

Group: Friday 9:30 AM with Sophia

Names: Miriam Brody, kambek karshe, Collins Kariuki, thomas mcconnell

1. On average the ratio is 8
 - a. $\theta(n^3)$
 - i. Work: $O(2n) = 2$, $O((2n)^2) = 4n^2$, $O((2n)^3) = 8n^3$
2. Ratio:
 - a. $\theta(1)$
 - i. We consider 8000 to be an anomaly, outside of 800 everything is constant
3. Ratio:
 - a. $\theta(n)$
 - i. A3 doubles thus it is linear
4. Ratio:
 - a. $\theta(\log(n))$
 - i. Increases slowly by about 100 each time
5. Ratio:
 - a. $\theta(n^2)$
 - i. The ratio is four, so a similar logic to 1. To how we got n^3 when the ratio was 8 we get n^2 when the ratio is 4
6. Ratio:
 - a. $\theta(n \log n)$
 - i. While the runtimes are roughly doubling (the n part), we are also adding a constant number to the doublings (the $\log n$ part).

Recurrences



$$T(n) = 2T(n/3) + d$$

$$T(n) = 7T(n/7) + n$$

if $f(n) = O(n^{\log_b a - \epsilon})$ for $\epsilon > 0$, then $T(n) = \Theta(n^{\log_b a})$

if $f(n) = \Theta(n^{\log_b a})$, then $T(n) = \Theta(n^{\log_b a} \log n)$

if $f(n) = \Omega(n^{\log_b a + \epsilon})$ for $\epsilon > 0$ and $af(n/b) \leq cf(n)$ for $c < 1$
then $T(n) = \Theta(f(n))$

$$T(n) = T(n-1) + \log n$$

$$T(n) = 8T(n/2) + n^3$$

$$T(n) = 8T\left(\frac{n}{2}\right) + n^3$$

$$\begin{aligned} &O(n^{3-\epsilon}) \\ &\Theta(n^3) \\ &\Omega(n^{3-\epsilon}) \end{aligned}$$

$$\begin{aligned} a &= 8 & n^{\log_b a} &= n^{\log_2 8} = \underline{\underline{n^3}} \\ b &= 2 \end{aligned}$$

$$f(n) = n^3$$

$$T(n) = \Theta(n^{\log_b a} \log n)$$

$$\text{is } n^3 = O(n^{3-\epsilon})$$

Ch

$$= \Theta(n^3 \log n)$$

$$k) = O(n \log(n)) \quad k \leq n$$

$$T(n) = T(n-1) + \log n \quad \text{Cost}$$

$$c \log n \quad \rightarrow c \log$$

$$c \log(n-1) \quad \rightarrow$$

$$c \log(n-2) \quad \rightarrow$$

$$c \log(n-(n-1)) \quad \rightarrow$$

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Verreys

$$\text{Total cost} = c \log(n) + c \log(n-1) + c \log(n-2) + \dots + 0$$

$$\sum_{i=1}^n c \log(i)$$

For $T(n) = T(n-1) + \log(n)$, we were not able to finish it up before the hour mark but we made an educated guess that the Theta complexity is $n \log(n)$ because for the total cost, despite the n decreasing we are adding up a significant amount of $\log(n)$'s multiple times, n times so we guessed that the recurrence should be big theta $n \log(n)$. We did not have enough time to validate our solution with the substitution method.

3. Everyone was at the group. Everyone felt comfortable to share and was included.