

# Computer Science Club

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## 0. INTRODUCTION

- a. Lots of clubs for specific languages
- b. Need to focus on solving problems
- c. Need more options to have healthy community
- d. Practice delivering complex discourse
- e. Video
  - i. Incentive for better presentation
  - ii. Accessible to more people
  - iii. Demonstrates competence of our local developers

## 1. MOTIVATION

- a. Why this problem
  - i. Fascinated me
  - ii. Didn't study it in depth before
- b. General introduction to how ingrained regular expressions are in languages
- c. Illustrate examples based on string metaphor
  - i. Constructors (/Colin/)
  - ii. Operators (=~)

## 2. BACKGROUND

- a. FSM
  - i. State diagram
    - 1. Vending machine diagram (machine metaphor)
    - 2. States
      - a. Start
      - b. Acceptor
    - 3. Transition function
  - ii. Idea of a language / alphabet we are using (coins, letters)
  - iii. Concepts of Finite and State

b. Finite Automata

- i. John Conway (game of life, etc.)
- ii. Simple example /colin/
  - 1. Break it up under concatenation
  - 2. Options with disjunctive union
  - 3. Add a cycle
  - 4. Add an epsilon
  - 5. Kleene star
    - a. Regular expression coined by Stephen Kleene
    - b. Regular expressions as in regular sets

c. General Properties

- i. DFA / NFA dichotomy
  - 1. Determined state = Deterministic Finite Automata (DFA)
  - 2. Indeterminate state = Non-deterministic Finite Automata (NFA)
  - 3. Complexity dictates determinism
    - a. Number of acceptor states
      - i. Single = DFA
      - ii. Multiple = NFA
    - b. Disjunctive unions
    - c. Epsilon
- ii. DFA and NFA are equivalent
- iii. Interests we need tonight
  - 1. Disjunctive union
  - 2. Concatenation
  - 3. Kleene star
- iv. Other properties that hold (more on this later)
  - 1. Conjunctive union
  - 2. Complement

3. REGULAR EXPRESSIONS

- a. An aside about things we are not talking about
  - i. Context free grammars
  - ii. Turing machines (Alan Turing)
  - iii. Recursive structures
  - iv. Pushdown automata
    - 1. Stacks
    - 2. Memory

- b. Definition: R is a regular expression if R is
- i.  $\emptyset$  (the null pattern)
  - ii.  $\epsilon$  (the null character)
  - iii. a (for some a in an alphabet)
  - iv. Concatenation  $R_1 R_2$  (where  $R_1$  and  $R_2$  are regular expressions)
  - v. Disjunctive union  $R_1 | R_2$  (where  $R_1$  and  $R_2$  are regular expressions)
  - vi. Kleene star  $R^*$  (where R is a regular expression)
  - vii. Notes
    1.  $R^+$  is simply  $RR^*$
    2. Generalized Regular Expressions
      - a. Conjunction union
      - b. Complement
- c. Brzozowski's derivatives (Janusz Brzozowski)
- i. Preliminaries
    1. Prescriptive generator
      - a. Recursive
      - b. Generates NFA
    2. Derivative idea is notional
  - ii. Define a function  $\delta$  such that
    1.  $\delta(\emptyset) = \emptyset$  (zero)
    2.  $\delta(\epsilon) = \epsilon$  (unit)
    3.  $\delta(c) = \emptyset$  (variable)
    4.  $\delta(R_1 R_2) = \delta(R_1) \delta(R_2)$  (multiplication)
    5.  $\delta(R_1 | R_2) = \delta(R_1) | \delta(R_2)$  (addition)
    6.  $\delta(R^*) = \epsilon$
  - iii. Example
    1.  $D_f(\text{foo|far})^* = (\text{oo|ar})(\text{foo|far})^*$
  - iv. Define derivative with respect to a character c recursively
    1.  $D_c(\emptyset) = \emptyset$
    2.  $D_c(\epsilon) = \emptyset$
    3.  $D_c(c) = \epsilon$
    4.  $D_c(c') = \emptyset$  if  $c \neq c'$
    5.  $D_c(R_1 R_2) = \delta(R_1) D_c(R_2) | D_c(R_1) R_2$  (think multiplication)
    6.  $D_c(R_1 | R_2) = D_c(R_1) | D_c(R_2)$  (think addition)
    7.  $D_c(R^*) = D_c(R) R^*$
  - v. Observations
    1. End states
      - a.  $\emptyset$  (no match)
      - b.  $\epsilon$  (has match)
    2. All other relations are recursive

#### 4. PROGRAMMING

##### a. OCaml benefits

###### i. Language

1. Originally for mathematical theorem proving
2. Polymorphic disjunctive union (sum) typing
3. Constructor matching
4. Hindley-Milner type checking
5. Modules and Functors
6. Object Orientation

###### ii. Execution

1. Bytecode compiler with forward and reverse step debugging
2. Fast native compiler
3. Various REPLs

###### iii. In the wild

1. Financial institutions (Jane Street Capital)
2. .Net CLR (F#)
3. Compilers
  - a. FFTW
  - b. Coq (More advanced formal proof language)

##### b. Walk through regex code