

Kotlin ▽

A Shape Safe eDSL for Differentiable Functional Programming

Breandan Considine

Université de Montréal

breandan.considine@umontreal.ca

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Overview

- 1 A Short Lesson on Computing Derivatives
- 2 Introduction and motivation
- 3 Architectural Overview
- 4 Usage
- 5 Future exploration

Differentiation

If we have a function, $P(x) : \mathbb{R} \rightarrow \mathbb{R}$, recall the derivative is defined as:

$$P'(x) = \lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h} = \frac{\Delta y}{\Delta x} = \frac{dP}{dx} \quad (1)$$

For $P(x_0, x_1, \dots, x_n) : \mathbb{R}^n \rightarrow \mathbb{R}$, the gradient is a vector of derivatives:

$$\nabla P = \left[\frac{\partial P}{\partial x_0}, \frac{\partial P}{\partial x_1}, \dots, \frac{\partial P}{\partial x_n} \right] \text{ where } \frac{\partial P}{\partial x_i} = \frac{dP}{dx_i} \quad (2)$$

For $\mathbf{P}(x_0, x_1, \dots, x_n) : \mathbb{R}^n \rightarrow \mathbb{R}^m$, the Jacobian is a vector of gradients:

$$\mathbf{J}_P = [\nabla P_0, \nabla P_1, \dots, \nabla P_n] \text{ or equivalently, } \mathbf{J}_{ij} = \frac{\partial P_i}{\partial x_j} \quad (3)$$

Type checking automatic differentiation

Suppose we have a program $P : \mathbb{R} \rightarrow \mathbb{R}$ where:

$$\mathbf{P}(p_0) = p_q \circ p_{q-1} \circ p_{q-2} \circ \cdots \circ p_1 \circ p_0 \quad (4)$$

From the chain rule of calculus, we know that:

$$\frac{dP}{dp_0} = \frac{dp_q}{dp_{q-1}} \frac{dp_{q-1}}{dp_{q-2}} \cdots \frac{dp_1}{dp_0} = \prod_{i=1}^n \frac{dp_i}{dp_{i-1}} \quad (5)$$

More generally, for $P : \mathbb{R}^n \rightarrow \mathbb{R}^m$, the chain rule also applies:

$$\mathbf{J_P} = \prod_{i=1}^q \mathbf{J_{p_i}} = \left(\left((\mathbf{J_{p_q}} \mathbf{J_{p_{q-1}}}) \cdots \mathbf{J_{p_2}} \right) \mathbf{J_{p_1}} \right) = \left(\mathbf{J_{p_q}} \left(\mathbf{J_{p_{q-1}}} \cdots (\mathbf{J_{p_2}} \mathbf{J_{p_1}}) \right) \right) \quad (6)$$

In order for \mathbf{P} to type check, what is the type signature of $p_{0 < i < n}$?

$$p_i : T_{out}(p_{i-1}) \rightarrow T_{in}(p_{i+1}) \quad (7)$$

Shape checking and inference

- Scalar functions implicitly represent shape as arity $f(1, 2) : \mathbb{R}^2 \rightarrow \mathbb{R}$
- To check matrix functions, we need a type-level encoding of shape
- Arbitrary matrix functions (e.g. convolution) require dependent types
- But parametric polymorphism will suffice for most matrix functions
- For arithmetical operations, we just need to check for equality

Math	Derivative	Code	Type Signature
$a(b)$	$\mathbf{J}_a \mathbf{J}_b$	<code>a(b)</code>	$(a : \mathbb{R}^\tau \rightarrow \mathbb{R}^\pi, b : \mathbb{R}^\lambda \rightarrow \mathbb{R}^\tau) \rightarrow (\mathbb{R}^\lambda \rightarrow \mathbb{R}^\pi)$
$a + b$	$\mathbf{J}_a + \mathbf{J}_b$	<code>a + b</code> <code>a.plus(b)</code> <code>plus(a, b)</code>	$(a : \mathbb{R}^\tau \rightarrow \mathbb{R}^\pi, b : \mathbb{R}^\lambda \rightarrow \mathbb{R}^\pi) \rightarrow (\mathbb{R}^\tau \rightarrow \mathbb{R}^\pi)$
ab	$\mathbf{J}_a b + \mathbf{J}_b b$	<code>a * b</code> <code>a.times(b)</code> <code>times(a, b)</code>	$(a : \mathbb{R}^\tau \rightarrow \mathbb{R}^{m \times n}, b : \mathbb{R}^\lambda \rightarrow \mathbb{R}^{n \times p}) \rightarrow (\mathbb{R}^\tau \rightarrow \mathbb{R}^{m \times p})$

- Abstract algebra can be useful when generalizing to new structures
- Helps us to easily translate between mathematics and source code
- Fields are a useful concept when computing over real numbers
 - A field is a set F with two operations $+$ and \times , with the properties:
 - Associativity: $\forall a, b, c \in F, a + (b + c) = (a + b) + c$
 - Commutivity: $\forall a, b \in F, a + b = b + a$ and $a \times b = b \times a$
 - Distributivity: $\forall a, b, c \in F, a \times (b + c) = (a \times b) + (a \times c)$
 - Identity: $\forall a \in F, \exists 0, 1 \in F$ s.t. $a + 0 = a$ and $a \times 1 = a$
 - $+$ inverse: $\forall a \in F, \exists (-a)$ s.t. $a + (-a) = 0$
 - \times inverse: $\forall a \neq 0 \in F, \exists (a^{-1})$ s.t. $a \times a^{-1} = 1$
- Extensible to other number systems (e.g. complex, dual numbers)
- What is a program, but a series of arithmetic operations?

Why Kotlin?

- Goal: To implement automatic differentiation in Kotlin
- Kotlin is a language with strong static typing and null safety
- Supports first-class functions, higher order functions and lambdas
- Has support for algebraic data types, via tuples sealed classes
- Extension functions, operator overloading other syntax sugar
- Offers features for embedding domain specific languages (DSLs)
- Access to all libraries and frameworks in the JVM ecosystem
- Multi-platform and cross-platform (JVM, Android, iOS, JS, native)



- Type system
 - Strong type system based on algebraic principles
 - Leverage the compiler for static analysis
 - No implicit broadcasting or shape coercion
 - Parameterized numerical types and arbitrary-precision
- Design principles
 - Functional programming and lazy numerical evaluation
 - Eager algebraic simplification of expression trees
 - Operator overloading and tapeless reverse mode AD
- Usage desiderata
 - Generalized AD with functional array programming
 - Automatic differentiation with infix and Polish notation
 - Partial derivatives and higher order derivatives and gradients
- Testing and validation
 - Numerical gradient checking and property-based testing
 - Performance benchmarks and thorough regression testing

How do we define algebraic types in Kotlin ▽?

// T: Group<T> is effectively a self type

```
interface Group<T: Group<T>> {  
    operator fun plus(f: T): T  
    operator fun times(f: T): T  
}
```

// Inherits from Group, default methods

```
interface Field<T: Field<T>>: Group<T> {  
    operator fun unaryMinus(): T  
    operator fun minus(f: T): T = this + -f  
    fun inverse(): T  
    operator fun div(f: T): T = this * f.inverse()  
}
```

Algebraic Data Types

```
class Var: Expr()
class Const(val num: Number): Expr()
class Sum(val e1: Expr, val e2: Expr): Expr()
class Prod(val e1: Expr, val e2: Expr): Expr()
```

```
sealed class Expr: Group {
    fun diff() = when(expr) {
        is Const -> Zero
        is Sum -> e1.diff() + e2.diff()
        is Prod -> e1.diff() * e2 + e1 * e2.diff()
        is Var -> One
    }
}
```

```
operator fun plus(e: Expr) = Sum(this, e)
operator fun times(e: Expr) = Prod(this, e)
}
```

Expression simplification

```
operator fun Expr.times(exp: Expr) = when {  
    this is Const && num == 0.0 -> Const(0.0)  
    this is Const && num == 1.0 -> exp  
    exp is Const && exp.num == 0.0 -> exp  
    exp is Const && exp.num == 1.0 -> this  
    this is Const && exp is Const -> Const(num*exp.num)  
    else -> Prod(this, e)  
}
```

```
// Sum(Prod(Const(2.0), Var()), Const(6.0))  
val q = Const(2.0) * Sum(Var(), Const(3.0))
```

Extension functions and contexts

```
class Expr<T: Group<T>>: Group<Expr<T>> {  
    // ...  
    operator fun plus(exp: Expr<T>) = Sum(this , exp)  
    operator fun times(exp: Expr<T>) = Prod(this , exp)  
}  
  
object DoubleContext {  
    operator fun Number.times(exp: Expr<DoubleReal>) =  
        Const(toDouble()) * exp  
}  
  
// Uses '*' operator in DoubleContext  
fun Expr<DoubleReal>.multiplyByTwo() =  
    with(DoubleContext) { 2 * this }
```

Automatic test case generation

```
val x = variable("x")
val y = variable("y")

val z = y * (sin(x * y) - x) // Function under test
val dz_dx = d(z) / d(x)      // Automatic derivative
val manualDx = y * (cos(x * y) * y - 1)

"dz/dx should be y * (cos(x * y) * y - 1)" {
    assertAll (NumGen) { cx, cy ->
        // Evaluate the results at a given seed
        val autoEval = dz_dx(x to cx, y to cy)
        val symbEval = manualDx(x to cx, y to cy)
        // Should pass if |adEval - manualEval| < eps
        autoEval shouldBeApproximately symbEval
    }
}
```

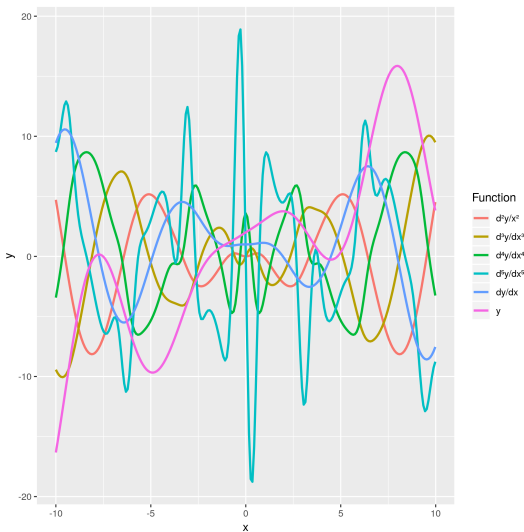
Usage: plotting higher derivatives of nested functions

```
with(DoublePrecision) { // Use double-precision
    val x = variable() // Declare an immutable variable
    val y = sin(sin(sin(x)))/x + sin(x) * x + cos(x) + x

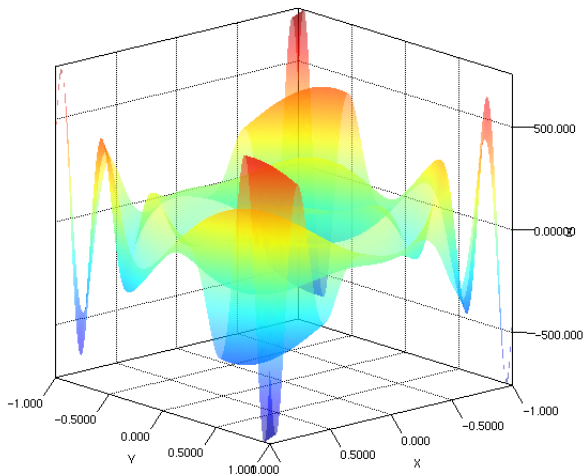
    // Lazily compute reverse-mode automatic derivatives
    val dy_dx = d(y) / d(x)
    val d2y_dx = d(dy_dx) / d(x)
    val d3y_dx = d(d2y_dx) / d(x)
    val d4y_dx = d(d3y_dx) / d(x)
    val d5y_dx = d(d4y_dx) / d(x)

    plot(-10..10, dy_dx, dy2_dx, d3y_dx, d4y_dx, d5y_dx)
}
```

$$y = \frac{\sin \sin \sin x}{x} + x \sin x + \cos x + x, \quad \frac{dy}{dx}, \quad \frac{d^2y}{dx^2}, \quad \frac{d^3y}{dx^3}, \quad \frac{d^4y}{dx^4}, \quad \frac{d^5y}{dx^5}$$



$$z = \sin(10(x^2 + y^2))/10, \quad \frac{\partial^3 z}{\partial^2 y \partial x}$$



Further directions to explore

- Theory Directions

- Generalization of types to higher order functions, vector spaces
- Dependent types via code generation to type-check convolution
- General programming operators and data structures
- Imperative define-by-run array programming syntax
- Program induction and synthesis, cf.
 - The Derivative of a Regular Type is its Type of One-Hole Contexts
 - The Differential Lambda Calculus (2003)
- Asynchronous gradient descent (cf. HogWild, YellowFin, et al.)

- Implementation Details

- Closer integration with Kotlin/Java standard library
- Encode additional structure, i.e. function arity into type system
- Vectorized optimizations for matrices with certain properties
- Configurable forward and backward AD modes based on dimension
- Automatic expression refactoring for numerical stability
- Primitive type specialization, i.e. `FloatVector <: Vector<T>?`

Learn more at:

<http://kg.ndan.co>

Liam Paull
Michalis Famelis



Université  de Montréal

