

Assignment 5 Design

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1 Introduction

In this assignment we will be using a public-key cryptography system to implement a cryptography program for encrypting and decrypting messages. More specifically, we will be implementing the Schmidt-Samoa (SS) algorithm for encryption and decryption. Our implementation of the algorithm will require us to be able to handle arbitrary precision integers. Towards that end, we will be making use of the GNU multi-precision arithmetic (GMP) library to handle large integers. The files we will be writing are:

- decrypt.c
- encrypt.c
- keygen.c
- numtheory.c
- randstate.c
- ss.c

2 Design and Psuedocode

The main programs we will be writing in this assignment (keygen, encrypt, and decrypt) require the usage of 2 libraries and a random state module. I will be using regular, pure C operators and arithmetic in my design, as GMP is hard to read and understand. However, it is easier to use if you just use it replace C operations and functions line by line with the GMP operations. All mentions of “mpz_t” in function paramters should be considered intgers.

2.1 randstate.c

The random state module will simply be used to initialize a global random “state” variable to use with GMP. It will seed both the GMP random function and the C library random function with the passed in seed.

```
func randstate_init(int seed):  
    seed gmp_randstate_ui  
    seed srandom()
```

```

gmp_randinit_mt(state) //state is a global extern variable

func randstate_clear(void):
    clear state with gmp_randclear()

```

2.2 numtheory.c

This file will contain the library for mathematic functions that are required for implementing the SS algorithm. For this program, we were provided psuedocode, so my design will closely resemble it.

```

//paramters: o is the output variable, a is the base that will
//be raised to exponent d, and n is the modulus
void pow_mod(mpz_t o, mpz_t a, mpz_t d, mpz_t n);
    int v = 1 //holds calculation results
    int p = a //p holds "a" to avoid changing "a"
    while exponent d > 0;
        if d is odd:
            v = (v * p) % n
        p = (p*p) % n
        d = floordiv(d/2)
    return v

```

For the primality testing function (the Miller-Rabin) test, we must evaluate the first step of the test: *write $n - 1 = 2^s r$ such that r is odd*. This equation can be rewritten as: $r = \frac{n-1}{2^s}$. This essentially means that we must divide $n - 1$ by 2 s times until it results in a positive number. We then hold on to the number of times required to reach an odd number as well as resultant odd number.

Secondly, we need to choose a random number a where $a \in \{2, 3, \dots, n - 1\}$. This is relatively simple. Since we need to exclude 0 and 1, we can just add 2 to the random number. Furthermore, gmp contains the function `mpz_urandomm(output, state, n)` which returns a random integer in the range $[0, n)$. Thus we need to just pass in $n-3$ and then add 2 to the output of the RNG to get a random number within the desired range. (the design uses $\%(n-4)$ due to how modulus in C works).

```

//in the SS implementation, the parity of the number will most likely
//be checked before being passed into is_prime. Thus, for now, assume
//that all arguments passed in are at least odd.
bool is_prime(mpz_t n, uint64_t iters);
    int r = (n-1)/2
    int s = 1 //first division done above. 2 divides all evens at least once
    while (is_odd(r) == false){
        r = r/2
    }

```

```

        s++;
    }

    int a
    for int i from 1 to iters:
        a = random()%(n-4)
        a += 2;
        int y = power_mod(a,r,n)
        if y != 1 AND y != n - 1:
            int j = 1
            while j <= (s-1) AND y != (n-1):
                y = power_mod(y,2,n)
                if y == 1:
                    return false
                j +=1
            if y != (n-1)
                return false
    return true;

//this function generate a random number, then uses a while loop
//to check and regenerate random numbers until a prime is found
//
void make_prime(mpz_t p, uint64_t bits, uint64_t iters);
    use urandomb(output, state, bit_cnt) to generate a random number
    bit_cnt should be set to the passed in "bits" argument
    while (is_prime(output, iters)):
        urandomb(output, state, bits)

```

The next 2 functions in the number theory library deal with modular inverses. However, to implement a modular inverse function we first need to create a function to find the greatest common divisor (GCD). Furthermore, mod-inverse as shown on the assignment spec contains a lot of parallel assignments. Since C cannot do parallel assignments, temporary variables will be needed to do the math.

```

//this function finds the greatest common divisor of a and b, and
//stores the result in d
func GCD(output, a, b):
    int t
    while b != 0;
        t = b
        b = a % b
        a = t
    return a

```

```

//compute the inverse i of a modulo n
func mod-inverse(inverse i, int a, modulo n):
    int r = n
    int r' = a
    int t = 0
    int t' = 1
    int q;

    //each parallel assignment will be spaced out to clarify the code
    while r' != 0:
        q = r/r'

        temp_r = r
        r = r'
        r' = temp_r - (q * r')

        temp_t = t
        t = t'
        t' = temp_t - (q * t')
    if r > 1:
        i = 0
    if t < 0:
        t = t + n
    return t

```

3 ss.c

This sections will cover functions pertaining to the actual logic of the Schmidt-Samoa algorithm. This file will contain the code for the creation of the public and private keys and will also contain functions to read and write to files. Furthermore, this file contains the functions to actually encrypt and decrypt text files. For the following functions, it should be noted that GMP uses pass by reference. Thus, all of the following functions are void and do not require return statements because changing variables inside the functions changes them outside the function scope too.

```

//
//makes components for a SS public key. Generates primes
//p and q, as well as the public modulus n. Takes in arguments
//nbits (for number of bits in n) and iters for number of iterations
//the miller-rabin test should go through to check primality
//

```

```

//First, we must determine the number of bits that go in p and q.
//for prime p, we have been given the bounds [nbits/5, (2*nbits)/5]
//we can use gmp_urandomm to obtain a number within those bounds.
//we only need to generate random numbers up to nbits/5 since
//we can just add nbits/5 again to ensure that the randomly
//generated number is within bounds
//
func ss_make_pub(int p, int q, int n, int nbits, int iters);

```

```

    int low = nbits/5 //will be added to randomly generated
    int rand_bits = urandomm(out, state, low + 1)
    rand_bits += nbits/5
    call make_prime twice with p and q as outputs:
        //both make_primes should accept iters as
        //an argument. make_prime for p should accept
        //the number of bits contained in rand_bits
        //while make_prime for q should accept 1 - nbits
        //bits.
    n = p^2 * q
    include assert statement to make sure n has at least nbits bits

```

To create a private key, we must find the least common multiple of the primes p and q. The formula for determining the least common multiple is $\text{lcm}(p,q) = \frac{p \cdot q}{\text{gcd}(p,q)}$.

```

func ss_make_priv(mpz_t d, mpz_t pq, mpz_t p, mpz_t q);
    compute public modulus (n = p^2*q)
    compute private modulus (pq = p*q)
    //below is computation of lcm (p-1,q-1)
    int lamda_pq = p*q
    lambda_pq = lambda_pq / gcd(p,q)

//the write to file functions are relatively simple and only
//require calls to gmp_fprintf
func ss_write_pub(mpz_t n, char username[], FILE *pbfile);
    write n to pbfile (gmp's fprintf function)
        //this should use the gmp format specifier %Zx, which specifies
        //a hex integer output
    on new line, write username to pbfile

func ss_write_priv(mpz_t pq, mpz_t d, FILE *pvfile);
    write pq to pvfile (gmp's fprintf function)
        //use gmp's hex specifier %Zx
    on new line, d to pvfile

```

```

//the functions to read the public and private keys will be very
//similar to the functions to write those keys.
func ss_read_pub(mpz_t n, char username[], FILE *pbfile);
    read formatted file using gmp's fscanf
    a single scan where n is set to the first line (private key)
    and username set to duplication (using strdup) of the second
    line should suffice

func ss_read_priv(mpz_t pq, mpz_t d, FILE *pvfile);
    read formatted file using fscanf
    a single scan, where we scan:
        line 1, pq
        line 2 d

```

The following functions require more work than the previous functions, as they pertain to the actual encryption and decryption of text. The functions that accomplish the encryption are `ss_encrypt` and `ss_decrypt`. These functions only require a singular call to power-mod because of how encryption and decryption are defined in SS. On the other hand, the functions to encrypt and decrypt entire files requires more work because of the limitations on the size of the text to be encrypted.

```

func ss_encrypt(mpz_t c, mpz_t m, mpz_t n);
     $c = m^n \bmod n = \text{pow\_mod}(c, m, n, n)$ 

func ss_decrypt(mpz_t m, mpz_t c, mpz_t d, mpz_t pq);
     $m = c^d \bmod pq = \text{pow\_mod}(m, c, d, pq)$ 

```

We are tasked with implementing file encryption and decryption in *blocks of text*. This is because of the modulus n , as the value of the text we are encrypting must be less than n . Furthermore, We need to insert a single byte to the front of each block since the value of a block cannot be 0 or 1. $\frac{\log_2(\sqrt{n})-1}{8}$

```

func ss_encrypt_file(FILE *infile, FILE *outfile, mpz_t n);
     $k = \sqrt[n]{n}$ 
     $k = \log_2(k) - 1$ 
     $k = k/8$ 
    allocate an array "arr" of type uint8_t to act as a block
    set index zero of arr to 0xFF (inserts full byte)
    while (file pointer not at EOF):
        use fread to read at most k-1 element of size char (1 byte)
        import above block into a mpz variable using mpz_import()
        encrypt mpz var with block with ss_encrypt()
        use gmp_fprintf to write encrypted output to file:

```

use the conversion Zx to convert mpz_t into hex

```
//this function will essentially reverse the logic of ss_encrypt_file()
func ss_decrypt_file(FILE *infile, FILE *outfile, mpz_t d, mpz_t pq);
    k =  $\frac{\log_2(\sqrt{pq})-1}{8}$  //see ss_encrypt_file() for breakdown of this equation
    allocate an array "arr" of type uint8_t to serve as a block
    while (file * not at EOF):
        use gmp_fscanf to read encrypted input, using conversion %Qd:
        read this input into var mpz_t c
        decrypt c into m using ss_decrypt()
        export block in c with mpz_export
        use fwrite to write j-1 bytes of block starting at index 1:
            //we start at index 1 since 0 contains prepended byte
```