



Theory NP

Week 17

Required Activities

- Check Announcements regularly (every 2-3 days)
- Read FA book: Chapter 9
- **Complete and submit Assignment 5 (due Feb 15)**

Tractable

- A problem is tractable if there exists a polynomial-bound algorithm that solves it
- Worst-case growth rate can be bounded by a polynomial
- Function of its input size
 - $P(n) = a_n n^k + \dots + a_1 n + a_0$ where k is a constant
 - $P(n) \in \theta(n^k)$
 - E.g. $2n$, $3n^3$, $n \log n$

Intractability

- **Dictionary Definition** of intractable: “difficult to treat or work.”
- **Computer Science:** problem is intractable if a computer has difficulty solving it
- A problem is intractable if it is not tractable
- Any algorithm with a growth rate not bounded by a polynomial
 - c^n , $c^{.01n}$, $n^{\log n}$, $n!$
- Property of the problem not the algorithm

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Three General Categories of Problems

- Problems for which **polynomial-time** algorithms have been found
- Problems that have been proven to be **intractable**
- Problems that have **not been proven** to be intractable, but for which polynomial-time algorithms have never been found

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Algorithms

- **Polynomial-time Algorithms**
 - $\Theta(n \log n)$ for sorting
 - $\Theta(\log n)$ for searching
 - $\Theta(n^3)$ for chained-matrix multiplication
- Proven to be **Intractable**
 - Unrealistic definition of the Problem (Hamiltonian Circuits)
 - Un-Decidable problems: The Halting Problem (proven un-decidable by Alan Turing).
 - Decidable intractable problems: researchers have shown some problems from automata and mathematical logic intractable
- **Not proven** to be intractable (there could exist polynomial time algorithm)
 - Traveling salesperson
 - 0-1 Knapsack
 - Graph coloring
 - Sum of subsets

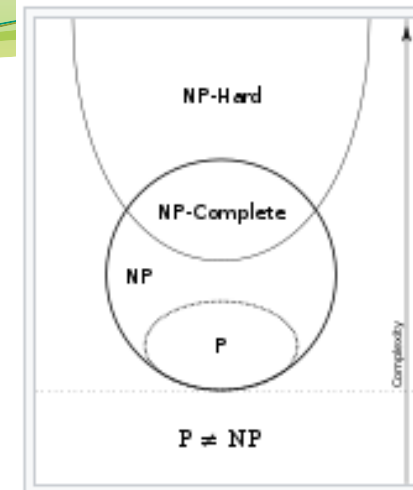
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Nondeterministic Polynomial time (NP)

- **Decision Problems** – output is “yes” or “no” (optimization problems are decision problems)
 - **Traveling Salesperson**: For a given positive number d , is there a tour having length $\leq d$?
 - **0-1 Knapsack**: For a given profit P , is it possible to load the knapsack such that total weight $\leq W$?
- **Deterministic** - given specific input, will produce same output
- **NP**
 - Decision problems solvable in polynomial time by a theoretical **non-deterministic** Turing machine (guess about solution in non-deterministic way then **verify** with deterministic algorithm to verify or reject guess)
 - Decision problems where “yes” have verifiable **proofs by deterministic** computations performed in **polynomial time**

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NP Problems



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- **P** – decision problems that can be **solved** by polynomial time algorithms
- **Polynomial-time nondeterministic algorithm** – **verification** stage is polynomial-time algorithm
- **NP** – nondeterministic polynomial time
 - Decision problems that can be “solved” by polynomial-time nondeterministic algorithm
 - Which means there is algorithm to **verify** in polynomial time
 - May not be an algorithm to actually solve problem in polynomial-time
- **NP-complete (NP-C)**
 - Every NP-complete is NP-hard but not vice versa
 - By definition: It is NP (there is verify polynomial algorithm) and every problem in NP can be reduced to it
 - In other words, solution can be **verified** quickly (polynomial time) but there is no known fast solution
 - In practice: New problem can be proven to be NP-complete if can reduce to known NP-complete problem
 - Usually solved using heuristic methods and approximation algorithms
- **NP-hard**
 - At least as hard as any NP problem to solve (can be harder)
 - Includes non-decision problems
 - Algorithm for “solving” it can be reduced into one for solving any other NP-complete in polynomial time

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Examples

- **NP-complete (also by definition NP and NP-hard)** when expressed as decision problem
 - Knapsack problem
 - Subset sum problem
 - Traveling salesman problem
 - Hamiltonian path problem
- **Halting problem** is NP-hard but not NP-complete (it is intractable)

Is P contained in NP?

- It has **not been proven** that there is a problem in NP that is not in P
- NP-P may be empty
- P=NP? One of the most important questions in CS
 - To show $P \neq NP$, find a problem in NP that is not in P
 - To show $P = NP$, find polynomial-time algorithm for each problem in NP

Questions ?

- Post in the discussions
- Send email to RMcFadden@HarrisburgU.edu
- Respond usually within 48hours