as a COMP2041 lecture example
// copy input to output using read/write system calls
// for each byte - very inefficient and Unix/Linux specific
#include <stdio.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
#include <stdlib.h>
#include <unistd.h>
copy_file_to_file(int in_fd, int out_fd) {
   while (1) {
        char c[1]:
        int bytes_read = read(in_fd, c, 1);
        if (bytes_read < 0) {
                perror("cp: "):
                exit(1);
        if (bytes read == 0)
        int bytes written = write(out fd. c. bytes read):
        if (bytes_written <= 0) {
            perror("cp: ");
            exit(1):
main(int argc, char *argv[]) {
    if (argc != 3) {
        fprintf(stderr, "cp <src-file> <destination-file>\n"):
        return 1:
    int in_fd = open(argv[1], O_RDONLY);
    if (in_fd < 0) {
        fprintf(stderr, "cp: %s: ", argv[1]):
        perror("");
        return 1;
    int out_fd = open(argv[2], 0_WRONLY|0_CREAT|0_TRUNC, S_IRWXU);
    if (out fd <= 0) {
        fprintf(stderr, "cp: %s: ", argv[2]);
        perror("");
        return 1:
   copy file to file(in fd. out fd):
    return 0:
```

Its suprisingly slow compared to /bin/cp

```
$ time /bin/cp input_file /dev/null
real 0m0.006s
user 0m0.000s
sys 0m0.004s
$ gcc cp0.c -o cp0
$ time ./cp0 input file /dev/null
real 0m6.683s
user 0m0.932s
sys 0m5.740s
$ gcc -03 cp0.c -o cp0
$ time ./cp0 input_file /dev/null
real 0m6.688s
user 0m0.900s
sys 0m5.776s
```

Notice that most execution time is spent executing system calls and as a consequence gcc -O3 is no help.

cp0.c is making 2 system calls for every byte copied - a huge overhead.

If it is modified to buffer its I/O into 8192 byte block it runs 1000x faster.

It now makes 2 system calls for every 8192 bytes copied

```
/home/cg2041/public html/code/performance/cp1.c
// Written by andrewt@cse.unsw.edu.au
// as a COMP2041 lecture example
// copy input to output using read/write system calls
// for every 4096 bytes - efficient but Unix/Linux specific
#include <stdio.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
#include <stdlib.h>
#include <unistd.h>
copy_file_to_file(int in_fd, int out_fd) {
   while (1) {
        char c[8192]:
        int bytes_read = read(in_fd, c, sizeof c);
        if (bytes read < 0) {
            perror("cp: ");
            exit(1):
        if (bytes read <= 0)
        int bytes_written = write(out_fd, c, bytes_read);
        if (bytes_written <= 0) {
            perror("cp: ");
            exit(1):
main(int argc, char *argv[]) {
     if (argc != 3) {
        fprintf(stderr, "cp <src-file> <destination-file>\n");
        return 1:
    int in_fd = open(argv[1], 0_RDONLY);
    if (in fd < 0) {
        fprintf(stderr, "cp: %s: ", argv[1]);
        perror(""):
        return 1;
    int out fd = open(argv[2], 0 WRONLY10 CREATIO TRUNC, S IRWXU):
    if (out_fd <= 0) {
        fprintf(stderr, "cp: %s: ", argv[2]);
        perror("");
        return 1;
   copy_file_to_file(in_fd, out_fd);
    return 0:
```

```
$ time ./cp1 input_file /dev/null
real     0m0.005s
user     0m0.000s
sys     0m0.008s
```

If we use the portable stdio library instead of a low-level copy, even a byte-byte by loop runs suprisingly fast, because stdio buffers the I/O behind the scenes.

```
/home/cx2041/public_html/code/performance/cp2.c
// Written by andrewt@cse.unsw.edu.au
// as a COMP2041 lecture example
// copy input to output using stdio functions
// stdio buffers reads & writes for you - efficient and portable
#include <stdio.h>
#include <stdlib.h>
copy file to file(FILE *in. FILE *out) {
    while (1) {
        int ch = fgetc(in);
        if (ch == EOF)
             break:
        if (fputc(ch. out) == EOF) {
            fprintf(stderr, "cp:"):
            perror("");
            exit(1);
main(int argc. char *argv□) f
    FILE *in. *out:
    if (argc != 3) {
        fprintf(stderr, "cp <src-file> <destination-file>\n"):
        return 1:
    3
    in = fopen(argv[1], "r");
    if (in == NULL) {
        fprintf(stderr, "cp: %s: ", argv[1]);
        perror("");
        return 1;
    out = fopen(argv[2], "w");
    if (out == NULL) {
        fprintf(stderr, "cp: %s: ", argv[2]);
        perror(""):
        return 1:
    copy file to file(in, out);
    return 0:
```

And with a little more complex code we get reasonable speed with portability:

```
/home/cs2041/public_html/code/performance/cp3.c
// Written by andrewt@cse.unsw.edu.au
// as a COMP2041 lecture example
#include <stdio.h>
#include <stdlib.h>
// copy input to output using stdio functions
// stdio buffers reads & writes for you - efficient and portable
copy_file_to_file(FILE *in, FILE *out) {
    char input[4096];
    while (1) {
        if(fgets(input, sizeof input, in) == NULL) {
        if (fprintf(out, "%s", input) == EOF) {
            fprintf(stderr, "cp:");
            perror("");
            exit(1);
main(int argc, char *argv[]) {
   FILE *in, *out;
    if (argc != 3) {
        fprintf(stderr, "cp <src-file> <destination-file>\n");
        return 1;
    in = fopen(argv[1], "r");
    if (in == NULL) {
        fprintf(stderr, "cp: %s: ", argv[1]);
        perror("");
        return 1;
   }
   out = fopen(argv[2], "w");
    if (out == NULL) {
        fprintf(stderr, "cp: %s: ", argv[2]);
        perror("");
        return 1;
   copy file to file(in. out):
    return 0:
```

For comparison Perl code which does a copy via an array of lines:

```
/home/cs2041/public html/code/performance/cp4.pl
#!/usr/bin/perl -w
# Simple cp implementation reading entire file into array
# Written by andrewt@cse.unsw.edu.au for COMP2041
die "Usage: cp <infile> <outfile>\n" if @ARGV != 2;
$infile = shift @ARGV;
$outfile = shift @ARGV;
open IN, '<', $infile or die "Cannot open $infile: $!\n";
open OUT, '>', $outfile or die "Cannot open $outfile: $!\n";
print OUT <IN>;
real 0m0.248s
      0m0.168s
user
       0m0.032s
SYS
```

And Perl code which unsets Perl's line terminator variable so a single read returns the whole file:

```
/home/cs2041/public_html/code/performance/cp5.pl
#!/usr/bin/perl -w
# Simple cp implementation reading entire file into array
# Written by andrewt@cse.unsw.edu.au for COMP2041
die "Usage: cp <infile> <outfile>\n" if @ARGV != 2;
$infile = shift @ARGV;
$outfile = shift @ARGV;
open IN, '<', $infile or die "Cannot open $infile: $!\n";
open OUT, '>', $outfile or die "Cannot open $outfile: $!\n";
undef $/;
print OUT <IN>;
real
        0m0.029s
user 0m0.008s
        0m0.020s
SYS
```

Here is a simple Perl program to calculate the n-th Fibonacci number:

```
/home/cs2041/public_html/code/performance/fib0.pl

#!/usr/bin/perl -w
sub fib($);
printf "fib(%d) = %d\n", $_, fib($_) foreach @ARGV;
sub fib($) {
    my ($n) = @_;
    return 1 if $n < 3;
    return fib($n-1) + fib($n-2);
}
```

It becomes slow near n=35.

```
$ time fib0.pl 35
fib(35) = 9227465
real 0m10.776s
user 0m10.729s
sys 0m0.016s
```

we can rewrite in C.

```
/home/cs2041/public html/code/performance/fib0.c
#include <stdlib.h>
#include <stdio.h>
int fib(int n) {
    if (n < 3) return 1;
    return fib(n-1) + fib(n-2);
}
int main(int argc, char *argv[]) {
    int i;
    for (i = 1; i < argc; i++) {
        int n = atoi(argv[i]);
        printf("fib(%d) = %d\n", n, fib(n));
    }
    return 0;
```

Faster but the program's complexity doesn't change and it becomes become slow near n=45.

```
$ gcc -03 -o fib0 fib0.c

$ time fib0 35

fib(45) = 1134903170

real 0m4.994s

user 0m4.976s

sys 0m0.004s
```

It is very easy to cache already computed results in a Perl hash.

```
/home/cs2041/public_html/code/performance/fib1.pl

#!/usr/bin/perl -w
sub fib($);
printf "fib(%d) = %d\n", $_, fib($_) foreach @ARGV;
sub fib($) {
    my ($n) = @_;
    return 1 if $n < 3;
    return $fib{$n} || ($fib{$n} = fib($n-1) + fib($n-2));
}
```

This changes the program's complexity from exponential to linear.

```
$ time fib1.pl 45
fib(45) = 1134903170
real 0m0.004s
user 0m0.004s
sys 0m0.000s
```

Now for Fibonanci we could also easily change the program to an iterative form which would be linear too.

But memoization is a general technique which can be employed in a variety of situations to improve perfrormnce. There is even a Perl package to support it.

```
/home/cs2041/public_html/code/performance/fib2.pl
#!/usr/bin/perl -w
use Memoize;
sub fib($);
memoize('fib');
printf "fib(%d) = %d\n", $_, fib($_) foreach @ARGV;
sub fib($) {
    my ($n) = @_;
    return 1 if $n < 3;
    return fib($n-1) + fib($n-2);
}</pre>
```