

Lectures Notes on Knowledge Representation and Processing

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2020

These notes were originally prepared for my CS course at University Erlangen-Nuremberg (FAU) given with Michael Kohlhase in Summer 2020. They are directed at 3rd semester CS undergraduates and master students but should be intelligible even for earlier students and could be interesting also for PhD students and for students from adjacent majors. The course is recommended both as a first course in the specialization area Artificial Intelligence as well as a one-off overview on knowledge representation.

The course was developed in Summer 2020 from scratch and materials were built along the way. It integrated current directions and recent results in research on knowledge representation pulling together materials in an entirely new and original way.

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Part I

Introduction

Chapter 1

Meta-Remarks

Important stuff that you should read carefully!

State of these notes I constantly work on my lecture notes. Therefore, keep in mind that:

- I am developing these notes in parallel with the lecture — they can grow or change throughout the semester.
- These notes are neither a subset nor a superset of the material discussed in the lecture. On the one hand, they may contain more details than mentioned in the lectures. On the other hand, important material such as background, diagrams, and examples may be part of the lecture but not mentioned in these notes.
- Unless mentioned otherwise, all material in these notes is exam-relevant (in addition to all material discussed in the lectures).

Collaboration on these notes I am writing these notes using LaTeX and storing them in a git repository on GitHub at <https://github.com/florian-rabe/Teaching>. As an experiment in teaching, I am inviting all of you to collaborate on these lecture notes with me. This would require familiarity with LaTeX as well as Git and GitHub — that is not part of this lecture, but it is an essential skill for you. Ask in the lecture if you have difficulty figuring it out on your own.

By forking and by submitting pull requests for this repository, you can suggest changes to these notes. For example, you are encouraged to:

- Fix typos and other errors.
- Add examples and diagrams that I develop on the board during lectures.
- Add solutions for the homeworks if I did not provide any (of course, I will only integrate solutions after the deadline).
- Add additional examples, exercises, or explanations that you came up or found in other sources. If you use material from other sources (e.g., by copying an diagram from some website), make sure that you have the license to use it and that you acknowledge sources appropriately!

I will review and approve or reject the changes. If you make substantial contributions, I will list you as a contributor (i.e., something you can put in your CV).

Any improvement you make will not only help your fellow students, it will also increase your own understanding of the material. Make sure your git commits carry a user name that I can connect to you.)

Other Advice I maintain a list of useful advice for students at https://github.com/florian-rabe/Teaching/blob/master/general/advice_for_students.pdf. It is mostly targeted at older students who work in individual projects with me (e.g., students who work on their BSc thesis). But much of it is useful for you already now or will become useful soon. So have a look.

Chapter 2

Fundamental Concepts

2.1 Abbreviations

knowledge representation and processing	KRP	the general area of this course
knowledge representation language	KRL	a languages used in KRP
knowledge representation tool	KRT	a tool implementing a KPL and processing algorithms for it

2.2 Motivation

2.2.1 Knowledge

Human knowledge pervades all sciences including computer science, mathematics, natural sciences and engineering. That is not surprising: “science” is derived from the Latin word “scire” meaning “to know”. Similarly, philosophy, from which all sciences derive, is named after the Greek words “philo” meaning loving and “sophia” meaning wisdom, and the for common ending “-logy” is derived from Greek “logos” meaning word (i.e., a representation of knowledge).

In regards to knowledge, computer science is special in two ways: Firstly, many branches of computer science need to understand KRP as a prerequisite for teaching computers to do knowledge-based tasks. In some sense, KRP is the foundation and ultimate goal of all artificial intelligence.¹ Secondly, modern information technology enables all sciences to apply computer-based KRP in order to vastly expand on the domain-specific tasks that can be automated. Currently all sciences are becoming more and more computerized, but most non-CS scientists (and many computer scientists for that matter) lack a systematic education and understanding of IT-KRP. That often leads to bad solutions when domain experts cannot see which KRP solutions are applicable or how to apply them.

2.2.2 Representation and Processing

It is no coincidence that this course uses the phrase “Representation and Processing”. In fact, this is an instance of a universal duality. Consider the following table of analogous concept pairs, which could be extended with many more examples:

Representation	Processing
Static	Dynamic
Situation	Change
Be	Become
Data Structures	Algorithms
Set	Function
State	Transition
Space	Time

¹Indeed, a major problem with the currently very successful machine learning-based AI technology is that it remains unclear when and how it does KRP. That can be dangerous because it leads to AI systems recommending decisions without being able to explain why that decision should be trusted.

Again and again, we distinguish a static concept that describes/represents what is a situation/state is and a dynamic concept that describes how it changes. If that change is a computer doing something with or acting on that representation, we speak of “processing”.

It is particular illuminating to contrast KRP to the standard CS course on Data Structures and Algorithms (DA).² Generally speaking, DA teaches the methods, and KRP teaches how to apply them. Data structures are a critical prerequisite for representing knowledge. But data structures alone do not capture what the data means (i.e., the knowledge) or if a particular representation makes any sense. Similarly, algorithms are the critical prerequisite for processing knowledge. But while algorithms can be systematically analyzed for efficiency, it is much harder to analyze if an algorithm processes knowledge correctly. The latter requires understanding what the input and output data means.

Capturing knowledge in computers is much harder than developing data structures and algorithms. It is ultimately the same challenge as figuring out if a computer system is working correctly — a problem that is well-known to be undecidable in general and very difficult in each individual case.

2.3 Components of Knowledge

2.3.1 Syntax and Semantics, Data and Knowledge

Four concepts are of particular relevance to understanding knowledge. They form a 2×2 -quadruple of concepts:

Syntax	Data
Semantics	Knowledge

All four concepts are primitive, i.e., they cannot be defined in simpler terms. All sciences have few carefully-chosen primitive on which everything builds. This is done most systematically in mathematics (where primitives include set or function). While mathematical primitives as well as some primitives in physics or CS are specified formally, the above four concepts can only be described informally, ultimately appealing to pre-existing human understanding. Moreover, this description is not standardized — different courses may use very different descriptions even they ultimately try to capture the same elusive ideas.

Data (in the narrow sense of computer science) is any object that can be stored in a computer, typically combined with the ability to input/output, transfer, and change the object. This includes bits, strings, numbers, files, etc.

Data by itself is useless because we would have no idea what to do with it. For example, the object $O = ((49.5739143, 11.0264941), "2020 - 04 - 21T16 : 15 : 00CEST")$ is useless data without additional information about its syntax and semantics. Similarly, a file is useless data unless we know which file format it uses.

Syntax is a system of rules that describes which data is **well-formed**. For O above the syntax could be “a pair of (a pair of two IEEE double precision floating point numbers) and a string encoding of an time stamp”. For a file, the syntax is often indicated by the file name extension, e.g., the syntax of an `html` file is given in Section 12 of the current HTML standard³.

Syntax alone is useless unless we know what the semantics, i.e., what the data means and thus how to correctly interpret and process the data. For example, the syntax of O allows to check that O is well-formed, i.e., indeed contains two numbers and a timestamp string. That allows rejecting ill-formed data such as $((49.5739143, 11.0264941), "foo")$. The HTML syntax allows us to check that a file conforms to the standard.

Semantics is a system of rules that determines the meaning of well-formed data. For example, ISO 8601 specifies that timestamp string refer to a particular date and time in a particular time zone. Further semantics for O might be implicit in the algorithms that produce and consume it: such as “the first component of the pair contains two numbers between 0 and 180 resp. 0 and 360 indicating latitude resp. longitude of a location on earth”. Semantics might be multi-staged, and further semantics about O might be that O indicates the location and time of the first lecture of this course. Similarly, Section 14 of the HTML standard specifies the semantics of well-formed HTML files by describing how they are to be rendered in a web browser.

Knowledge is the combining of some data with its syntax and semantics. That allows applying the semantics to obtain the meaning of the data (if syntactically well-formed and signaling an error otherwise). In computer systems,

²The course is typically called “Algorithms and Data Structures”, but that is arguably awkward because algorithms can exist if there are data structure to work with. Compare my notes on that course in this repository, where I emphasize data structures much more than is commonly done in that course.

³<https://html.spec.whatwg.org/multipage/>

- data is represented using primitive data (ultimately the bits provided by the hardware) and encodings of more complex data (bytes, arrays, strings, etc.) in terms of simpler ones,
- syntax is theoretically specified using grammars and practically implemented in programming languages using data structures,
- semantics is represented using algorithms that process syntactically well-formed data,
- knowledge is elusive and often emerges from executing the semantics, e.g., rendering of an HTML file.

2.3.2 Semantics as Syntax Transformation

In order to capture knowledge better in computer systems, we often use two syntax levels: one to represent the data itself and another to represent the knowledge. These can be seen as input and output data. In that case, semantics is a function that translates from the data syntax to the knowledge syntax, and knowledge is the pair of the data and the result of applying the semantics. The following table gives some examples.

Data syntax	Semantics function	Knowledge syntax
SPARQL query	evaluation	result set
SQL query	evaluation	result table
program	compiler	binary code
program expression	interpreter	result value
logical formula	interpretation in a model	mathematical object
HTML document	rendering	graphical representation

Thus, the role of syntax vs. semantics may depend on the context: just like one function's output can be another function's input, one interpretation's knowledge can be another one's syntax. For example, we can first compile a program into binary and then execute it to return its value.

Such hierarchies of evaluation levels are very common in computer systems. In fact, most state-of-the-art compilers are subdivided into multiple phases each further interpreting the output of the previous one. Thus, if knowledge is represented in computers, it is invariably data itself but relative to a different syntax.

2.3.3 Heterogeneity of Semantics and Knowledge

While it is easy to design languages to represent data in general, it is very difficult to designing KRLs that capture the human-level quality of knowledge. Over the last few decades, the KRP area in computer science has diversified into different subareas that approach this research problem in fundamentally different ways. In fact, KRP in the very general sense of this course is usually not even studied by itself — instead the subareas are so different, specialized, and large that they all sustain their respective university courses and research conferences.

This is related to the fact the data naturally comes in fundamentally different forms such as graphs, arrays, tables in the sense of relational databases, programs in a programming language, logical formulas, or natural language texts. We speak of **heterogeneous** data. These different forms of data are supported by highly specialized KPTs: graph databases, array databases, relational databases, package databases for programming languages, theorem databases for logics (e.g., the Isabelle Archive of Formal Proofs), databases of research papers (such as the arXiv), and so on. All of these are very successful for their respective kind of data. And all of them include specifications of semantics and KP algorithms that implement this semantics. But it can vary massively how the semantics is specified and implemented. This has caused major practical problems for tool interoperability: many projects require data in multiple formats and algorithms from multiple tools. But the respective tools are often islands that assume that all data is represented in the tool's language and users do not use outside tools. Therefore, the import/export capabilities of the tools are often limited.

Moreover, transporting data across systems is usually ignorant of the semantics: while each tool takes relatively good care to implement the semantics correctly, there is much less certainty that the semantics is preserved when exchanging data across tools. For a trivial example, consider a tool that measures length in inches vs. a tool that uses centimeters, both using floating point numbers for the data: if they exchange the data, i.e., just the numbers, they may mis-communicate the semantics.⁴

This problem is not easy to fix though. The heterogeneity of data and semantics is so extreme that it is, in some cases, an open theoretical problem how knowledge can be shared at all across tools. The basic idea — exchange the

⁴Problems like this have been involved in major disasters such as the Mars Climate Orbiter.

data in a way that preserves semantics — can be difficult to implement if both tools use entirely different paradigms to specify semantics.

2.4 The Tetrapod Model of Knowledge

The Tetrapod model of knowledge is an ongoing research project by the instructors of this course. A first publication was made in [CFKR20]. The structure of this course will draw heavily on the Tetrapod model to get an overview of the different approaches to KPR and their interoperability problems.

2.4.1 Five Aspects of Knowledge

The Tetrapod model distinguishes five basic **aspects** of knowledge and KPR as described below. For each aspect, there is a variety dedicated KRLs supported by highly optimized KPTs as indicated in the following table:

Aspect	KRLs (examples)	KPTs (examples)
ontologization	ontology languages (OWL), description logics (ALC)	reasoners, SPARQL engines (Virtuoso)
concretization	relational databases (SQL, JSON)	databases (MySQL, MongoDB)
computation	programming languages (C)	interpreters, compilers (gcc)
deduction	logics (HOL)	theorem provers (Isabelle)
narration	document languages (HTML, LaTeX)	editors, viewers

Ontologization focuses on developing and curating a coherent and comprehensive ontology of concepts. This focuses on identifying the central concepts in a domain and their relations. For example, a medical ontology would define concepts for every symptom, disease, and medication and then define relations for which symptoms and medications are related to which disease.

Ontologies typically abstract from the knowledge: they standardize identifiers for the concepts and spell out some properties and relations but do not try to capture all details of the knowledge. Well-designed ontologies can capture exactly that different KPTs must share and can thus serve as interoperability layers between them.

While organization can use ontology languages such as OWL or RDF, the inherent complexity of formal objects in computer science and mathematics usually requires going beyond general purpose ontology languages (similar to how the programming languages underlying computer algebra systems usually go beyond general purpose programming languages).

Concretization uses languages based on numbers, strings, lists, and records to obtain concrete representations of datasets in order to store and query their properties efficiently. Because concrete objects are so simple and widely used, it is possible and common to build concrete datasets on top of general purpose data representation languages and tools such as JSON or SQL.

Computation uses specification and programming languages to represent algorithmic knowledge.

Deduction uses logics and theorem provers to obtain verifiable correctness.

Narration uses natural language to obtain texts are easy to understand for humans. Because narrative languages are not well-standardized (apart from general purpose languages such as free text or \LaTeX), it is common to develop narrative libraries on top of ad-hoc languages that impose some formal structure on top of informal text, such as a fixed tree structure whose leafs are free text or a particular set of \LaTeX macros that must be used. Narrative libraries can be classified based on whether entries are derived from publications (e.g., one abstract per paper in zbMATH) or mathematical concepts (e.g., one page per concept in $n\text{Lab}$).

2.4.2 Relations between the Aspects

The aspects can be visualized as the corners of tetrahedron with ontologization in the center and edges and faces representing solutions that mix two or three aspects as seen in Figure ??.

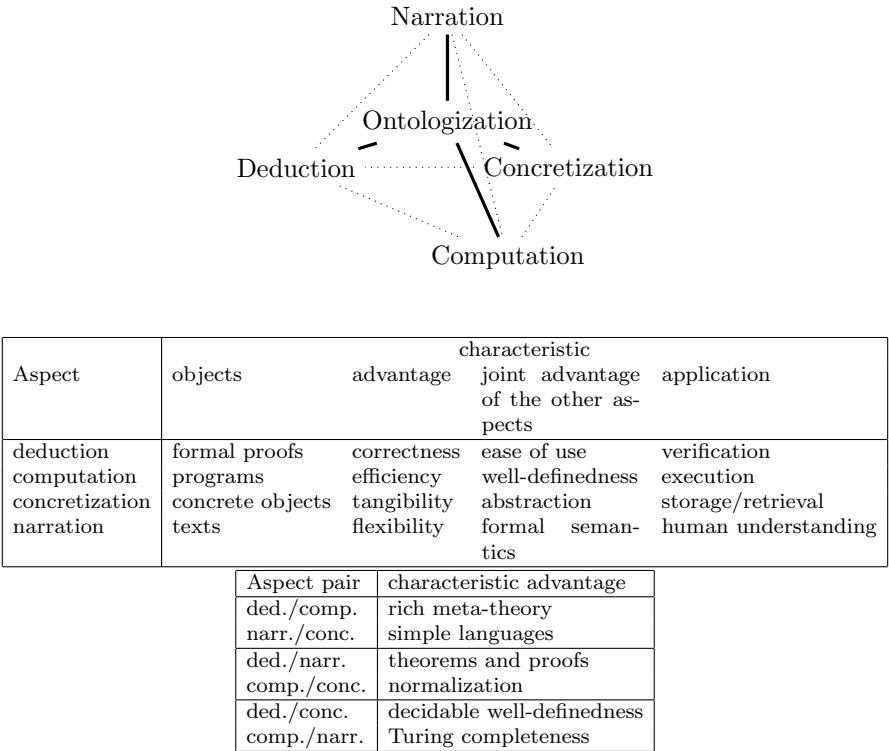


Figure 2.1: Shared properties and advantages of aspects

Most approaches try to incorporate all or multiple aspects. But all languages and tools tend to be heavily biased towards and optimized for a single one of the four corner aspects. This is not due to ignorance but because each aspect provides characteristic advantages that are extremely hard to capture at once. In fact, every combination of aspects shares characteristic advantages and disadvantages as sketched in Figure 2.1. For example, deductive and narrative definitions of a function involved well-definedness arguments, and a function defined by a concrete table is trivially well-defined, but a computational definition of a function may throw exceptions when running; but only the latter can store and compute functions efficiently. Consequently, dedicated and mostly disjoint communities have evolved that have produced large aspect-specific datasets.

Chapter 3

Overview of This Course

3.1 Structure

The subsequent *parts* of this course follow the Tetrapod model with one part per aspect. Each of these will describe the concepts, languages, and tools of the respective aspect as well as their relation to other aspects.

The aspects of the Tetrapod are typically handled in individual courses, which describe highly specialized languages and tools in depth. On the contrary, the overall goal of this course will be seeing all of them as different approaches to semantics and knowledge representation. The course will focus on universal principles and their commonalities and differences as well as their advantages and disadvantages.

The subsequent *chapters* of this first part will be dedicated to aspect-independent material. These will not necessarily be taught in the order in which they appear in these notes. Instead, some of them will be discussed in connection to how they are relevant in individual aspects.

3.2 Exercises and Running Example

Typical practical projects, e.g., the ones that a strong CS graduate might be put in charge of, involve heterogeneous data and knowledge that must be managed using a variety of optimized aspect-specific languages and tools. Interoperability between these is often a major source of inefficiency and bugs.

The exercises accompanying the course will mimic this situation: they will be designed around a single large project that requires choosing and integrating methods, languages, and tools from all aspects.

Concretely, this project will be the development of a univis-like system for a university. It will involve heterogeneous data such as course and program descriptions, legal texts, websites, grade tables, and transcript generation code.

Over the course of the semester students will implement a completely functional system applying the lessons of the course. This is very unusual and often impossible for other courses: as any university course must teach many different things from a wide area, it is rarely possible to find a project that requires many and only lessons from a single course. Here KRP is special because its material pervades all aspects of system development.

Chapter 4

Representing Syntax and Semantics

4.1 Context-Free Grammars

4.2 Inductive Data Types

4.3 Semantics as a Recursive Function

4.4 Context-Sensitive Syntax

Chapter 5

Encoding Data

5.1 Data Representation Languages

5.2 Typed Data

5.3 Encoding Typed Data in Untyped Representation Languages

Part II

Ontological Knowledge

Chapter 6

Ontologies

6.1 General Principles

Motivation An ontology is an abstract representation of the main concepts in some domain. Here *domain* refers to any area of the real world such as mathematics, biology, diseases and medications, human relationships, etc. Many examples can be found at <https://bioportal.bioontology.org/>, including the Gene ontology one of the biggest.

Contrary to the other four aspects, ontological knowledge representations do not aim at capturing the entire semantics of the domain objects. Instead, they focus on defining unique identifiers for the those objects and describing some of their properties and relations to each other.

We use the word **ontologization** to refer to the process of organizing the knowledge of a domain in ontologies.

Ontologies are most valuable when they are *standardized* (either sanctioned through a formal body or a quasi-standard because everyone uses it). A standard ontology allows everybody in the domain to use the identifiers defined by the ontology in a way that avoids misunderstandings. Thus, in the simplest form, an ontology can be seen as a dictionary defining the technical terms of a domain. For example, the Gene ontology defines identifier G0:0000001 to have the formal name "mitochondrion inheritance" and the informal definition "The distribution of mitochondria, including the mitochondrial genome, into daughter cells after mitosis or meiosis, mediated by interactions between mitochondria and the cytoskeleton."

Ontology Languages An ontology is written in **ontology language**. Common ontology languages are

- description logics such as ALC,
- the W3C ontology language OWL, which is the standard ontology languages of the semantic web,
- the entity-relationship model, which focuses on modeling rather than formal syntax,
- modeling languages like UML, which is the main ontology language used in software engineering.

Ontology languages are not committed to a particular domain — in the Tetrapod model, they correspond to programming languages and logics, which are similarly uncommitted. Instead, an ontology language is a formal language that standardizes the syntax of how ontologies can be written as well as their semantics.

Ontologies The details of the syntax vary between ontology languages. But as a general rule, every **ontology** declares

- **individuals** — concrete objects that exist in the real world, e.g., "Florian Rabe" or "WuV"
- **concepts** — abstract groups of individuals, e.g., "instructor" or "course"
- **relations** — binary relations between two individuals, e.g., "teaches"
- **properties** — binary relations between an individuals and a concrete value (a number, a date, etc.), e.g., "has-credits"
- **concept assertions** — the statement that a particular individual is an instance of a particular concept
- **relation assertions** — the statement that a particular relation holds about two individuals
- **property assertions** — the statement that a particular individual has a particular value for a particular property
- **axioms** — statements about relations between concepts, typically in the form subconcept of statements like

"instructor" \sqsubseteq "person"

All assertions can be understood and spoken as subject-predicate-object **triples** as follows:

Assertion	Triple		
	Subject	Predicate	Object
concept assertion	"Florian Rabe"	is-a	"instructor"
relation assertion	"Florian Rabe"	"teaches"	"WuV"
property assertion	"Florian Rabe"	"has credits"	7.5

This uses a special relation **is-a** between individuals and concepts. Some languages group **is-a** with the other binary relations between individuals for simplicity although it is technically a little different.

The possible values of properties must be fixed by the ontology language. Typically, it includes at least standard types such as integers, floating point numbers, and strings. But arbitrary extensions are possible such as dates, RGB-colors, lists, etc. In advanced languages, it is possible that the ontology even introduces its own basic types and values.

Ontologies are often divided into two parts:

- The **abstract** part contains everything that holds in general independent of which individuals: concepts, relations, properties, and axioms. It describes the general rules how the worlds works without committing to a particular set of inhabitants of the world. This part is commonly called the **TBox** (T for terminological).
- The **concrete** part contains everything that depends on the choice of individuals: individuals and assertions. It populates the world with inhabitants. This part is commonly called the **ABox** (A for assertional).

Synonyms Because these principles pervade all formal languages, many competing synonyms are used in different domains. Common synonyms are:

Here	OWL	Description logics	ER model	UML	semantics via logics
individual	instance	individual	entity	object, instance	constant
concept	class	concept	entity-type	class	unary predicate
relation	object property	role	role	association	binary predicate
property	data property	(not common)	attribute	field of base type	binary predicate

In particular, the individual-concept relation occurs everywhere and is known under many names:

domain	individual	concept
type theory, logic	constant, term	type
set theory	element	set
database	row	table
philosophy ¹	object	property
grammar	proper noun	common noun

6.2 A Basic Ontology Language

We could study practical ontology languages like ALC or OWL now. But those feature a lot of other details that can block the view onto the essential parts. Therefore, we first define a basic ontology language ourselves in order to have full control over the details.

6.2.1 Syntax

Definition 6.1 (Syntax of BOL). A BOL-ontology is given by the grammar in Fig. 6.1. It is well-formed if

- no identifier is declared twice,
- every property assertion assigns a value of the type required by the property declaration,
- every reference to an atomic individual/concept/relation/property is declared as such.

The above grammar exhibits some general structure that we find throughout formal KR languages. In particular, an ontology consists of **named declarations** of four different kinds as well as some axioms. Each declaration

clarifies which kind it is (in our case by starting with a keyword) and introduces a new identifier. For each kind, there are complex expression. These are anonymous and built inductively; their base cases are references to the corresponding identifiers. Sometimes (in our case: individuals and properties), the references are the only expressions of the kind. Sometimes (in our case: concepts and relations), there can be many productions for complex expressions. The complex expressions are used to build axioms; in our case, these are the three kinds of assertions and other formulas.

6.2.2 Semantics

We give the semantics of BOL as an example of a semantics by translation. We fix one language that we have already understood and define an interpretation function that maps all complex expression of the syntax into the semantic language.

For simple ontology languages like BOL, ALC, OWL, etc., it is common to use first-order logic (FOL) as the semantic language. More specifically, we use SFOL, the typed variant of FOL with

Definition 6.2 (Semantic of BOL). Every BOL-ontology O is interpreted as a FOL-theory $\llbracket O \rrbracket$ according to Fig. 6.2 where we assume that $\llbracket O \rrbracket$ contains

- a type ι (for individuals),
- additional types and constants corresponding to base types and values of BOL.

Like with the syntax, we can observe some general principles. Every BOL-declaration is translated to a FOL declaration for the same name, and ontologies are translated declaration-wise. For every kind of complex expression, there is one inductive function mapping BOL-expressions to FOL-expressions. The base cases of references to declared identifiers are translated to themselves, i.e., to the identifiers of the same name declared in the FOL theory.

Ontologies	
$O ::= D^*$	
Declarations	
$D ::=$ individual ID concept ID relation ID property ID : T I is-a C I R I I P V F	atomic individual atomic concept atomic relation atomic property concept assertion relation assertion property assertion other axioms
Formulas	
$F ::=$ $C \equiv C$ $C \sqsubseteq C$	concept equality concept subsumption
Individual expressions	
$C ::=$ ID	atomic individuals
Concept expressions	
$C ::=$ ID $C \sqcup C$ $C \sqcap C$ $\text{dom}R$ $\text{rng}R$	atomic concepts union of concepts intersection of concepts domain of a relation range of a relation
Relation expressions	
$R ::=$ ID $R; R$ $R \cup R$ $R \cap R$ R^* Δ_C	atomic relations composition of relations union of relations intersection of relations transitive closure of a relation identity relation of a concept
Property expressions	
$P ::=$ ID	atomic properties
Identifiers	
ID $::=$ alphanumeric string	
Basic types and values	
$T ::=$ int float bool string	types
$T ::=$ (omitted)	values

Figure 6.1: Grammar of BOL

BOL Syntax X	Semantics $\llbracket X \rrbracket$ in FOL
ontology D_1, \dots, D_n	FOL theory $\llbracket D_1 \rrbracket, \dots, \llbracket D_n \rrbracket$
BOL declaration individual i concept i relation i property $i : T$ $I \text{ is-a } C$ $I_1 R I_2$ $I P V$ F	FOL declaration nullary function symbol $i : \iota$ unary predicate symbol $i \subseteq \iota$ binary predicate symbol $i \subseteq \iota \times \iota$ binary predicate symbol $i \subseteq \iota \times T$ axiom $\llbracket C \rrbracket(\llbracket I \rrbracket)$ axiom $\llbracket R \rrbracket(\llbracket I_1 \rrbracket, \llbracket I_2 \rrbracket)$ axiom $\llbracket P \rrbracket(\llbracket I \rrbracket, \llbracket V \rrbracket)$ axiom $\llbracket F \rrbracket$
Formula $C_1 \equiv C_2$ $C_1 \sqsubseteq C_2$	Formula $\forall x : \iota. \llbracket C_1 \rrbracket(x) \Leftrightarrow \llbracket C_2 \rrbracket(x)$ $\forall x : \iota. \llbracket C_1 \rrbracket(x) \Rightarrow \llbracket C_2 \rrbracket(x)$
Individual i	Terms of type ι i
Concept i $C_1 \sqcup C_2$ $C_1 \sqcap C_2$ $\text{dom} R$ $\text{rng} R$	Formula with free variable $x : \iota$ $i(x)$ $\llbracket C_1 \rrbracket(x) \vee \llbracket C_2 \rrbracket(x)$ $\llbracket C_1 \rrbracket(x) \wedge \llbracket C_2 \rrbracket(x)$ $\exists y : \iota. \llbracket R \rrbracket(x, y)$ $\exists y : \iota. \llbracket R \rrbracket(y, x)$
Relation i $R_1; R_2$ $R_1 \cup R_2$ $R_1 \cap R_2$ R^* Δ_C	Formula with free variables $x : \iota, y : \iota$ $i(x, y)$ $\exists m : \iota. \llbracket R_1 \rrbracket(x, m) \wedge \llbracket R_2 \rrbracket(m, y)$ $\llbracket R_1 \rrbracket(x, y) \vee \llbracket R_2 \rrbracket(x, y)$ $\llbracket R_1 \rrbracket(x, y) \wedge \llbracket R_2 \rrbracket(x, y)$ (tricky, omitted) $x = y \wedge \llbracket C \rrbracket(x)$
Property of type T i	Formula with free variables $x : \iota, y : T$ $i(x, y)$

Figure 6.2: Interpretation Function for BOL

Part III

Concretized Knowledge

Part IV

Computational Knowledge

Part V

Deductive Knowledge

Various methods have been developed to represent and perform inferences. We structure our presentation by how each method relates to computation, the aspect most whose integration with inference has drawn the most attention. In general, the ubiquity of underspecified function symbols and quantified variables means that logical expressions usually do not normalize to unique values. At best, computations like $y := f(x)$ can be represented as open-ended conjectures where different options for y are produced, each together with a proof of the respective equality. Therefore, inference systems usually sacrifice computation or at least its efficiency.

Proof assistants sit at the extreme end of this spectrum. They employ strong logics and high-level declarations to provide a convenient way to formalize domain knowledge and reason about it. The reasoning is usually interactive in order to represent inferences that are too difficult to be fully automated. Most proof assistants integrate at least some of the other methods to overcome this weakness.

Further along the spectrum, *automated theorem provers* use simpler logics than interactive proof assistants. They are fully automatic and much faster, but can handle much fewer theorems, and typically do not check their proofs. *Satisfiability checkers* continue this progression by aiming at decidable automation support, whereas theorem proving is usually an semi-decidable search problem. That limits them to propositional logic or specific theories of more expressive logics (usually of first-order logic) that are complete, i.e., where every formula can be proved or disproved. In the special cases, where satisfiability checkers are applicable, they come close to verified computation systems.

Orthogonal to the above triplet, there are several methods for realizing Turing-complete computation naturally inside a logic. Here imperative and object-oriented computation are usually avoided in favor of other programming paradigms that are easier to reason about. *Rewriting* aims at optimizing the $f(x) \rightsquigarrow y$ progression, allowing users to mark specific transformations as rewrite steps. *Terminating recursion* is the method of adding recursive functions to a logic in order to make it a pure functional programming language. Finally, *logic programming* restricts attention to theorems of a special form, for which proof search is simple and predictable so that users can represent computations by supplying axioms that guide the proof search.

Part VI

Narrative Knowledge

Part VII

Conclusion

Bibliography

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