

1 Syntax

1.1 Language E of Expressions

Lan- guage	Types	Expressions	Comments
E (Ch.4)	num str	x $\text{num}[n]$ $\text{str}[s]$ $e_1 + e_2$ $e_1 * e_2$ $e_1 \wedge e_2$ $\text{let } x \text{ be } e_1 \text{ in } e_2$	
ED (Ch.8.1)		$\text{fun } f(x : \tau_1) : \tau_2 = e_1 \text{ in } e_2$ $e_1(e_2)$	Limited extension, superceded by next: First-order functions, with their names are from a different variable supply here.
EF (Ch.8.2)	$\tau_1 \rightarrow \tau_2$	$\lambda(x : \tau)e$ $e_1(e_2)$	Full functions as first-class citizens, with variable names.

1.2 Language T of Gödel total functions

♠JB: ...TODO♠

1.3 Language family PCF of (general) recursive functions

Lan- guage	Types	Expressions	Comments
PCF (Ch.19)	nat $\tau_1 \rightarrow \tau_2$	x z $s(e)$ $\text{ifz } e (e_0, x.e_1)$ $\lambda(x : \tau)e$ $e_1(e_2)$ $\text{fix } (x : \tau) \text{ is } e$	
FPC (Ch.20)	t rect is τ	$\text{fold}_{t,\tau}(e)$ $\text{unfold}(e)$	Full functions as first-class citizens, with variable names.

2 Typing

Lan- guage	Rules	Comments
E (Ch.4)	$\overline{\Gamma, x : \tau \vdash x : \tau} \quad \overline{\Gamma \vdash \mathbf{str}[s] : \mathbf{str}} \quad \overline{\Gamma \vdash \mathbf{num}[n] : \mathbf{num}}$ $\frac{\Gamma \vdash e_1 : \mathbf{num} \quad \Gamma \vdash e_2 : \mathbf{num}}{\Gamma \vdash \mathbf{plus}(e_1, e_2) : \mathbf{num}} \quad \frac{\Gamma \vdash e_1 : \mathbf{num} \quad \Gamma \vdash e_2 : \mathbf{num}}{\Gamma \vdash \mathbf{times}(e_1, e_2) : \mathbf{num}}$ $\frac{\Gamma \vdash e_1 : \mathbf{str} \quad \Gamma \vdash e_2 : \mathbf{str}}{\Gamma \vdash \mathbf{cat}(e_1, e_2) : \mathbf{str}} \quad \frac{\Gamma \vdash e : \mathbf{str}}{\Gamma \vdash \mathbf{len}(e) : \mathbf{num}}$ $\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma, x : \tau_1 \vdash e_2 : \tau_2}{\Gamma \vdash \mathbf{let } x \mathbf{ be } e_1 \mathbf{ in } e_2 : \tau_2}$	Typing axiom and atoms num operations conversions local binding
ED (Ch.8.1)	...	♠JB: TODO♠