

# Introduction to Number Theory and Algorithms

Christian J. Rudder

August 2024

## Contents

<b>Contents</b>	<b>1</b>
<b>1 Asymptotic Notation</b>	<b>4</b>
1.1 Asymptotic Notation . . . . .	4

*This page is left intentionally blank.*

## Prerequisites

## Asymptotic Notation

### 1.1 Asymptotic Notation

Asymptotic analysis is a method for describing the limiting behavior of functions as inputs grow infinitely.

#### Definition 1.1: Asymptotic

Let  $f(n)$  and  $g(n)$  be two functions. As  $n$  grows, if  $f(n)$  grows closer to  $g(n)$  never reaching, we say that “ $f(n)$  is **asymptotic** to  $g(n)$ .”

We call the point where  $f(n)$  starts behaving similarly to  $g(n)$  the **threshold**  $n_0$ . After this point  $n_0$ ,  $f(n)$  follows the same general path as  $g(n)$ .

#### Definition 1.2: Big-O: (Upper Bound)

Let  $f$  and  $g$  be functions.  $f(n)$  our function of interest, and  $g(n)$  our function of comparison.

Then we say  $f(n) = O(g(n))$ , “ $f(n)$  is **big-O** of  $g(n)$ ,” if  $f(n)$  grows no faster than  $g(n)$ , up to a constant factor. Let  $n_0$  be our asymptotic threshold. Then, for all  $n \geq n_0$ ,

$$0 \leq f(n) \leq c \cdot g(n)$$

Represented as the ratio  $\frac{f(n)}{g(n)} \leq c$  for all  $n \geq n_0$ . Analytically we write,

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} < \infty$$

Meaning, as we chase infinity, our numerator grows slower than the denominator, bounded, never reaching infinity.

**Examples:**

(i.)  $3n^2 + 2n + 1 = O(n^2)$

(ii.)  $n^{100} = O(2^n)$

(iii.)  $\log n = O(\sqrt{n})$

**Proof 1.1:**  $\log n = O(\sqrt{n})$ 

We setup our ratio:

$$\lim_{n \rightarrow \infty} \frac{\log n}{\sqrt{n}}$$

Since  $\log n$  and  $\sqrt{n}$  grow infinitely without bound, they are of indeterminate form  $\frac{\infty}{\infty}$ . We apply L'Hopital's Rule, which states that taking derivatives of the numerator and denominator will yield an evaluateable limit:

$$\lim_{n \rightarrow \infty} \frac{\log n}{\sqrt{n}} = \lim_{n \rightarrow \infty} \frac{\frac{d}{dn} \log n}{\frac{d}{dn} \sqrt{n}}$$

Yielding derivatives,  $\log n = \frac{1}{n}$  and  $\sqrt{n} = \frac{1}{2\sqrt{n}}$ . We substitute these back into our limit:

$$\lim_{n \rightarrow \infty} \frac{\frac{1}{n}}{\frac{1}{2\sqrt{n}}} = \lim_{n \rightarrow \infty} \frac{2\sqrt{n}}{n} = \lim_{n \rightarrow \infty} \frac{2}{\sqrt{n}} = 0$$

Our limit approaches 0, as we have a constant factor in the numerator, and a growing denominator. Thus,  $\log n = O(\sqrt{n})$ , as  $0 < \infty$ . ■