

## Chapter 2-[Logic](#)

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3:20 PM

### **Logik:**

$\wedge$

is shorthand for "and" (equivalent to the AND Operator)

A **Proposition** is a statement which is true or false (but never both!). (It can't be a question. It also can't contain variables, and a claim needs to be stated in the proposition) *Not as useful in Proofing.*

$\vee$

Is shorthand for "or" (equivalent to the OR Operator)

$\neg p$

Is shorthand for "Not p" (equivalent to the NOT Operator)

$p \rightarrow q$

Is shorthand for "If p then q"

$p \leftrightarrow q$

Is shorthand for "if p then q, and the converse is also true"

$\neg q \rightarrow \neg p$

Is shorthand for the contrapositive of if p then q, and **is** equivalent to if p then q.

Ordering in Mathematical Shorthand: apply the "not" operators first, then the "and" and "or". Then you take the results and do the implication operations. (similar to high school algebra rules)

Logical Equivalence:

A frequently useful fact is that  $p \rightarrow q$  is logically equivalent to  $\neg p \vee q$ .

De Morgan's Law:

$\text{not (A and B)} \rightarrow \text{not A or not B}$

$\text{not (A or B)} \rightarrow \text{not A and not B}$

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$\text{not (A and B)} \neq \text{not A and not B}$

T and F are special constant propositions with no variables that are, respectively, always true and always false.

$$p \wedge \neg p \equiv F$$

The equal operator = can only be applied to **objects** such as numbers. When comparing **logical expressions** that return true/false values, you must use  $\equiv$ .

Distributive Rules:

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

(you can distribute either operator over the other)

The Conditional to OR/AND Rule:

$$p \rightarrow q \equiv \neg p \vee q$$

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

(Recall this is true because in  $p \rightarrow q$ , when  $p$  is false, the statement is always true, and when  $p$  is true, the statement is true only when  $q$  is true)

A **Predicate** is a statement that becomes true or false if you substitute in values for its variables.

Quantifiers:

There are only 3, and the first 2 are the most common:

$\forall$

*For All* (universal quantifier),

$\exists$

*There Exists* (existential quantifier),

$\exists!$

*There exists a unique ...* (unique existence, used when there is only 1 object that matches some requirements)

Examples:

$$\forall x \in \mathbb{R}, x^2 + 3 \geq 0$$

$$\exists y \in \mathbb{R}, y = \sqrt{2}$$

$$\exists! x \in \mathbb{R}, x^2 = 0$$

"Such that" is sometimes abbreviated "s.t."

*(Mathematicians are lazy)*

Contrapositives also exists for statements with quantifiers.

For example:

$$\forall x, \text{ if } p(x), \text{ then } q(x)$$

And its contrapositive:

$$\forall x, \text{ if } \neg q(x), \text{ then } \neg p(x)$$

The quantifier stays the same: only transform the if/then part.

Representing 2D points:

Either refer to 2 single variables and pair them up, like so:

$$\forall x, y \in \mathbb{Z}$$

And refer to them as a pair later on, or

Treat the pair as a single variable, like so:

$$\exists (x, y) \in \mathbb{R}^2, x^2 + y^2 = 1$$

Or be even more abstract:

$$\exists p \in \mathbb{R}^2, p \text{ is on the unit circle}$$

(and then later define what "p" and "the unit circle" means)

### **Predicate Logic "De Morgan's Laws":**

$$\neg(\forall x, P(x)) \equiv \exists x, \neg P(x)$$

"for all x in A, P(x)" is false when (is equivalent) to "there is some value x in A such that P(x) is false"

$$\neg(\exists x, P(x)) \equiv \forall x, \neg P(x)$$

"there exists x in A such that P(x) is true" is false when (is equivalent to) "for all x in A, P(x) is false"

( $P(x)$  is a statement like  $x > 0$  or whatever, and  $A$  can be some set like real numbers or whatever)

Example:

$$\forall x, P(x) \rightarrow (Q(x) \wedge R(x))$$

Its **negation** is:

$$\begin{aligned}\neg(\forall x, P(x) \rightarrow (Q(x) \wedge R(x))) &\equiv \exists x, \neg(P(x) \rightarrow (Q(x) \wedge R(x))) \\ &\equiv \exists x, P(x) \wedge \neg(Q(x) \wedge R(x)) \\ &\equiv \exists x, P(x) \wedge (\neg Q(x) \vee \neg R(x))\end{aligned}$$

### **Bounding and Scope:**

A quantifier **binds** to a variable, but if a variable hasn't been bound by a quantifier, (or otherwise given a value or a set of replacement values), it is called a **Free Variable**.

**Free variables** don't have a defined truth value, so they **cannot** be a step in a proof.

The "bound" variable in a quantification is only defined for a limited time, called the **Scope** of the binding.

Qualifiers define the scope of the variable. For example, the  $i$  in

$$\sum_{i=0}^n \frac{1}{i}$$

is only defined while you are still inside the summation.