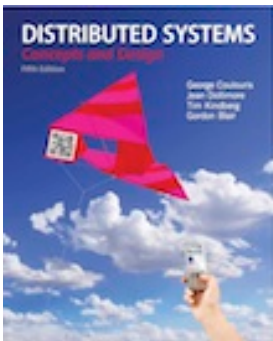


# Slides for Chapter 5: Operating System support

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*From* Coulouris, Dollimore, Kindberg and Blair

Distributed Systems:  
Concepts and Design

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Figure 5.1  
System layers

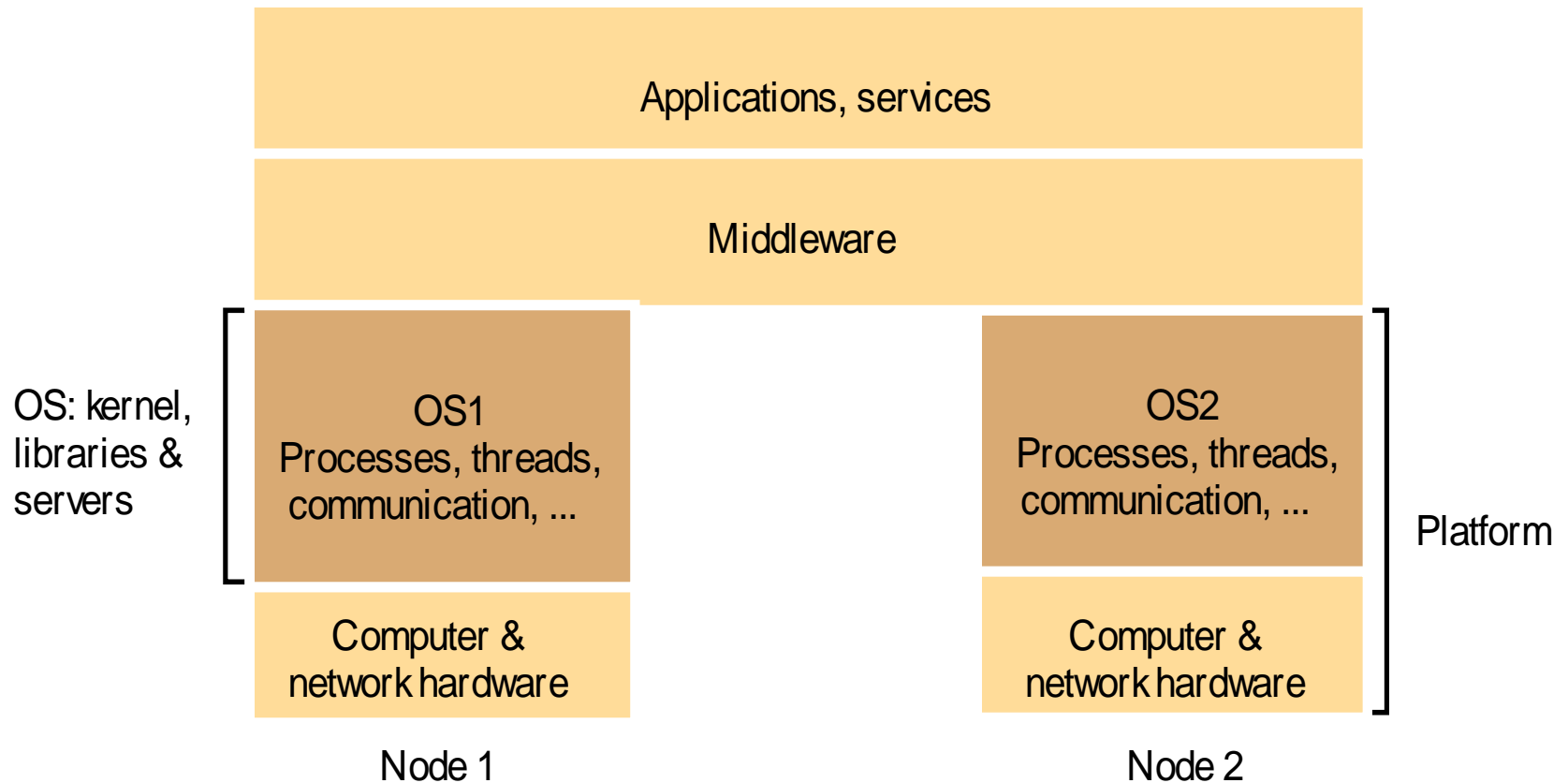


Figure 5.2  
Core OS functionality

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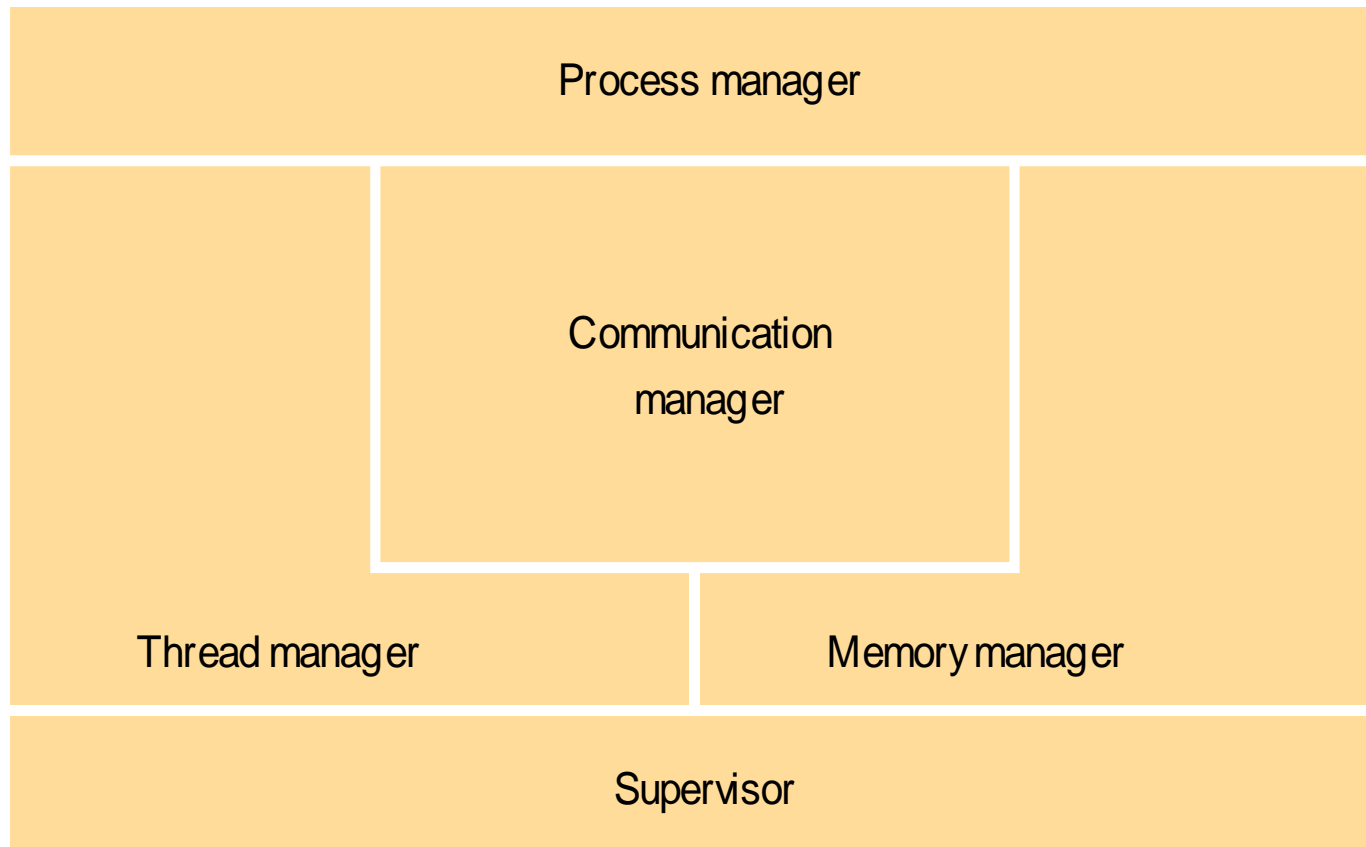


Figure 5.3  
Address space

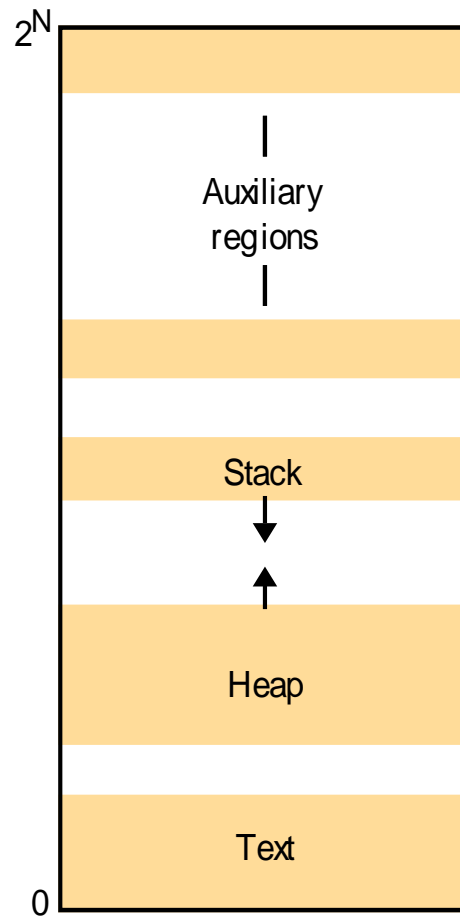


Figure 5.4  
Copy-on-write

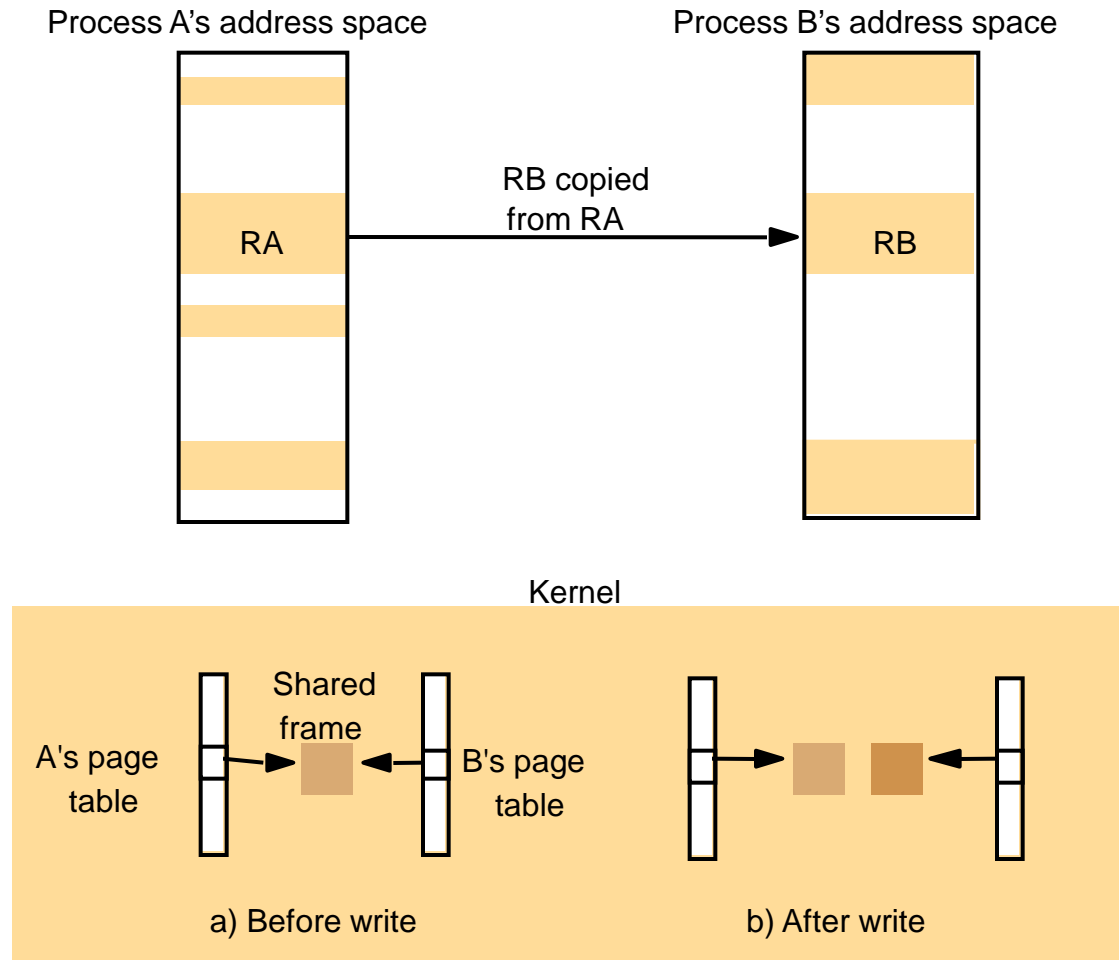


Figure 5.5  
Client and server with threads

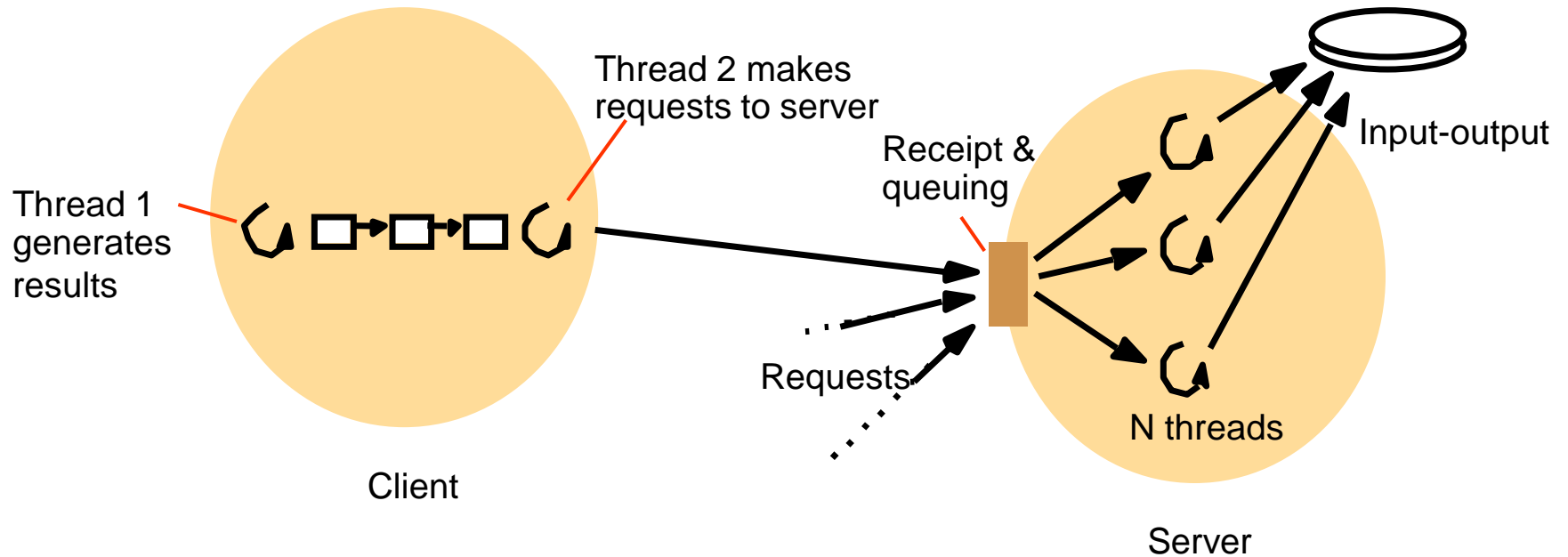
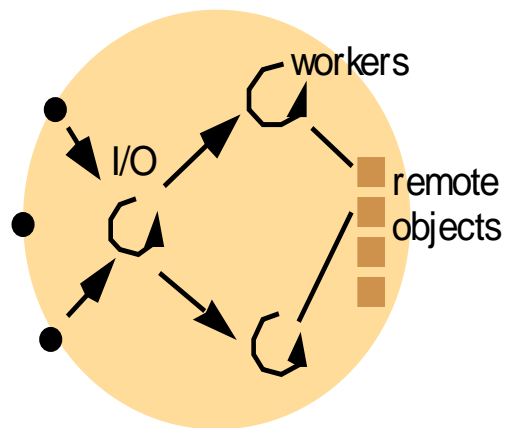
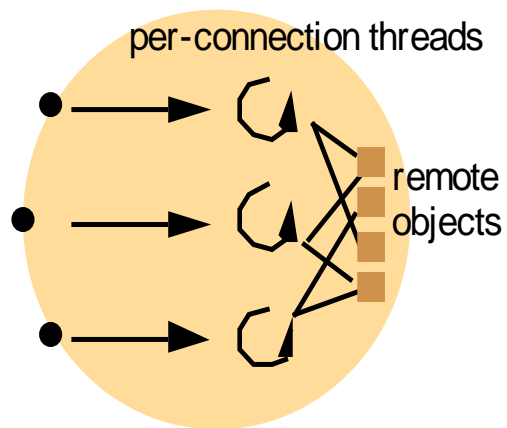


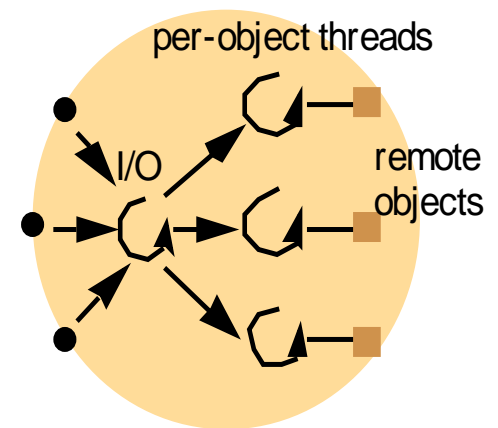
Figure 5.6  
Alternative server threading architectures (see also Figure 5.5)



a. Thread-per-request



b. Thread-per-connection



c. Thread-per-object

## Figure 5.7

### State associated with execution environments and threads

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<i>Execution environment</i>	<i>Thread</i>
Address space tables	Saved processor registers
Communication interfaces, open files	Priority and execution state (such as <i>BLOCKED</i> )
Semaphores, other synchronization objects	Software interrupt handling information
List of thread identifiers	Execution environment identifier
Pages of address space resident in memory; hardware cache entries	

---



## Figure 5.8

### Java thread constructor and management methods

*Thread(ThreadGroup group, Runnable target, String name)*

Creates a new thread in the *SUSPENDED* state, which will belong to *group* and be identified as *name*; the thread will execute the *run()* method of *target*.

*setPriority(int newPriority), getPriority()*

Set and return the thread's priority.

*run()*

A thread executes the *run()* method of its target object, if it has one, and otherwise its own *run()* method (*Thread* implements *Runnable*).

*start()*

Change the state of the thread from *SUSPENDED* to *RUNNABLE*.

*sleep(int millisecs)*

Cause the thread to enter the *SUSPENDED* state for the specified time.

*yield()*

Causes the thread to enter the *READY* state and invoke the scheduler.

*destroy()*

Destroy the thread.

## Figure 5.9

### Java thread synchronization calls

---

*thread.join(int millisecs)*

Blocks the calling thread for up to the specified time until *thread* has terminated.

*thread.interrupt()*

Interrupts *thread*: causes it to return from a blocking method call such as *sleep()*.

*object.wait(long millisecs, int nanosecs)*

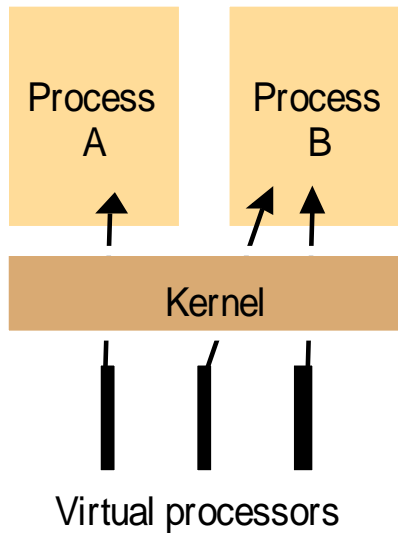
Blocks the calling thread until a call made to *notify()* or *notifyAll()* on *object* wakes the thread, or the thread is interrupted, or the specified time has elapsed.

*object.notify()*, *object.notifyAll()*

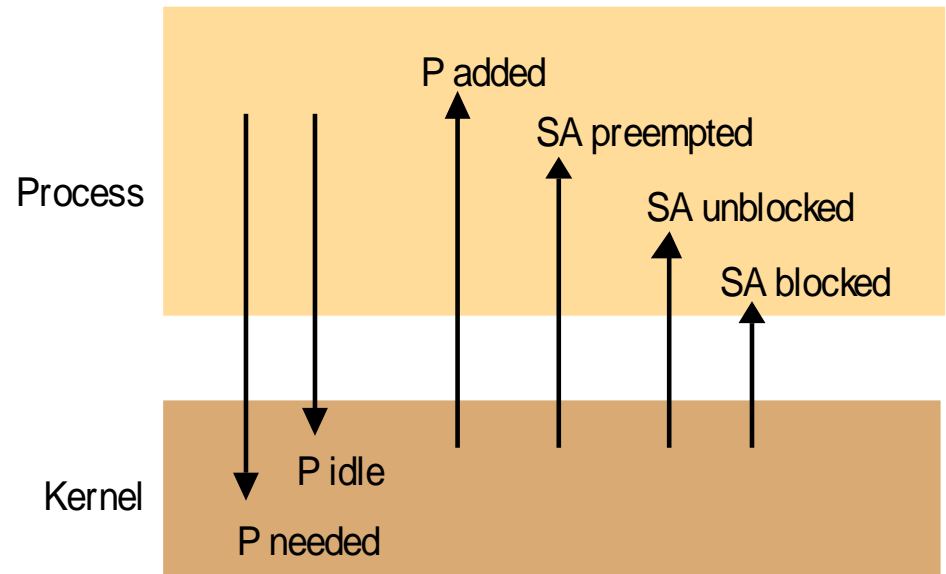
Wakes, respectively, one or all of any threads that have called *wait()* on *object*.

## Figure 5.10

### Scheduler activations



A. Assignment of virtual processors to processes



B. Events between user-level scheduler & kernel  
Key. P = processor; SA = scheduler activation

## Figure 5.11

### Invocations between address spaces

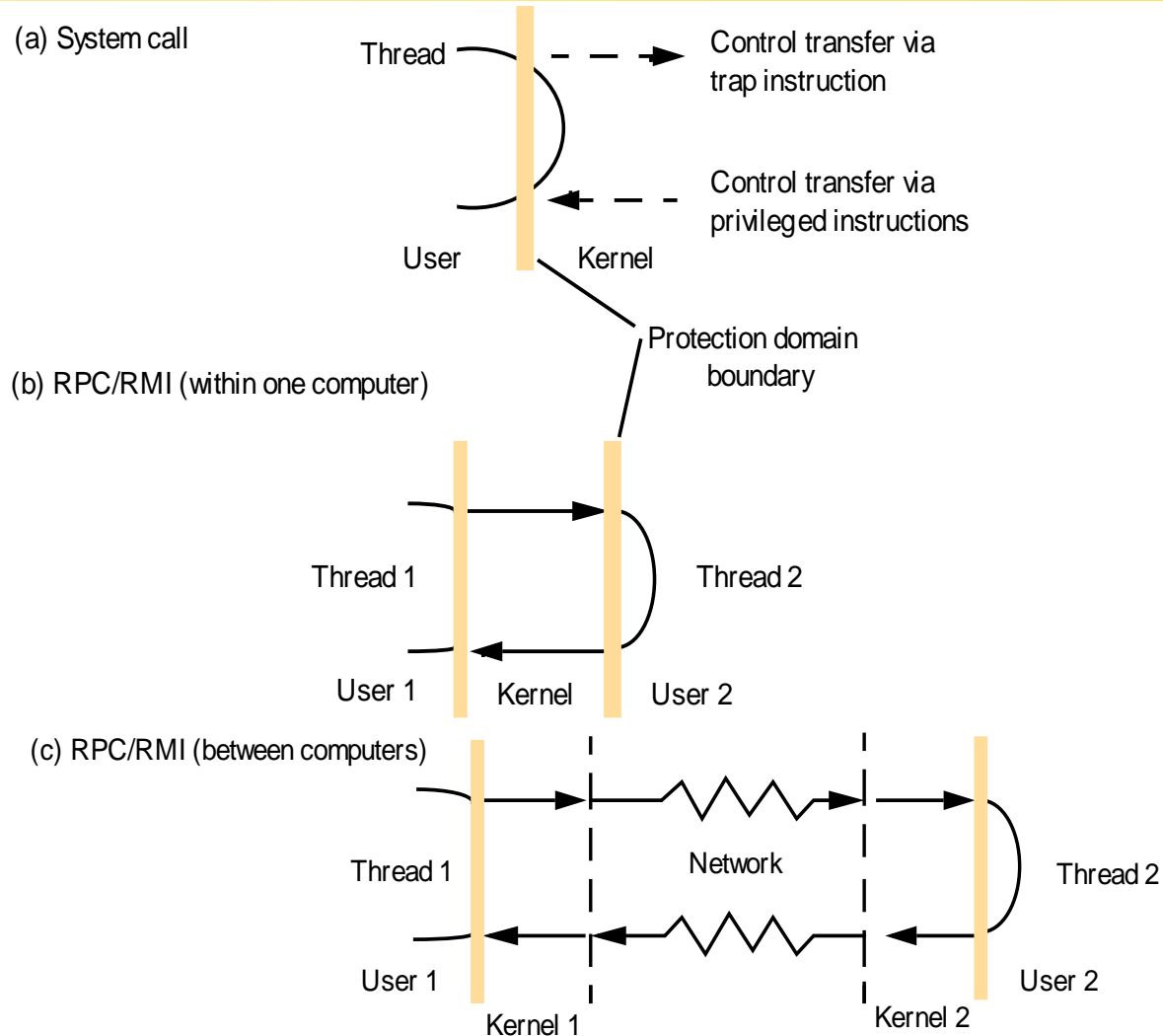


Figure 5.12  
RPC delay against parameter size

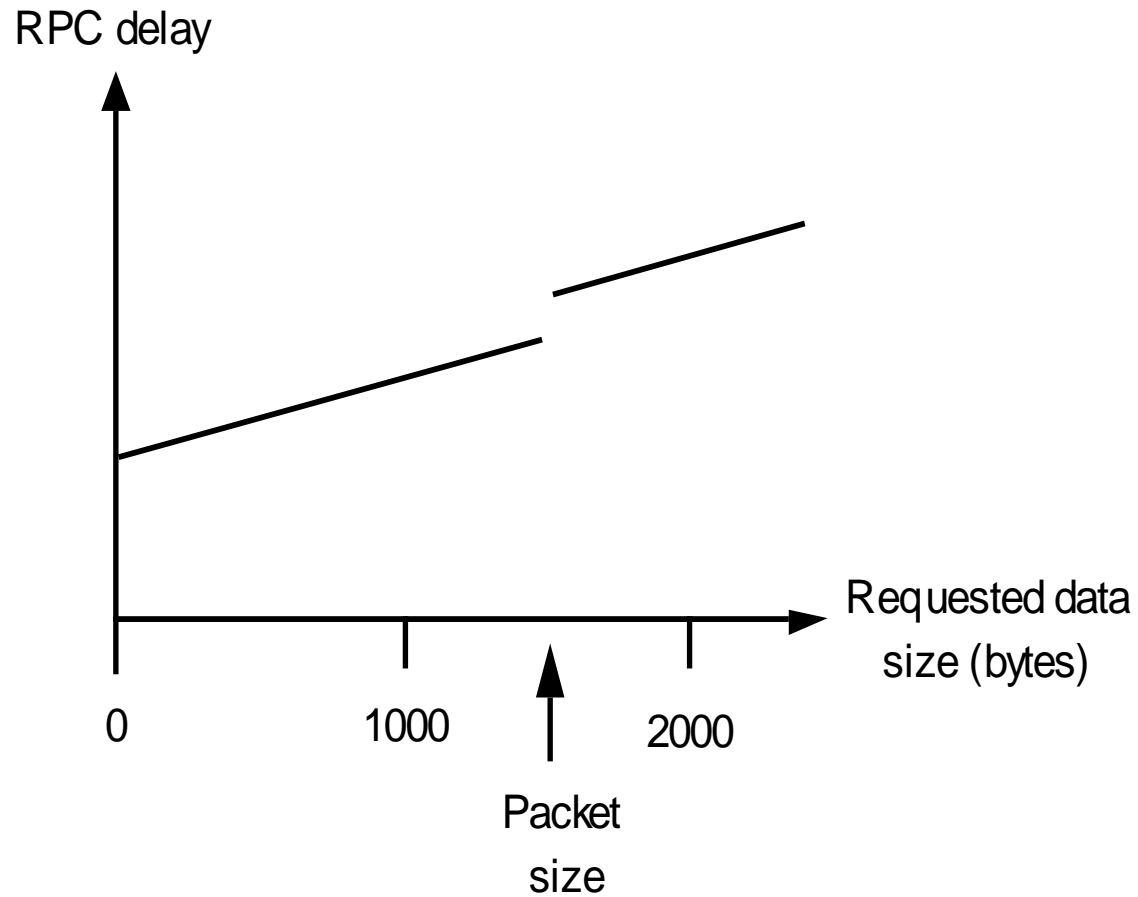


Figure 5.13  
A lightweight remote procedure call

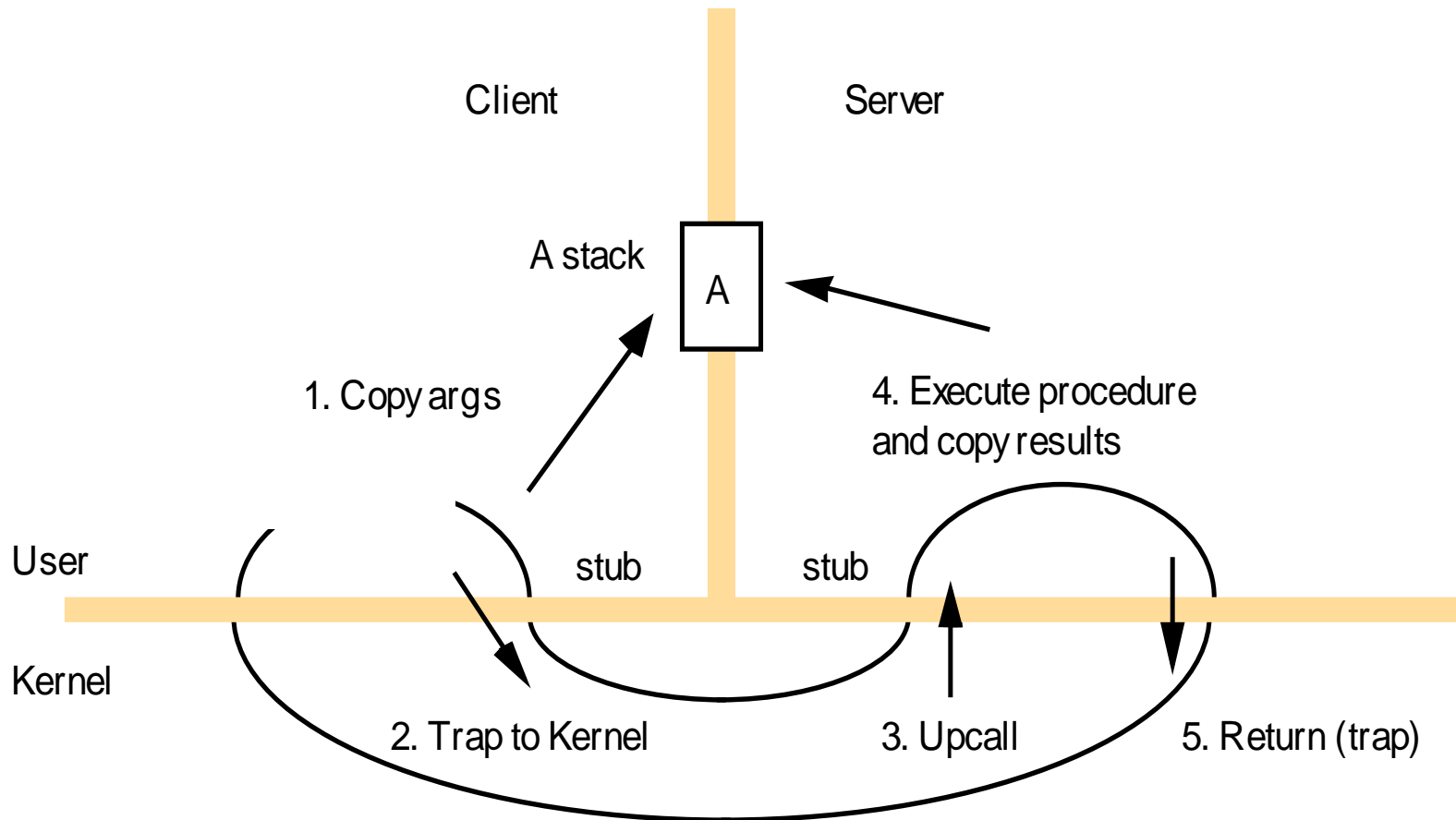


Figure 5.14  
Times for serialized and concurrent invocations

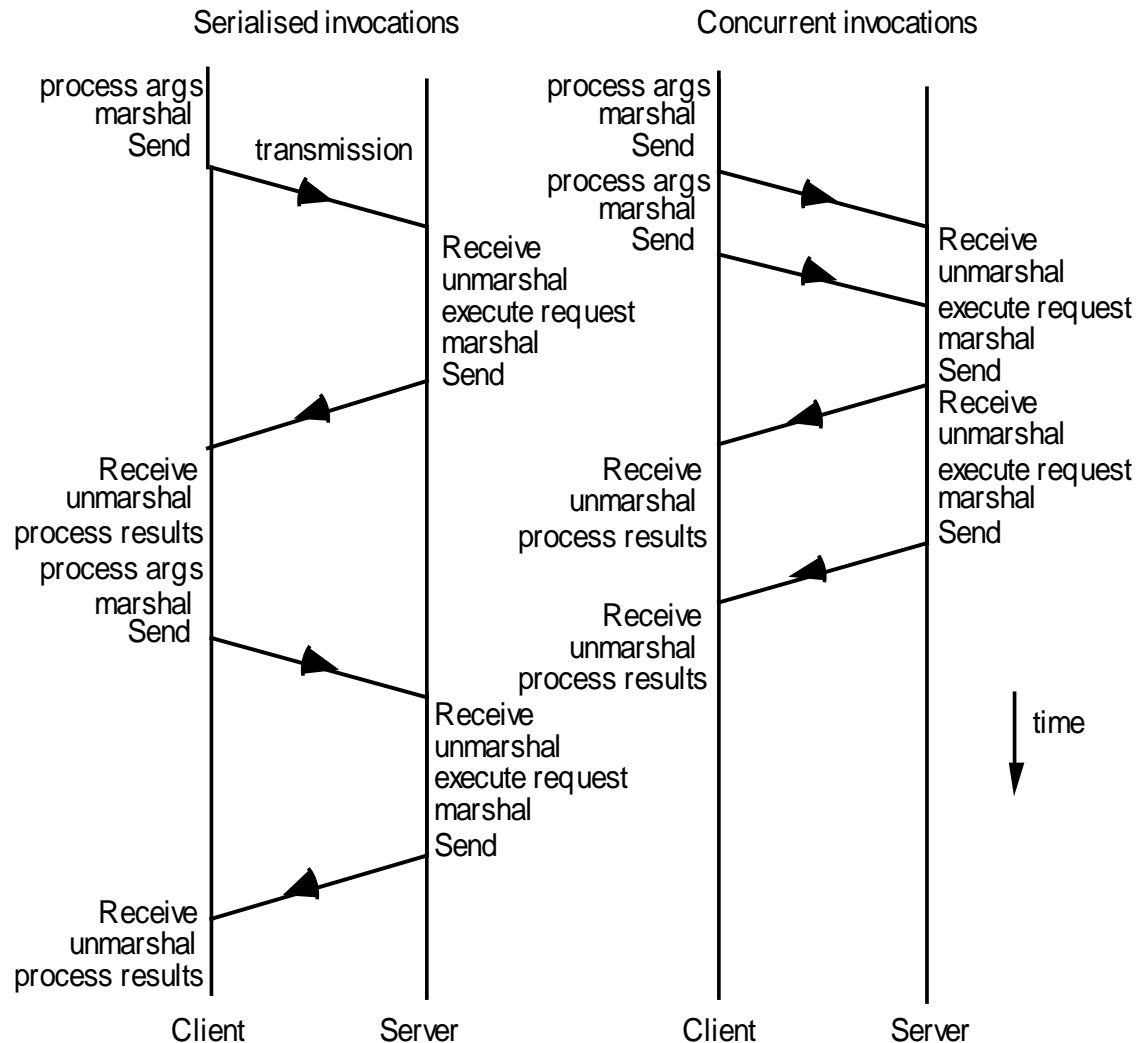


Figure 5.15  
Monolithic kernel and microkernel

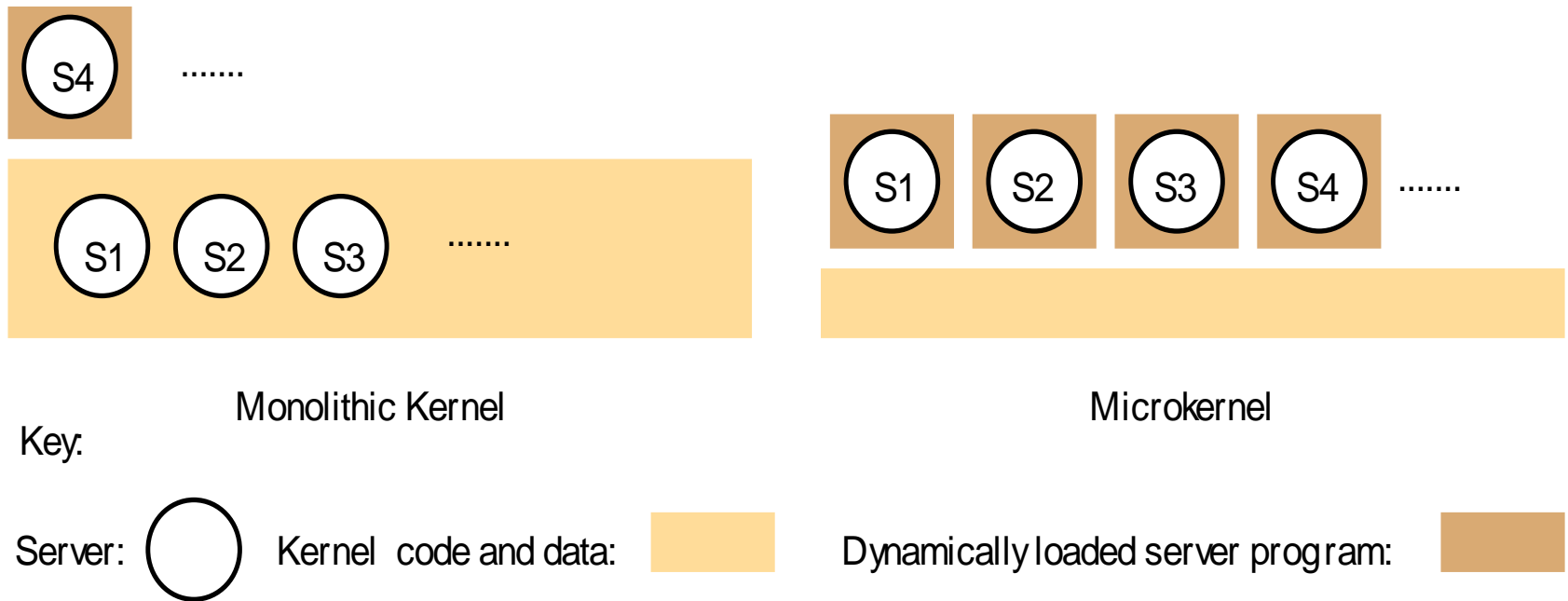
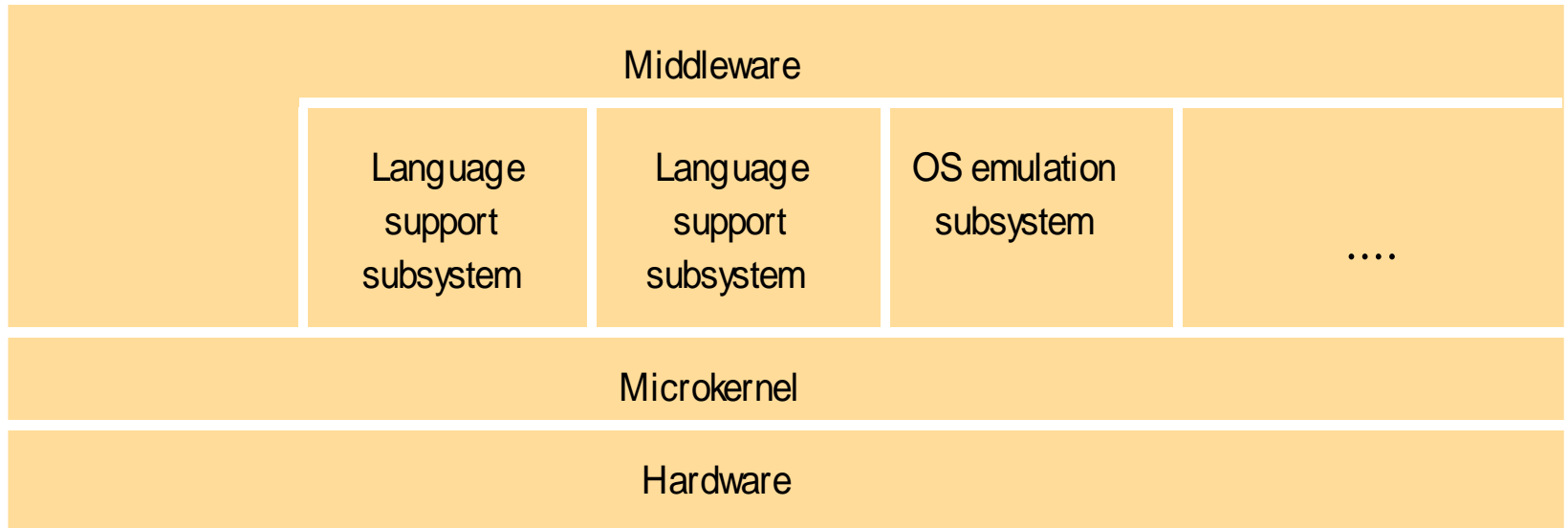




Figure 5.16  
The role of the microkernel



The microkernel supports middleware via subsystems

Figure 5.17  
The architecture of Xen

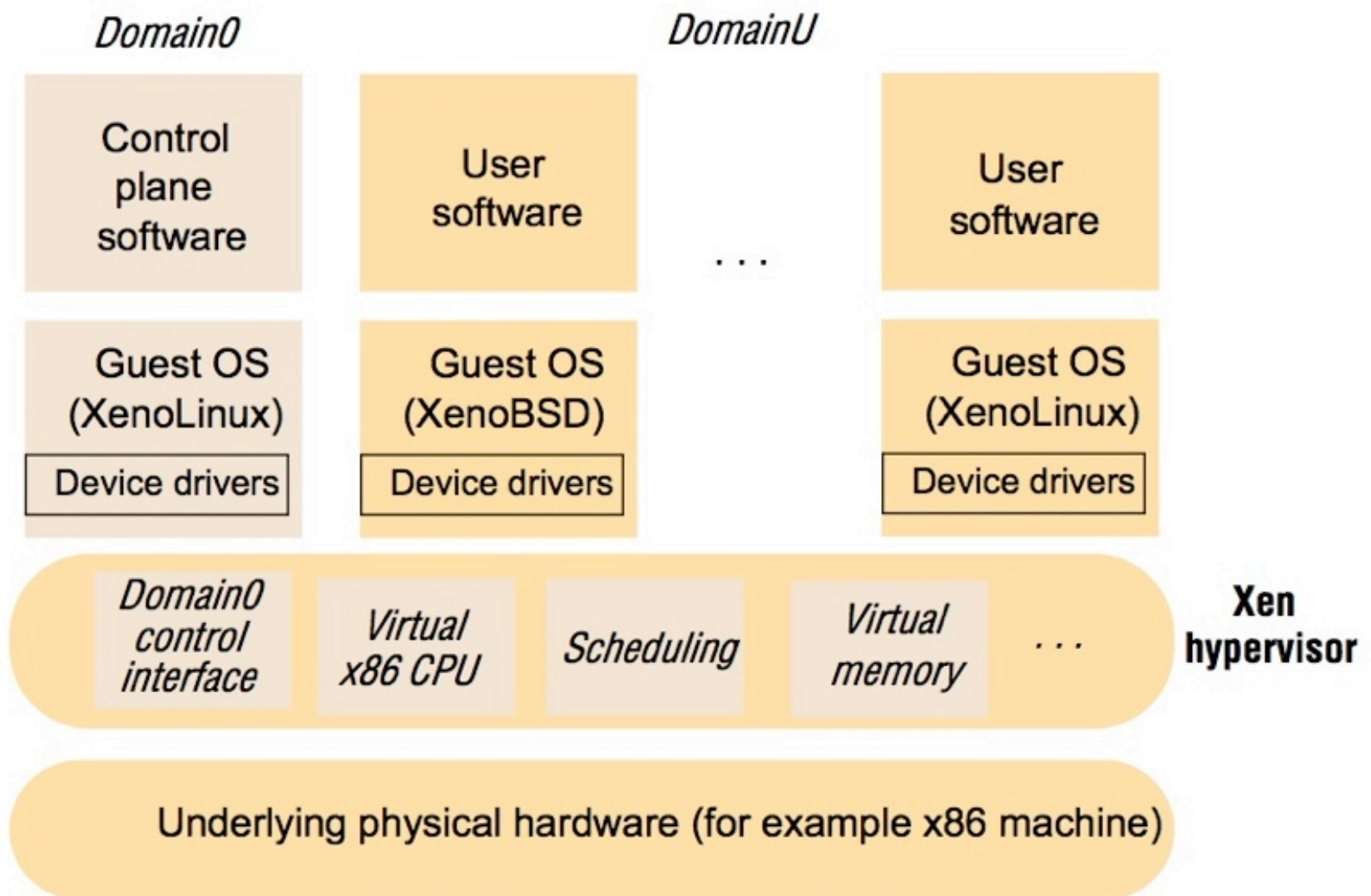


Figure 5.18  
Use of rings of privilege

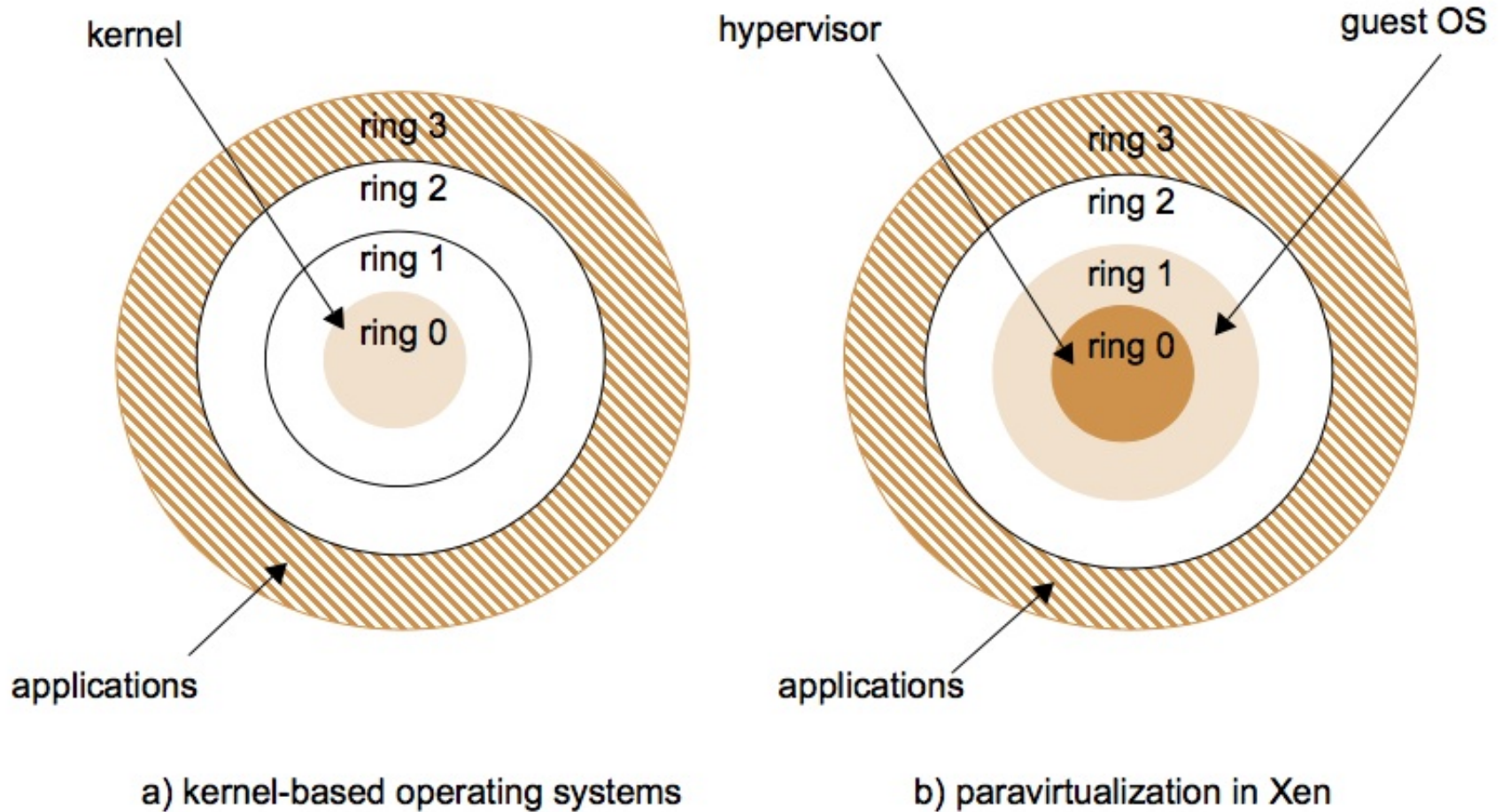


Figure 5.19  
Virtualization of memory management

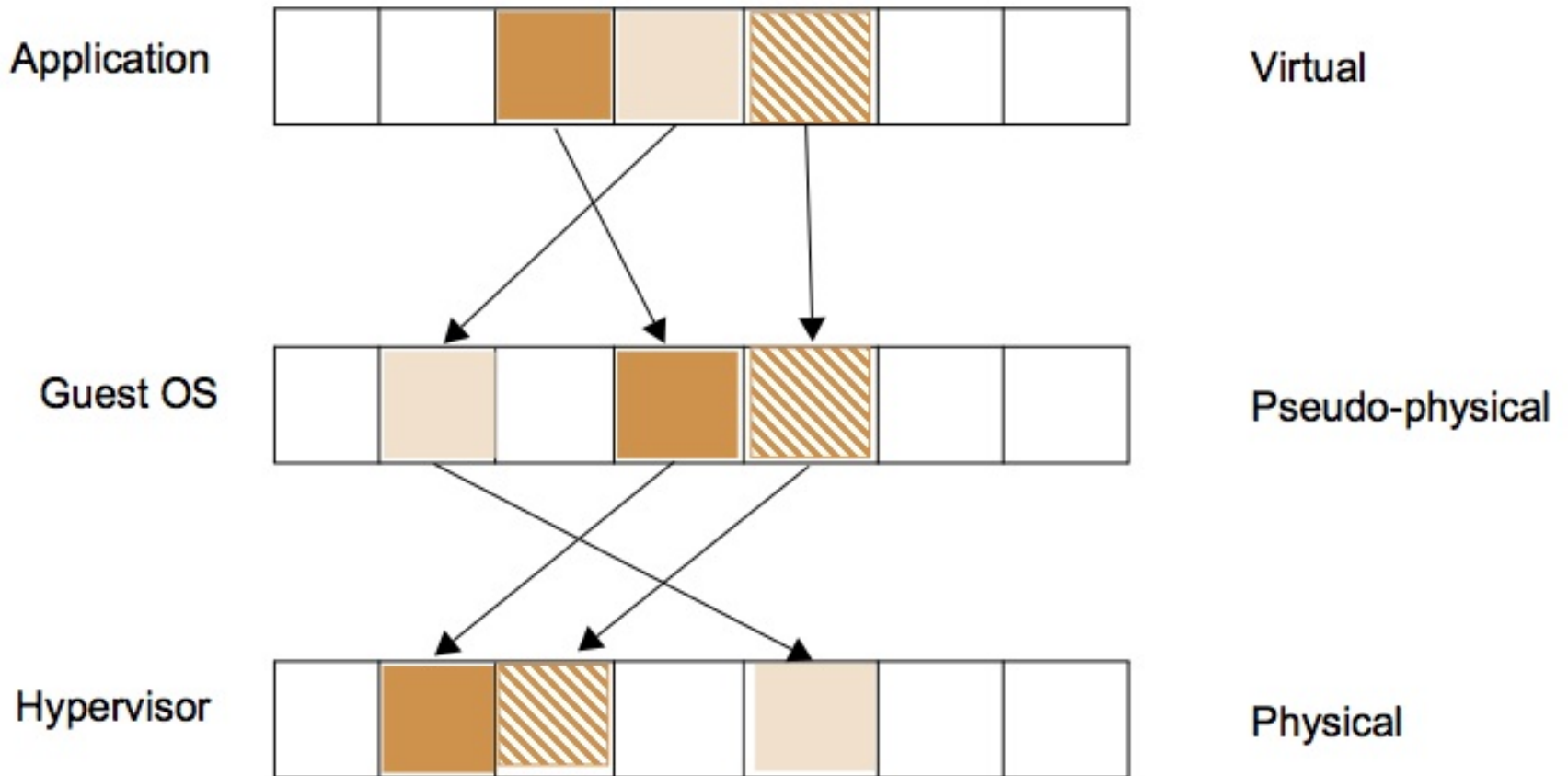


Figure 5.20  
Split device drivers

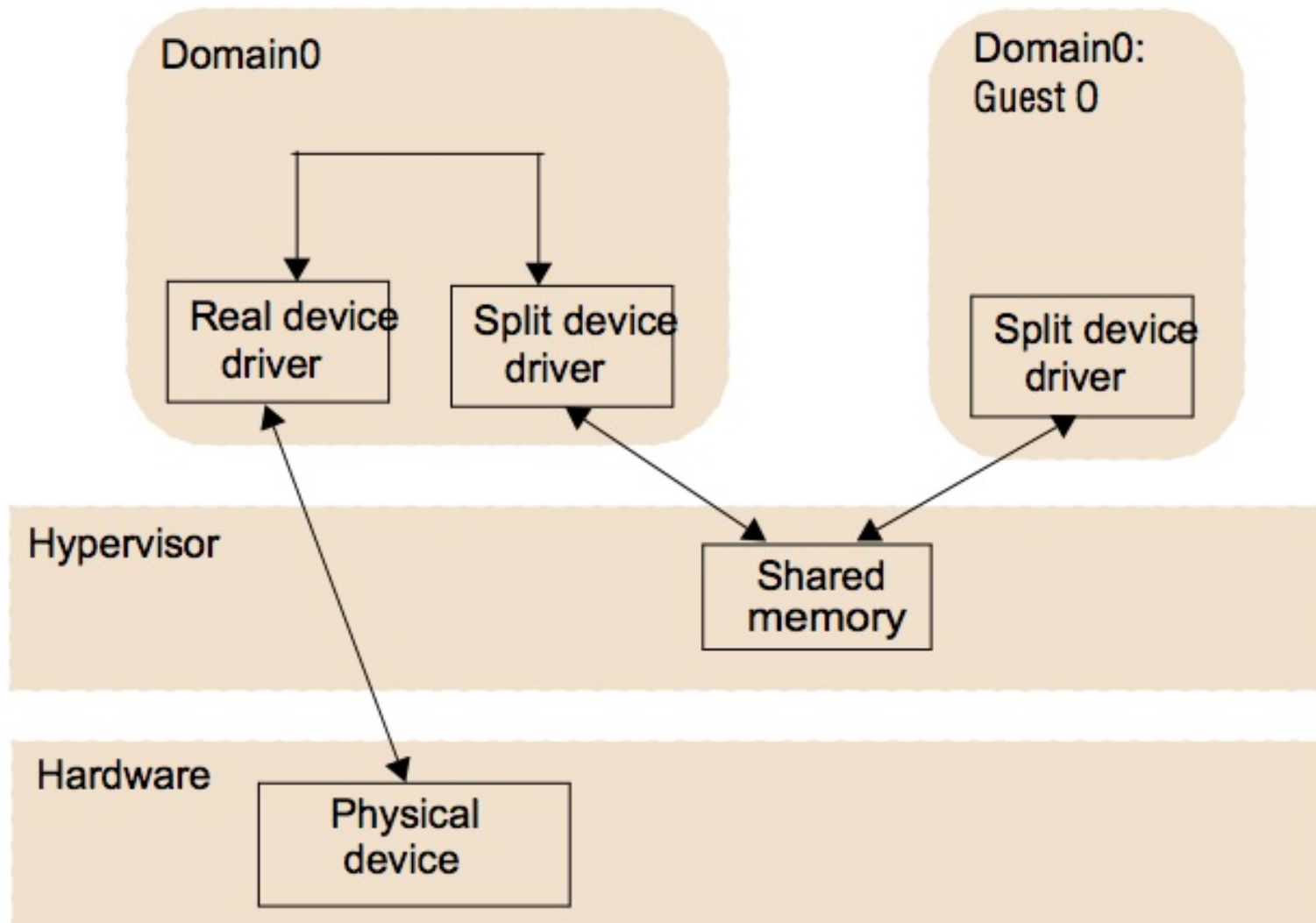


Figure 5.21  
I/O rings

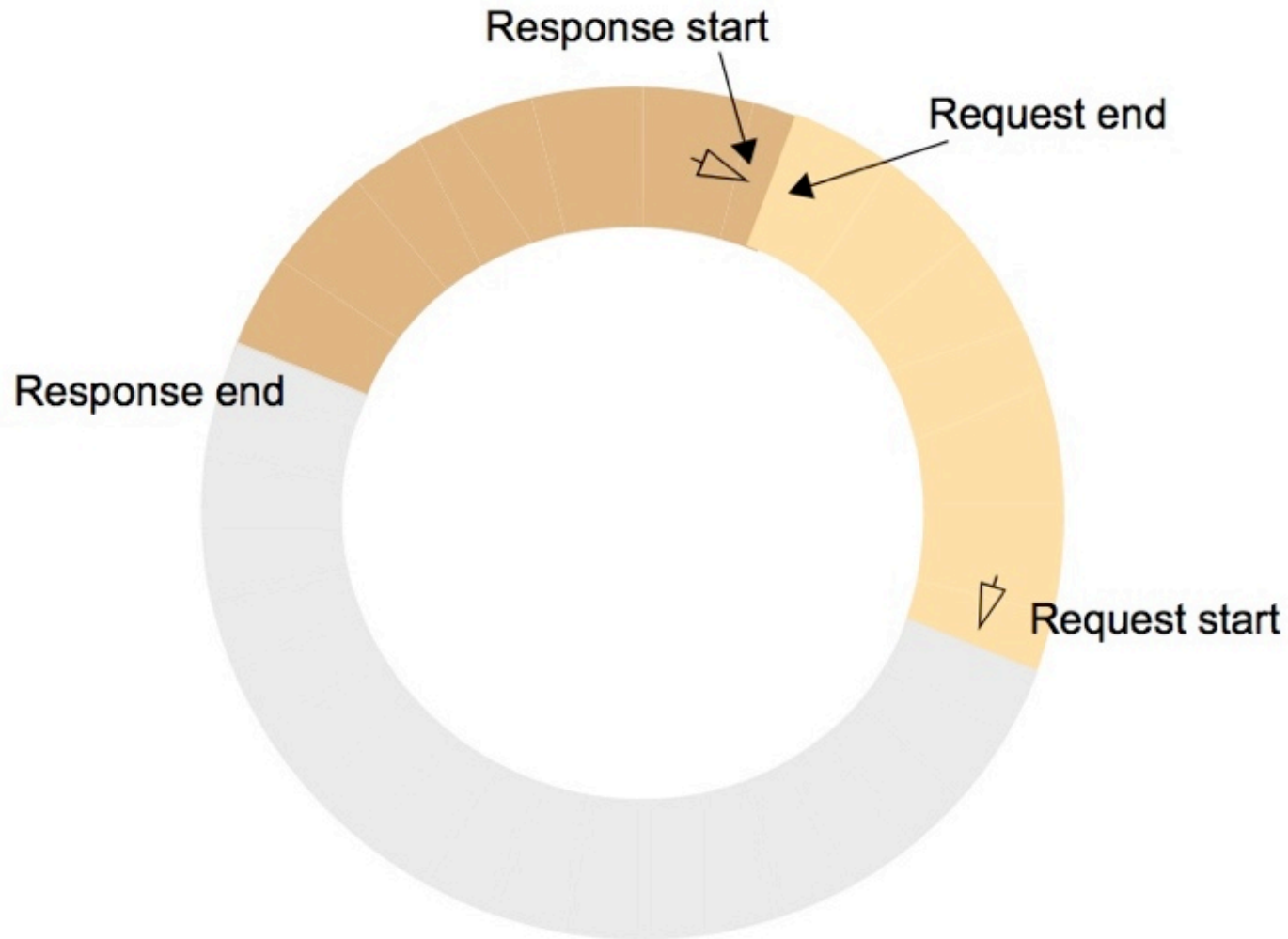


Figure 5.22  
The XenoServer Open Platform Architecture

