

Analyses of various algorithms

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Abstract

algorithm goes brrrr.

1 Introduction

1 Hi, we are students. How are you?(see Cormen et al., *Introduction to Algorithms*)

2 Theory

2 Provide a brief description of the algorithms you will be investigating, including pseudocode for the algorithms. Describe in particular the expected runtime of algorithms in terms of problem size. Use a separate subsection for each algorithm.

2.1 Algo 1: Bubble sort

3

Algorithm 1 Insertion sort algorithm from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

4 BUBBLESORT(A)
4 1 **for** $i = 1$ **to** $A.length - 1$
2 **for** $j = A.length$ **downto** $i + 1$
3 **if** $A[j] < A[j - 1]$
4 exchange $A[j]$ with $A[j - 1]$

Pseudocode for the first algorithm is shown in Listing 1. Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (1)$$

It is achieved for correctly sorted input data.

2.2 Algo 2: Insertion sort

Algorithm 2 Insertion sort algorithm from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

INSERTION-SORT(A)
1 **for** $j = 2$ **to** $A.length$
2 $key = A[j]$
3 $i = j - 1$
4 **while** $i > 0$ and $A[i] > key$
5 $A[i + 1] = A[i]$
6 $i = i - 1$
7 $A[i + 1] = key$

Pseudocode for the second algorithm is shown in Listing 2. Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (2)$$

It is achieved for correctly sorted input data.

2.3 Algo 3: Quicksort

Algorithm 3 Quicksort algorithm from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

```

QUICKSORT( $A, p, r$ )
1  if  $p < r$ 
2       $q = \text{PARTITION}(A, p, r)$ 
3      QUICKSORT( $A, p, q - 1$ )
4      QUICKSORT( $A, q + 1, r$ )

```

Algorithm 4 Partition from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

```

PARTITION( $A, p, r$ )
1   $x = A[r]$ 
2   $i = p - 1$ 
3  for  $j = p$  to  $r - 1$ 
4      if  $A[j] \leq x$ 
5           $i = i + 1$ 
6          exchange  $A[i]$  with  $A[j]$ 
7  exchange  $A[i + 1]$  with  $A[r]$ 
8  return  $i + 1$ 

```

Pseudocode for the second algorithm is shown in Listing 4. Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (3)$$

It is achieved for correctly sorted input data.

2.4 Algo 4: Quicksort Median of Three

Algorithm 5 Quicksort median of three algorithm from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

```

INSERTION-SORT( $A$ )
1  for  $j = 2$  to  $A.length$ 
2       $key = A[j]$ 
3       $i = j - 1$ 
4      while  $i > 0$  and  $A[i] > key$ 
5           $A[i + 1] = A[i]$ 
6           $i = i - 1$ 
7       $A[i + 1] = key$ 

```

Pseudocode for the second algorithm is shown in Listing 5. Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (4)$$

It is achieved for correctly sorted input data.

2.5 Algo 5: Quicksort Insertion sort hybrid

Algorithm 6 Insertion sort algorithm from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

```

INSERTION-SORT( $A$ )
1  for  $j = 2$  to  $A.length$ 
2       $key = A[j]$ 
3       $i = j - 1$ 
4      while  $i > 0$  and  $A[i] > key$ 
5           $A[i + 1] = A[i]$ 
6           $i = i - 1$ 
7       $A[i + 1] = key$ 

```

Pseudocode for the second algorithm is shown in Listing 6. Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (5)$$

It is achieved for correctly sorted input data.

2.6 Algo 6: Mergesort

why are you like this

Algorithm 7 Mergesort algorithm from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

```

MERGE-SORT( $A, p, r$ )
1  if  $p < r$ 
2       $q = \lfloor (p + r) / 2 \rfloor$ 
3      MERGE-SORT( $A, p, q$ )
4      MERGE-SORT( $A, q + 1, r$ )
5      MERGE( $A, p, q, r$ )

```

Algorithm 8 Merge from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

```

MERGE( $A, p, q, r$ )
1   $n_1 = q - p + 1$ 
2   $n_2 = r - q$ 
3  let  $L[1..n_1 + 1]$  and  $R[1..n_2 + 1]$  be new arrays
4  for  $i = 1$  to  $n_1$ 
5       $L[i] = A[p + i - 1]$ 
6  for  $j = 1$  to  $n_2$ 
7       $R[j] = A[q + j]$ 
8   $L[n_1 + 1] = \infty$ 
9   $R[n_2 + 1] = \infty$ 
10  $i = 1$ 
11  $j = 1$ 
12 for  $k = p$  to  $r$ 
13     if  $L[i] \leq R[j]$ 
14          $A[k] = L[i]$ 
15          $i = i + 1$ 
16     else  $A[k] = R[j]$ 
17          $j = j + 1$ 

```

Pseudocode for the second algorithm is shown in Listing 8. Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (6)$$

It is achieved for correctly sorted input data.

2.7 Algo 7: Mergesort Insertion sort hybrid

Algorithm 9 Insertion sort algorithm from Cormen et al., *Introduction to Algorithms*, Ch. 2.1.

```

INSERTION-SORT( $A$ )
1  for  $j = 2$  to  $A.length$ 
2       $key = A[j]$ 
3       $i = j - 1$ 
4      while  $i > 0$  and  $A[i] > key$ 
5           $A[i + 1] = A[i]$ 
6           $i = i - 1$ 
7       $A[i + 1] = key$ 

```

Pseudocode for the second algorithm is shown in Listing 9. Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (7)$$

It is achieved for correctly sorted input data.

2.8 Algo 8: Python sort

Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (8)$$

It is achieved for correctly sorted input data.

2.9 Algo 9: Numpy sort

Best case runtime for this algorithm is

$$T(n) = \Theta(n) . \quad (9)$$

It is achieved for correctly sorted input data.

3 Method used

In this section we will show our workflow and automation of tests.(brady_haran_death_2020)

4 Benchmarks

In this section we will analyse the performance of the different algorithms on exponentially growing test data.

5 Discussion

We will in this section discuss our findings, and thoughts¹

6 Acknowledgements

We acknowledge that we may, or may not be made out of bread.

References

Cormen, Thomas H. et al. *Introduction to Algorithms*. 3rd. Computer science. USA: McGraw-Hill, 2009. 1292 pp. ISBN: 978-0-262-03384-8. URL: <https://books.google.no/books?id=aefUBQAAQBAJ>.

¹Although these thoughts are subjective, we will present them as facts. For our thoughts feel real to us.