Algorithms and Data Structures Jacobs University Bremen Dr. Florian Rabe Quiz 7 given: 2017-05-16

You have 20 minutes.

Problem 1 Points: 5+2+3

Consider the following recursive function:

```
\begin{aligned} & \mathbf{fun} \ product(x:\mathbb{Z}):\mathbb{Z} = \\ & \mathbf{if} \ x == 0 \\ & 1 \\ & \mathbf{else} \\ & x \cdot product(x-1) \end{aligned}
```

- 1. Convert it to a dynamic program by completing the blanks below.
- 2. Give the Θ -class of the space complexity of the resulting dynamic program.
- 3. Explain in which situations a dynamic program performs better than the recursive variant. Does it matter in this case?

```
\begin{aligned} & \mathbf{fun} \ product(x:\mathbb{Z}):\mathbb{Z} = \\ & results := \mathbf{new} \ Array[\mathbb{Z}](x) \\ & \mathbf{for} \ i \ \mathbf{in} \ 0, \dots, n \end{aligned}
```

Problem 2 Points: 3+2+5

Assume a fixed list D of denominations of coins, sorted decreasingly. We want to find the smallest set of coins whose denominations add up to n.

We use the following greedy-style algorithm:

```
\begin{aligned} & \textbf{fun } greedyCoinChange(D:List[\mathbb{Z}],\,n:\mathbb{Z}):List[\mathbb{Z}] = \\ & solution := Nil \\ & \textbf{for } x \textbf{ in } D \\ & \textbf{while } x + sumList(solution) \leq n \\ & x := prepend(x, solution) \\ & solution \end{aligned} \begin{aligned} & \textbf{fun } sumList(l:List[\mathbb{Z}]):\mathbb{Z} = \\ & \text{return the sum of all elements in } l \end{aligned}
```

- 1. Give the result for running greedyCoinChange([50, 20, 10, 5, 2, 1], 14).
- 2. For fixed D, the algorithm's run time is $\Theta(n^2)$. Give a simple improvement that makes the run time $\Theta(n)$.
- 3. The algorithm is not correct in general. Give an example where greedyCoinChange(D, n) returns a non-optimal result.