Group: Friday 9:30 AM with Sophia

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- 1. On average the ratio is 8
  - a.  $\theta(n^3)$ 
    - i. Work: O(2n) = 2,  $O((2n)^2) = 4n^2$ ,  $O((2n)^3) = 8n^3$
- 2. Ratio:
  - a.  $\theta(1)$ 
    - We consider 8000 to be an anomaly, outside of 800 everything is constant
- 3. Ratio:
  - a.  $\theta(n)$ 
    - A3 doubles thus it is linear
- Ratio:
  - a.  $\theta(\log(n))$ 
    - Increases slowly by about 100 each time
- 5. Ratio:
  - a.  $\theta(n^2)$ 
    - i. The ratio is four, so a similar logic to 1. To how we got n^3 when the ratio was 8 we get n^2 when the ratio is 4
- 6. Ratio:
  - a.  $\theta(nlogn)$ 
    - i. While the runtimes are roughly doubling (the n part), we are also adding a constant number to the doublings (the log n part).

## Recurrences



$$T(n) = 2T(n/3) + d$$
  $T(n) = 7T(n/7) + n$ 

if 
$$f(n) = O(n^{\log_b a - \varepsilon})$$
 for  $\varepsilon > 0$ , then  $T(n) = \Theta(n^{\log_b a})$   
if  $f(n) = \Theta(n^{\log_b a})$  then  $T(n) = \Theta(n^{\log_b a} \log_b a)$ 

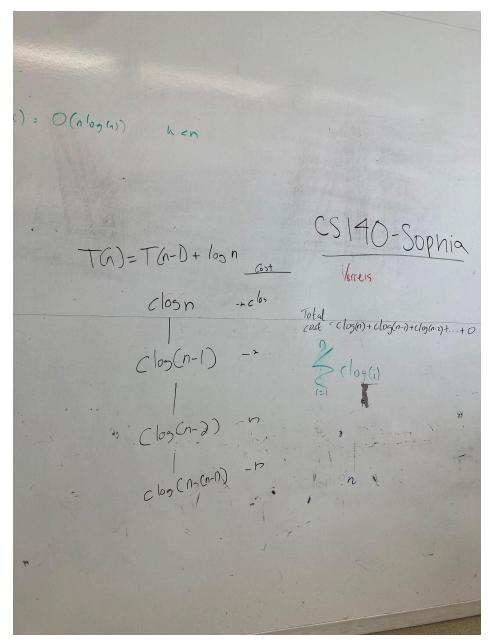
if 
$$f(n) = \Theta(n^{\log_b a})$$
, then  $T(n) = \Theta(n^{\log_b a} \log n)$ 

if 
$$f(n) = \Omega(n^{\log_b a + \varepsilon})$$
 for  $\varepsilon > 0$  and  $af(n/b) \le cf(n)$  for  $c < 1$   
then  $T(n) = \Theta(f(n))$ 

$$T(n) = T(n-1) + \log n$$
  $T(n) = 8T(n/2) + n^3$ 

$$T(n) = 8T(\frac{n}{2}) + n^{2}$$

$$O(n^{3}, 0)$$



For  $T(n) = T(n-1) + \log(n)$ , we were not able to finish it up before the hour mark but we made an educated guess that the Theta complexity is  $n\log(n)$  because for the total cost, despite the n decreasing we are adding up a significant amount of  $\log(n)$ 's multiple times, n times so we guessed that the recurrence should be big theta  $n\log(n)$ . We did not have enough time to validate our solution with the substitution method.

3. Everyone was at the group. Everyone felt comfortable to share and was included.