Multiprocessing: processes; fork

COSC 208, Introduction to Computer Systems, 2021-11-08

Announcements

• Project 2 Part B due tomorrow at 11pm

Outline

- Warm-up
- · System calls
- Creating processes: actual code & fork

Warm-up:

Q1: True or False

- 1. Code stored on secondary storage (e.g., a solid state drive) is called a process -- false; it is a program
- 2. Each process has its own code, heap, stack, and register values -- true
- 3. The CPU is in user mode when executing application code, and kernel mode when executing OS code -- true
- 4. A process can directly execute instructions on the CPU -- true; user mode uses Limited Restricted Execution (OS isn't an intermediary until privileged instructions)
- 5. A process can directly access input and output ports -- false; indirect, need a system call

System calls

- Invoked via a special assembly instruction: trap (generic) or SVC (on ARM)
 - Example program

```
#include <stdio.h>
#include <unistd.h>
int user() {
   int uid = getuid();
   return uid;
}
int main() {
   int u = user();
   printf("User %d is running this process\n", u);
}
```

Assembly code

```
efourquet@ulna:~$ cd demo/
efourquet@ulna:~/demo$ clang -o process.o -static process.c
efourquet@ulna:~/demo$ objdump -d process.o >process.s
efourquet@ulna:~/demo$ strace -i ./process
echo $UID
```

SVC

- Functions in the C standard library that involve a privileged operation (e.g., printf) put the system call number in a register and invoke a trap instruction programmer doesn't have to worry about these details; they can just call the appropriate function in the C standard library
- svc makes system calls; 0xae 174 call number; go to kernel, save registers
- When svc is executed
 - 1. CPU saves registers to the kernel stack kernel stack is at a fixed location in memory
 - Why do we need to save the registers? so we can return to user when __getuid is done
 - 2. CPU switches to kernel mode
 - 3. CPU uses system call number to index into table of trap handlers
 - At boot, initialize table of trap handlers with pointers into OS code for handling each type of syscall
 - 4. Branch and link to trap handler code
 - 5. CPU restores registers from the kernel stack
 - 6. CPU switches to user mode
 - 7. Resume execution after svc
- What should we do if a process tries to perform a privileged operation without making a system call?
 - Let the code keep running code may assume privileged operation was successful
 - Kill the process

Creating processes

- int fork()
 - Creates an exact copy of the running process, except for the return value from fork return ∅ to child (i.e., new) process; return child's process ID to parent process (i.e., process that called fork)
 - Both child and parent resume execution from place where fork was called
- Q2: What does the following code output?

```
int main(int argc, char **argv) {
    printf("Before fork\n");
    int pid = fork();
    printf("After fork\n");
    return 0;
}
```

```
Before fork
After fork
After fork
```

• Q3: What does the following code output (assuming the new process has PID 1819)?

```
int main(int argc, char **argv) {
    printf("Before fork");
    int pid = fork();
    if (pid == 0) {
        printf("Child gets %d\n", pid);
    } else {
        printf("Parent gets %d\n", pid);
    }
    return 0;
}
```

```
Before fork
Child gets 0
Parent gets 1819
```

OR

```
Before fork
Parent gets 1819
Child gets 0
```

• Creating an exponentially increasing number of processes (known as a fork bomb)

```
int main() {
    while(1) {
       fork();
    }
```

}	