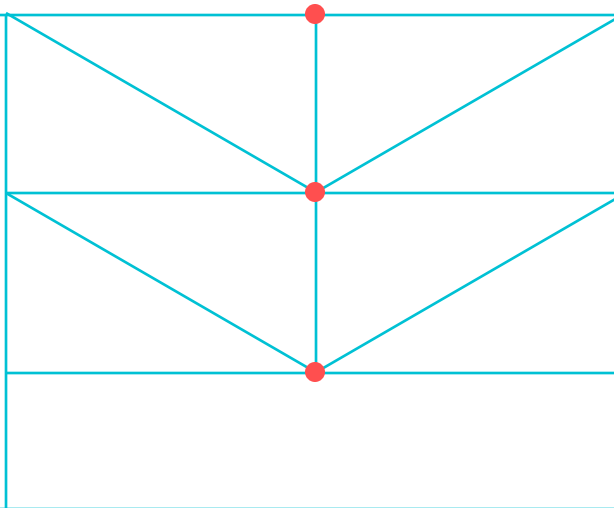


# Introduction to IEEE 802.15.4



**TUHH**  
Institute of  
Communication  
Networks



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**31st of March 2025**

# Course Contents

## Online lectures

- ~~Lecture 1 – Introduction~~
- Lecture 2 – IEEE 802.15.4
- Lecture 3 – IETF 6TiSCH

## Physical Meeting 13th to 17th of April

- Lecture 4 – Theoretical Analysis
- Cooja Simulations and Experiments
- Industrial visit and a talk
- Team Presentations on self learned material
- Lecture 5 – Research Project Results

More details on our padlet,

<https://tuhh.padlet.org/c00zll01/enabling-industry-4-0-j88rkh1i3j7rzmV3>

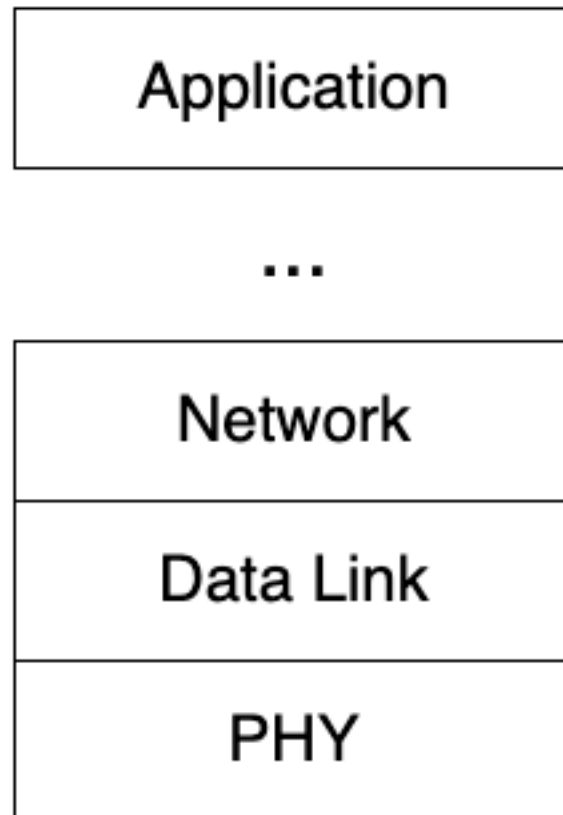
# Wireless Communication for Deterministic Networks

- 2006: Dust Networks Time Sync Mesh Protocol (TSMP)
- 2008: WirelessHART
  - TDMA with fixed time slots
  - Vendor-driven
- 2011: ISA 100.11a
  - TDMA (variable time slots) + CSMA
- 2012: **IEEE 802.15.4 TSCH** (Time Slotted Channel Hopping) mode
  - Configurable TDMA – predictable latency, avoids idle listening and extends device lifetime
  - Channel hopping – improves the reliability in the presence of narrowband interference

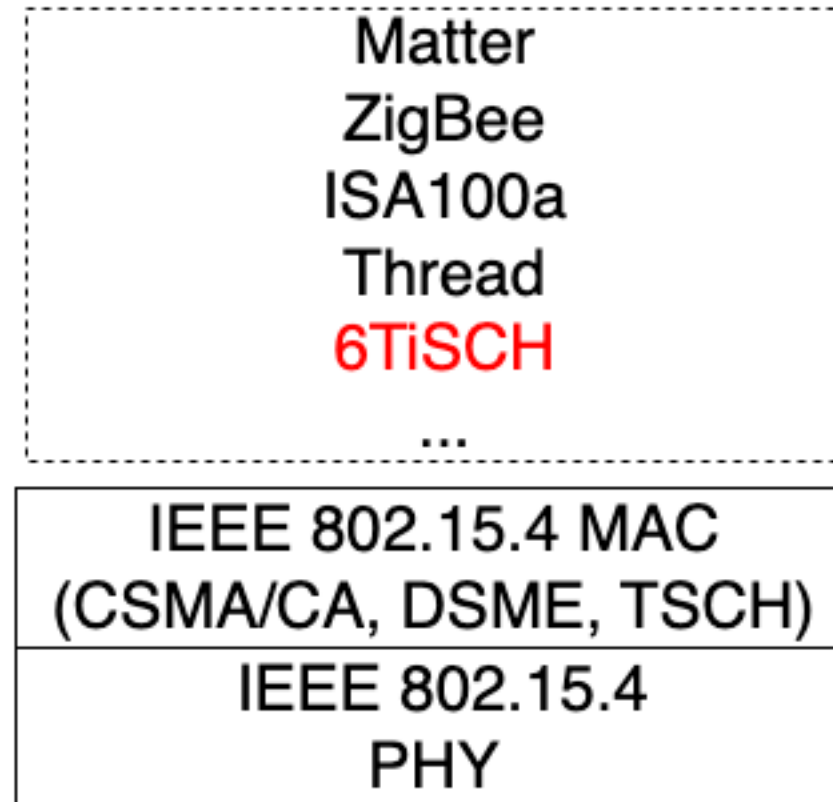
**Any other wireless protocols that you heard about for industry 4.0?**

# Protocol Stack for Industry 4.0 Applications

## TCP/IP Protocol Stack

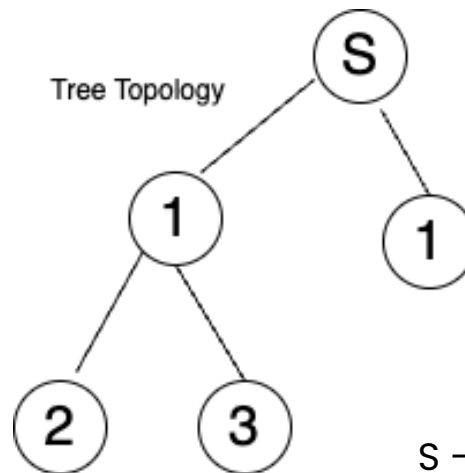
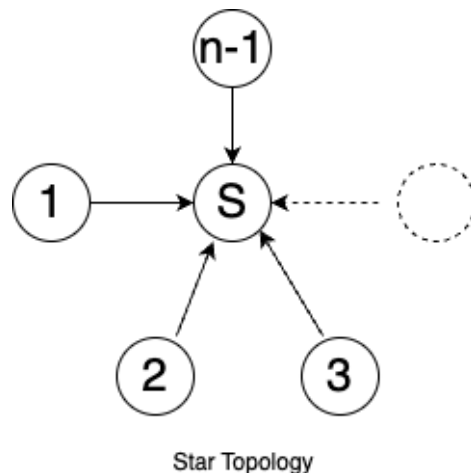


## Protocol stacks for Low Power applications



# Types of MAC – IEEE 802.15.4

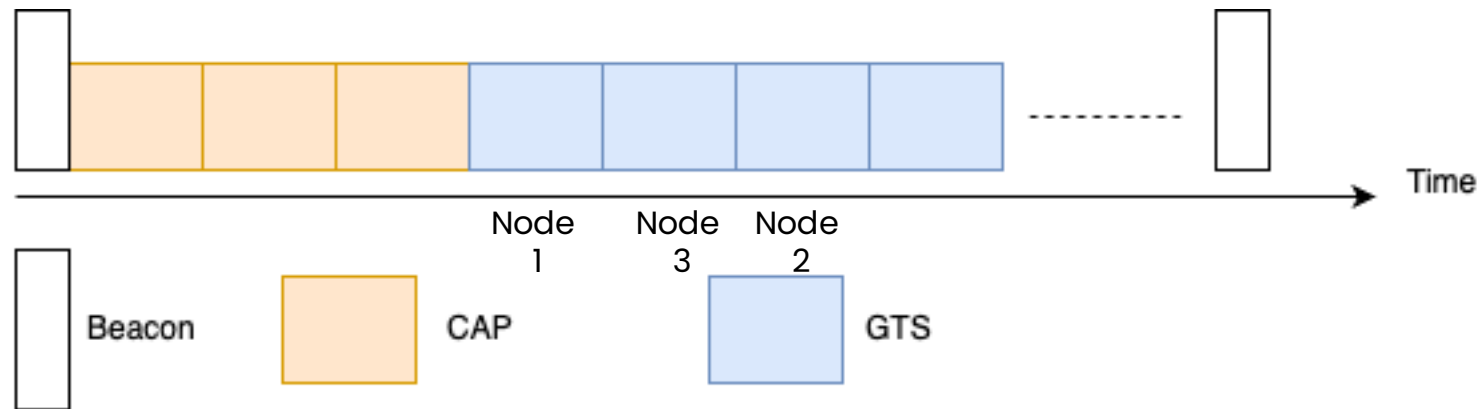
- CSMA/CA – Carrier Sense Multiple Access / Collision Avoidance
  - No guaranteed access
- DSME – Deterministic and Synchronous Multichannel Extension
  - TDMA slots with periodic CSMA/CA access for slot negotiation, broadcast traffic
- TSCH – Time Slotted Channel Hopping
  - TDMA combined with FDMA, scheduled communication *cells* (shared or dedicated)



## 802.15.4 MAC Modes – DSME

### DSME (Deterministic and Synchronous Multi-Channel Extension)

- Uses **beacon-enabled mode** with **multi-channel communication**.
- Supports **two phases**:
  - **CAP (Contention Access Period)**: Nodes compete for access using CSMA/CA.
  - **CFP (Contention-Free Period)**: Time slots are allocated for deterministic access.



# IEEE 802.15.4 TSCH Mode

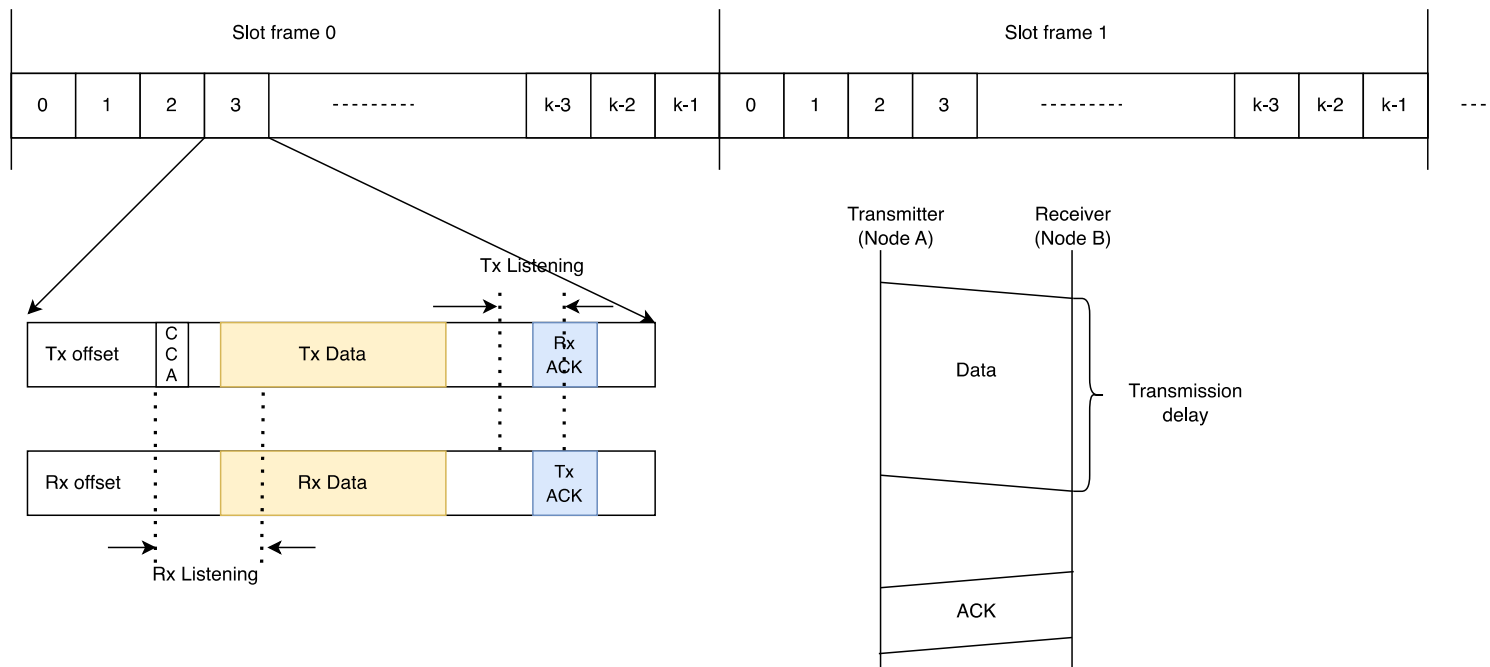
These slides were prepared using the following references:

- [1] IEEE standards for Low-Rate Wireless Networks, <https://ieeexplore.ieee.org/stamp/stamp.jsp?tp=&arnumber=9144691>
- [2] Appendix A of RFC 7554 - TSCH Protocol Highlights, <https://datatracker.ietf.org/doc/html/rfc7554#page-15>
- [3] RFC 8480 - 6TiSCH Operation Sublayer (6top) Protocol (6P), <https://datatracker.ietf.org/doc/html/rfc8480>



# IEEE 802.15.4 TSCH Mode

- Slotframe consists of  $k$  time slots
  - Structure repeats over time
- A pair of nodes communicate within a time slot
  - Actions that are performed are **transmit**, **receive**, or **sleep**



**Any factors that impact the duration of a timeslot?**



# Time Slot Duration

- Example: IEEE 802.15.4 in 2.4 GHz frequency band
    - Max Transmission delay ~ 4 ms (127 bytes @ 250 kbps)
    - a shorter ACK takes about 1 ms
- => 10ms slot (a typical duration)
- leaves 5 ms for radio turnaround, packet processing and security operations

# Slotframe

- Continuously repeats over time
- Multiple slotframes are possible (sorted by priority)
  - Control, application traffic, etc.
  - Resulting schedule is a superposition
- Slotframe size depends on the application needs (from 10 to 1000 of time slots)
  - Shorter vs. longer slotframes?
    - Data transmission rate vs. power consumption

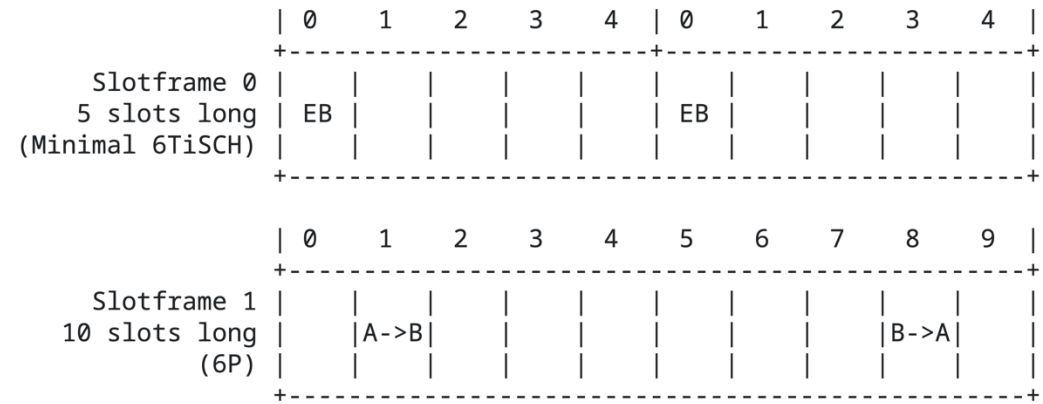
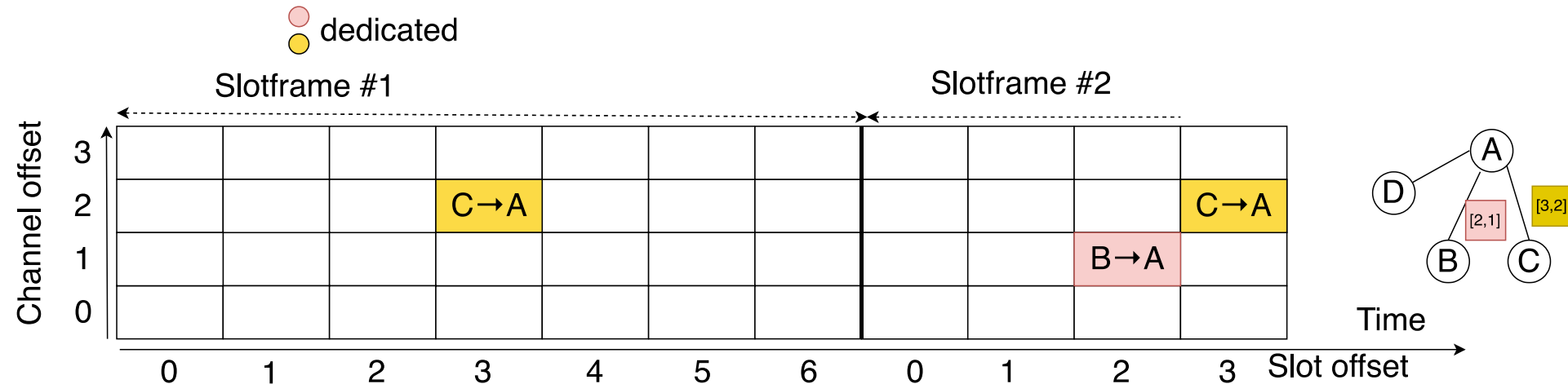


Figure 1. 2-Slotframe Structure when using 6P alongside the Minimal 6TiSCH Configuration<sup>1</sup>

# Channel Offset vs. Slot Offset

$[\text{slotOffset}, \text{channelOffset}]$  is called a *cell*, either transmit (TX), receive (RX) or both.



1. C uses cell  $[3,2]$  to communicate with A. Does node C use the same channel every slotframe?
2. How to support broadcast traffic (e.g. routing information)?
3. How to negotiate/install a dedicated cell between two nodes?

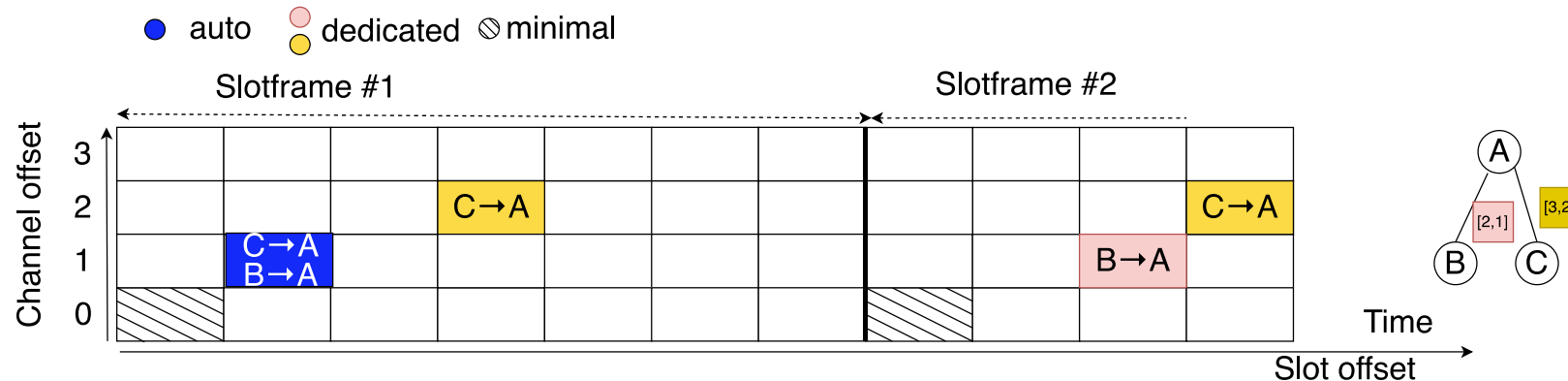
# TSCH Cells

- **Dedicated**

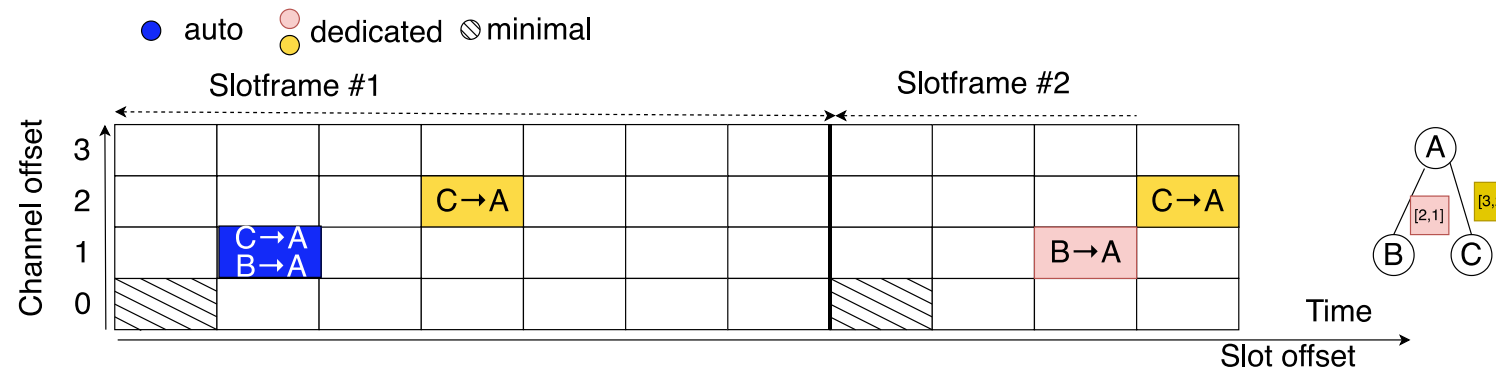
- Contention-free
- Communication between a *single* pair of nodes
- Either transmit or receive mode

- **Shared**

- Contention-based (Slotted ALOHA)
- Communication from/to *multiple* nodes
  - **Auto cell** among neighbors
  - **Minimal cell** [0, 0] is for everything - broadcast traffic (beacons, routing) and joining the network



- Enables shared communication between neighbors
- Maintained autonomously by each node without negotiation
  - Auto RX Cell: [slotOffset, channelOffset] computed as a hash of the 64-bit Extended Unique Identifier (EUI-64) of the node itself
  - Auto TX Cell: [slotOffset, channelOffset] computed as a hash of EUI-64 destination address (MAC) in the unicast frame to be transmitted

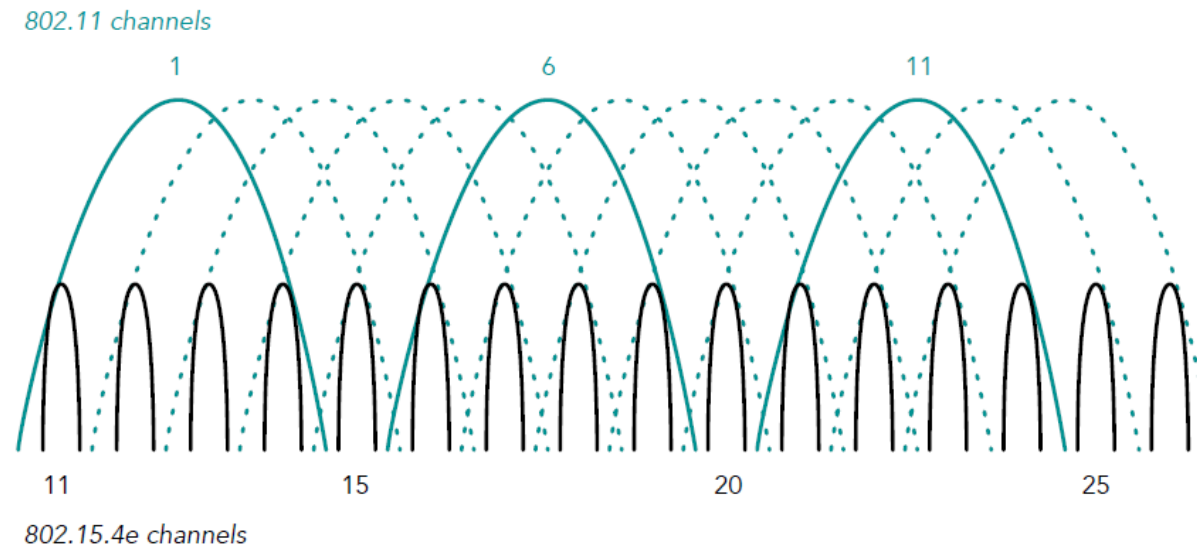


- Example implementation of the hash function see RFC 9033 (Appendix A).<sup>2</sup>

2. <https://datatracker.ietf.org/doc/html/rfc9033#name-example-implementation-of-t>

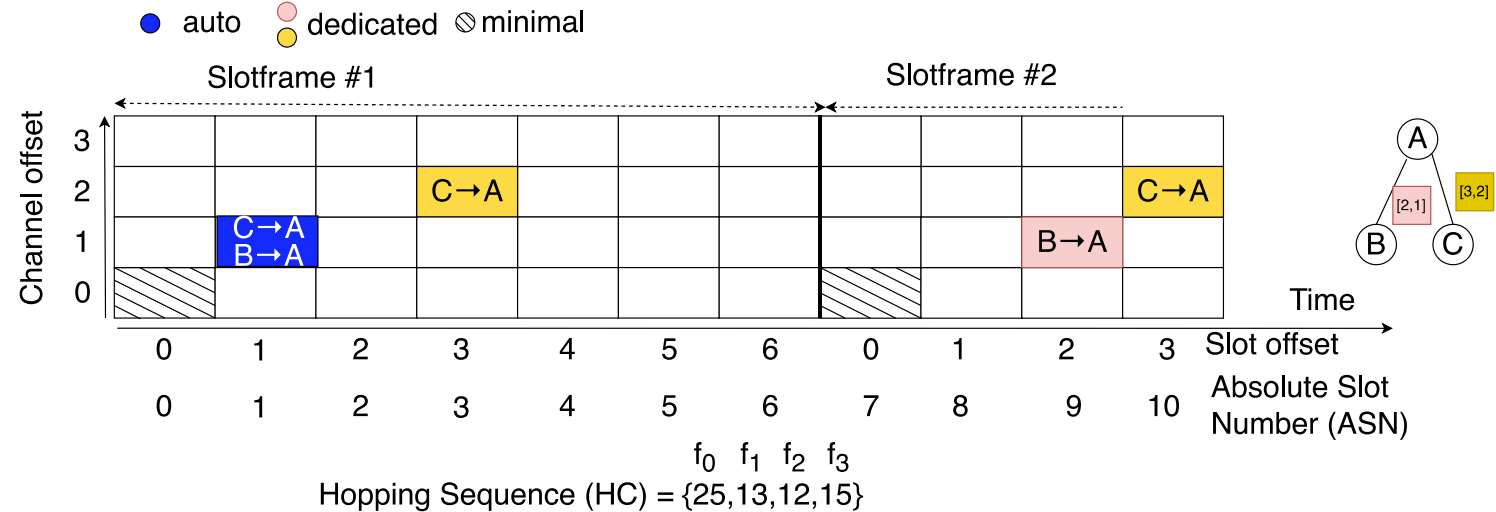
# Channel Hopping

- Hopping Sequence (HS) is a configurable, unordered list of channels to be used for communication
  - Advertised via Enhanced Beacons (EBs) on the minimal cell
  - IEEE 802.15.4 has 16 Channels (11 to 26)



# Channel Hopping

- C has a transmit cell to node A on a channel offset 2, translated by both nodes into a frequency  $f$



$$f = HC [(ASN + channelOffset) \bmod |HC|]$$

- When a new network is created, the ASN is initialized to 0 and increments throughout the network lifetime

**Which channels are used for data transmission from C to A in slotframe #1 and slotframe #2?**

## Channel Hopping contd.

- For a given cell
  - same channel offset, but with different ASN
  - this results in different frequency from the HC
- How to ensure all frequencies from the HC are used?
  - HC and slotframe lengths are coprime
- Time Synchronization among nodes is crucial (out of the scope of this lecture)

**Is TSCH channel hopping robust enough?**



# Blind Channel Hopping

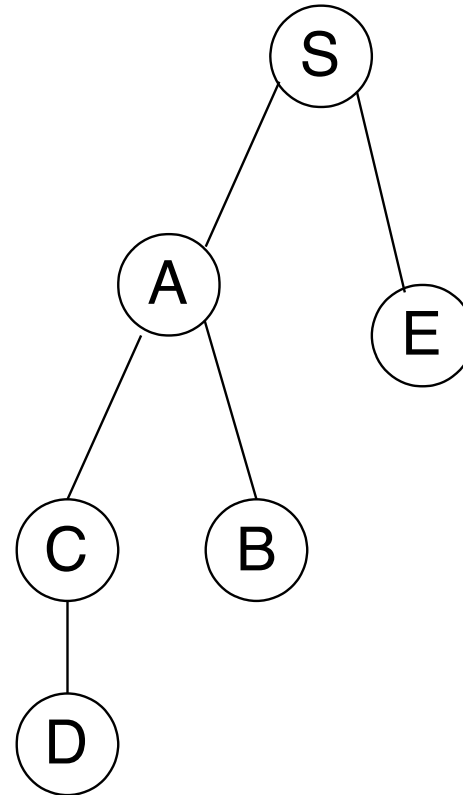
- Mainly designed to work in a **distributed** manner with a very simple algorithm
- Compared to Cellular networks
  - Full control is done by the base station (**centralized** approach)
    - Channels are monitored continuously
    - Dynamic reallocation of cells

# TSCH Open Issues

- **Scheduling:** trade off between throughput, delay and power consumption
  - Centralized vs. distributed
  - Time slot length, slotframe size
  - Spectrum sensing and blacklisting of channels
  - Number of cells per link
    - Avoiding internal collisions
    - Adaptive cell allocation (traffic monitoring)
- Synchronization
- Scalability
- Deployment
- ...

# End-to-End Delay

- Transmission
- Propagation
- Processing
- Queuing

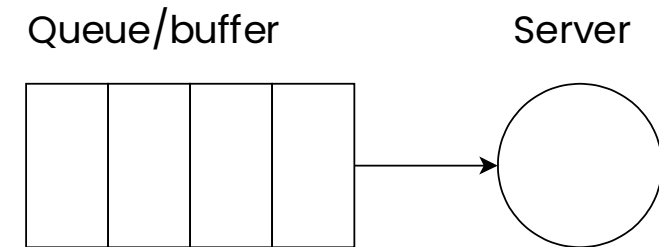


# Queue Model

- Single server
- Unlimited buffer
- Arrival rate  $\lambda$  in packets per slotframe (pkt/sf)
  - Periodic arrivals
- Service rate  $\mu$  in cells per slotframe

## Are these scenarios valid?

- $\lambda = 1.3, \mu = 2$
- $\lambda = 2.3, \mu = 2$
- $\lambda = 1.3, \mu = 1.6$

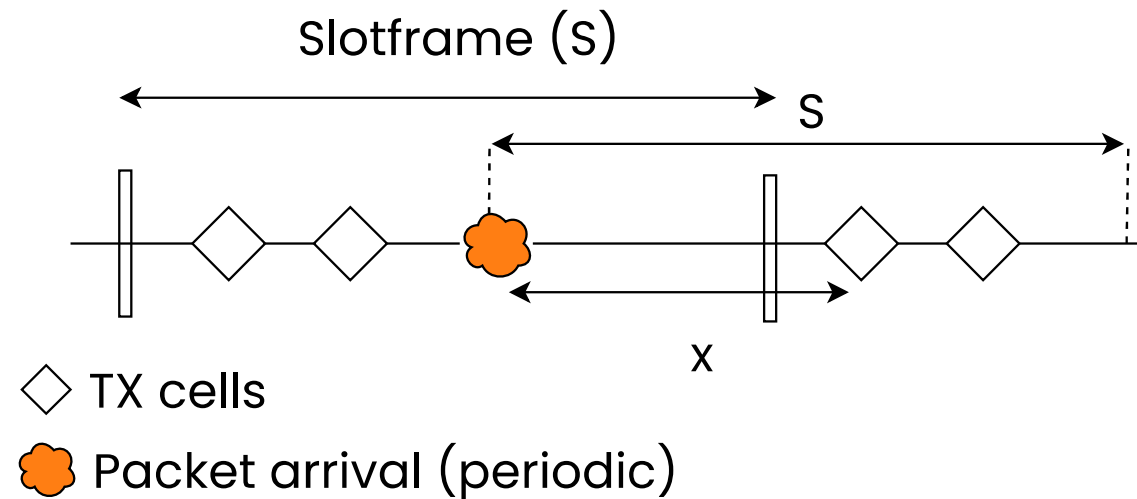


# Service Time Modeling

- Packet arrives in an empty queue
- The Cumulative Distribution Function (CDF) of the service time is

$$F_W = (W < x) = 1 - \left(1 - \frac{x}{S}\right)^\mu,$$

- $S$  is the slotframe length in time slots
- $0 < W < S$



**Run cell A1 of the Jupyter Notebook.**