Efficiency: memory hierarchy; locality

COSC 208, Introduction to Computer Systems, 2021-10-25

Outline

- Memory hierarchy
- Locality
- Optimizing for locality

No warm-up — Happy Monday!

Memory hierarchy

- Compares various forms of storage in terms of
    - Access latency
    - Capacity
    - Cost
    - Volatility
- Access latency
    - Let's consider a 1hz CPU, which means 1 cycle = 1 second
    - Registers — 1 cycle = 1 second
    - Caches — ~10 cycles = ~10 seconds
    - Main memory — ~100 cycles = ~2 minutes
    - Solid-state drive — ~1 million cycles = ~11.5 days
    - Hard (i.e., traditional) disk drive — ~10 million cycles = ~115 days
    - Remote (i.e., network) storage — ~20ms = ~2 years
- Storage capacity
    - Let's assume 1 byte = 1mL
    - Registers — 30 * 8B = ~250mL = ~1 cup
    - Caches (Core i7 in MacBook Pro)
        - L1 — 32KB + 32KB = 64L = ~1 tank of gas
        - L2 — 512KB * 4 cores = 2048L = ~7 bathtubs
    - Main memory = 32GB (in MacBook Pro) = ~13 olympic swimming pools
    - SSD = 1TB (in MacBook Pro) = ~Lake Moraine
- Cost
    - 2 x 16GB DRAM = ~$100 = $3.12 per GB
    - 1TB SSD = $80 = $0.08 per GB
    - 2TB HDD = $60 = $0.03 per GB
- Volatility
    - Primary storage (registers, caches, and main memory) — volatile (i.e., data is lost if power is lost)
    - Secondary storage (SSD, HDD, network storage) — non-volatile (i.e., data is preserved if power is lost)

Data movement

- Recall: *How does data move between the CPU, main memory, and secondary storage in the von Neumann Architecture?* — bus
- *Why does data move to/from secondary storage?* — data stored in primary storage is lost when a machine looses power
- *Why does data move between registers and main memory?* — not enough room in registers to store all values used by a program at runtime
- *How can we move less data?*
    - Make better use of registers — i.e., eliminate unnecessary loads/stores
    - Add additional memory to the CPU — i.e., a cache
- *How do we decide what/when to move data between the registers, cache, and CPU?* — based on locality

Temporal vs. spatial locality

- *What is temporal locality?*
  - Access the same data repeatedly
  - E.g., for loop variable
- *What is spatial locality?*
  - Access data with a similar scope
  - E.g., next item in array
  - E.g., local variables/parameters, which are stored in the same stack frame

Optimizing assembly code for locality

- Q1: *Cross-out redundant loads and stores from the assembly code*

```
000000000000088c <interest_due>:
    88c:    sub sp, sp, #0x20
    890:    str w0, [sp, #12]
    894:    str w1, [sp, #8]
    898:    ldr w0, [sp, #12]
    89c:    ldr w1, [sp, #8]
    8a0:    mul w0, w1, w0
    8a4:    str w0, [sp, #20]
    8a8:    mov w0, #0x4b0
    8ac:    str w0, [sp, #24]
    8b0:    ldr w1, [sp, #20]
    8b4:    ldr w0, [sp, #24]
    8b8:    sdiv    w0, w1, w0
    8bc:    str w0, [sp, #28]
    8c0:    ldr w0, [sp, #28]
    8c4:    add sp, sp, #0x20
    8c8:    ret

00000000000008cc <make_payment>:
    8cc:    stp x29, x30, [sp, #-48]!
    8d0:    mov x29, sp
    8d4:    str w0, [sp, #28]
    8d8:    str w1, [sp, #24]
    8dc:    str w2, [sp, #20]
    8e0:    ldr w1, [sp, #20]
    8e4:    ldr w0, [sp, #28]
    8e8:    bl  88c <interest_due>
    8ec:    str w0, [sp, #40]
    8f0:    ldr w1, [sp, #24]
    8f4:    ldr w0, [sp, #40]
    8f8:    sub w0, w1, w0
    8fc:    str w0, [sp, #44]
    900:    ldr w1, [sp, #44]
    904:    ldr w0, [sp, #28]
    908:    cmp w1, w0
    90c:    b.le    918 <make_payment+0x4c>
    910:    str wzr, [sp, #28]
    914:    b   928 <make_payment+0x5c>
    918:    ldr w1, [sp, #28]
    91c:    ldr w0, [sp, #44]
    920:    sub w0, w1, w0
    924:    str w0, [sp, #28]
    928:    ldr w0, [sp, #28]
    92c:    ldp x29, x30, [sp], #48
    930:    ret
```