

CSCE 221 Assignment 6 Cover Page

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Type of sources		
People	Peer Teacher	Dr. Teresa Leyk
Web pages (provide URL)	http://www.cplusplus.com/forum/beginner/89137/ http://www.cplusplus.com/reference/list/list/erase/	http://programmers.stackexchange.com/
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Other Sources		

I certify that I have listed all the sources that I used to develop the solutions/codes to the submitted work.
On my honor as an Aggie, I have neither given nor received any unauthorized help on this academic work.

Your Name Chris Comeaux Date 4/24/2016

Assignment 6 Description

In programming Assignment 6, students had to represent any directed graph as a directed acyclic graph (DAG). To do this, students had to extract the strongly connected components of the graph and represent them as a single vertex. In the first part of the program, students had to read in an adjacency list, create a graph from the list, and print the graph to the screen (as an adjacency list). In the second part, students had to run a depth first search on the graph to find the search order, create a transpose of the graph, and run the depth first search on the transpose. Running the depth first search on the transpose will return the strongly connected components. From there, students had to use the strongly connected components found in the last step and create the final acyclic graph.

Implementation

To implement the notion of a graph I used a STL vector of vertices. Each vertex comprised of an edgeList and a label. The edgeList was a STL list of edges. An edge contained a starting label, end label, a weight, and a visited bit. The starting label represented the starting node of a directed edge and the end label represented the ending node of a directed edge. To display the graph to the screen, I outputted the graph as an adjacency list in the displayGraph() member function of the Graph class. To transpose the graph (reverseGraph(Graph& T)), I visited every edge in every vertices' edge list, created a new edge in the corresponding transpose graph's vertices edgeList, and reversed the start and finish label. The depth first search was implemented through two functions: getDFSOrder and getDFSOrder. These functions utilized a stack and two vectors. One vector held the visited nodes while the other vector held the order of nodes starting from the lowest end time stamp to the highest end time stamp. The stack was used to keep track of the nodes in the current path. To create the final DAG, makeSCG(Graph graph, vector<int> order) was called on the transpose graph. Like the two functions listed above, makeSCG(Graph graph, vector<int> order), utilized 2 vectors and one stack implemented in the same manor. However, this function had an additional vector of vectors of ints. The inner vector each represented one strongly connected component. This vector of vector of ints was then used to extract the strongly connected components from the original graph and create a new graph with the vertices. Finally, the function deleted all edges that did not belong in the DAG, checked for any loops, and displayed the graph to the screen.

C++ Features Used

I did not use any aspects of generic programming in programming assignment 6. However, I did use some object oriented programming features such as abstraction. I created 3 objects, a Graph object, a Vertex object, and a Edge object. The implementation of Edge was hidden behind the Vertex object and the implementation of Vertex was hidden behind the Graph object. Each object knew about the class below it but not above it. For example, the Graph class knew what a Vertex and Edge was, the Vertex class knew what a Edge was but did not know anything about the Graph class, and the Edge class did not know anything about the Vertex or Graph classes. Another C++ feature was that I used 3 standard library containers in my implementation: stack, vector, and list.

Assumptions on Input Data

There were a few assumptions made on the input data. The first assumption was that the input was in the form of an adjacency list. Furthermore, this adjacency list had to be terminated by a -1 and on a separate line. Finally, an assumption was made that each number in the list was separated by a space. A visual representation of these assumptions would be: # # # # -1, where # represents any number and is not limited to the number of # listed.

Testing

I used three test cases to test my program listed below. NOTE: the adjacency list listed under GRAPH is identical to the input data.

Test1:

```
[cmc236]@build ~/csce221/PA6> (19:21:28 04/24/16)
:: ./main input.data
GRAPH
1:  2 4 5 -1
2:  3 4 7 -1
3:  4 -1
4:  6 7 -1
5:  4 -1
6:  5 -1
7:  6 -1

TRANSPOSE GRAPH
1:  -1
2:  1 -1
3:  2 -1
4:  1 2 3 5 -1
5:  1 6 -1
6:  4 7 -1
7:  2 4 -1

ACYCLIC GRAPH
1:  2 4 -1
2:  3 4 -1
3:  4 -1
4:  -1
```

Test 2:

```
[cmc236]@build ~/csce221/PA6> (19:21:36 04/24/16)
:: ./main input1.data
GRAPH
1:  2 -1
2:  3 5 6 -1
3:  4 7 -1
4:  3 8 -1
5:  1 6 -1
6:  7 -1
7:  6 -1
8:  4 7 -1

TRANSPOSE GRAPH
1:  5 -1
2:  1 -1
3:  2 4 -1
4:  3 8 -1
5:  2 -1
6:  2 5 7 -1
7:  3 6 8 -1
8:  4 -1

ACYCLIC GRAPH
1:  2 -1
2:  3 -1
3:  -1
4:  3 -1
```

Test 3:

```
[cmc236]@build ~/csce221/PA6> (19:24:29 04/24/16)
:: ./main input2.data
GRAPH
1:  2 4 -1
2:  4 -1
3:  2 -1
4:  5 -1
5:  6 -1
6:  2 3 -1

TRANSPPOSE GRAPH
1:  -1
2:  1 3 6 -1
3:  6 -1
4:  1 2 -1
5:  4 -1
6:  5 -1

ACYCLIC GRAPH
1:  2 -1
2:  -1
```

Running Time Functions

GRAPH:

buildGraph: $O(n * m)$ where n is the number of lines in the file and m is the number of characters in each line. This is because this function reads in one line at a time then does work on each character in each line.

displayGraph: $O(n * m)$. This is because this function loops through every edge and vertex in the graph. n is the number of vertices and m is the number of edges.

reverseGraph: $O(n * m)$ Like the function above, this function loops through every edge and every vertex. n is the number of vertices and m is the number of edges.

findLabel: $O(n)$ where n is the number of vertices. This function loops through each every vertex and return the index of the vertex with the label that was passed to it.

getDFSOrder: $O(1)$ This function is constant because it simply initializes the depth first search. It first finds the index of the first vertex, add that vertex to the the visited vector and vertex stack and then calls vertex traversal.

vertexTraversal: $O(n)$ This function is recursive and continues to call itself as long as all vertices have not been visited. Since it visits all vertices it is $O(n)$ where n is the number of vertices.

makeSCG: $O(n * m)$ This function, like the other listed above, visits every vertex and every edge in the output graph. even though it has many loops, its can be bounded by $O(n * m)$ where n is the number of vertices in the graph and m is the nuber of edges in the graph

VERTEX

reverseEdge: $O(n)$ where n is the number of edges in the edge list. This is because this function loops though every node in the edgeList and reverses the edge

isVisited: $O(n)$ where n is the number of vertices in the visited vector. This function compares the passed value to every number in the vector too see if has been visited or not.

isVertex: $O(n)$ where n is the number of vertices in the graph call on. This vertex compares the passed label to every label in the vertices vector to see if there is a corresponding vertex with that label.

isLoop: $O(n * m)$ where n is the number of vertices in the graph and m is the number of edges in the graph. This function goes through every vertex's edges and sees if there is any type of loop.

Real Life Application

One application that this program can be used is when calculating the maximum execution time of a set. To do this we need to have a DAG of the set in memory. Since my program can be used to create a DAG from any directed graph, you could pass the original graph to my program, extract the DAG and then pass it to another algorithm that calculates the maximum execution time of the set.

Conclusion

Programming assignment 6 taught students about graphs and graph traversals, especially depth first search. To complete this assignment students had to understand graphs enough to represent a graph in c++ and then write a depth first search to traverse the graph. Furthermore, the students had to understand what strongly connected components were and how to recognize them in graphs to create a DAG from any directed graph.