

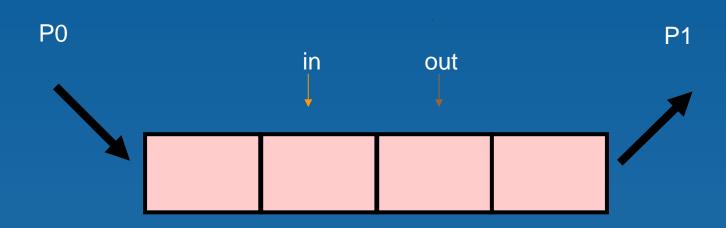
## 经典同步问题 ( Classical Problem s )

#### 经典同步问题

- ◆生产者 消费者问题
  Producer-Consumer Problem
- ◆读者 写者问题 Readers and Writers Problem
- ◆哲学家就餐问题
  Dining-Philosophers Problem

#### 生产者-消费者问题

- ◆ 生产者进程"生产"出信息,存储在缓冲区, 供消费者进程"消费"
  - ∞ unbounded-buffer 缓冲区的容量无限
  - ∞ bounded-buffer 缓冲区的容量有限



#### 生产者-消费者问题的一种解决方案

◆假设一个拥有 N个 items 的缓冲区

- ◆定义信号量 mutex ,初始值为 1
- ◆定义信号量 full , 初始值为 0
- ◆定义信号量 empty , 初始值为 N

#### 生产者-消费者问题的一种解决万案 (续)

◆生产者进程 producer while (true) { // produce an item wait (empty); wait (mutex); // add the item to the buffer signal (mutex); signal (full);

# 生产者-消费者问题的一种解决万案(续)

◆消费者进程 consumer while (true) { wait (full); wait (mutex); // remove an item from buffer signal (mutex); signal (empty); // consume the removed item

### 死锁 (Deadlock) vs 饥饿 (Starvation)

◆如果 wait 操作排序不当,会怎样?

```
Producer
                                 Consumer
wait (empty);
                                  wait (mutex);
wait (mutex);
                                  wait (full);
signal (mutex);
                                  signal (mutex);
signal (full);
                                  signal (empty);
```

#### 死锁 vs 饥饿

- ◆死锁 (Deadlock) two or more processes are waiting indefinitely for an event that can be caused by only one of the waiting processes
- ◆饥饿 (Starvation) indefinite blocking. A process may never be removed from the semaphore queue in which it is suspended
- ◆关于死锁的概念,会在后续单元详述

## 读者 - 写者 (Readers-Writers) 问题

- ◆场景:一组并发进程共享一个数据集,其中
  - ☞读者 (Readers) 进程 只对数据集实施 "读"操作,从不试图更新数据集
  - ™写者 (Writers) 进程 对数据集既实施 "读"操作,也实施"写"操作
- ◆问题

允许若干读者进程同时读取数据集;在任何时刻,只允许单个写者进程访问数据集。

◆如何保证?

#### 第一类读者 - 写者问题 (First Readers-Writers)

- No readers will be kept waiting unless a writer has already obtained permission to use the shared object.
- ◆读者进程不会被要求等待,除非已经有一个 个写者进程获准访问共享数据集
- In other words, no reader should wait for other readers to finish simply because a writer is waiting.

### 第一类读者-写者问题的一个解决方案

- ◆共享数据集
- ◆定义
  - ☆信号量 mutex ,初始值为 1
  - ☆信号量 wrt ,初始值为 1
  - 輕整型变量 readcount ,初始值为 0

### 第一类读者-写者问题的一个解决方案

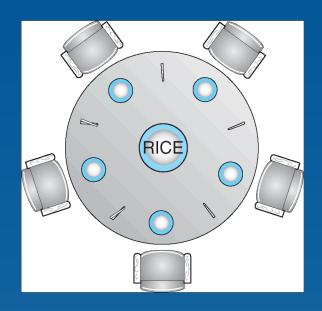
◆写者进程 while (true) { wait (wrt); // writing is performed signal (wrt);

## 第一类读者-写者问题的一个解决方案

◆ 读者进程

```
while (true) {
     wait (mutex);
     readcount ++;
     if (readercount == 1) wait (wrt);
     signal (mutex)
          // reading is performed
      wait (mutex);
      readcount --;
      if (redacount == 0) signal (wrt);
      signal (mutex);
```

## 哲学家就餐 (Dining-Philosophers) 问题



- ◆ 共享数据集: 碗里的食物
- 定义信号量数组 chopstick [5] , 初始值为 1

### 哲学家就餐 (Dining-Philosophers) 问题

◆ 哲学家进程 *i* While (true) { wait ( chopstick[i] ); wait (chopStick[(i + 1) % 5]); // eat signal (chopstick[i]); signal (chopstick[ (i + 1) % 5]); // think 此方案有缺陷: 死锁!

