# Multiprocessing: scheduling

COSC 208, Introduction to Computer Systems, 2022-04-26

#### **Outline**

- · Scheduling processes
- · First In First Out (FIFO) scheduling
- Shortest Job First (SJF) scheduling
- Preemption
- Shortest Time-to-Completion First (STCF) scheduling
- Round Robin (RR) scheduling

### Scheduling processes

- OS decides which process to run and for how long
- What factors should the OS consider when making these decisions?
  - Time for process to complete
  - Overhead of context switching
  - Fairness
  - User interaction
- For now, consider one of these metrics: turnaround time
  - $\circ$   $T_{turnaround} = T_{complete} T_{arrive}$
- For now, assume a process starts and runs to completion—i.e., no I/O and no preemption

## First In First Out (FIFO) scheduling

Process	Arrival time	Duration
А	0	15
В	5	15
С	10	15

•	What is the average	turnaround	time for the	processes i	usina FIFO?

• What happens if A's duration is 60?		
• What happens if C's duration is 60?		

• How can we change the schedule so the average turnaround time when A's duration is 60 is more like the average turnaround time when C's duration is 60?

## Shortest Job First (SJF) scheduling

_	Process	Arrival time	Duration
	Α	0	60
	В	0	15
	С	0	15

• What is the average turnaround time for the above processes using SJF?

Process	Arrival time	Duration
Α	0	60
В	5	15
С	10	15

• What is the average turnaround time for the above processes using SJF?

• We're back to FIFO—What happened!? How can we fix this?

### Preemption

- OS only regains control when a system call occurs—e.g., read/write file, yield
  - o Syscalls may occur infrequently, or never, due to program design, bugs, or malicious behavior
- How does an OS forcibly regain control?—set a timer that raises an interrupt
  - Interrupt causes a trap instruction to be executed
  - Interrupts can also be raised by devices—e.g., Network Interface Card (NIC)
- What must the OS do if it decides to run another process?—perform a context switch
  - Save the machine state associated with the process that was running—in particular, the contents of all registers are saved in the process's control structure
  - Restore the machine state associated with the process that should run—again, the contents of all registers are loaded from the process's control structure

## Shortest Time-to-Completion First (STCF) scheduling

- · Allow preemption
- If a process arrives that has less computation remaining than the currently running process, then preempt the current process and run the new process
- Also known as Preemptive Shortest Job First (PSJF)

Process	Arrival time	Duration		
Α	0	60		
В	5	15		
С	10	15		

Now consider re	sponse time: T <sub>respon</sub>	<sub>se</sub> = T <sub>first_run</sub> - T <sub>arrive</sub>		
What is the ave	age response time fo	r the same processes	s using STCF scheduling	?
Now consider w	ait time: $T_{wait} = \sum (T_s)$	tart_run - T <sub>become_read</sub>	<sub>dy</sub> )	

What major assumption have we made thus far that is impractical in a real system?

## Round Robin (RR)

• Let each process run for a small amount of time, then switch to the next process; when you get to the last process, then start again with the first process and repeat

Process	Arrival time	Duration
А	Just before 0	60
В	Just before 5	15
С	Just before 10	15

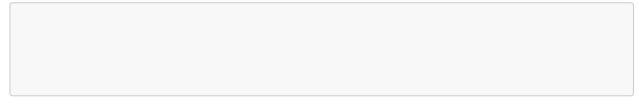
• What is the average turnaround time and response time for the above processes assuming we let a process run for 5 seconds before switching processes?



- Average Turnaround =
- Average Response =
- Average Wait =
- In practice, there is a queue of processes that are in the ready state, resulting in the following schedule:

	Α	В	A	С	В	A	С	В	A	С		
0		5	10	15	20	25 3	30	35	40	45	50	90

• Determine the schedule for the above process with a time quantum of 10.



- Average Turnaround =
- Average Response =
- Average Wait =
- What happens to average response time as we increase the time quantum?