

# LO-RISC

**Learning Optimized Reduced Instruction Set Computer**

A minimal Instruction Set Architecture designed for speed and simplicity

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COA lab Final Project

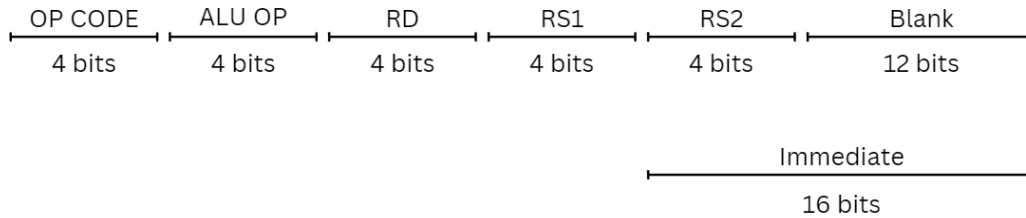
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# Instruction Set

Each Instruction in LO-RISC is 32 bits long, divided into 8, 4-bit chunks representable by hexadecimal.

**OP CODE** is connected to the control unit and essentially selects the instruction while **ALU OP** is directly connected to the ALU and selects its function. Three register operand addresses (4-bit each) follow, followed by 12 trailing blank (0) bits. For Immediate instructions **RS2** and Blank Bits are replaced by the **immediate** field.



OP code		ALU op		Register	
Hex	Instruction	Hex	Operation	Hex	function
0*	ALU	0	Add	0	Read Only (0)
1*	ALU Immediate	1	Subtract	1-9	General Purpose
20	Load	2	SLT	A	Function Arguments
30	Store	3	SGT	B	Return Address
40	Branch	4	AND	C	System Reserved
50	Branch LT	5	OR	D	Debug Register
60	Branch GT	6	XOR	E	Stack Pointer
70	Branch on 0	7	NOT	F	Function Output
80	Move	8	NOR	*Apart from given function, all registers can be used as GPR	
90	conditional move	9	LUI		
A0	Jump Register	A	SLA		
FF	Halt	B	SRL	*7 segment display constantly displays debug register value (Nexys A7 only)	
EF	No Operation	C	SRA		
* Denotes ALU OP		D	Increment	*Stack pointer initializes to memory address 1023	
		E	Decrement		
		F	HAM		

## Examples:

add \$1 \$2 \$3 ► 0x00123000 ►  $R1 = R2 + R3$

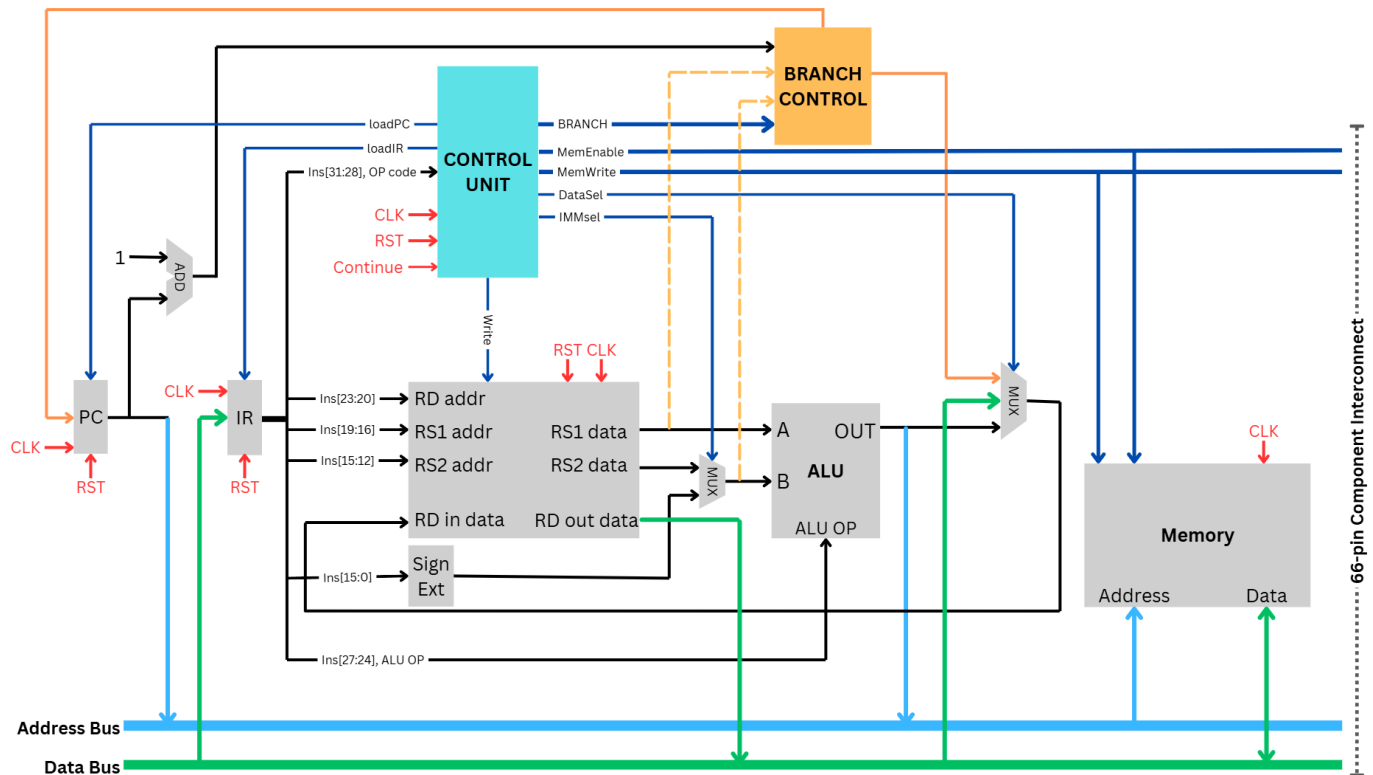
slai \$5 \$7 1 ► 0x1A570001 ►  $R5 = R7 \ll 1$

ld \$3 8(\$6) ► 0x20360008 ►  $R3 = \text{Mem}[R6+8]$

br #10 ► 0x4000000A ►  $PC = PC + 10$

bmi \$5 \$2 ► 0x50050020 ►  $PC \leq PC + 30$  if ( $R5 < 0$ )

# Data Path



The CPU uses Von Neuman Architecture for memory access. Instructions and Data share 4 KB of system memory.

32-bit **Address** and **Data** buses along with control signals, **Memory Enable** and **Memory Write Enable** make up the 66-pin (32+32+1+1) **Component Interconnect (CI)**. Components such as memory or I/O devices can be connected to the CPU via this interconnect. Tri-State buffers are used to control access to the address and data buses. A **UART I/O module** is used for Input/Output through serial UART and is connected to the CPU via the component interconnect.

All Verilog modules can be found in '**Verilog Assets/Sources**' directory.

# Input/Output

An I/O module is used for Serial communications through UART.

- A Serial terminal application such as **Tera Term** or **Minicom** can be used for interacting with the system. Recommended settings for Tera Term are present in '**Programs/system**'.
- It operates at a baud rate of 460800 leading to a max effective bandwidth of 368.64 Kbits/sec Kbits/sec.
- Memory addresses 4096 and 4097 are reserved for this module. It has two registers, the **command register** and the **data register** at reserved memory addresses **4096** and **4097** respectively.
- The module starts at idle state, the registers can only be written to in this state.
- The module is controlled through the command register. When command register is set to:
  - 1, it transmits the first byte of data register as an ASCII code
  - 2, it listens for an ASCII transmission and puts the transmission into the data register
  - 3, it transmits the entire data register as a decimal integer
- After the transmission/reception is complete, the command register resets to 0 and the module is idle again.
- The command register can be polled to check the state of transmission/reception.

# Assembly

A LO-RISC assembly file is composed of two parts, data section preceded by `.data` and instruction section preceded by `.text` along with macros of the form: `num = 4242`

## Data:

Data entries are composed of the label followed by data type and the corresponding data. Data is placed in memory after the instructions, sequentially in the order of data entries.

```
myvar: .int 42
myarr: .arr {3,4,5,7}
mychar: .char 'k'
mystr: .str "Hello"
```

## Instructions (*case insensitive*):

**Labels** (Eg: `Label_1:`) denote specific points in a program used for calculating the effective address for branching.

**The following instructions are available:**

**a)** Arithmetic and logic instructions: ADD, SUB, AND, OR, XOR, NOR, NOT, SL, SRL, SRA, INC, DEC, SLT, SGT, LUI, HAM. There are corresponding immediate addressing versions with a suffixing "I" (like ADDI, SUBI, etc.). Assume that all shift instructions can have either 0 (no shift) or 1 (1-bit shift) as operand. Some example uses are as follows:

```
add $1 $2 $3 #R1 = R2 + R3
slai $5 $7 1 #R5 = R7 << 1
```

**b)** Load and store instructions: LD, ST (all load and stores are 32-bits) and use register indexed addressing (any of the registers R1..R15 can be used). Some example uses are as follows:

```
ld $1 myvar #r3 = Mem[location of myvar]
ld $3 myarr($2) #R3 = Mem[Location of myarr[$2]]
```

**c)** Branch instructions: BR, BMI, BPL, BZ. Some example uses are as follows:

```
br loop #branch to loop
bz $5 lab #Branch to lab if R5 = 0
bmi $5 lab #branch to lab if R5 < 0
bpl $5 lab #branch to lab if R5 > 0
jr $ra #branch to address at RA
```

**d)** Register to register transfer: MOVE, CMOV. Some example uses are as follows:

```
move $4 $6 #R4 = R6
cmov $1 $2 $3 #R1 = (R2 < R3) ? R2 : R3
```

## Pseudo instructions

**a)** LA loads memory address into register

```
la $1 program
#equivalent to:
lui $1 (upper 16 bits of program)
ori $1 $1 (lower 16 bits of program)
```

**b)** LI loads value into register

```
li $1 12 #equivalent to: addi $1 $0 12
```

**b)** JAL to be used for function calls

```
jal func
# equivalent to:
addi $ra $0 (current address)
br func
```

# Dasmon

Inspired in part by Steve Wozniak's Wozmon, Dasmon is the system software for this implementation of the LO-RISC CPU. Source assembly file for Dasmon is present in '**Programs/system**'.

Dasmon allows the user to load and run programs as well as inspect memory locations:

- Enter instructions or data in hexadecimal separated by newline, each will be loaded in successive memory locations starting at address **271**, current memory location will be shown by the prompt, example:  
**271:\$>** indicates the user is writing the instruction/data to address 271.
- Press **shift+C** to clear the screen, **shift+R** to run the program.
- Enter ':' followed by an address in decimal to see its contents (*\*output may overflow*)

Dasmon provides subroutines for certain common operations including I/O. The assembler will assume the program is to be used with Dasmon and thus, will provide the following labels for calling these subroutines

- **jal booth\_mul** will store the product of **0(\$a)** and **1(\$a)** at **\$fo**.
- **jal printc** will print the first byte of **\$a** register as an ASCII character.
- **jal prints** will expect pointer to a string at **\$a** register and print the string.
- **jal printi** will print the value of **\$a** register as a positive decimal integer.
- **jal getchar** will listen for a transmission and put the received ASCII character at **\$fo** register.
- **jal getint** will listen for a transmission and put the received positive integer at **\$fo** register.

These labels are reserved and thus cannot be defined anywhere within the program.

the Assembler is present as portable executable (**Programs/asm.exe**) as well as python source code (**Programs/source/asm.py**). Usage is as follows:

**./asm program.s** or **python source/asm.py program.s**

Output will be **program.out** containing machine code instructions in hexadecimal separated by newline. Contents of this file can then be copy-pasted into a serial terminal (eg: tera term) to load the program. Dasmon will take the program as input and load it into memory. Press **shift+R** to subsequently, run the program.

Example programs are present in '**Programs**' directory.

To assemble dasmon itself, use **python source/asm-sys.py system/dasmon.s**, this will output dasmon.coe.

This coefficient file can then be loaded into the system memory with Block Memory Generator.