

IE 521 Convex Optimization

Lecture 22: Constrained Subgradient Methods

Niao He

1st May 2019

Outline

Constrained
Convex Problems

Constrained Convex Problems

Constrained
Subgradient
Method

Constrained Subgradient Method

Algorithm
Convergence
Proof

Algorithm

Convergence

Proof

Constrained Level
Method

Constrained Level Method

Lower Bound
Complexity

Lower Bound Complexity

Recap: Simple Constrained Convex Problems

$$\begin{array}{ll} \min & f(x) \\ \text{s.t.} & x \in X \end{array}$$

Subgradient Method

$$x_{t+1} = \Pi_X(x_t - \gamma_t g_t), \quad g_t \in \partial f(x_t)$$

Problem Class	Stepsize	Convergence
Convex Lipschitz	$O(\frac{1}{\sqrt{t}})$	$O(\frac{D_X M_f}{\sqrt{t}})$
Strongly Convex Lipschitz	$O(\frac{1}{\mu t})$	$O(\frac{M_f^2}{\mu t})$

Bundle Methods

- ▶ Kelley method
- ▶ Level-set method

General Constrained Convex Problems

We will focus on the general convex problem:

$$\begin{aligned} \min_{x \in X} \quad & f_0(x) \\ \text{s.t.} \quad & f_i(x) \leq 0, i = 1, \dots, m \end{aligned}$$

Assumptions

- ▶ X is simple and admits easy-to-compute projections
- ▶ First-order oracles for $f_0(x), f_i(x)$ are available

Note $f_0(x), f_i(x)$ are not necessarily differentiable or smooth

Problem Reformulation

The problem can be rewritten as

$$\begin{array}{ll} \min_{x \in X} & f_0(x) \\ \text{s.t.} & f(x) \leq 0 \end{array} \quad (\text{P})$$

where

$$f(x) = \max_{1 \leq i \leq m} f_i(x).$$

Remark. We can easily compute the subgradient for $f(x)$.

Problem Reformulation

The problem can be rewritten as

$$\begin{array}{ll} \min_{x \in X} & f_0(x) \\ \text{s.t.} & f(x) \leq 0 \end{array} \quad (P)$$

where

$$f(x) = \max_{1 \leq i \leq m} f_i(x).$$

Remark. We can easily compute the subgradient for $f(x)$.

So how to solve (P) given access to $f'_0(f)$, $f'(x)$?

Constrained Subgradient Method

0. Initialize $x_1 \in X$

1. For $t \geq 1$, compute

$$g_t = \begin{cases} f'_0(x_t), & \text{if } f(x_t) < \gamma_t \|f'(x_t)\|_2 \\ f'(x_t), & \text{if } f(x_t) \geq \gamma_t \|f'(x_t)\|_2 \end{cases}$$

$$x_{t+1} = \Pi_X(x_t - \gamma_t \frac{g_t}{\|g_t\|_2})$$

Convergence of Constrained Subgradient Method

IE 521 Convex
Optimization

Niao He

Constrained
Convex Problems

Constrained
Subgradient
Method

Algorithm
Convergence
Proof

Constrained Level
Method

Lower Bound
Complexity

Theorem. Let f_0 and f be Lipschitz continuous with constant $M_{f_0} > 0$ and $M_f > 0$. Let D_X be the diameter of the set X such that $\forall x, y \in X, \|x - y\|_2 \leq D_X$. Set the stepsize $\gamma_t = \frac{D_X}{\sqrt{t+0.5}}$. We have, for any $T \geq 3$,

$$\min_{1 \leq t \leq T} f_0(x_t) - f_0^* \leq \frac{\sqrt{3}D_X M_{f_0}}{\sqrt{T-1.5}}$$

$$\min_{1 \leq t \leq T} f(x_t) \leq \frac{\sqrt{3}D_X M_f}{\sqrt{T-1.5}}$$

Proof of Convergence

- First, we have

$$\|x_{t+1} - x^*\|_2^2 \leq \|x_t - x^*\|_2^2 - 2\gamma_t \underbrace{\frac{g_t^T(x_t - x^*)}{\|g_t\|_2^2}}_{\nu_t} + \gamma_t^2$$

- Define

$$I = \{t \in [\frac{T}{3}, T] : g_t = f'_0(x_t)\}$$

$$I^c = \{t \in [\frac{T}{3}, T] : g_t = f'(x_t)\}$$

- For $t \in I^c$, $\nu_t \geq \gamma_t$. This is because

$$f'(x_t)^T(x_t - x^*) \geq f(x_t) - f(x^*) \geq f(x_t) \geq \gamma_t \|f'(x_t)\|_2$$

Proof of Convergence (cont'd)

► **Claim:** There exists $t^* \in I$, such that $\nu_{t^*} < \gamma_{t^*}$.

► Otherwise, suppose $\nu_t \geq \gamma_t, \forall t \in I$. Then

$$\|x_{t+1} - x^*\|_2^2 \leq \|x_t - x^*\|_2^2 - \gamma_t^2, \forall t \in [\frac{T}{3}, T].$$

► This implies that

$$\sum_{t=\lfloor T/3 \rfloor}^T \gamma_t^2 \leq D_X^2.$$

► On the other hand,

$$\sum_{t=\lfloor T/3 \rfloor} \gamma_t^2 \geq D_X^2 \int_{T/3}^{T-1} \frac{1}{t-0.5} dt = D_X^2 \ln 3 > D_X^2.$$

Proof of Convergence (cont'd)

- Since $\nu_{t^*} < \gamma_{t^*} \leq \gamma_{\lfloor T/3 \rfloor}$, we have

$$f_0(x_{t^*}) - f_0^* \leq f_0'(x_{t^*})^T (x_{t^*} - x^*) \leq \nu_{t^*} \|g_{t^*}\|_2 \leq \frac{\sqrt{3}D_X M_{f_0}}{\sqrt{T - 1.5}}.$$

- Since $t^* \in I$, we have

$$f(x_{t^*}) \leq \gamma_{t^*} \|f'(x_{t^*})\|_2 \leq \frac{\sqrt{3}D_X M_f}{\sqrt{T - 1.5}}$$

Constrained Convex Problem

Constrained Convex Problem:

$$\begin{aligned} \min_{x \in X} \quad & f_0(x) \\ \text{s.t.} \quad & f(x) \leq 0 \end{aligned} \tag{P}$$

with optimal value being $\alpha^* := \text{Opt}(P)$.

Parametric Function: Define

$$\Phi(x; \alpha) = \max\{f_0(x) - \alpha, f(x)\}$$

$$\phi(\alpha) = \min_{x \in X} \Phi(x; \alpha)$$

Note that $\phi(\alpha)$ is a one-dimensional function.

Properties

Lemma. The following results hold:

1. $\phi(\alpha^*) = 0$
2. $\phi(\alpha) \leq 0, \forall \alpha \geq \alpha^*$.
3. $\phi(\alpha) > 0, \forall \alpha < \alpha^*$
4. $\phi(\alpha) - \beta \leq \phi(\alpha + \beta) \leq \phi(\alpha), \forall \beta \geq 0$.

Remark.

- ▶ $\phi(\alpha)$ is non-increasing, 1-Lipschitz continuous, and convex.
- ▶ The smallest root of function $\phi(\alpha)$ corresponds to the optimal value of (P) .
- ▶ Procedures for finding the root: Bisection, Newton's method, etc.

Generic Two-stage Scheme

0. Initialize α_1

1. For $i = 1, \dots, N$

► (Approximately) solve the subproblem:

$$\min_{x \in X} \left\{ \Phi(x; \alpha_i) := \max\{f_0(x) - \alpha_i, f(x)\} \right\}$$

through some subgradient routines (e.g., subgradient method, bundle methods)

► Update $\alpha_i \rightarrow \alpha_{i+1}$ through some route-finding procedure (e.g., approximate Newton step)

Remark. Total number of complexity is expected to be

$$O\left(\frac{1}{\epsilon^2} \log\left(\frac{1}{\epsilon}\right)\right).$$

Constrained Level Method

Model:

$$f_{0,t}(x) = \max_{1 \leq i \leq t} \left\{ f_0(x_i) + f'_0(x_i)^T (x - x_i) \right\}$$

$$f_t(x) = \max_{1 \leq i \leq t} \left\{ f(x_i) + f'(x_i)^T (x - x_i) \right\}$$

$$\Phi_t(x; \alpha) = \max \{ f_{0,t}(x) - \alpha, f_t(x) \}$$

$$\phi_t(\alpha) = \min_{x \in X} \Phi_t(x; \alpha)$$

Model's Smallest Root:

$$\alpha_t^* := \min \{ \alpha : \phi_t(\alpha) = 0 \} = \min_{x \in X} \{ f_{0,t}(x) : f_t(x) \leq 0 \}$$

Remark.

- ▶ $\phi_1(\alpha) \leq \phi_2(\alpha) \leq \cdots \leq \phi(\alpha)$
- ▶ $\alpha_1^* \leq \alpha_2^* \leq \cdots \leq \alpha^*$

Constrained Level Method

Constrained
Convex Problems

Constrained
Subgradient
Method

Algorithm
Convergence
Proof

Constrained Level
Method

Lower Bound
Complexity

0. Initialize $\alpha_1 < \alpha^*, \kappa \in (0, 1/2)$

1. For $i = 1, \dots, N$

- ▶ Run Level-set Method to solve $\min_{x \in X} \Phi(x; \alpha_i)$ and generate $\{x_\tau\}_{\tau=1}^t \in X$ until

$$\underbrace{\min_{x \in X} \Phi_t(x; \alpha_i)}_{\text{model's best value}} \geq (1 - \kappa) \underbrace{\min_{1 \leq \tau \leq t} \Phi_t(x_\tau; \alpha_i)}_{\text{model's record value}}$$

- ▶ Set $t(i) = t$
- ▶ Terminate if $\min_{1 \leq \tau \leq t} \Phi_t(x_\tau; \alpha_i) \leq \epsilon$
- ▶ Otherwise, set $\alpha_{i+1} = \alpha_{t(i)}^*$

Complexity of Constrained Level Method

Constrained
Convex ProblemsConstrained
Subgradient
MethodAlgorithm
Convergence
ProofConstrained Level
MethodLower Bound
Complexity

- ▶ Number of master steps is at most

$$N(\epsilon) \leq \frac{1}{\ln(2 - 2\kappa)} \ln \frac{t^* - t_0}{(1 - \kappa)\epsilon}$$

- ▶ Number of Level Method iterations at each step is at most

$$t(i) \leq \frac{D_X^2 M_f^2}{\kappa^2 \epsilon^2 (1 - \alpha)^2 \alpha (2 - \alpha)}$$

Lower Bound Complexity

Theorem. (Nemirovski & Yudin, 1979) For any $1 \leq t \leq n$, $x_1 \in \mathbb{R}^n$, there exists a M -Lipschitz continuous convex function f and a convex set X with diameter D_X , such that for any first-order algorithm that generates:

$$x_{t+1} \in x_1 + \text{span}(g_1, \dots, g_t)$$

where $g_i \in \partial f(x_i)$, we always have

$$\min_{1 \leq s \leq t} f(x_s) - f_* \geq \frac{D_X M_f}{(1 + \sqrt{t})}$$

Lower Bound Complexity

Theorem. (Nemirovski & Yudin, 1979) For any $1 \leq t \leq n$, $x_1 \in \mathbb{R}^n$, there exists a μ -strongly convex, M -Lipschitz continuous function f and a convex set X , for any first-order method that generates

$$x_{t+1} \in x_1 + \text{span}(g_1, \dots, g_t)$$

where $g_i \in \partial f(x_i)$, we always have

$$\min_{1 \leq s \leq t} f(x_s) - f_* \geq \frac{M^2}{8\mu t}$$

Worse-case Example

Consider the problem $\min_{x \in X} f(x)$ where

$$f(x) = C \cdot \max_{1 \leq i \leq t} x_i + \frac{\mu}{2} \|x\|_2^2,$$

$$X = \{x \in \mathbb{R}^n : \|x\|_2 \leq R := D_X/2\}$$

- The subdifferential set of function f :

$$\partial f(x) = \mu x + C \cdot \text{Conv}\{e_i : i \text{ that } x_i = \max_{1 \leq j \leq t} x_j\}$$

- **Oracle example:** Given an input x , it returns $g = C \cdot e_i + \mu x$, with $i = \min\{i : x_i = \max_{1 \leq j \leq t} x_j\}$.

- The optimal solution and optimal value:

$$x_{*i} = \begin{cases} -\frac{C}{\mu t} & 1 \leq i \leq t \\ 0 & t < i \leq n \end{cases} \quad \text{and} \quad f_* = -\frac{C^2}{2\mu t}.$$

Worst-case Example

- ▶ W.l.o.g., set $x_1 = 0$.
- ▶ By induction, we can show that $x_t \in \text{span}(e_1, \dots, e_{t-1})$.
- ▶ This implies for $1 \leq s \leq t$, $f(x_s) \geq 0$.

$$\min_{1 \leq s \leq t} f(x_s) - f_* \geq \frac{C^2}{2\mu t}.$$

- ▶ If $C = \frac{M\sqrt{t}}{1+\sqrt{t}}$, $\mu = \frac{M}{R(1+\sqrt{t})}$, then $f(x)$ is M -Lipschitz continuous. Moreover,

$$\min_{1 \leq s \leq t} f(x_s) - f_* \geq \frac{C^2}{2\mu t} = \frac{MD_X}{4(1+\sqrt{t})}$$

- ▶ If $C = \frac{M}{2}$, $\mu = \frac{M}{2R}$, then $f(x)$ is M -Lipschitz continuous and μ -strongly convex. Moreover,

$$\min_{1 \leq s \leq t} f(x_s) - f_* \geq \frac{C^2}{2\mu t} = \frac{M^2}{8\mu t}$$

References

Constrained
Convex Problems

Constrained
Subgradient
Method

Algorithm
Convergence
Proof

Constrained Level
Method

Lower Bound
Complexity

- ▶ Nesterov (2004), Chapter 3.2-3.3