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CS 4412

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Project One Report

1. Working Examples:

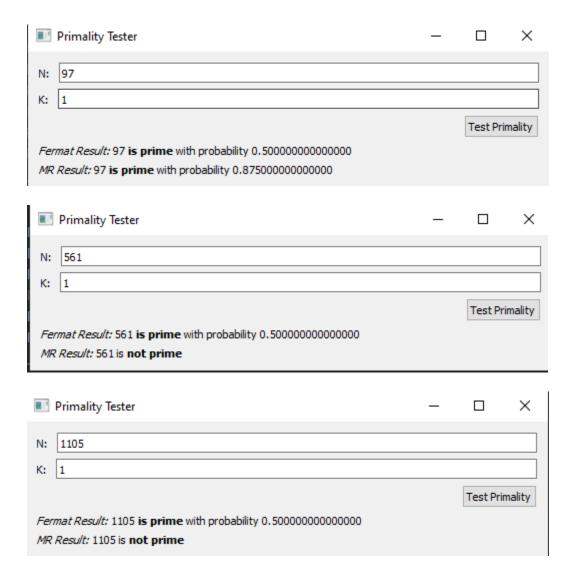
■ Primality Tester	_		×
N: 291			
K: 4			
		Test Prin	mality
Fermat Result: 291 is not prime			
MR Result: 291 is not prime			
Drimality Tostor	_		×
■ Primality Tester		Ш	^
N: 37			
K: 3			
		Test Prin	mality
Fermat Result: 37 is prime with probability 0.87500000000000		Test Pri	mality
		Test Pri	mality
Fermat Result: 37 is prime with probability 0.87500000000000000000000000000000000000		Test Prir	
Fermat Result: 37 is prime with probability 0.87500000000000	_	Test Prin	mality
Fermat Result: 37 is prime with probability 0.87500000000000000000000000000000000000	_		
Fermat Result: 37 is prime with probability 0.87500000000000000 MR Result: 37 is not prime Primality Tester	_		
Fermat Result: 37 is prime with probability 0.87500000000000 MR Result: 37 is not prime Primality Tester N: 23	_		×
Fermat Result: 37 is prime with probability 0.87500000000000 MR Result: 37 is not prime Primality Tester N: 23	_		×

ı	■ P	rimality Tester	_		×
	N: 2	23			
	K: 7	7			
ı				Test Pri	mality
		at Result: 23 is prime with probability 0.992187500000000 esult: 23 is prime with probability 0.998046875000000			

Discussion experimentation.

There were several bugs that I almost didn't catch. The Miller-Rabin method was confusing and I did not read some edge case text properly at first. I had the number 293 causing problems for a while, but now it should all be working properly.

The following images demonstrate requirements passing. I chose 1105 for my unique Carmichael number.



Complexity:

The complexity and space of the probability calculations was O(1) for both functions.

The complexity of the modular exponentiation function was O(log(N)), and the space requirements was O(n) where n was the number of operations that the function completed. Normally space requirements for this equation can be O(1) but I decided to store intermediate values instead of run cleanup operations that removed nodes from the end of my list.

Note that I decided to use an iterative approach instead of a recursive approach because I don't like recursion and thought it would be interesting.

The time complexity of the Fermat function was O(k) * the complexity of the modular exponential function, where k was the number of trials. The space complexity was O(k) as I stored an array of random numbers instead of generating as needed.

The complexity of the Miller-Rabin function was O(k) * O(log(N)) where k was the number of trials and O(Log(N)) represents that there was a repeated square root operation happening where N was the size of the prime input, ie. The bigger the input the proportionally more time the function took.

Probability equations:

I documented both of my probability equations. In short, the Miller-Rabin method captured ¾ of the numbers that Fermat wouldn't. In short, the Miller-Rabin method was 4x more likely to solve a given number, and so its probability of failure was 25% of what the Fermat's formula was.

The base equation for each was $(1/2^k)$ where k is the number of tests.