

Assembly: Tracing; conditionals

COSC 208, Introduction to Computer Systems, 2022-03-24

Outline

- Warm-up
- Tracing assembly code
- Conditionals

Warm-up

- Q1: Write the C code equivalent for each line of assembly, treating registers as if they were variable names. For example, the C code equivalent for `sub sp, sp, #0x20` is `sp = sp - 0x20`

```
0000000000400544 <sum>:
400544: d10083ff      sub sp, sp, #0x20    sp = sp - 0x20
400548: b9001fe0      str w0, [sp, #28]    *(sp + 28) = w0
40054c: f9000be1      str x1, [sp, #16]    *(sp + 16) = x1
400550: f9400be8      ldr x8, [sp, #16]    x8 = *(sp + 16)
400554: b9400109      ldr w9, [x8]         w9 = *x8
400558: b9000fe9      str w9, [sp, #12]    *(sp + 12) = w9
40055c: b9401fe9      ldr w9, [sp, #28]    w9 = *(sp + 28)
400560: b9400fea      ldr w10, [sp, #12]   w10 = *(sp + 12)
400564: 0b0a0129      add w9, w9, w10      w9 = w9 + w10
400568: b9000be9      str w9, [sp, #8]     *(sp + 8) = w9
40056c: b9400be0      ldr w0, [sp, #8]     w0 = *(sp + 8)
400570: 910083ff      add sp, sp, #0x20    sp = sp + 0x20
```

Tracing assembly code

- Q2: The assembly corresponds to the following C code. Label each line of assembly code with the line number of the line of C code from which the assembly instruction was derived.

- C code

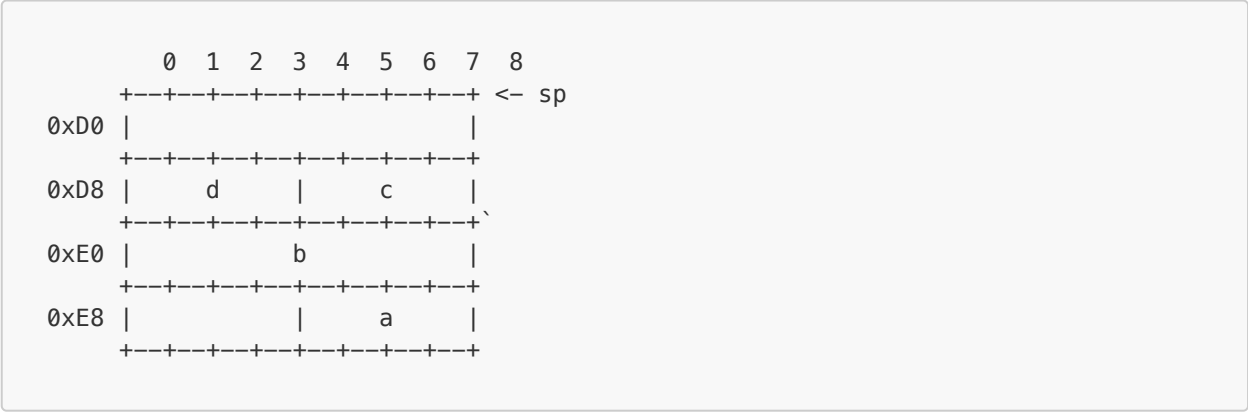
```
1 int sum(int a, int *b) {
2     int c = *b;
3     int d = a + c;
4     return d;
5 }
```

- Assembly code

```
0000000000400544 <sum>:
400544: d10083ff    sub sp, sp, #0x20    // Line 1
400548: b9001fe0    str w0, [sp, #28]    // |
40054c: f9000be1    str x1, [sp, #16]    // V
400550: f9400be8    ldr x8, [sp, #16]    // Line 2
400554: b9400109    ldr w9, [x8]         // |
400558: b9000fe9    str w9, [sp, #12]    // V
40055c: b9401fe9    ldr w9, [sp, #28]    // Line 3
400560: b9400fea    ldr w10, [sp, #12]   // |
400564: 0b0a0129    add w9, w9, w10      // |
400568: b9000be9    str w9, [sp, #8]     // V
40056c: b9400be0    ldr w0, [sp, #8]     // Line 4
400570: 910083ff    add sp, sp, #0x20    // V
```

- Q3: Place in the stack below the parameters *a*, *b* and local variables *c* and *d* (before executing last assembly instruction; and assuming *sp* = 0xF0 initially)

- Stack (before executing last assembly instruction; assume *sp* = 0xF0 initially)



Conditionals

- Q4: The following C code was compiled into assembly. Label each line of assembly code with the line number of the line of C code from which the assembly instruction was derived.

```
1 int divide(int numerator, int denominator) {  
2     int result = -1;  
3     result = numerator / denominator;  
4     return result;  
5 }
```

```
0000000000400544 <divide>:  
400544: d10043ff    sub sp, sp, #0x10      // Line 1  
400548: 12800008    mov w8, #0xffffffff    // Line 2  
40054c: b9000fe0    str w0, [sp, #12]      // Line 1  
400550: b9000be1    str w1, [sp, #8]       // V  
400554: b90007e8    str w8, [sp, #4]       // Line 2  
400558: b9400fe8    ldr w8, [sp, #12]      // Line 3  
40055c: b9400be9    ldr w9, [sp, #8]       // |  
400560: 1ac90d08    sdiv w8, w8, w9        // |  
400564: b90007e8    str w8, [sp, #4]       // V  
400568: b94007e0    ldr w0, [sp, #4]       // Line 4  
40056c: 910043ff    add sp, sp, #0x10      // |  
400570: d65f03c0    ret                    // V
```

- Why is `#0xffffffff` being stored in `w8`? — this is the two's complement representation of -1
- When might this function cause an error? — when denominator is 0

- How would you modify the C code to avoid an error?

```

1 int divide_safe(int numerator, int denominator) {
2     int result = -1;
3     if (denominator != 0) {
4         result = numerator / denominator;
5     }
6     return result;
7 }

```

Conditional assembly code

```

0000000000000076c <divide_safe>:                                // Line
76c:    d10083ff    sub     sp, sp, #0x20    // 1
770:    b9000fe0    str     w0, [sp, #12]   // 1
774:    b9000be1    str     w1, [sp, #8]    // 1
778:    12800000    mov     w0, #0xffffffff // 2
77c:    b9001fe0    str     w0, [sp, #28]   // 2
780:    b9400be0    ldr     w0, [sp, #8]     // 3
784:    7100001f    cmp     w0, #0x0        // 3
788:    540000a0    b.eq    79c <divide_safe+0x30> // 3
78c:    b9400fe1    ldr     w1, [sp, #12]   // 4
790:    b9400be0    ldr     w0, [sp, #8]    // 4
794:    1ac00c20    sdiv    w0, w1, w0      // 4
798:    b9001fe0    str     w0, [sp, #28]   // 4
79c:    b9401fe0    ldr     w0, [sp, #28]   // 6
7a0:    910083ff    add     sp, sp, #0x20   // 6
7a4:    d65f03c0    ret                                // 6

```

- What does the *cmp* instruction do? — compares a register's value to another value
- What does the *b.eq* instruction do? — "jumps" (i.e., branches) to a different instruction when the compared values are equal
- Why does the assembly check if *w0 == 0* when the C code contains *!= 0*? — the C code checks for the condition that must be true to execute the if body, whereas the assembly code checks for the condition that must be true to **skip over** the if body
- How would we express this in C code? — using an if statement and a *goto* statement

```

1 int divide_safe_goto(int numerator, int denominator) {
2     int result = -1;
3     if (denominator == 0)
4         goto after;
5     result = numerator / denominator;
6 after:
7     return result;
8 }

```

More Practice with conditionals

- What happens if the code includes an `else` statement? — if condition is true, execute the if body and skip over the else body; if condition is false, skip over the if body and execute the else body

```
1 int flip(int bit) {
2     int result = -1;
3     if (bit == 0) {
4         result = 1;
5     }
6     else {
7         result = 0;
8     }
9     return result;
10 }
```

- Q3: The above C code was compiled into assembly (using `gcc`). Label each line of assembly code with the line number of the line of C code from which the assembly instruction was derived.

```
0000000000000071c <flip>:                                // Line
71c:    d10083ff    sub     sp, sp, #0x20    // 1
720:    b9000fe0    str     w0, [sp, #12]   // 1
724:    12800000    mov     w0, #0xffffffff // 2
728:    b9001fe0    str     w0, [sp, #28]   // 2
72c:    b9400fe0    ldr     w0, [sp, #12]   // 3
730:    7100001f    cmp     w0, #0x0        // 3
734:    54000081    b.ne    744 <flip+0x28> // 3
738:    52800020    mov     w0, #0x1        // 4
73c:    b9001fe0    str     w0, [sp, #28]   // 4
740:    14000002    b       748 <flip+0x2c> // 5
744:    b9001fff    str     wzr, [sp, #28]  // 7
748:    b9401fe0    ldr     w0, [sp, #28]   // 9
74c:    910083ff    add     sp, sp, #0x20   // 9
750:    d65f03c0    ret                                // 9
```

- Q4: Write a function called `flip_goto` that behaves the same as `flip` but matches the structure of the assembly code that will be generated for `flip`. (Hint: you'll need two `goto` statements.)

```
int flip_goto(int bit) {
    int result = -1;
    if (bit != 0)
        goto else_body;
    result = 1;
    goto after_else;
else_body:
    result = 0;
after_else:
    return result;
}
```

Extra practice

- QA: Write a function called `adjust_goto` that behaves the same as `adjust` but matches the structure of the assembly code that will be generated for `adjust`. (Hint: you'll need two `goto` statements.)

```
int adjust(int value) {
    if (value < 10) {
        value = value * 10;
    }
    else {
        value = value / 10;
    }
    return value;
}
```

```
int adjust_goto(int value) {
    if (value >= 10)
        goto else_body;
    value = value * 10;
    goto after_if;
else_body:
    value = value / 10;
after_if:
    return value;
}
```